

# Packing Boxes with Score Constraints

Asyl L. Hawa

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## 1 Example

### 1.1 Creating problem instance

In this example, we will attempt to find a feasible alignment of nine boxes, with score widths ranging between 1mm and 70mm. The minimum score separation constraint, or “threshold”, will be set at 70mm, the industry standard.

Initially, we create 18 random values between 1 and 70, which will be our score widths. We also need to add an extra box with two score of widths 70mm, which will be our “dominating scores”, giving us a total of  $n = 10$  boxes and 20 scores. These dominating scores will eventually be discarded.

We then sort the scores in non-decreasing order, and place them in a vector. The scores will now be addressed by their indices throughout the rest of the algorithm (see Table 1).

Index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
Scores	3	4	12	22	25	35	36	37	38	45	49	54	54	55	55	64	65	66	70	70

Table 1: allScores vector.

Since every box consists of two scores, each score is randomly assigned a “mate”, so that one score (the smaller score) represents the left-hand side of the box, and the other score represents the right-hand side of the box. The only exceptions to the random allocation of mates are the dominating scores - they must be assigned to one another. We can then assign measurements to each pair of mates to represent the width of each individual box.

Mates		Box Width
0	6	314
1	10	372
2	16	297
3	11	220
4	17	959
5	9	738
7	12	622
8	15	635
13	14	859
18	19	0

Table 2: Mates.

Next, we create an adjacency matrix. This adjacency matrix contains information regarding threshold constraint. If the sum of two scores is greater than or equal to the threshold, the adjacency matrix contains “1” in the relevant cell. If two scores are mates, the adjacency matrix contains “2”. Else, the adjacency matrix contains “0”.

## 1.2 Matching scores

Once we have created the problem instance, we then have to attempt to match the scores together. If two scores can be matched, it means that the boxes can be placed next to each other, and the total width of the two scores from the boxes is equal to or exceeds the threshold, and can therefore be scored using the knives in the scoring machine.

The matching is performed as follows: starting from the smallest score (0), we attempt to match it with the largest possible score it is adjacent to ( $= 1$  in the adjacency matrix), and that has not yet been matched (in this case, 19). We continue until all scores have been matched. If a score cannot be matched due to the largest score being its mate, we match the score with the second largest score available. As we can see in Table 3, score 7 was unable to be matched with the largest possible score, 12, as they are mates. Instead, 7 was matched with the next largest score, 11.

Once the matching algorithm is complete, we then examine the size of the matching list. For the instance to be feasible, all of the scores must be matched. If not all of the scores are matched, i.e. there are less than  $n$  matchings in the matching list, then there are not enough matching edges to create a path, and the instance is infeasible.

Match	
0	19
1	18
2	17
3	16
4	15
5	14
6	13
7	11
8	12
9	10

Table 3: Match List.

## 1.3 Mate-Induced Structure

Cycle 0	0	6	13	14	15	9	10	1	18	19
Cycle 1	2	16	3	11	7	12	8	15	4	17

Table 4: Mate-Induced Cycles.

	Edges	
cycle 1	1	2
cycle 2	4	5
cycle 3	6	8

Table 5: T-Matrix.

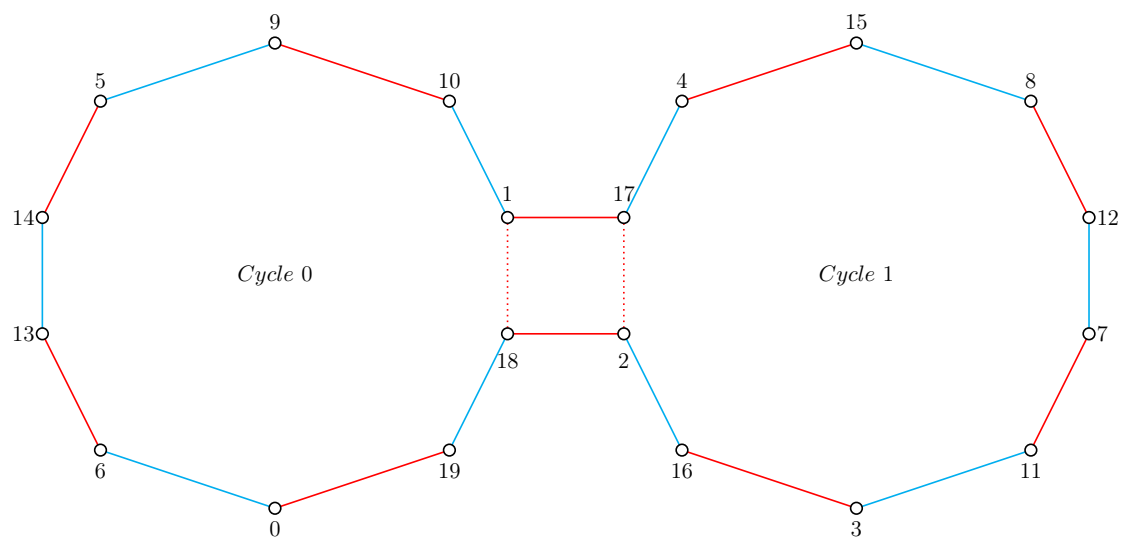


Figure 1: Patching MIS blue = mates, red = matching, explain dotted lines.