

Practical C++ Decompilation

Igor Skochinsky Hex-Rays

Recon 2011 Montreal

Outline

- Class layouts
- Virtual tables
- Methods, constructors and destructors
- RTTI and alternatives
- Dealing with C++ in IDA and Hex-Rays decompiler

Class layouts

- Class fields are generally placed in memory in the order of declaration
- Equivalent structure can be produced by removing all methods
- Example:

Single inheritance

- With simple inheritance, fields of the derived class are placed after the base class
- Example:

Multiple inheritance

- With multiple inheritance, first the base classes are laid out, then the fields of the derived class
- Example:

```
class C: public A,
                    00000000 C
                                  struc
public B
                                  dd?
                    00000000 a1
                                  dd?
                    00000004 a2
 int c4;
                    00000008 a1
                                  dd?
                                  dd?
}
                    0000000C a2
                                  dd?
                    0000010 b3
                                  dd?
                    0000014 c4
                                  ends
                    0000018 C
```

Virtual inheritance

- In case of virtual inheritance, the place of a virtual base class is not fixed and can change in future derived classes
- The compiler has to track offset of a virtual base in each specific class inheriting from it
- MSVC implements it by producing a virtual base table (vbtable) with offsets to each of the virtual bases
- GCC puts offsets to virtual bases into the virtual function table (vftable)

Virtual inheritance example

```
class B : public virtual A {
                              00000000 B struc
                              00000000 _vbptr
public:
                                             dd?
                              00000004 b3
 int b3;
                                             dd?
};
                              00000008 a1 dd?
                              0000000C a2 dd?
                              00000010 B ends
class C : public virtual A, public
                              00000000 C
                                             struc
                              00000000 _vbptr
B {
                                             dd?
public:
                              00000004 b3
                                             dd?
                              00000008 c4 dd?
 int c4;
};
                              0000000C a1 dd?
                              00000010 a2 dd?
                              00000014 C ends
```

Virtual inheritance example

MSVC	GCC
<pre>const B::`vbtable': dd 0 dd 8</pre>	<pre>`vtable for'B: dd 8 dd 0 dd offset `typeinfo for'B dd 0</pre>
const C::`vbtable': dd 0 dd 0Ch	`vtable for'C: dd 0Ch dd 0 dd offset `typeinfo for'C dd 0

Virtual tables

- If class has virtual methods, compiler creates a table of pointers to those methods
- The pointer to the table is placed into a hidden field
- Methods are usually arranged in the order of declaration
- When inheriting, overridden methods are replaced and new ones are added at the end
- If inherited class does not override or add new virtual methods, the table can be reused
- MSVC uses separate tables for virtual functions and virtual bases, GCC combines them

Method calls

- Standard method calls take a hidden this parameter
- In MSCV x86, ecx is traditionally used (__thiscall)
- In other compilers, it's usually inserted as the first parameter

```
mov ecx, [edi+0Ch]
call CUser::IsCachedLogon(void)
```

```
cached = m_pUser->IsCachedLogon()
```

Static method calls

- Static methods do not need a class instance
- Because of that, they behave as standard functions
- Can't be easily distinguished from standalone functions

```
push eax ; nLengthNeeded
  push edi ; hObj
  call CSession::GetDesktopName(HDESK___
*,ushort * *)
```

CSession::GetDesktopName(hDesk, &name)

Virtual method calls

- Virtual methods can be overridden in derived classes
- The call address has to be calculated dynamically
- The address is loaded from the virtual table
- Virtual methods also expect this pointer passed

```
mov eax, [ebp+pUnk]
mov ecx, [eax]
push eax
call dword ptr [ecx+8]

pUnk->Release();
```

Constructors

- First code executed during an object's lifetime
- Usually performs the following actions:
 - a) call constructors of base classes
 - b) initialize vfptrs if has virtual functions
 - c) call constructors of complex members
 - d) run initialization list actions
 - e) execute the constructor body
- In optimized code, some steps can be shuffled and some calls inlined
- Calling convention same as normal methods (hidden this pointer)
- MSVC returns the this pointer from constructors

Constructor example

```
public: thiscall CMachine::CMachine(void) proc near
 mov edi, edi
 push esi
 mov esi, ecx
 call CDataStoreObject::CDataStoreObject(void)
 and dword ptr [esi+24h], 0
 or dword ptr [esi+28h], 0FFFFFFFh
 and dword ptr [esi+2Ch], 0
 or dword ptr [esi+30h], 0FFFFFFFh
 mov dword ptr [esi], offset const CMachine::`vftable'
 mov eax, esi
 pop esi
 retn
public: thiscall CMachine::CMachine(void) endp
```

Constructor calls I

- Objects can be constructed in different ways
- Local (automatic) variables usually get memory allocated on stack

```
push [ebp+arg_0]
 lea ecx, [ebp+var_7A]
 call CFolderIdString::CFolderIdString(_GUID
const &)
```

