

DU_FlareBlitz Codebook

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1 Numerical algorithms

1.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a;
// computes lcm(a,b)
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                 if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
b >>= 1;
        return ret;
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
int yy = x = 1;
        while (b) {
                 int q = a / b;
                int t = b; b = a%b; a = t;
t = xx; xx = x - q*xx; x = t;
                 t = yy; yy = y - q*yy; y = t;
        return a;
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
        int x, y;
        VI ret,
        int g = extended_euclid(a, n, x, y);
        if (!(b%g)) {
                 x = mod(x*(b / g), n);
                 for (int i = 0; i < g; i++)
                         ret.push_back(mod(x + i*(n / g), n));
        return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
        int x, y;
        int g = extended_euclid(a, n, x, y);
if (g > 1) return -1;
        return mod(x, n);
```

```
// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = 1 \text{cm} (\text{m1, m2}).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
        int g = extended_euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make_pair(0, -1);
        return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2 / g);
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On // failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
        PII ret = make_pair(r[0], m[0]);
        for (int i = 1; i < m.size(); i++) {
                ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);
                if (ret.second == -1) break;
        return ret:
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
        if (!a && !b)
                if (c) return false:
                x = 0; y = 0;
                return true;
        if (!a)
                if (c % b) return false;
                x = 0; y = c / b;
                return true;
        if (!b)
                if (c % a) return false;
x = c / a; y = 0;
                return true;
        int g = gcd(a, b);
        if (c % g) return false;
        x = c / g * mod_inverse(a / g, b / g);
        y = (c - a * x) / b;
        return true;
int main() {
        // expected: 2
        cout << gcd(14, 30) << endl;</pre>
        // expected: 2 -2 1
        int x, y;
int g = extended_euclid(14, 30, x, y);
        cout << g << " " << x << " " << y << endl;
        VI sols = modular_linear_equation_solver(14, 30, 100);
        for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";
        cout << endl;
        // expected: 8
        cout << mod_inverse(8, 9) << endl;</pre>
        // expected: 23 105
        PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
        cout << ret.first << " " << ret.second << endl;</pre>
        ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
        cout << ret.first << " " << ret.second << endl;</pre>
        if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;</pre>
        cout << x << " " << y << endl;
        return 0;
```

Systems of linear equations, matrix inverse, determi-

// Gauss-Jordan elimination with full pivoting.

```
nant
```

```
// Uses:
    (1) solving systems of linear equations (AX=B)
    (2) inverting matrices (AX=I)
    (3) computing determinants of square matrices
// Running time: O(n^3)
// INPUT: a[][] = an nxn matrix
             b[][] = an nxm matrix
// OUTPUT: X
                   = an nxm matrix (stored in b[][])
             A^{-1} = an nxn matrix (stored in a[][])
             returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
 const int n = a.size():
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  for (int i = 0; i < n; i++) {
   int pj = -1, pk = -1;
    for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
      for (int k = 0; k < n; k++) if (!ipiv[k])
   ipiv | pk | ++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
   if (pj != pk) det *= -1;
irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] *= c;
    for (int p = 0; p < m; p++) b[pk][p] *= c;
    for (int p = 0; p < n; p++) if (p != pk) {
     c = a[p][pk];
     a[p][pk] = 0;
     for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c; for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
   for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);
  return det;
int main() {
 const int n = 4:
  const int m = 2:
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};
  double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
  VVT a(n), b(n);
  for (int i = 0; i < n; i++) {
   a[i] = VT(A[i], A[i] + n);
   b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
  // expected: 60
  cout << "Determinant: " << det << endl;</pre>
  // expected: -0.233333 0.166667 0.133333 0.0666667
               0.166667 0.166667 0.333333 -0.333333
               0.233333 0.833333 -0.133333 -0.0666667
  cout << "Inverse: " << endl;</pre>
  for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++)
     cout << a[i][j] << ' ';
    cout << endl;
```

2 Graph algorithms

2.1 Bellman-Ford shortest paths with negative edge weights (C++)

```
// This function runs the Bellman-Ford algorithm for single source
// shortest paths with negative edge weights. The function returns
// false if a negative weight cycle is detected. Otherwise, the
// function returns true and dist[i] is the length of the shortest
// path from start to i.
// Running time: O(|V|^3)
     INPUT: start, w[i][j] = cost of edge from i to j
    OUTPUT: dist[i] = min weight path from start to i
              prev[i] = previous node on the best path from the
                         start node
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std:
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool BellmanFord (const VVT &w, VT &dist, VI &prev, int start) {
 int n = w.size();
  prev = VI(n, -1);
  dist = VT(n, 1000000000);
dist[start] = 0;
  for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {</pre>
        if (dist[j] > dist[i] + w[i][j]){
          if (k == n-1) return false;
          dist[j] = dist[i] + w[i][j];
          prev[j] = i;
  return true:
```

2.2 Dijkstra and Floyd's algorithm (C++)

```
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;

typedef double T;
typedef vector<T> VT;
```

```
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
// This function runs Dijkstra's algorithm for single source
// shortest paths. No negative cycles allowed!
// Running time: O(|V|^2)
    INPUT: start, w[i][j] = cost of edge from i to j
    OUTPUT: dist[i] = min weight path from start to i
              prev[i] = previous node on the best path from the
                        start node
void Dijkstra (const VVT &w, VT &dist, VI &prev, int start) {
 int n = w.size();
  VI found (n);
  prev = VI(n, -1);
  dist = VT(n, 1000000000);
  dist[start] = 0;
  while (start != -1) {
    found[start] = true;
    int best = -1:
    for (int k = 0; k < n; k++) if (!found[k]) {
     if (dist[k] > dist[start] + w[start][k]){
       dist[k] = dist[start] + w[start][k];
        prev[k] = start;
      if (best == -1 || dist[k] < dist[best]) best = k;</pre>
    start = best;
// This function runs the Floyd-Warshall algorithm for all-pairs
// shortest paths. Also handles negative edge weights. Returns true
// if a negative weight cycle is found.
// Running time: O(|V|^3)
    INPUT: w[i][j] = weight of edge from i to j
OUTPUT: w[i][j] = shortest path from i to j
             prev[i][j] = node before j on the best path starting at i
bool FloydWarshall (VVT &w, VVI &prev) {
  int n = w.size();
  prev = VVI (n, VI(n, -1));
  for (int k = 0; k < n; k++) {
   w[i][j] = w[i][k] + w[k][j];
prev[i][j] = k;
  // check for negative weight cycles
  for (int i=0; i<n; i++)</pre>
   if (w[i][i] < 0) return false;</pre>
  return true;
```

2.3 Fast Dijkstra's algorithm

