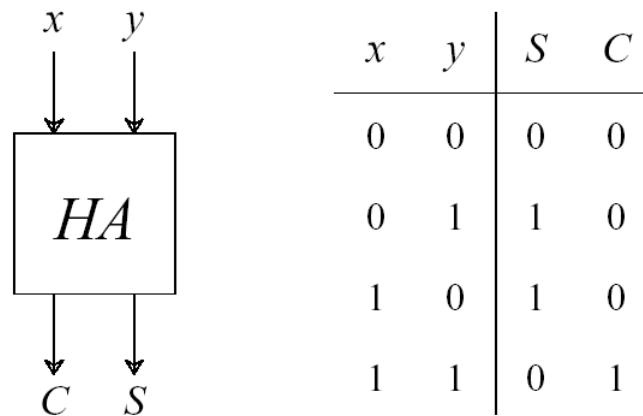


LAB 1

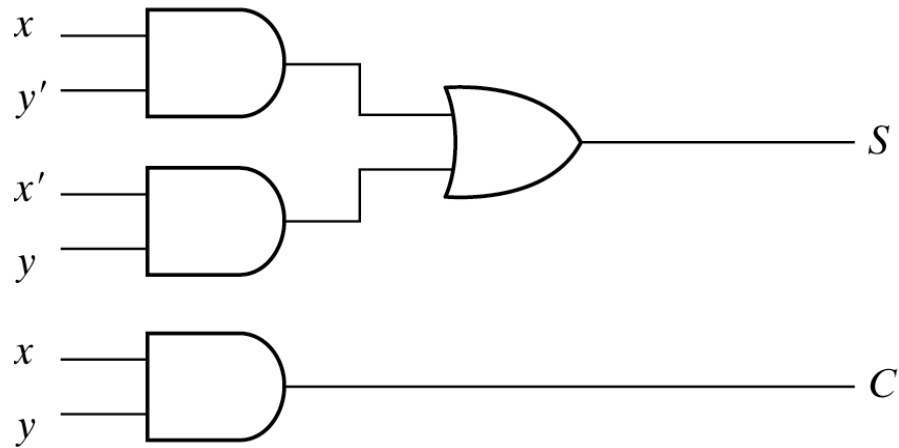
# DESIGN OF ARITHMETIC CIRCUIT

# Adder

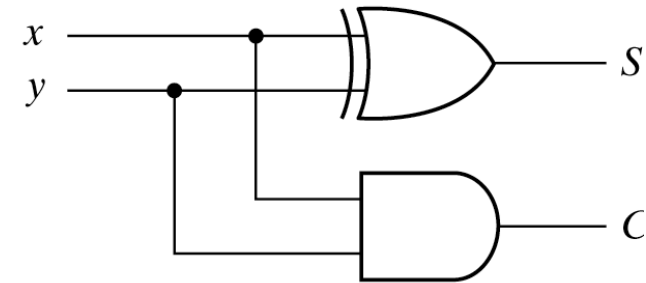
- ❑ The most basic arithmetic operation is the addition of 2 bits. A combinational circuit that performs this operation is called a *half-adder*.
- ❑ A combinational circuit that performs the addition of 3 bits is called a *full-adder*, which can be implemented by 2 half-adders.



