Remote Memory Operations

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Materials are taken from the book, Using MPI-2: Advanced Features of the Message-Passing Interface by William Gropp, Ewing Lusk, and Rajeev Thakur.

Topic Overview

- Introduction to Remote Memory Operations
- A Non-Remote Memory Operation code
- Remote Memory Operations

Introduction to Remote Memory Operations

Send-Recv:

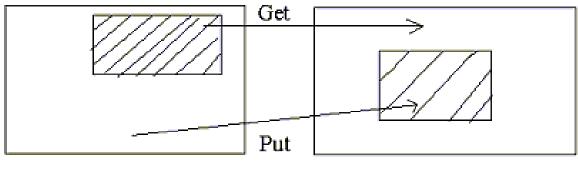
Specify what memory locations can be modified. Specify when this can happen.

Advantages:

- 1. Ensuring the correctness of the data.
- 2. Changes in memory of the receiver can happen only when the receiver allows.

• Disadvantages:

- 1. Expressiveness and low performance.
- 2. What if data that is updated frequently.
- 3. Communication costs might be big.



Address space of Process 0

Address space of Process 1

- One single routine that specifies both where the data is coming from and where it is going to.
- They are called one-sided communication routines.
- Three main steps when using the RMA.
 - 1. Define the memory of the process that can be used for RMA operations.

```
MPI_Win_create
```

2. Specify the data to be moved and where to move it.

```
MPI_Put, MPI_Get, MPI_Accumulate
```

3. Specify how we know that the data is available.

```
MPI_Win_fence
```

A Non-RMA code

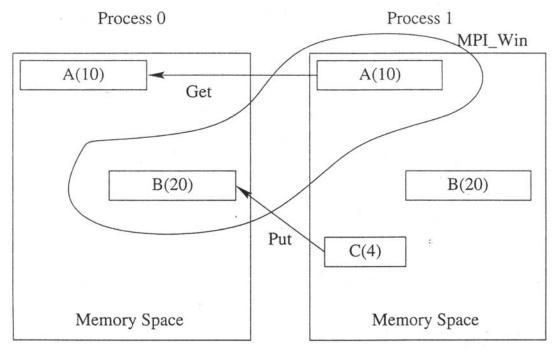
$$\pi = \int_0^1 \frac{4}{1+x^2} \, dx.$$

• Compute the Riemann sum.

- *base:Initial address of window.
- size: Size of window in bytes.
- disp_unit: Local unit size for displacements, in bytes.
 Unit size in different processors maybe different.
- info: Info argument.
- *win: Window object returned by the call.
- Need to be called by all processors in comm which contribute to this window.
- A "window" refers to a region of memory within a single process.
- A "window object" describes the collection of windows that are the input to the MPI_Win_create call.

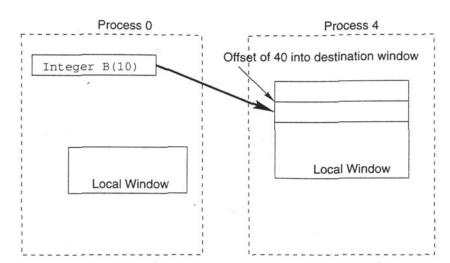
- nwin consists one integer from cpu0, none from the rest.
- piwin consists one double from cpu0, none from the rest.

Remote Memory Access (RMA)-MPI_Win_create



Remote Memory Access (RMA)-MPI_Put

- *origin_addr: Initial address of origin buffer.
- origin_count: Number of entries in origin buffer.
- target_rank: Rank of target.
- target_disp: Displacement from window start to the beginning of the target buffer.
- target_count: Number of entries in target buffer.



Remote Memory Access (RMA)-MPI_Get

- *origin_addr: Address of the buffer in which to receive the data.
- origin_count: Number of entries in origin buffer.
- target_rank: Rank of target.
- target_disp: Displacement from window start to the beginning of the target buffer.
- target_count: Number of entries in target buffer.

Remote Memory Access (RMA)-MPI_Win_fence

- What if two processors try to access the same data at the same time?
- Answer: Use MPI_Win_fence.

```
int MPI_Win_fence(int assert, MPI_Win win)
```

- assert: Program assertion. Provide information about the fence that can be used by some MPI implementation to provide better performance.
- win: Window object.
- MPI_Win_fence completes any RMA operations that started since the last call to MPI_Win_fence.
- MPI_Win_fence is collective over all processes in the group associated with the window object.

- The easy rules:
 - 1. Do not overlap accesses on window.
 - 2. Separate non-RMA access from RMA access with MPI_Win_fence.
- It is possible to have several MPI window objects whose local windows overlap.
- BUT, because of the performance issues and the complexities of the rules for correct use, we recommend avoiding the use of overlapping window.
- MPI_Accumulate.

```
if (myid == 0) {
 printf("Enter the number of intervals: (0 quits) ");
 printf("\n");
 scanf("%d",&n);
 pi = 0.0;
MPI Win fence(0, nwin);
if (myid != 0)
   MPI_Get(&n, 1, MPI_INT, 0, 0, 1, MPI_INT, nwin);
MPI_Win_fence(0, nwin);
if (n == 0)
   break;
else{
   h = 1.0 / (double) n;
    sum = 0.0;
   for (i = myid + 1; i \le n; i += numprocs) {
         x = h * ((double)i - 0.5);
         sum += (4.0 / (1.0 + x*x));
    mypi = h * sum;
   MPI_Win_fence(0, piwin);
    MPI_Accumulate(&mypi, 1, MPI_DOUBLE, 0, 0, 1, MPI_DOUBLE,
                    MPI_SUM, piwin);
   MPI_Win_fence(0, piwin);
 }
```

Remote Memory Access (RMA)-Example

```
MPI-Win_create(A, ..., &win);
MPI Win fence (0, win);
if (rank == 0)
* Process 0 puts data into many local windows */
   MPI Put( ... , win );
/* This fence completes the MPI_Put operations initiated
   by process 0 */
MPI_Win_fence( 0, win );
/* All processes initiate access to some window to extract data */
MPI_Get( ... , win );
/* The following fence completes the MPI_Get operations */
MPI Win fence (0, win);
/* After the fence, processes can load and store
into A, the local window */
A[rank] = 4;
printf( "A[%d] = %d\n", 0, A[0] );
MPI_Win_fence(0, win);
/* We need a fence between stores and RMA operations */
MPI_Put( ..., win );
/* The following fence completes the preceding Put */
MPI_Win_fence( 0, win );
```

- Homework: Use RMA to compute variance of n numbers.
- Homework: Use RMA to implement the Jacobi method for solving linear system.