

Computer Sciences Department
University of Wisconsin-Madison
CS/ECE 552 – Introduction to Computer Architecture
In-Class Exercise (04/23)

Answers to all questions should be uploaded on Canvas.

1. [2 points] From the textbook:

- a. [1 point] (6.10) So far in class we have discussed the 3C's for classifying cache misses (cold/compulsory, capacity, conflict). However, in multiprocessor systems there is a 4th 'C' – coherence misses. Name at least one scenario where a coherence miss occurs.

Different/multiple cores try to access the same data. They each keep a copy in its own cache but that copy does not match with each other.

- b. [1 point] (6.10) What mechanism determines when a write will be seen by read from a different core?
Invalidation-based protocols

2. [8 points] Consider a program with four threads running in parallel on a quad-core multiprocessor. Assume that all memory contents are **initially 0**:

```
# thread A
addi $t0, $zero, 4096
addi $t1, $zero, 1
sw    $t1, 0($t0)
```

```
# thread B
addi $t0, $zero, 4096
addi $t2, $zero, 1
sw    $t2, 4($t0)
```

```
# thread C
addi $t0, $zero, 4096
lw    $s0, 0($t0)
lw    $s1, 4($t0)
```

```
# thread D
addi $t0, $zero, 4096
```

```
lw    $s2, 4($t0)
lw    $s3, 0($t0)
```

Fill in the table below, specifying whether or not each program outcome (i.e., set of register contents after executing all threads) is **possible** if the multiprocessor enforces **sequential consistency**.

Register Contents After Execution		Possible With Sequential Consistency (Y/N)?
Thread C	Thread D	
[\$s0] = 0 [\$s1] = 0	[\$s2] = 0 [\$s3] = 0	Y
[\$s0] = 0 [\$s1] = 0	[\$s2] = 0 [\$s3] = 1	Y
[\$s0] = 0 [\$s1] = 1	[\$s2] = 0 [\$s3] = 1	Y
[\$s0] = 1 [\$s1] = 0	[\$s2] = 0 [\$s3] = 0	Y
[\$s0] = 1 [\$s1] = 0	[\$s2] = 1 [\$s3] = 0	N
[\$s0] = 1 [\$s1] = 1	[\$s2] = 0 [\$s3] = 0	Y
[\$s0] = 1 [\$s1] = 1	[\$s2] = 0 [\$s3] = 1	Y
[\$s0] = 1 [\$s1] = 1	[\$s2] = 1 [\$s3] = 0	Y
[\$s0] = 1 [\$s1] = 1	[\$s2] = 1 [\$s3] = 1	Y