```
// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
//
// Running time: O(|E| log |V|)

#include <queue>
#include <cstdio>
using namespace std;
const int INF = 2000000000;
typedef pair<int, int> PII;
int main() {
    int N, s, t;
    scanf("%d%d%d", &N, &s, &t);
    vector<vector<PII>> edges(N);
    for (int i = 0; i < N, i++) {</pre>
```

```
int M;
                 scanf("%d", &M);
for (int j = 0; j < M; j++) {
                          int vertex, dist;
                          scanf("%d%d", &vertex, &dist);
                          edges[i].push_back(make_pair(dist, vertex)); // note order of arguments here
         // use priority queue in which top element has the "smallest" priority
        priority_queue<PII, vector<PII>, greater<PII> > Q;
vector<int> dist(N, INF), dad(N, -1);
        Q.push(make_pair(0, s));
         dist[s] = 0;
        while (!Q.empty()) {
    PII p = Q.top();
                 Q.pop();
                 int here = p second;
                 if (here == t) break;
                 if (dist[here] != p.first) continue;
                 for (vector<PII>::iterator it = edges[here].begin(); it != edges[here].end(); it++) {
                          if (dist[here] + it->first < dist[it->second]) {
                                   dist[it->second] = dist[here] + it->first;
                                   dad[it->second] = here;
                                   Q.push(make_pair(dist[it->second], it->second));
        printf("%d\n", dist[t]);
        if (dist[t] < INF)</pre>
                 for (int i = t; i != -1; i = dad[i])
                          printf("%d%c", i, (i == s ? '\n' : ' '));
        return 0;
Sample input:
5 0 4 2 1 2 3 1
2 2 4 4 5
3 1 4 3 3 4 1
2 1 5 2 1
Expected:
4 2 3 0
```

2.4 Strongly connected components

```
#include<memory.h>
struct edge{int e, nxt;};
int V. E.
edge e[MAXE], er[MAXE];
int sp[MAXV], spr[MAXV];
int group_cnt, group_num[MAXV];
bool v[MAXV];
int stk[MAXV];
void fill_forward(int x)
  int i
  for(i=sp[x];i;i=e[i].nxt) if(!v[e[i].e]) fill_forward(e[i].e);
  stk[++stk[0]]=x;
void fill_backward(int x)
  int i;
  v[x]=false;
  group_num[x]=group_cnt;
  for(i=spr[x];i;i=er[i].nxt) if(v[er[i].e]) fill_backward(er[i].e);
void add_edge(int v1, int v2) //add edge v1->v2
  e [++E].e=v2; e [E].nxt=sp [v1]; sp [v1]=E;
  er[ E].e=v1; er[E].nxt=spr[v2]; spr[v2]=E;
void SCC()
  int i:
  stk[0]=0;
  memset(v, false, sizeof(v));
  for(i=1;i<=V;i++) if(!v[i]) fill_forward(i);</pre>
  group_cnt=0;
```

```
for(i=stk[0];i>=1;i--) if(v[stk[i]]){group_cnt++; fill_backward(stk[i]);}
```

2.5 Eulerian path

```
typedef list<Edge>::iterator iter;
        int next_vertex;
        iter reverse_edge;
        Edge(int next_vertex)
                :next_vertex(next_vertex)
                { }
};
const int max vertices = :
int num vertices;
list<Edge> adj[max vertices];
                                        // adjacency list
vector<int> path;
void find_path(int v)
        while(adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().reverse_edge);
                adj[v].pop_front();
                find path(vn):
        path.push_back(v);
void add_edge(int a, int b)
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
        adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse_edge = itb;
        itb->reverse_edge = ita;
```

2.6 Kruskal's algorithm

```
Uses Kruskal's Algorithm to calculate the weight of the minimum spanning
forest (union of minimum spanning trees of each connected component) of
a possibly disjoint graph, given in the form of a matrix of edge weights (-1 if no edge exists). Returns the weight of the minimum spanning forest (also calculates the actual edges - stored in T). Note: uses a
disjoint-set data structure with amortized (effectively) constant time per
union/find. Runs in O(E*log(E)) time.
#include <iostream>
#include <vector>
#include <algorithm>
#include <queue>
using namespace std;
typedef int T;
struct edge
  int u, v;
  T d;
};
struct edgeCmp
  int operator()(const edge& a, const edge& b) { return a.d > b.d; }
int find(vector <int>& C, int x) { return (C[x] == x) ? x : C[x] = find(C, C[x]); }
T Kruskal (vector <vector <T> >& w)
  int n = w.size();
```

```
T weight = 0;
  vector <int> C(n), R(n);
  for(int i=0; i<n; i++) { C[i] = i; R[i] = 0; }</pre>
  vector <edge> T;
  priority_queue <edge, vector <edge>, edgeCmp> E;
  for(int i=0; i<n; i++)</pre>
    for (int j=i+1; j<n; j++)</pre>
      if(w[i][j] >= 0)
        e.u = i; e.v = j; e.d = w[i][j];
         E.push(e);
  while (T.size() < n-1 && !E.empty())
    edge cur = E.top(); E.pop();
    int uc = find(C, cur.u), vc = find(C, cur.v);
    if(uc != vc)
      T.push_back(cur); weight += cur.d;
      if(R[uc] > R[vc]) C[vc] = uc;
else if(R[vc] > R[uc]) C[uc] = vc;
      else { C[vc] = uc; R[uc]++; }
  return weight:
int main()
  int wa[6][6] = {
    { 0, -1, 2, -1, 7, -1 }, 
{ -1, 0, -1, 2, -1, -1 },
    { 2, -1, 0, -1, 8, 6 },
{ -1, 2, -1, 0, -1, -1 },
    { 7, -1, 8, -1, 0, 4 },
{ -1, -1, 6, -1, 4, 0 } };
  vector <vector <int> > w(6, vector <int>(6));
  for(int i=0; i<6; i++)
    for (int j=0; j<6; j++)</pre>
      w[i][j] = wa[i][j];
  cout << Kruskal(w) << endl;
  cin >> wa[0][0];
```

2.7 Minimum spanning trees