# Half Adder

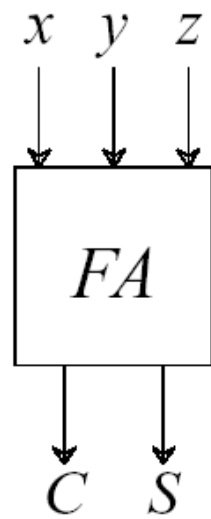


(a)  $S = xy' + x'y$   
 $C = xy$



(b)  $S = x \oplus y$   
 $C = xy$

# Full Adder



$x$	$y$	$z$	$S$	$C$
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

# Full Adder

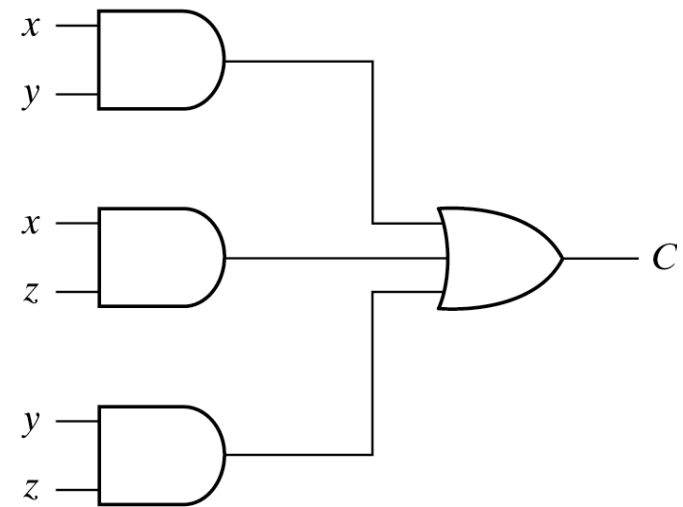
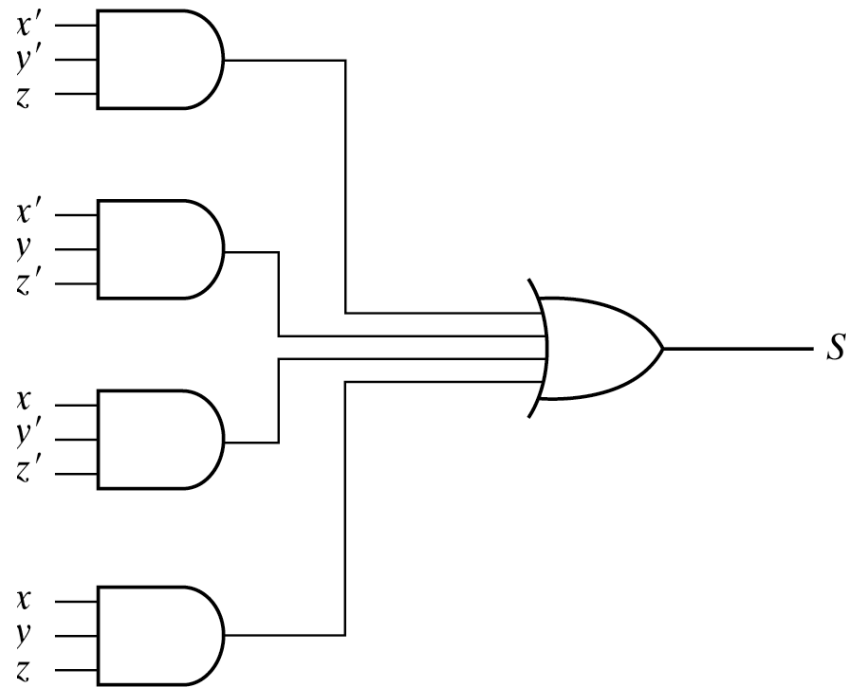
		$yz$			$y$	
		0 0	0 1	1 1	1 0	
$x$	0		1		1	
$x$	1	1		1		
		$z$				

