

# Data structures and Algorithms

## Lecture 2b: stacks and queues [GT 2.1]

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## Stacks and queues

These ADTs are restricted forms of List, where insertion and removal happen only in particular locations:

- stacks follow last-in-first-out (LIFO)
- queues follows first-in-first-out (FIFO)

So why should we care are a less general ADT?

- operations names are part of computing culture
- numerous applications
- simpler/more efficient implementations than Lists

# Stack ADT



Main stack operations:

- **push**(e): inserts an element, e
- **pop**(): removes and returns the last inserted element

Auxiliary stack operations:

- **top**(): returns the last inserted element without removing it
- **size**(): returns the number of elements stored
- **isEmpty**(): indicates whether no elements are stored

## Stack Example

operation	returns	stack
push(5)	-	[5]
push(3)	-	[5, 3]
size()	2	[5, 3]
pop()	3	[5]
isEmpty()	False	[5]
pop()	5	[]
isEmpty()	True	[]
push(7)	-	[7]
push(9)	-	[7, 9]
top()	9	[7, 9]
push(4)	-	[7, 9, 4]
pop()	4	[7, 9]

# Stack Applications

## Direct applications

- Keep track of a history that allows undoing such as Web browser history or undo sequence in a text editor
- Chain of method calls in a language supporting recursion
- Context-free grammars

## Indirect applications

- Auxiliary data structure for algorithms
- Component of other data structures

# Method Stacks

The runtime environment keeps track of the chain of active methods with a stack, thus allowing **recursion**

When a method is called, the system pushes on the stack a frame containing

- Local variables and return value
- Program counter

When a method ends, we pop its frame and pass control to the method on top

```
main() {  
    int i = 5;  
    foo(i);  
}  
  
foo(int j) {  
    int k;  
    k = j+1;  
    bar(k);  
}  
  
bar(int m) {  
    ...  
}
```

bar  
PC = 1  
m = 6

foo  
PC = 3  
j = 5  
k = 6

main  
PC = 2  
i = 5

---

## Parentheses Matching

Each “(”, “{”, or “[” must be paired with a matching “)”, “}”, or “]”

- correct: ( )(( )){([ ( ))}
- correct: ((( )(( )){([ ( ))}
- incorrect: )(( )){([ ( ))}
- incorrect: ({ [ ]})
- incorrect: (

Scan input string from left to right:

- If we see an opening character, push it to a stack
- If we see a closing character, pop character on stack and check that they match

## Stack implementation based on arrays

- A simple way of implementing the Stack ADT uses an array
- We add elements from left to right
- A variable keeps track of the index of the top element

Algorithm *size()*

return  $t + 1$

Algorithm *pop()*

if *isEmpty()* then

return null

else

$t \leftarrow t - 1$

return  $S[t + 1]$





## Stack implementation based on arrays

- The array storing the stack elements may become full.
- A push operation will then either grow the array or signal an error.

**Algorithm** *push(o)*  
if  $t = S.length - 1$  then  
    signal *stack overflow error*  
else  
     $t \leftarrow t + 1$   
     $S[t] \leftarrow o$



## Stack implementation based on arrays

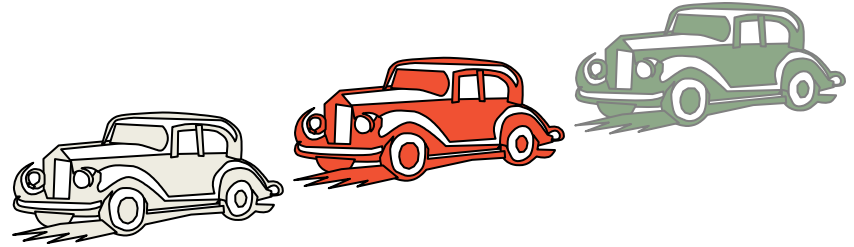
### Performance

- Let  $n$  be the number of elements in the stack
- The space used is  $O(n)$
- Each operation runs in time  $O(1)$

### Qualifications

- Trying to push a new element into a full stack causes an implementation-specific exception or
- Pushing an item on a full stack causes the underlying array to double in size, which implies each operation runs in  $O(1)$  amortized time.