```
// This function runs Prim's algorithm for constructing minimum
// weight spanning trees.
// Running time: O(|V|^2)
     INPUT: w[i][j] = cost \ of \ edge \ from \ i \ to \ j
              NOTE: Make sure that w[i][j] is nonnegative and
              symmetric. Missing edges should be given -1
              weight.
    OUTPUT: edges = list of pair<int,int> in minimum spanning tree
             return total weight of tree
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
```

```
T Prim (const VVT &w, VPII &edges) {
  int n = w.size();
  VI found (n);
  VI prev (n, -1);
  VT dist (n, 1000000000);
  int here = 0;
  dist[here] = 0;
  while (here != -1) {
    found[here] = true;
    int best = -1;
    for (int k = 0; k < n; k++) if (!found[k]) {
   if (w[here][k] != -1 && dist[k] > w[here][k]) {
        dist[k] = w[here][k];
prev[k] = here;
       if (best == -1 || dist[k] < dist[best]) best = k;</pre>
  T tot_weight = 0;
  for (int i = 0; i < n; i++) if (prev[i] != -1) {
    edges.push_back (make_pair (prev[i], i));
    tot_weight += w[prev[i]][i];
  return tot weight:
int main() {
  int ww[5][5] = {
    {0, 400, 400, 300, 600},
     \{400, 0, 3, -1, 7\},\
     {400, 3, 0, 2, 0},
     {300, -1, 2, 0, 5},
    {600, 7, 0, 5, 0}
  VVT w(5, VT(5));
  for (int i = 0; i < 5; i++)
for (int j = 0; j < 5; j++)
w[i][j] = ww[i][j];</pre>
  VPII edges;
  cout << Prim (w, edges) << endl;</pre>
  for (int i = 0; i < edges.size(); i++)</pre>
    cout << edges[i].first << " " << edges[i].second << endl;</pre>
```

3 Data structures

3.1 Suffix array

```
// Suffix array construction in O(L log^2 L) time. Routine for
// computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
            of substring s[i...L-1] in the list of sorted suffixes.
            That is, if we take the inverse of the permutation suffix[],
            we get the actual suffix array.
#include <vector>
#include <iostream>
#include <string>
using namespace std;
struct SuffixArray
 const int L:
  string s;
  vector<vector<int> > P;
  vector<pair<int,int>,int> > M;
  SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<int>(L, 0)), M(L) {
   for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
      P.push_back(vector<int>(L, 0));
```

```
for (int i = 0; i < L; i++)</pre>
        M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
      sort(M.begin(), M.end());
      for (int i = 0; i < L; i++)
        P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i;
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
    int len = 0;
    if (i == j) return L - i,
for (int k = P size() - 1, k >= 0 && i < L && j < L, k--) {
      if (P[k][i] == P[k][j]) {
        len += 1 << k;
    return len;
};
// BEGIN CUIT
// The following code solves UVA problem 11512: GATTACA.
#define TESTING
#ifdef TESTING
int main() {
 int T;
  cin >> T;
  for (int caseno = 0; caseno < T; caseno++) {
    string s;
    cin >> s;
    SuffixArray array(s);
    vector<int> v = array.GetSuffixArray();
int bestlen = -1, bestpos = -1, bestcount = 0;
    for (int i = 0; i < s.length(); i++) {</pre>
      int len = 0, count = 0;
      for (int j = i+1; j < s.length(); j++) {
        int 1 = array.LongestCommonPrefix(i, j);
        if (1 >= len) {
          if (1 > len) count = 2; else count++;
          len = 1;
      if (len > bestlen || len == bestlen && s.substr(bestpos, bestlen) > s.substr(i, len)) {
        bestlen = len;
        bestcount = count:
        bestpos = i;
    if (bestlen == 0) {
      cout << "No repetitions found!" << endl;</pre>
    } else {
      cout << s.substr(bestpos, bestlen) << " " << bestcount << endl;</pre>
// END CUT
  // bobocel is the O'th suffix
  // obocel is the 5'th suffix
       bocel is the 1'st suffix
       ocel is the 6'th suffix
         cel is the 2'nd suffix
         el is the 3'rd suffix
           l is the 4'th suffix
  SuffixArray suffix("bobocel");
  vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
  for (int i = 0; i < v.size(); i++) cout << v[i] << " ";
  cout << endl;
  cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
// BEGIN CUT
#endif
// END CUT
```

3.2 Binary Indexed Tree

#include <iostream>

```
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
  while (x \le N) {
   tree |x| += v;
    x += (x & -x);
// get cumulative sum up to and including x
int get(int x) {
  while(x) {
    \mathbf{x} -= (\mathbf{x} \& -\mathbf{x});
  return res:
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
 int idx = 0, mask = N;
  while (mask && idx < N)
    int t = idx + mask;
    if(x >= tree[t]) {
      idx = t;
      x -= tree[t];
    mask >>= 1;
  return idx;
```

3.3 Union-find set

```
#include <iostream>
#include vector>
using namespace std;
struct UnionFind {
    vector<int> C;
    UnionFind(int n) : C(n) { for (int i = 0; i < n; i++) C[i] = i; }
    int find(int x) { return (C[x] == x) ? x : C[x] = find(C[x]); }
    void merge(int x, int y) { C[find(x)] = find(y); }
};
int main() {
    int n = 5;
    UnionFind uf(n);
    uf. merge(0, 2);
    uf. merge(1, 0);
    uf. merge(3, 4);
    for (int i = 0; i < n; i++) cout << i << " " << uf.find(i) << endl;
    return 0;
}</pre>
```

3.4 Lazy segment tree