$$S = x'y'z + x'yz' + xy'z' + xyz$$

		$yz$			$y$	
		0 0	0 1	1 1	1 0	
$x$	0			1		
$x$	1		1	1	1	
		$z$				

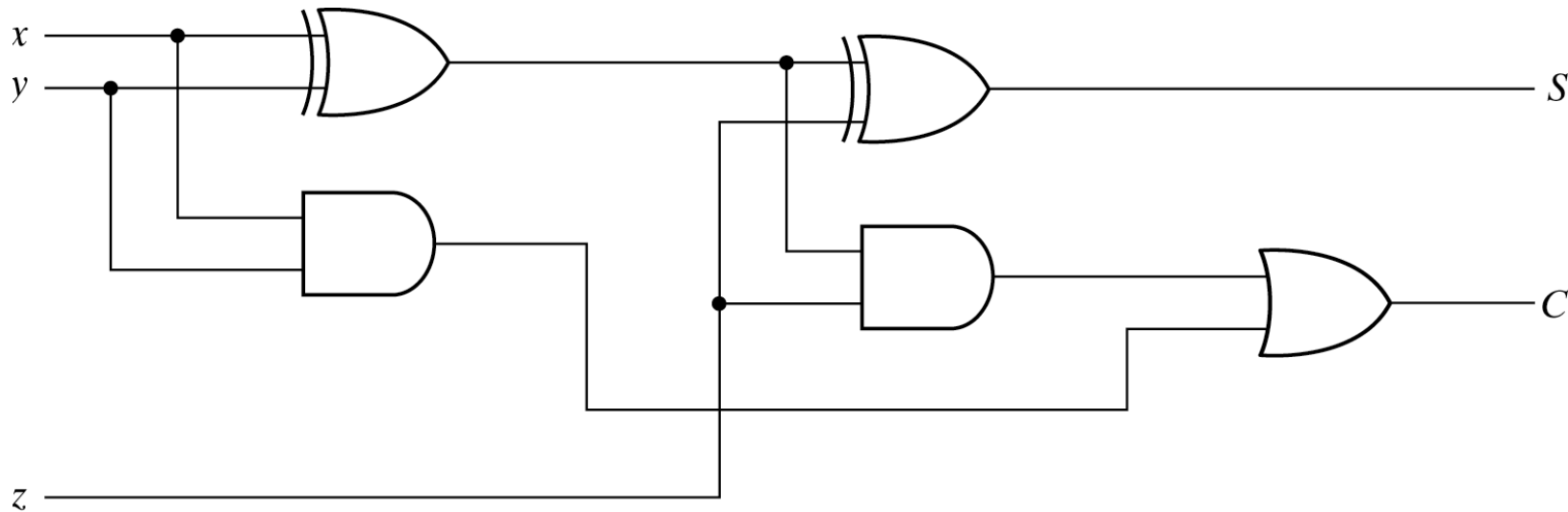
$$C = xy + xz + yz$$

# Implementation of FA



# Implementation of FA

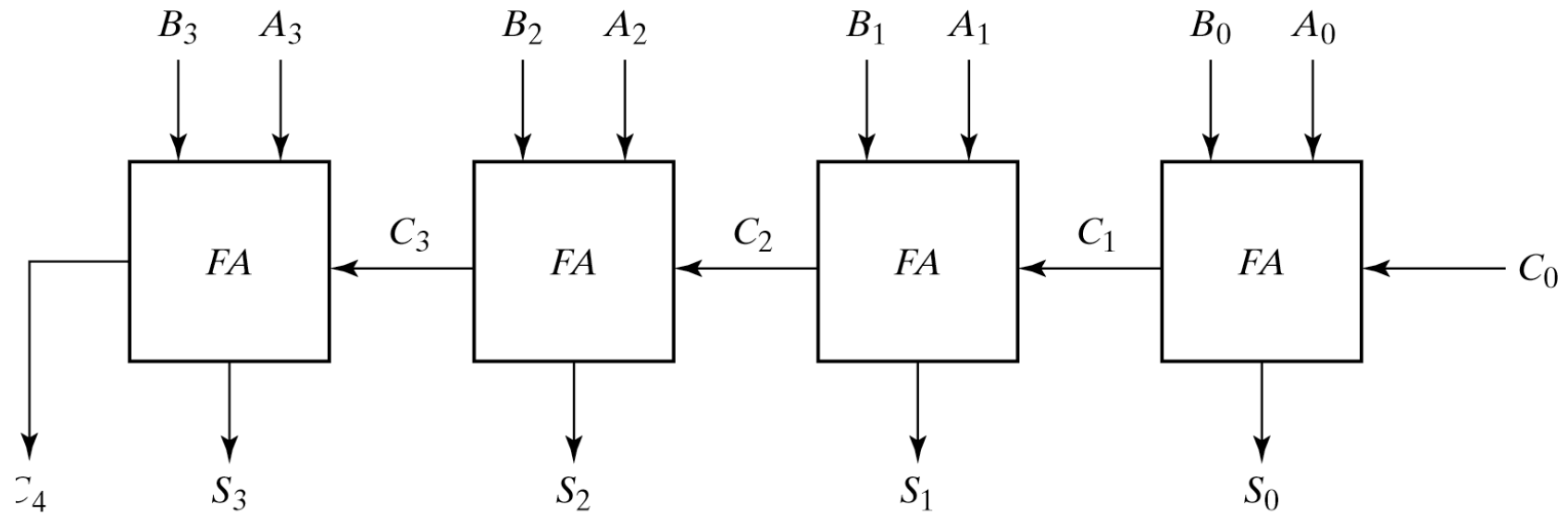
(with two half adder)