# Queue ADT



Main queue operations:

- **enqueue(e)**: inserts an element, e, at the end of the queue
- **dequeue()**: removes and returns element at the front of the queue

Auxiliary queue operations:

- **first()**: returns the element at the front without removing it
- **size()**: returns the number of elements stored
- **isEmpty()**: indicates whether no elements are stored

Boundary cases:

- Attempting the execution of dequeue or first on an empty queue signals an error or returns null

## Queue Example

<b>Operation</b>		<b>Output</b>	<b>Q</b>
enqueue(5)	—	(5)	
enqueue(3)	—	(5, 3)	
dequeue()	5	(3)	
enqueue(7)	—	(3, 7)	
dequeue()	3	(7)	
first()	7	(7)	
dequeue()	7	()	
dequeue()	<i>null</i>	()	
isEmpty()	<i>true</i>	()	
enqueue(9)	—	(9)	
enqueue(7)	—	(9, 7)	
size()	2	(9, 7)	
enqueue(3)	—	(9, 7, 3)	
enqueue(5)	—	(9, 7, 3, 5)	
dequeue()	9	(7, 3, 5)	

# Queue applications

Buffering packets in streams, e.g., video or audio

## Direct applications

- Waiting lists, bureaucracy
- Access to shared resources (e.g., printer)
- Multiprogramming

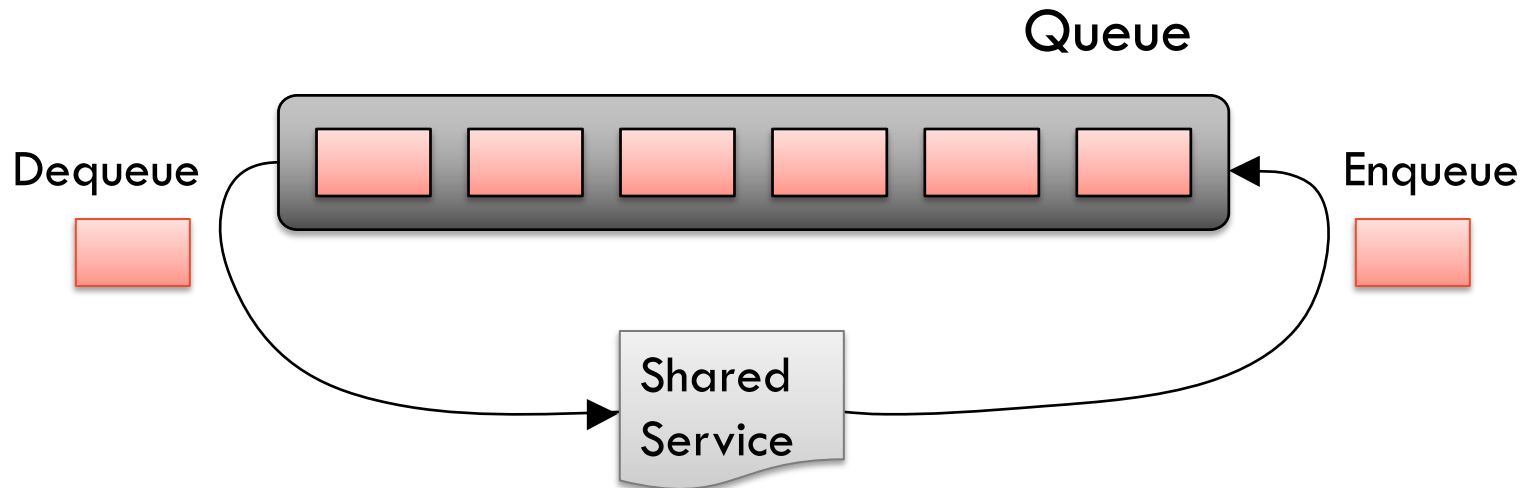
## Indirect applications

- Auxiliary data structure for algorithms
- Component of other data structures

## Queue application: Round Robin Schedulers

Implement a round robin scheduler using a queue  $Q$  by repeatedly performing the following steps:

1.  $e = Q.dequeue()$
2. Service element  $e$
3.  $Q.enqueue(e)$



## Queue implementation based on arrays

Use an array of size  $N$  in a circular fashion

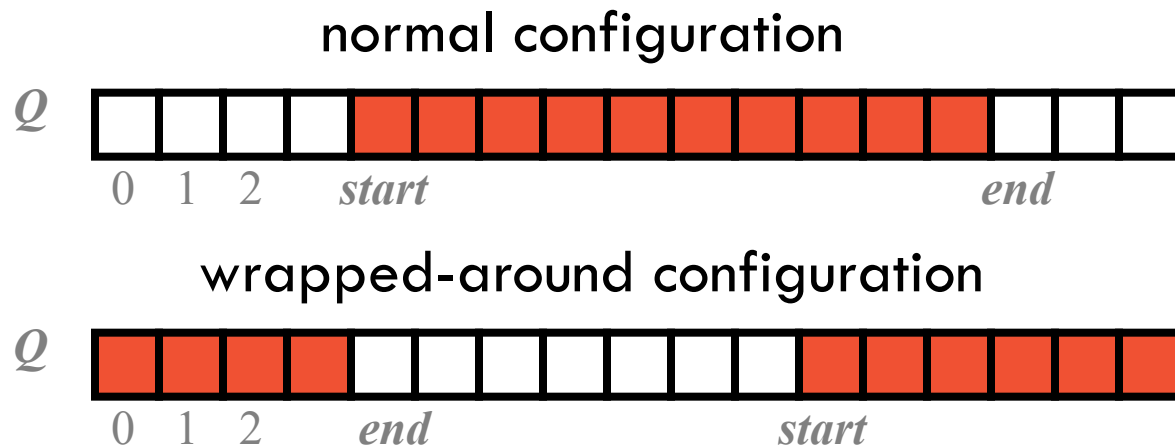
Two variables keep track of the front and size

