

$$\begin{array}{c}
(4) \ ((\lambda x. \forall @x) @ \text{BOY}) \vee (y = z) \\
\frac{\frac{}{t \rightarrow (t \rightarrow t)} \quad \frac{}{(e \rightarrow t) \rightarrow t} \quad \frac{x \quad \text{BOY}}{e \rightarrow t} \quad \frac{}{e \rightarrow (e \rightarrow t)} \quad \frac{y \quad z}{e \rightarrow e}}{e \rightarrow t} @ \\
\frac{}{t} @ \\
\hline
\text{Typing fails} \times
\end{array}$$

$$\begin{array}{c}
(5) \ ((\text{IN}(\text{utrecht})) \wedge (\text{WITH}(\text{ashley}))) \\
\frac{\frac{}{t \rightarrow (t \rightarrow t)} \quad \frac{}{e \rightarrow (e \rightarrow t)} \quad \frac{\text{utrecht}}{e}}{e \rightarrow t} @ \quad \frac{\frac{}{e \rightarrow (e \rightarrow t)} \quad \frac{\text{ashley}}{e}}{e \rightarrow t} @ \\
\hline
\text{Typing fails} \times
\end{array}$$

$$\begin{array}{c}
(6) \ \lambda x. \lambda y. ((\text{WINE}(x)) \vee (y @ \text{DANCE})) \\
\text{interpreting } y \text{ as being of type } (e \rightarrow t) \rightarrow t. \\
\frac{x \quad y \quad \wedge \quad \text{WINE } x \quad y \quad \text{DANCE}}{e \quad (e \rightarrow t) \rightarrow t \quad t \rightarrow (t \rightarrow t) \quad e \rightarrow t \quad e \quad (e \rightarrow t) \rightarrow t \quad e \rightarrow t} @ \\
\frac{}{t} @ \quad \frac{}{t} @ \\
\frac{}{t \rightarrow t} @ \\
\frac{}{t} @ \\
\frac{}{((e \rightarrow t) \rightarrow t) \rightarrow t} \lambda. \\
\frac{}{e \rightarrow (((e \rightarrow t) \rightarrow t) \rightarrow t)} \lambda.
\end{array}$$

3 Recipes for semantic composition exercise 3.2

$$\begin{array}{c}
\begin{array}{ccccc}
\text{No} & \text{dog} & \text{is} & \text{in} & \text{a} & \text{room} \\
\hline
\text{NP/N} & \text{N} & (\text{S}\backslash\text{NP})/\text{PP} & \text{PP/NP} & \text{NP/N} & \text{N} \\
\hline
[[\text{No}]] & [[\text{dog}]] & [[\text{is}]] & [[\text{in}]] & [[\text{a}]] & [[\text{room}]]
\end{array} \\
\frac{}{NP} > \quad \frac{}{NP} > \\
\frac{}{[[\text{No}]] @ [[\text{dog}]]} & \frac{}{[[\text{a}]] @ [[\text{room}]]} > \\
\frac{}{PP} > \\
\frac{}{[[\text{in}]] @ ([[a]] @ [[room]])} > \\
\frac{}{S\backslash NP} > \\
\frac{}{[[is]] @ ([[in]] @ ([[a]] @ [[room]]))} < \\
\frac{}{S} < \\
(([[No]] @ [[dog]]) @ ([is] @ ([in] @ ([a] @ [room]))))
\end{array}$$

4 Computing semantics exercise 4.2

Semantic lexicon
$\llbracket \text{no} \rrbracket = \lambda PQ. \forall x (P(x) \rightarrow \neg Q(x))$
$\llbracket \text{dog} \rrbracket = \text{DOG}$
$\llbracket \text{is} \rrbracket = \lambda PQ. Q @ P$
$\llbracket \text{in} \rrbracket = \lambda Px. P @ (\lambda y. \text{IN}(y, x))$
$\text{IN}(x, y)$ is interpreted as “ x is in y ”.
$\llbracket \text{a} \rrbracket = \lambda PQ. \exists x (Q(x) \wedge P(x))$
$\llbracket \text{room} \rrbracket = \text{ROOM}$