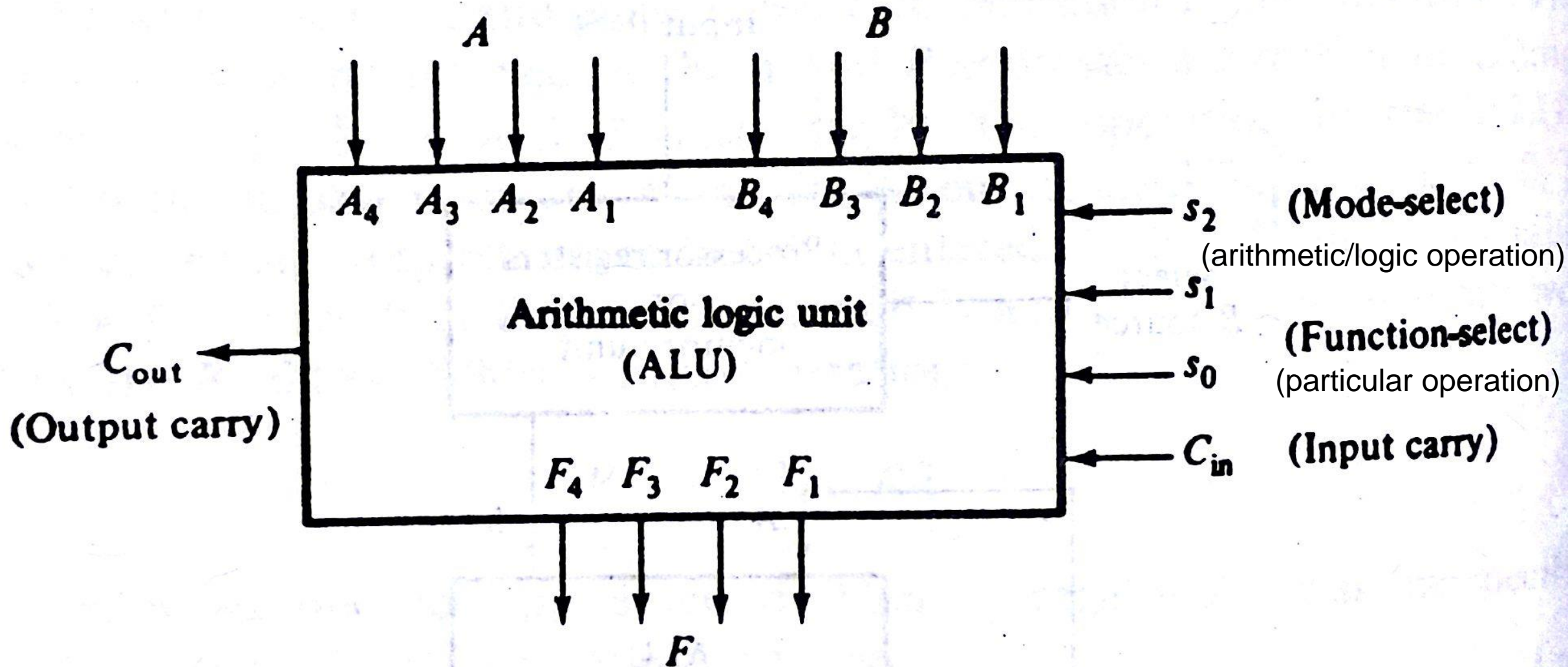
$$\begin{aligned} S &= xy'z' + x'yz' + xyz + x'y'z \\ &= z'(xy' + x'y) + z(xy + x'y') \\ &= z'(xy' + xy') + z(xy' + x'y)' \\ &= z \oplus (x \oplus y) \end{aligned}$$