```
public void update(int begin, int end, int val) {
        update(1,0,origSize-1,begin,end,val);
public void update(int curr, int tBegin, int tEnd, int begin, int end, int val)
        if(tBegin >= begin && tEnd <= end)</pre>
                  update[curr] += val;
                  leaf[curr] += (Math.min(end,tEnd)-Math.max(begin,tBegin)+1) * val;
                  int mid = (tBegin+tEnd)/2;
                 if(mid >= begin && tBegin <= end)
    update(2*curr, tBegin, mid, begin, end, val);</pre>
                 if(tEnd >= begin && mid+1 <= end)</pre>
                          update(2*curr+1, mid+1, tEnd, begin, end, val);
public long query(int begin, int end) {
        return query (1, 0, origSize-1, begin, end);
public long query(int curr, int tBegin, int tEnd, int begin, int end) {
        if(tBegin >= begin && tEnd <= end)
                 if(update[curr] != 0) {
                           leaf[curr] += (tEnd-tBegin+1) * update[curr];
                           if(2*curr < update.length){</pre>
                                   update[2*curr] += update[curr];
update[2*curr+1] += update[curr];
                           update[curr] = 0;
                  return leaf[curr]:
                  leaf[curr] += (tEnd-tBegin+1) * update[curr];
                  if(2*curr < update.length){</pre>
                          update[2*curr] += update[curr];
                           update[2*curr+1] += update[curr];
                 update[curr] = 0;
int mid = (tBegin+tEnd)/2;
                 long ret = 0;
                 if(mid >= begin && tBegin <= end)
    ret += query(2*curr, tBegin, mid, begin, end);</pre>
                  if (tEnd >= begin && mid+1 <= end)
                         ret += query(2*curr+1, mid+1, tEnd, begin, end);
                  return ret:
```

3.5 Lowest common ancestor

```
const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max_nodes];
                                           // children[i] contains the children of node i
int A[max_nodes][log_max_nodes+1];
                                           // A[i][i] is the 2^i-th ancestor of node i, or -1 if that
      ancestor does not exist
int L[max_nodes];
                                            // L[i] is the distance between node i and the root
// floor of the binary logarithm of n
int lb(unsigned int n)
    if(n==0)
        return -1;
    int p = 0;
    if (n >= 1<<16) { n >>= 16; p += 16;
    if (n >= 1<< 8) { n >>= 8; p += 8; }
if (n >= 1<< 4) { n >>= 4; p += 4; }
if (n >= 1<< 2) { n >>= 2; p += 2; }
    if (n >= 1<< 1) {
    return p;
void DFS(int i, int 1)
    for(int j = 0; j < children[i].size(); j++)</pre>
        DFS(children[i][j], 1+1);
int LCA(int p, int q)
    // ensure node p is at least as deep as node q
    if(L[p] < L[q])
        swap(p, q);
    // "binary search" for the ancestor of node p situated on the same level as q
    for(int i = log_num_nodes; i >= 0; i--)
        if(L[p] - (1<<i) >= L[q])
```

```
p = A[p][i];
    if(p == q)
        return p;
    // "binary search" for the LCA
    for(int i = log_num_nodes; i >= 0; i--)
        if(A[p][i] != -1 && A[p][i] != A[q][i])
            p = A[p][i];
            q = A[q][i];
    return A[p][0];
int main(int argc,char* argv[])
    // read num_nodes, the total number of nodes
    log_num_nodes=1b(num_nodes);
    for(int i = 0; i < num_nodes; i++)</pre>
        // read p, the parent of node i or -1 if node i is the root
        A[i][0] = p;
        if (p != -1)
            children[p] push_back(i);
        else
            root = i;
    // precompute A using dynamic programming
    for(int j = 1; j <= log_num_nodes; j++)
    for(int i = 0; i < num_nodes; i++)</pre>
            if(A[i][j-1] != -1)
                A[i][j] = A[A[i][j-1]][j-1];
            else
                 A[i][j] = -1;
    // precompute L
    DFS (root, 0);
    return 0;
```

4 Geometry

4.1 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is
// #defined.
// Running time: O(n log n)
     INPUT: a vector of input points, unordered.
    OUTPUT: a vector of points in the convex hull, counterclockwise, starting
              with bottommost/leftmost point
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
// BEGIN CUT
#include <map>
// END CUT
using namespace std;
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT (
 Тх, у;
  PT() {}
 PT(T x, T y) : x(x), y(y) {}
bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(rhs.y,rhs.x); }</pre>
 bool operator==(const PT &rhs) const { return make_pair(y,x) == make_pair(rhs.y,rhs.x); }
```

```
T cross(PT p, PT q) { return p.x+q.y-p.y+q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
bool between (const PT &a, const PT &b, const PT &c) {
  return (fabs(area2(a,b,c)) < EPS && (a.x-b.x) \star(c.x-b.x) <= 0 && (a.y-b.y) \star(c.y-b.y) <= 0);
#endif
void ConvexHull(vector<PT> &pts) {
  sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
    while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >= 0) up.pop_back();
while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();</pre>
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  pts = dn;
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE_REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear():
  dn.push_back(pts[0]);
  dn.push back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {</pre>
    if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
  pts = dn;
#endif
// BEGIN CUT
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)
int main() {
 int t
  scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {</pre>
    int n
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);</pre>
    vector<PT> h(v);
    map<PT.int> index:
    for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
    double len = 0:
    for (int i = 0; i < h.size(); i++) {</pre>
     double dx = h[i].x - h[(i+1)%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
    printf("%.2f\n", len);
for (int i = 0; i < h.size(); i++) {
   if (i > 0) printf(" ");
      printf("%d", index[h[i]]);
    printf("\n");
// END CUT
```

4.2 Miscellaneous geometry

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = le100;
double EPS = le-12;
```

struct PT { double x, y, PT() {} PT(double x, double y) : x(x), y(y) {} PT(const PT &p) : x(p.x), y(p.y) PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); PT operator * (double c) const { return PT(x*c, y*c); PT operator / (double c) const { return PT(x/c, y/c); } double dot (PT p, PT q) { return p.x*q.x+p.y*q.y; } double dist2(PT p, PT q) { return dot(p-q,p-q); } double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; } ostream &operator<<(ostream &os. const PT &p) {
 return os << "(" << p.x << "," << p.y << ")"; // rotate a point CCW or CW around the origin PT RotateCCW90 (PT p) { return PT(-p.y,p.x); }
PT RotateCW90 (PT p) { return PT(p.y,-p.x); } PT RotateCCW(PT p, double t) { return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t)); // project point c onto line through a and b // assuming a != b PT ProjectPointLine(PT a, PT b, PT c) { return a + (b-a) *dot(c-a, b-a) /dot(b-a, b-a); // project point c onto line segment through a and b PT ProjectPointSegment (PT a, PT b, PT c) { double r = dot(b-a,b-a); if (fabs(r) < EPS) return a; r = dot(c-a, b-a)/r;if (r < 0) return a; if (r > 1) return b; return a + (b-a) *r; // compute distance from c to segment between a and b
double DistancePointSegment (PT a, PT b, PT c) { return sqrt(dist2(c, ProjectPointSegment(a, b, c))); // compute distance between point (x,y,z) and plane ax+by+cz=d double DistancePointPlane(double x, double y, double z, double a, double b, double c, double d) return fabs(a*x+b*v+c*z-d)/sgrt(a*a+b*b+c*c); // determine if lines from a to b and c to d are parallel or collinear bool LinesParallel(PT a, PT b, PT c, PT d) { return fabs(cross(b-a, c-d)) < EPS; bool LinesCollinear (PT a, PT b, PT c, PT d) { return LinesParallel(a, b, c, d) && fabs(cross(a-b, a-c)) < EPS && fabs(cross(c-d, c-a)) < EPS; // determine if line segment from a to b intersects with // line segment from c to d bool SegmentsIntersect(PT a, PT b, PT c, PT d) { if (LinesCollinear(a, b, c, d)) {
 if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
 dist2(b, c) < EPS || dist2(b, d) < EPS) return true; if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)return false; return true; if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false; if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false; // compute intersection of line passing through a and b // with line passing through c and d, assuming that unique // intersection exists; for segment intersection, check if // seaments intersect first PT ComputeLineIntersection(PT a, PT b, PT c, PT d) { b=b-a; d=c-d; c=c-a; assert (dot (b, b) > EPS && dot (d, d) > EPS); return a + b*cross(c, d)/cross(b, d); // compute center of circle given three points

PT ComputeCircleCenter(PT a, PT b, PT c) {

