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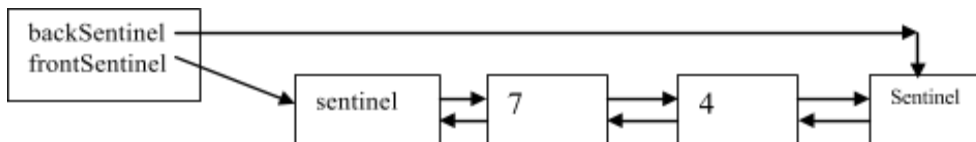
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## Worksheet 19: Linked List Deque

**In Preparation:** Read Chapter 7 to learn more about the Deque data type. If you have not done so already, you should complete worksheets 17 and 18 to learn about the basic features of a linked list.

In this lesson we continue looking at variations on the theme of linked lists, this time including double links and sentinels on both the front and the back of the list. This will be the first of several lessons that will develop a very general purposed linked list abstraction. In this lesson we will emphasize the deque aspects of the list. In the following lesson we will add more operations.



### Conceptual Deque operations

```
void addFront (TYPE e)
```

```
void addBack (TYPE e)
```

```
TYPE front ()
```

```
TYPE back () void
```

```
removeFront ()
```

```
void removeBack ()
```

```
int isEmpty ()
```

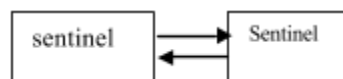
A deque, you recall, allows insertions at both the beginning and the end of the container. The conceptual interface is shown at right. (Recall that the actual functions may differ from the conceptual interface in two ways. First, the names are likely differ to accommodate the C restriction that all functions have unique names. Second, the actual functions will pass the underlying data area as an argument).

Our linkedList data structure will have two

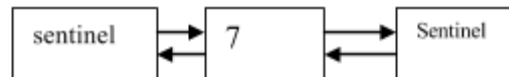
major variations from the linked list stack and queue you have examined earlier. First, an integer data field will remember the number of elements (that is, the size) in the container. Second, we will use sentinels at both the front and back of the linked list.

Sentinels ensure the list is never completely empty. They also mean that all instructions can be described as special cases of more general routines. The internal method `_addBefore` will add a link before the given location. The second routine `_removeLink`, will remove its argument link. Both of these use the underscore convention to indicate they are internal methods. When a value is removed from a list make sure you free the associated link fields. Use an assertion to check that allocations of new links are successful, and that you do not attempt to remove a value from an empty list. The following instructions will help you understand how to implement these operations.

1. Draw a picture of an initially empty LinkedList, including the two sentinels.



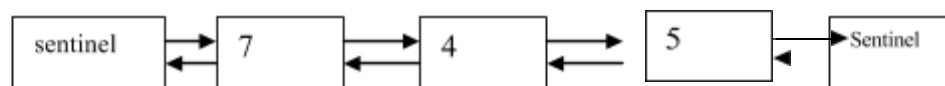
2. Draw a picture of the LinkedList after the insertion of one value.



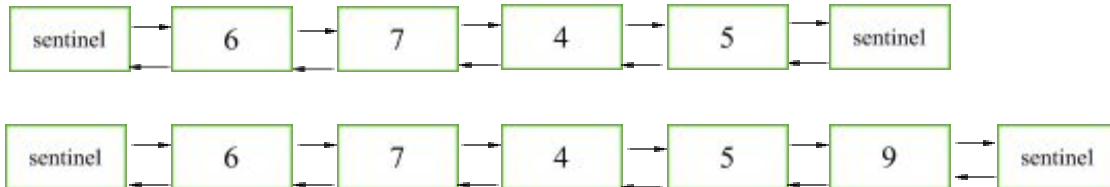
3. Based on the previous two pictures, can you describe what property characterizes an empty collection?

If the sentinels link to each other, the list is empty.

4. Draw a picture of a LinkedList with three or more values (in addition to the sentinels).



5. Draw a picture after a value has been inserted into the front of the collection. Notice that this is between the front sentinel and the following element. Draw a picture showing an insertion into the back. Notice that this is again between the last element and the ending sentinel. Abstracting from these pictures, what would the function `addBefore` need to do, where the argument is the link that will follow the location in which the value is inserted.



`addBefore` would need to create a link, assign the pointer to the next link of the last list element to the new link, and assign the sentinel's pointer to the previous link to the new link.

6. Draw a picture of a `linkedList` with three or more values, then examine what is needed to remove both the first and the last element. Can you see how both of these operations can be implemented as a call on a common operation, called `_removeLink`?

See picture in question 4. Both operations would require changing the previous and next pointers of the links adjacent to the removed link. Instead of pointing to the removed link, they will point to each other.

7. What is the algorithmic complexity of each of the deque operations?

The algorithmic complexity of each operation is  $O(1)$ .

```
struct dlink {
    TYPE value;
    struct dlink * next;
    struct dlink * prev;
};

struct linkedList {
    int size;
    struct dlink * frontSentinel;
```

```

    struct dlink * backSentinel;
};

/* these functions are written for you */
void LinkedListInit (struct linkedList *q) {
    q->frontSentinel = malloc(sizeof(struct dlink));
    assert(q->frontSentinel != 0);
    q->backSentinel = malloc(sizeof(struct dlink));
    assert(q->backSentinel);
    q->frontSentinel->next = q->backSentinel;
    q->backSentinel->prev = q->frontSentinel;
    q->size = 0;
}

void linkedListFree (struct linkedList *q) {
    while (q->size > 0)
        linkedListRemoveFront(q);
    free (q->frontSentinel);
    free (q->backSentinel);
    q->frontSentinel = q->backSentinel = null;
}

void LinkedListAddFront (struct linkedList *q, TYPE e)
    { _addBefore(q, q->frontSentinel->next, e); }

void LinkedListAddback (struct linkedList *q, TYPE e)
    { _addBefore(q, q->backSentinel, e); }

void linkedListRemoveFront (struct linkedList *q) {
    assert(! linkedListIsEmpty(q));
    _removeLink (q, q->frontSentinel->next);
}

void LinkedListRemoveBack (struct linkedList *q) {
    assert(! linkedListIsEmpty(q));
    _removeLink (q, q->backSentinel->prev);
}

int LinkedListIsEmpty (struct linkedList *q) {
    return q->size == 0;
}

/* write addLink and removeLink. Make sure they update the size field correctly */
/* _addBefore places a new link BEFORE the provide link, lnk */
void _addBefore (struct linkedList *q, struct dlink *lnk, TYPE e) {
    struct dlink *newLink = malloc(sizeof(struct dlink));

```

```
    newLink->value = e;
    newLink->prev = lnk->prev;
    newLink->next = lnk;

    lnk->prev->next = newLink;
    lnk->prev = newLink;
    q->size++;
}

void _removeLink (struct linkedList *q, struct dlink *lnk) {
    lnk->prev->next = lnk->next;
    lnk->next->prev = lnk->prev;

    free(lnk);
    lnk = 0;
    q->size--;
}

TYPE LinkedListFront (struct linkedList *q) {
    assert(LinkedListIsEmpty(q) == 0);

    return q->frontSentinel->next->value;
}

TYPE LinkedListBack (struct linkedList *q) {
    assert(LinkedListIsEmpty(q) == 0);

    return q->backSentinel->prev->value;
}
```