$$\begin{aligned} C &= x'yz + xy'z + xyz' + xyz \\ &= x'yz + xy'z + xy \\ &= z(x'y + xy') + xy \\ &= z(x \oplus y) + xy \end{aligned}$$

# Binary Adder







**Figure 9-5** Block diagram of a 4-bit ALU

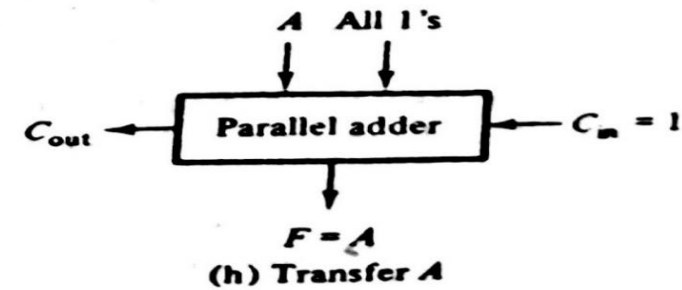
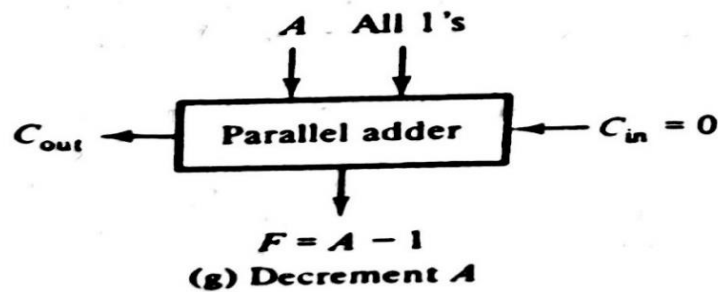