```
b = (a+b)/2;
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// \ {\it integer \ arithmetic \ by \ taking \ care \ of \ the \ division \ appropriately}
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {
  int j = (i+1) %p.size();</pre>
    if ((p[i].y <= q.y && q.y < p[j].y ||
  p[j].y <= q.y && q.y < p[i].y) &&</pre>
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
      return true:
    return false:
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret;
  b = b-a;
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C:
  if (D < -EPS) return ret;
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R | | d+min(r, R) < max(r, R)) return ret;</pre>
  double x = (d*d-R*R+r*r)/(2*d);
  double v = sart(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (v > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0:
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
 PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {
  int j = (i+1) % p.size();</pre>
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
```

```
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++)
    for (int k = i+1; k < p.size(); k++) {
      int j = (i+1) % p.size();
      int 1 = (k+1) % p.size();
      if (i == 1 || j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
  return true;
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5,2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
  cerr << ProjectPointSegment (PT(-5,-2), PT(10,4), PT(3,7)) << " "
       << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "
       << ProjectPointSegment (PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
  cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
       << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  // expected: 0 0 1
  << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  // expected: 1 1 1 0
  cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "</pre>
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
  // expected: (1,2)
  cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
  // expected: (1.1)
  cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
  vector<PT> v:
  v.push_back(PT(0,0));
  v.push_back(PT(5,0));
  v.push_back(PT(5,5));
  v.push_back(PT(0,5));
  // expected: 1 1 1 0 0
  cerr << PointInPolygon(v, PT(2,2)) << " "
       << PointInPolygon(v, PT(2,0)) << " "
       << PointInPolygon(v, PT(0,2)) << " "
       << PointInPolygon(v, PT(5,2)) << " "
       << PointInPolygon(v, PT(2,5)) << endl;
  // expected: 0 1 1 1 1
  cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
       << PointOnPolygon(v, PT(2,0)) << " "
       << PointOnPolygon(v, PT(0,2)) << " "
       << PointOnPolygon(v, PT(5,2)) << " "
       << PointOnPolygon(v, PT(2,5)) << endl;
  // expected: (1,6)
               (5,4) (4,5)
               blank line
               (4,5) (5,4)
               blank line
               (4.5) (5.4)
  vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
  u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; ce
                                                        ": cerr << endl:
  u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
  for (int i = 0; i < u.size(); i++) cerr << u[i] <<
                                                        "; cerr << endl;
  u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
  u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
```

```
for (int i = 0, i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt (2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;

// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;
return 0;</pre>
```

4.3 Slow Delaunay triangulation

```
// Slow but simple Delaunay triangulation. Does not handle
// degenerate cases (from O'Rourke, Computational Geometry in C)
// Running time: O(n^4)
             x[] = x-coordinates
              y[] = y-coordinates
// OUTPUT: triples = a vector containing m triples of indices
                         corresponding to triangle vertices
#include<vector>
using namespace std;
typedef double T;
struct triple (
    int i, j, k;
    triple() {}
    triple(int i, int j, int k) : i(i), j(j), k(k) {}
vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
        int n = x.size();
        vector<T> z(n);
        vector<triple> ret;
        for (int i = 0; i < n; i++)
z[i] = x[i] * x[i] + y[i] * y[i];</pre>
        for (int i = 0; i < n-2; i++) {
   for (int j = i+1; j < n; j++) {
      for (int k = i+1; k < n; k++) {</pre>
                      if (j == k) continue;
                      double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])*(z[j]-z[i]);
                      double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])*(z[k]-z[i]);
                      double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])*(y[j]-y[i]);
                      bool flag = zn < 0;
                      for (int m = 0; flag && m < n; m++)</pre>
                          flag = flag && ((x[m]-x[i])*xn +
                                            (y[m]-y[i])*yn +
                                            (z[m]-z[i])*zn <= 0);
                      if (flag) ret.push_back(triple(i, j, k));
            }
        return ret;
int main()
    T \times s[] = {0, 0, 1, 0.9};
    vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
    vector<triple> tri = delaunayTriangulation(x, y);
    //expected: 0 1 3
    int i;
    for(i = 0; i < tri.size(); i++)</pre>
        printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
```

5 Combinatorial optimization

5.1 Dense max-flow