**start** : index of the front element

**end** : index past the last element

**size** : number of stored elements

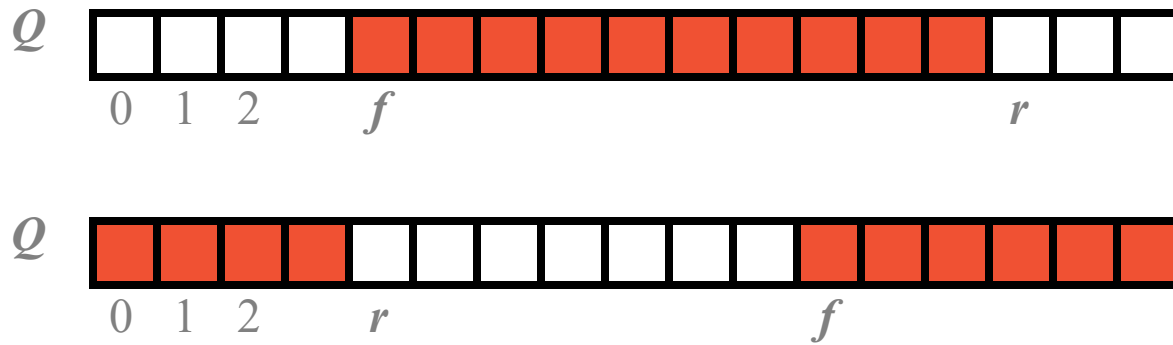
These are related as follows  $\text{end} = (\text{start} + \text{size}) \bmod N$ ,  
so we only need two



## Queue Operations: Enqueue

Return an error if the array is full. Alternatively, we could grow the underlying array as dynamic arrays do

```
def enqueue(o)
  if size() =  $N - 1$  then
    return "queue full"
  else
    last  $\leftarrow$  (first + size) mod  $N$ 
     $Q[r]$   $\leftarrow$  o
    size  $\leftarrow$  (size + 1)
```



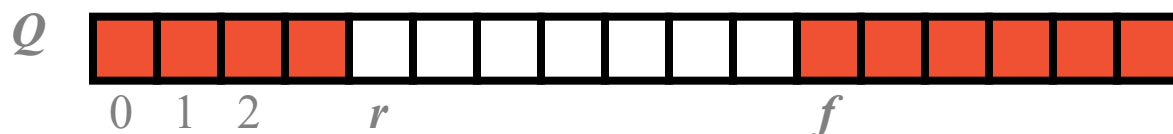


## Queue Operations: Dequeue

Note that operation dequeue returns error if the queue is empty

One could alternatively signal an error

```
def dequeue()  
    if isEmpty() then  
        return "queue empty"  
    else  
        o ← Q[f]  
        f ← (f + 1) mod N  
        s ← (s - 1)  
        return o
```



## Double-ended queues: Deques

- A linear structure that allows insertions and deletions at both ends

Method	Time
size, isEmpty	$O(1)$
getFirst, getLast	$O(1)$
addFirst, addLast	$O(1)$
removeFirst, removeLast	$O(1)$

**Table 5.4:** Performance of a deque realized by a doubly linked list.

## Double-ended queue operations

The deque abstract data type is richer than both the stack and the queue ADTs. The fundamental methods of the deque ADT are as follows:

- `addFirst( $e$ )`: Insert a new element  $e$  at the head of the deque.
- `addLast( $e$ )`: Insert a new element  $e$  at the tail of the deque.
- `removeFirst()`: Remove and return the first element of the deque; an error occurs if the deque is empty.
- `removeLast()`: Remove and return the last element of the deque; an error occurs if the deque is empty.

Additionally, the deque ADT may also include the following support methods:

- `getFirst()`: Return the first element of the deque; an error occurs if the deque is empty.
- `getLast()`: Return the last element of the deque; an error occurs if the deque is empty.
- `size()`: Return the number of elements of the deque.
- `isEmpty()`: Determine if the deque is empty.