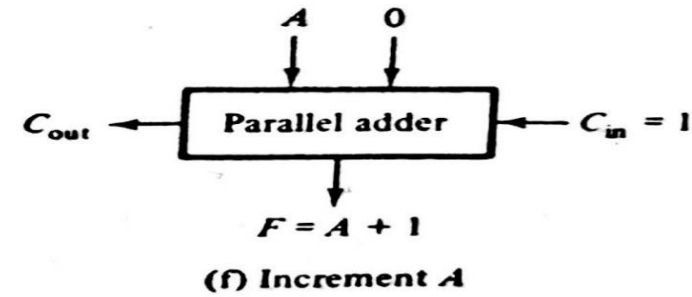
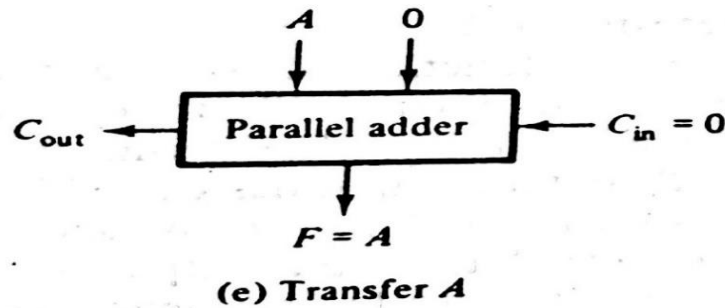
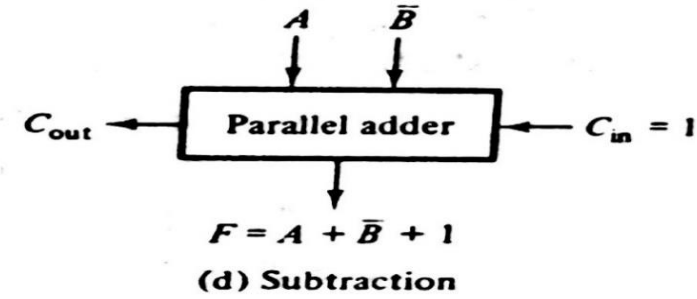
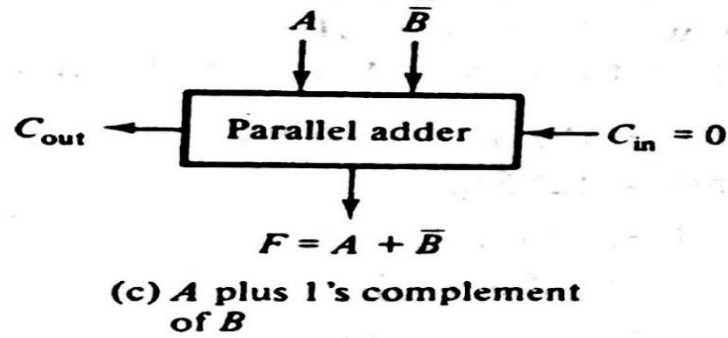
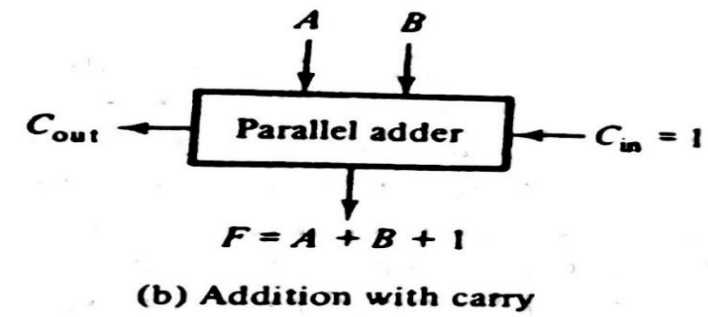
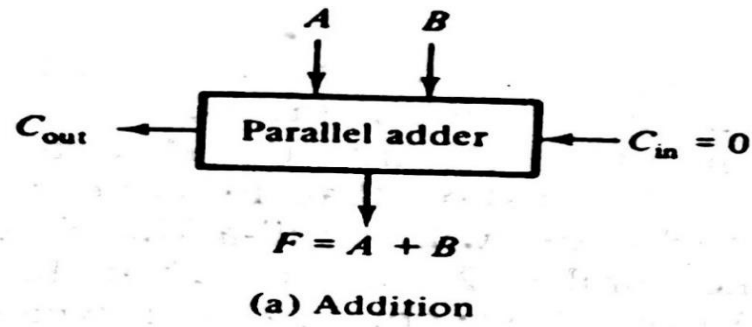