```
// Adjacency matrix implementation of Dinic's blocking flow algorithm.
// Running time:
        0(|V|^4)
// INPUT:
        - graph, constructed using AddEdge()
        - source
        - sink
         - maximum flow value
         - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
struct MaxFlow {
  int N;
  VVI cap, flow;
  VI dad, Q;
    N(N), cap(N, VI(N)), flow(N, VI(N)), dad(N), Q(N) {}
  void AddEdge(int from, int to, int cap) {
  this->cap[from][to] += cap;
  int BlockingFlow(int s, int t) {
    fill(dad.begin(), dad.end(), -1);
     dad[s] = -2;
     int head = 0, tail = 0;
    Q[tail++] = s;
while (head < tail) {</pre>
       int x = Q[head++];
       for (int i = 0; i < N; i++) {
  if (dad[i] == -1 && cap[x][i] - flow[x][i] > 0) {
            dad[i] = x;
            Q[tail++] = i;
     if (dad[t] == -1) return 0;
     int totflow = 0;
     for (int i = 0; i < N; i++) {
  if (dad[i] == -1) continue;</pre>
       int amt = cap[i][t] - flow[i][t];
for (int j = i; amt && j != s; j = dad[j])
    amt = min(amt, cap[dad[j]][j] - flow[dad[j]][j]);
if (amt == 0) continue;
       flow[i][t] += amt;
       flow[t][t] r= amt;
flow[t][i] r= amt;
for (int j = i; j != s; j = dad[j]) {
  flow[dad[j]][j] += amt;
          flow[j][dad[j]] -= amt;
       totflow += amt;
     return totflow;
  int GetMaxFlow(int source, int sink) {
    int totflow = 0;
     while (int flow = BlockingFlow(source, sink))
      totflow += flow;
     return totflow;
};
```

```
int main() {
  MaxFlow mf(5);
  mf.AddEdge(0, 1, 3);
  mf.AddEdge(0, 2, 4);
  mf.AddEdge(0, 3, 5);
  mf.AddEdge(0, 4, 5);
  mf.AddEdge(1, 2, 2);
 mf.AddEdge(2, 3, 4);
  mf.AddEdge(2, 4, 1);
 mf.AddEdge(3, 4, 10);
  // should print out "15"
  cout << mf.GetMaxFlow(0, 4) << endl;</pre>
// The following code solves SPOJ problem #203: Potholers (POTHOLE)
#ifdef COMMENT
int main() {
 int t;
  cin >> t:
  for (int i = 0; i < t; i++) {</pre>
    int n;
    cin >> n:
    MaxFlow mf(n);
    for (int j = 0; j < n-1; j++) {
     int m:
      cin >> m:
      for (int k = 0; k < m; k++) {
       int p:
        cin >> p;
        int cap = (j == 0 || p == n-1) ? 1 : INF;
        mf.AddEdge(j, p, cap);
    cout << mf.GetMaxFlow(0, n-1) << endl;</pre>
  return 0:
#endif
// END CUT
```

5.2 Min-cost max-flow

```
//\ {\tt Implementation\ of\ min\ cost\ max\ flow\ algorithm\ using\ adjacency}
// matrix (Edmonds and Karp 1972). This implementation keeps track of
// forward and reverse edges separately (so you can set cap[i][j] !=
// \operatorname{cap}[j][i]). For a regular max flow, set all edge costs to 0.
// Running time, O(|V|^2) cost per augmentation
     max flow:
       min cost max flow: O(|V|^4 * MAX_EDGE_COST) augmentations
      - graph, constructed using AddEdge()
      - source
       - (maximum flow value, minimum cost value)
      - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow (
  int N
  VVL cap, flow, cost;
  VI found,
  VL dist, pi, width;
```

```
VPII dad;
  MinCostMaxFlow(int N) :
    N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
    found(N), dist(N), pi(N), width(N), dad(N) {}
  void AddEdge(int from, int to, L cap, L cost) {
    this->cap[from][to] = cap;
    this->cost[from][to] = cost;
  void Relax(int s, int k, L cap, L cost, int dir) {
  L val = dist[s] + pi[s] - pi[k] + cost;
  if (cap && val < dist[k]) {</pre>
     dist[k] = val;
dad[k] = make_pair(s, dir);
      width[k] = min(cap, width[s]);
  L Dijkstra(int s, int t) {
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
      int best = -1;
      found[s] = true;
      for (int k = 0; k < N; k++) {
        if (found[k]) continue;
        Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
Relax(s, k, flow[k][s], -cost[k][s], -1);
        if (best == -1 || dist[k] < dist[best]) best = k;</pre>
      s = best;
    for (int k = 0; k < N; k++)
pi[k] = min(pi[k] + dist[k], INF);</pre>
    return width[t];
  pair<L, L> GetMaxFlow(int s, int t) {
    L totflow = 0, totcost = 0;
    while (L amt = Dijkstra(s, t)) {
      totflow += amt;
      for (int x = t; x != s; x = dad[x].first) {
        if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
          totcost += amt * cost[dad[x].first][x];
        | else {
          flow[x][dad[x].first] -= amt;
          totcost -= amt * cost[x][dad[x].first];
      }
    return make_pair(totflow, totcost);
};
// The following code solves UVA problem #10594: Data Flow
int main() {
  int N, M;
  while (scanf("%d%d", &N, &M) == 2) {
    VVL v(M, VL(3));
    for (int i = 0; i < M; i++)
      scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[i][2]);
    LDK
    scanf("%Ld%Ld", &D, &K);
    MinCostMaxFlow mcmf(N+1);
    for (int i = 0; i < M; i++) {
      mcmf.AddEdge(int(v[i][0]), int(v[i][1]), K, v[i][2]);
      mcmf.AddEdge(int(v[i][1]), int(v[i][0]), K, v[i][2]);
    mcmf.AddEdge(0, 1, D, 0);
    pair<L, L> res = mcmf.GetMaxFlow(0, N);
    if (res.first == D) {
      printf("%Ld\n", res.second);
    } else {
      printf("Impossible.\n");
```

return 0;

// END CUT

6 Miscellaneous

6.1 Longest increasing subsequence

```
// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
// Running time: O(n log n)
    INPUT: a vector of integers
    OUTPUT: a vector containing the longest increasing subsequence
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
VI LongestIncreasingSubsequence(VI v) {
  VPII best;
  VI dad(v.size(), -1);
  for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASIG
    PII item = make_pair(v[i], 0);
    VPII::iterator it = lower_bound(best.begin(), best.end(), item);
    item second = i;
#else
    PII item = make_pair(v[i], i);
    VPII::iterator it = upper_bound(best.begin(), best.end(), item);
#endif
    if (it == best.end()) {
     dad[i] = (best.size() == 0 ? -1 : best.back().second);
      best.push_back(item);
    } else {
     dad[i] = it == best.begin() ? -1 : prev(it)->second;
      *it = item:
  VI ret;
  for (int i = best.back().second; i >= 0; i = dad[i])
   ret.push_back(v[i]);
  reverse(ret.begin(), ret.end());
  return ret:
```

6.2 Prime numbers