We are inserting correct lexical terms in
 $(\llbracket \text{No} \rrbracket @ \llbracket \text{dog} \rrbracket) @ (\llbracket \text{is} \rrbracket @ (\llbracket \text{in} \rrbracket @ (\llbracket \text{a} \rrbracket @ \llbracket \text{room} \rrbracket)))$.

$$\llbracket \text{a} \rrbracket @ \llbracket \text{room} \rrbracket = \lambda PQ. \exists x (Q(x) \wedge P(x)) @ \text{room} \rightarrow_{\beta} \lambda Q. \exists x (Q(x) \wedge \text{ROOM}(x))$$

$$\llbracket \text{in} \rrbracket (\llbracket \text{a} \rrbracket @ \llbracket \text{room} \rrbracket) = \lambda Px. P @ (\lambda y. \text{IN}(y, x)) @ \lambda Q. \exists x (Q(x) \wedge \text{ROOM}(x)) \rightarrow_{\beta}$$

$$\begin{aligned} &= \lambda x. (\lambda Q. (\exists x (Q(x) \wedge \text{ROOM}(x))) @ (\lambda y. \text{IN}(x, y))) \rightarrow_{\alpha} \\ &= \lambda x. (\lambda Q. (\exists z (Q(z) \wedge \text{ROOM}(z))) @ (\lambda y. \text{IN}(x, y))) \rightarrow_{\beta} \\ &= \lambda x. (\exists z ((\lambda y. \text{IN}(x, y))(z) \wedge \text{ROOM}(z))) \rightarrow_{\beta} \end{aligned}$$

$$\begin{aligned} \llbracket \text{is} \rrbracket (\llbracket \text{in} \rrbracket (\llbracket \text{a} \rrbracket @ \llbracket \text{room} \rrbracket)) &= (\lambda PQ. Q @ P) @ (\lambda x. \exists z (\text{IN}(x, z) \wedge \text{ROOM}(z))) \rightarrow_{\beta} \\ &= (\lambda Q. Q @ (\lambda x. \exists z (\text{IN}(x, z) \wedge \text{ROOM}(z)))) \rightarrow_{\beta} \end{aligned}$$

$$\begin{aligned} \llbracket \text{No} \rrbracket @ \llbracket \text{dog} \rrbracket &= \lambda PQ. \forall x (P(x) \rightarrow \neg Q(x)) @ \text{DOG} \rightarrow_{\beta} \\ \llbracket \text{No} \rrbracket @ \llbracket \text{dog} \rrbracket &= \lambda Q. \forall x (\text{DOG}(x) \rightarrow \neg Q(x)) \end{aligned}$$

$$\begin{aligned} &(\llbracket \text{is} \rrbracket @ (\llbracket \text{in} \rrbracket @ (\llbracket \text{a} \rrbracket @ \llbracket \text{room} \rrbracket))) @ (\llbracket \text{No} \rrbracket @ \llbracket \text{dog} \rrbracket) = \\ &(\lambda Q. Q @ (\lambda x. \exists z (\text{IN}(x, z) \wedge \text{ROOM}(z)))) @ \lambda Q. \forall x (\text{DOG}(x) \rightarrow \neg Q(x)) \rightarrow_{\beta} \\ &= \lambda Q. \forall x (\text{DOG}(x) \rightarrow \neg Q(x)) @ (\lambda x. \exists z (\text{IN}(x, z) \wedge \text{ROOM}(z))) \rightarrow_{\beta} \\ &= \forall x (\text{DOG}(x) \rightarrow \neg (\lambda x. \exists z (\text{IN}(x, z) \wedge \text{ROOM}(z))) @ x) \rightarrow_{\beta} \\ &= \forall x (\text{DOG}(x) \rightarrow \neg \exists z (\text{IN}(x, z) \wedge \text{ROOM}(z))) \end{aligned}$$

5 Natural tableau reasoning

We tried to find a conclusive tableau solution but we were unable to. Our hypothesis is that it is neutral.