FIGURE 9-6 Operations obtained by controlling one set of inputs to a parallel adder

$F = A + 0 = A$ , which transfers input  $A$  into output  $F$ . Adding 1 through  $C_{in}$  as in Fig. 9-6(f), we obtain  $F = A + 1$ , which is the increment operation.

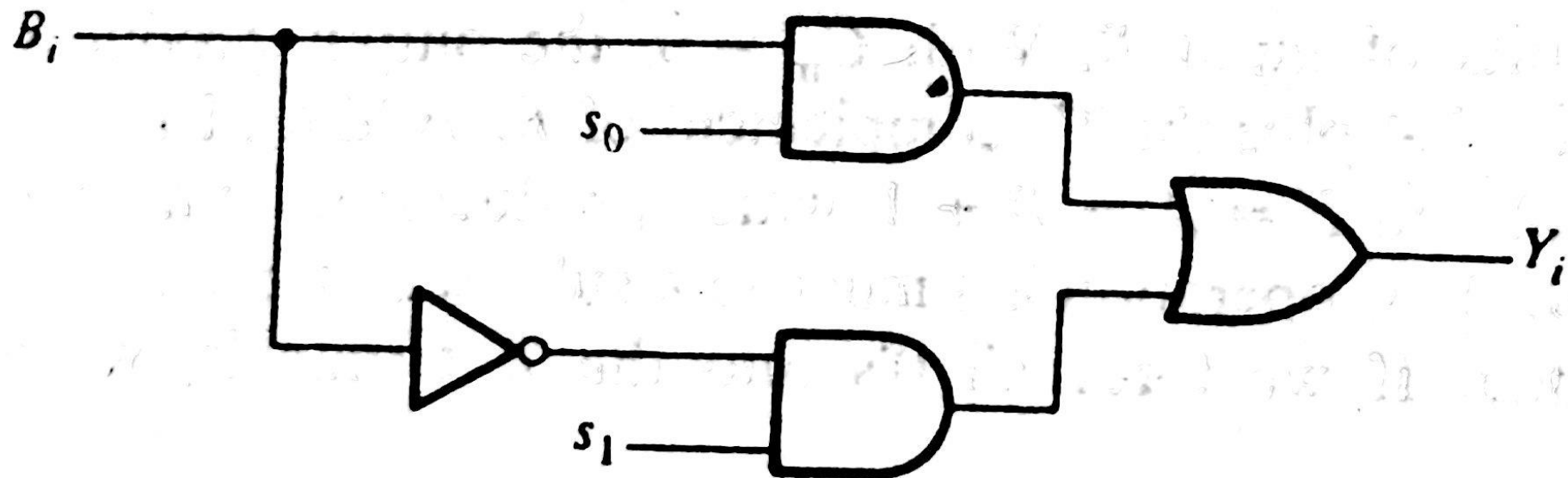
The condition illustrated in Fig. 9-6(g) inserts all 1's into the  $B$  terminals. This produces the decrement operation  $F = A - 1$ . To show that this condition is indeed a decrement operation, consider a parallel adder with  $n$  full-adder circuits. When the output carry is 1, it represents the number  $2^n$  because  $2^n$  in binary consists of a 1 followed by  $n$  0's. Subtracting 1 from  $2^n$ , we obtain  $2^n - 1$ , which in binary is a number of  $n$  1's. Adding  $2^n - 1$  to  $A$ , we obtain  $F = A + 2^n - 1 = 2^n + A - 1$ . If the output carry  $2^n$  is removed, we obtain  $F = A - 1$ .

To demonstrate with a numerical example, let  $n = 8$  and  $A = 9$ . Then:

$$\begin{aligned} A &= 0000 \ 1001 = (9)_{10} \\ 2^n &= 1 \ 0000 \ 0000 = (256)_{10} \\ 2^n - 1 &= 1111 \ 1111 = (255)_{10} \\ A + 2^n - 1 &= 1 \ 0000 \ 1000 = (256 + 8)_{10} \end{aligned}$$

Removing the output carry  $2^n = 256$ , we obtain  $8 = 9 - 1$ . Thus, we have decremented  $A$  by 1 by adding to it a binary number with all 1's.