```
// O(sgrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x)
  if(x<=1) return false;</pre>
  if(x<=3) return true;</pre>
  if (!(x%2) || !(x%3)) return false;
  LL s=(LL) (sqrt((double)(x))+EPS);
  for(LL i=5;i<=s;i+=6)
    if (!(x%i) || !(x%(i+2))) return false;
  return true:
// Primes less than 1000:
                               59
                                     61
                                                              79
                                                                    83
                              109
                                                             139
                                                                   149
```

```
229
     283
                 307
                                                     347
                                                           349
                                                                       359
                       383
                             389
                                               409
                                                     419
                                                           421
                                                                       433
     439
           443
                       457
                             461
                                   463
                                         467
                                               479
                                                     487
                                                          491
     509
                       541
                             547
                                               569
                                                           577
                                                                 587
     599
           601
                 607
                       613
                             617
                                   619
                                         631
                                               641
                                                     643
                                                           647
     661
           673
                       683
                             691
                                         709
                                               719
                                                                       743
                       769
                                   787
                                         797
                                               809
                                                     811
                                                          821
                                                                823
                                                                      827
                                         877
     829
           839
                 8.5.3
                       8.57
                             8.59
                                   863
                                               881
                                                    883
                                                          887
                                                                907
                                                                      911
     919
           929
                       941
                             947
                                   9.5.3
                                         967
                                               971
                                                          983
// Other primes:
     The largest prime smaller than 10 is 7.
     The largest prime smaller than 100 is 97.
     The largest prime smaller than 1000 is 997
     The largest prime smaller than 10000 is 9973.
     The largest prime smaller than 100000 is 99991.
     The largest prime smaller than 1000000 is 999983
     The largest prime smaller than 10000000 is 9999991.
     The largest prime smaller than 100000000 is 99999989.
     The largest prime smaller than 1000000000 is 999999937.
     The largest prime smaller than 10000000000 is 9999999967.
     The largest prime smaller than 10000000000 is 99999999977.
     The largest prime smaller than 100000000000 is 99999999989. The largest prime smaller than 1000000000000 is 999999999971.
     The largest prime smaller than 1000000000000 is 9999999999973.
     The largest prime smaller than 10000000000000 is 999999999999999.
     The largest prime smaller than 100000000000000 is 99999999999997.
     The largest prime smaller than 100000000000000 is 999999999999997.
```

6.3 C++ input/output

```
#include <iostream>
#include <iomanip>
using namespace std;
int main()
    // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);</pre>
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed):
    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl;
    cout.unsetf(ios::showpoint);
    // Output a '+' before positive values
    cout.setf(ios::showpos);
cout << 100 << " " << -100 << endl;</pre>
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal
    cout << hex << 100 << " " << 1000 << " " << 10000 << dec << endl;
```

6.4 Knuth-Morris-Pratt

```
/*
Finds all occurrences of the pattern string p within the
text string t. Running time is O(n + m), where n and m
are the lengths of p and t, respecitively.

*/
#include <iostream>
#include <string>
#include <vector>
using namespace std;

typedef vector<int> VI;

void buildPi(string& p, VI& pi)
{
   pi = VI(p.length());
   int k = -2;
   for(int i = 0; i < p.length(); i++) {
     while(k >= -1 && p[k+1] != p[i])
     k = (k == -1) ? -2 : pi[k];
}
```

```
pi[i] = ++k;
int KMP (string& t, string& p)
  VI pi:
  buildPi(p, pi);
  int k = -1;
  for(int i = 0; i < t.length(); i++) {
  while(k >= -1 && p[k+1] != t[i])
      k = (k == -1) ? -2 : pi[k];
    if(k == p.length() - 1) {
     // p matches t[i-m+1, ..., i]
cout << "matched at index " << i-k << ": ";
      cout << t.substr(i-k, p.length()) << endl;</pre>
      k = (k == -1) ? -2 : pi[k];
  return 0;
int main()
  string a = "AABAACAADAABAABA", b = "AABA";
 KMP(a, b); // expected matches at: 0, 9, 12
 return 0:
```

6.5 Topological sort (C++)

```
// This function uses performs a non-recursive topological sort.
// Running time: O(|V|^2). If you use adjacency lists (vector<map<int> >),
                  the running time is reduced to O(|E|).
     INPUT: w[i][j] = 1 if i should come before j, 0 otherwise
     OUTPUT: a permutation of 0, \ldots, n-1 (stored in a vector)
               which represents an ordering of the nodes which
               is consistent with w
// If no ordering is possible, false is returned.
#include <iostream>
#include <queue>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool TopologicalSort (const VVI &w, VI &order) {
  int n = w.size();
  VI parents (n);
  queue<int> q;
  order.clear();
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++)
    if (w[j][i]) parents[i]++;</pre>
      if (parents[i] == 0) q.push (i);
  while (q.size() > 0){
   int i = q.front();
   q.pop();
    order push_back (i);
    for (int j = 0; j < n; j++) if (w[i][j]) {
     parents[j]--;
      if (parents[j] == 0) q.push (j);
  return (order.size() == n);
```

6.6 Random STL stuff

```
// Example for using stringstreams and next_permutation
#include <algorithm>
#include <iostream>
#include <sstream>
#include <vector>
using namespace std;
int main (void) (
 vector<int> v;
  v.push back(1);
 v.push back(2);
  v.push_back(3);
  v.push_back(4);
  // Expected output: 1 2 3 4
                       4 3 2 1
  do [
   ostringstream oss;
oss << v[0] << " " << v[1] << " " << v[2] << " " << v[3];
    // for input from a string s,
    // istringstream iss(s);
    // iss >> variable;
    cout << oss.str() << endl;</pre>
  } while (next_permutation (v.begin(), v.end()));
  v.push_back(1);
  v.push_back(2);
  v.push back(1);
  v.push_back(3);
  // To use unique, first sort numbers. Then call
  // unique to place all the unique elements at the beginning
  // of the vector, and then use erase to remove the duplicate
  // elements.
  sort(v.begin(), v.end());
  v.erase(unique(v.begin(), v.end()), v.end());
  // Expected output: 1 2 3
  for (size_t i = 0; i < v.size(); i++)
  cout << v[i] << " ";</pre>
  cout << endl;
```

6.7 Fast exponentiation

```
Uses powers of two to exponentiate numbers and matrices. Calculates
n^k in O(\log(k)) time when n is a number. If A is an n x n matrix,
calculates A^k in O(n^3*log(k)) time.
#include <iostream>
#include <vector>
using namespace std:
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T power (T x, int k) {
  T ret = 1;
  while(k) {
   if(k & 1) ret *= x;
   k >>= 1; x *= x;
  return ret:
VVT multiply(VVT& A, VVT& B) {
 int n = A.size(), m = A[0].size(), k = B[0].size();
```