$s_1$	$s_0$	$Y_i$
0	0	0
0	1	$B_i$
1	0	$B_i'$
1	1	1

**Figure 9-7 True/complement, one/zero circuit**



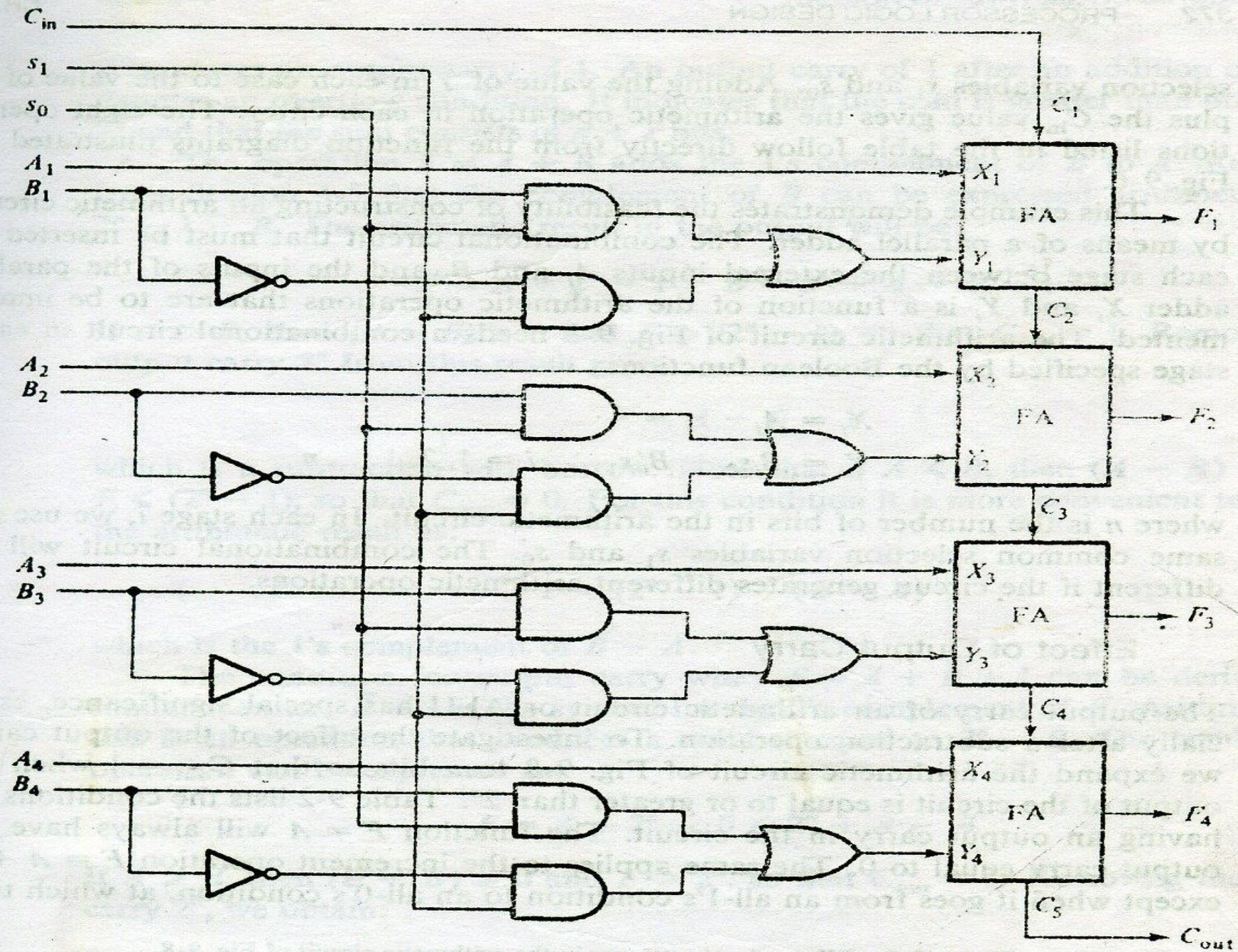


Figure 9-8 Logic diagram of arithmetic circuit



Function select			Y equals	Output equals	Function
$s_1$	$s_0$	$C_{in}$			
0	0	0	0	$F = A$	Transfer $A$
0	0	1	0	$F = A + 1$	Increment $A$
0	1	0	$B$	$F = A + B$	Add $B$ to $A$
0	1	1	$B$	$F = A + B + 1$	Add $B$ to $A$ plus 1
1	0	0	$\bar{B}$	$F = A + \bar{B}$	Add 1's complement of $B$ to $A$
1	0	1	$\bar{B}$	$F = A + \bar{B} + 1$	Add 2's complement of $B$ to $A$
1	1	0	All 1's	$F = A - 1$	Decrement $A$
1	1	1	All 1's	$F = A$	Transfer $A$

Enough Talk, Let's Work 😊