```
VVT C(n, VT(k, 0));
  for(int i = 0; i < n; i++)</pre>
    for(int j = 0; j < k; j++)
for(int 1 = 0; 1 < m; 1++)</pre>
        C[i][j] += A[i][1] * B[1][j];
VVT power(VVT& A, int k) {
  int n = A.size();
  VVT ret(n, VT(n)), B = A;
  for(int i = 0; i < n; i++) ret[i][i]=1;</pre>
    if(k & 1) ret = multiply(ret, B);
    k >>= 1; B = multiply(B, B);
  return ret;
int main()
  /* Expected Output:
     2.37^48 = 9.72569e+17
     376 264 285 220 265
     550 376 529 285 484
     484 265 376 264 285
     285 220 265 156 264
      529 285 484 265 376 */
  double n = 2.37;
  int k = 48;
  cout << n << "^" << k << " = " << power(n, k) << endl;
  double At [5] [5] = {
    { 0, 0, 1, 0, 0 },
    { 1, 0, 0, 1, 0 },
    { 0, 0, 0, 0, 1 },
    { 1, 0, 0, 0, 0 },
    { 0, 1, 0, 0, 0 } };
  vector <vector <double> > A(5, vector <double>(5));
  for(int i = 0; i < 5; i++)
for(int j = 0; j < 5; j++)</pre>
      A[i][j] = At[i][j];
  vector <vector <double> > Ap = power(A, k);
  cout << endl:
  for(int i = 0; i < 5; i++) {</pre>
    for(int j = 0; j < 5; j++)
     cout << Ap[i][j] << " ";
    cout << endl:
```

6.8 Longest common subsequence

```
Calculates the length of the longest common subsequence of two vectors.
Backtracks to find a single subsequence or all subsequences. Runs in
\mathcal{O}\left(m\star n\right) time except for finding all longest common subsequences, which
may be slow depending on how many there are.
#include <iostream>
#include <vector>
#include <set>
#include <algorithm>
using namespace std;
typedef int T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
void backtrack(VVI& dp, VT& res, VT& A, VT& B, int i, int j)
  if(A[i-1] == B[j-1]) { res.push_back(A[i-1]); backtrack(dp, res, A, B, i-1, j-1); }
```

```
if(dp[i][j-1] >= dp[i-1][j]) backtrack(dp, res, A, B, i, j-1);
    else backtrack(dp, res, A, B, i-1, j);
void backtrackall(VVI& dp, set<VT>& res, VT& A, VT& B, int i, int j)
  if(!i || !j) { res.insert(VI()); return; }
  if(A[i-1] == B[j-1])
    set < VT > tempres;
    backtrackall(dp, tempres, A, B, i-1, j-1);
    for(set<VT>::iterator it=tempres.begin(); it!=tempres.end(); it++)
      VT temp = *it;
     temp.push_back(A[i-1]);
      res.insert(temp);
  else
    if(dp[i][j-1] >= dp[i-1][j]) backtrackall(dp, res, A, B, i, j-1);
    if(dp[i][j-1] <= dp[i-1][j]) backtrackall(dp, res, A, B, i-1, j);
VT LCS (VT& A, VT& B)
 VVI dp;
 int n = A.size(), m = B.size();
  dp.resize(n+1);
  for (int i=0; i<=n; i++) dp[i].resize(m+1, 0);</pre>
  for (int i=1; i<=n; i++)</pre>
    for(int j=1; j<=m; j++)</pre>
      if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
      else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
  VT res;
 backtrack(dp, res, A, B, n, m);
  reverse(res.begin(), res.end());
 return res;
set<VT> LCSall (VT& A, VT& B)
  int n = A.size(), m = B.size();
  dp.resize(n+1);
  for(int i=0; i<=n; i++) dp[i].resize(m+1, 0);</pre>
  for(int i=1; i<=n; i++)
    for (int j=1; j<=m; j++)</pre>
      if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
     else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
  backtrackall(dp, res, A, B, n, m);
  return res;
 int a[] = { 0, 5, 5, 2, 1, 4, 2, 3 }, b[] = { 5, 2, 4, 3, 2, 1, 2, 1, 3 };
 VI A = VI(a, a+8), B = VI(b, b+9);
 VI C = LCS(A, B);
  for(int i=0; i<C.size(); i++) cout << C[i] << " ";</pre>
  cout << endl << endl;
  set <VI> D = LCSall(A, B);
  for(set<VI>::iterator it = D.begin(); it != D.end(); it++)
    for(int i=0; i<(*it).size(); i++) cout << (*it)[i] << " ";</pre>
    cout << endl;
```

6.9 Miller-Rabin Primality Test (C)

```
// Randomized Primality Test (Miller-Rabin):
// Error rate: 2^(-TRIAL)
// Almost constant time. srand is needed
```

```
#include <stdlib.h>
#define EPS 1e-7
typedef long long LL;
LL ModularMultiplication (LL a, LL b, LL m)
        LL ret=0, c=a;
        while(b)
                if(b&1) ret=(ret+c)%m;
b>>=1; c=(c+c)%m;
        return ret;
LL ModularExponentiation(LL a, LL n, LL m)
        LL ret=1, c=a;
        while(n)
                if(n&1) ret=ModularMultiplication(ret, c, m);
                n>>=1; c=ModularMultiplication(c, c, m);
        return ret;
bool Witness(LL a, LL n)
        LL u=n-1;
  int t=0;
        while(!(u&1)){u>>=1; t++;}
```

```
LL x0=ModularExponentiation(a, u, n), x1;
    for(int i=1; i<=1;i++)
    {
        x1=ModularMultiplication(x0, x0, n);
        if(x1==1 && x0!=1 && x0!=n-1) return true;
        x0=x1;
    }
    if(x0!=1) return true;
    return false;
}
LL ret=rand(); ret*=32768;
    ret+=rand(); ret*=32768;
    ret+=rand(); ret*=32768;
    ret+=rand(); ret*=32768;
    return ret*n;
}
bool IsPrimeFast(LL n, int TRIAL)
{
    while(TRIAL--)
    {
        LL a=Random(n-2)+1;
        if(Witness(a, n)) return false;
    }
    return true;
}</pre>
```