CFolderIdString folderString(guid);

Constructor calls II

- Objects can be constructed with the new operator
- The compiler calls generic or class-specific operator new
- Allocated memory is passed to the constructor

```
34h
   push
  call
          operator new(uint)
   pop
          ecx
  test eax, eax
  jz
                                     CMachine *machine =
          short @@nomem
                                          new CMachine();
  mov
          ecx, eax
  call
          CMachine::CMachine(void)
   jmp
          short @@ok
@@nomem:
  xor
          eax, eax
@@ok:
```

Constructor calls III

- Global objects are constructed at the program start
- Usual implementation uses a table of compiler-generated functions that call constructor with memory reservered in the data section
- MSCV adds destructors using atexit() calls
- GCC uses a separate table of destructors
- The table is handled in the runtime start-up code

Global objects: MSVC I

```
push offset ___xc_z
push offset ___xc_a
call initterm
  _xc_a dd 0
  dd offset sub_42FAB74B
  dd offset sub_42FAB765
  dd offset sub_42FAB7E1
  _{\mathsf{xc}}\mathsf{z} dd 0
```

Global objects: MSVC II

```
void cdecl `dynamic initializer for
'g PrivateProfileCache''(void) proc near
   mov ecx, offset g_PrivateProfileCache
   call CPrivateProfileCache::CPrivateProfileCache(void)
    push offset `dynamic atexit destuctor for
'g PrivateProfileCache''(void); void ( cdecl *)()
   call _atexit
   pop ecx
   retn
void cdecl `dynamic initializer for
'g PrivateProfileCache''(void) endp
```

Global objects: MSVC III

```
void __cdecl `dynamic atexit destuctor for
'g_PrivateProfileCache''(void) proc near
   mov ecx, offset g_PrivateProfileCache
   jmp CPrivateProfileCache::~CPrivateProfileCache(void)
void __cdecl `dynamic atexit destuctor for
'g_PrivateProfileCache''(void) endp
```

Global objects: GCC

- .ctors section contains pointers to "global constructor" functions
- .dtors contains pointers to "global destructor" functions
- Sometimes two global arrays
 (__CTOR_LIST__/_DTOR_LIST__) are used instead
- Both types call a common "initialization_and_destruction" function which either constructs or destructs globals for a module

```
void __static_initialization_and_destruction_0(int
_initialize_p, int _priority)
```

Constructor calls IV

- Static objects are initialized on first use
- Common way is to use a guard variable

```
mov eax, ds:_guard_aa
and eax, 1 static A aa;
jnz short @@skip
mov ecx, ds:_guard_aa
or ecx, 1
mov ds:_guard, ecx
mov ecx, offset aa
call A::A(void)
```

 GCC uses ABI-specified helper functions cxa guard acquire/ cxa guard release.

Array construction

- Each element of the array has to be constucted separately
- If any of the constructors throws an exception, all previous elements must be destructed
- MSVC uses a helper function, "vector constructor iterator"
- very useful because in one place we get instance size, constructor and (in case of EH iterator) destructor

```
push offset ATL::CComTypeInfoHolder::stringdispid::stringdispid(void)
push esi
push 0Ch
push eax
call `vector constructor iterator'(void *,uint,int,void * (*)(void *))
strings = new ATL::CComTypeInfoHolder::stringdispid[count];
```

Destructors

- Unlike constructors, a class can have only one destructor
- Takes a pointer to instance and reverses actions of the constructor:
 - a) initialize vfptrs if has virtual functions
 (this is done so that any virtual calls in the body use the methods of the current class)
 - b) execute the destructor body
 - c) call destructors of complex class members
 - d) call destructors of base classes
- Simple destructors can be inlined, so you can often see the vfptr reloaded many times in the same function

Destructor example

```
virtual thiscall CMruLongList::~CMruLongList(void) proc near
 mov edi, edi
 push esi
 mov esi, ecx
 mov eax, [esi+30h]
        dword ptr [esi], offset const CMruLongList::`vftable'
 mov
 test eax, eax
 jz short loc 42D93239
 push eax
                            : hMem
 call ds:LocalFree(x)
 and dword ptr [esi+30h], 0
loc 42D93239:
 mov ecx, esi
 pop esi
 virtual thiscall CMruLongList::~CMruLongList(void) endp
```

Virtual destructors

- When deleting object by pointer, a proper operator delete must be called
- It can be different for different classes in hierarchy
- Compiler has to make sure the correct operator regardless of the pointer type
- MSVC uses a helper function (deleting destructor) which is placed into the virtual table instead of the actual destructor
- It calls the actual destructor and then operator delete
- GCC emits multiple destructors (in-charge, not-in-charge, incharge deleting) and calls the corresponding one

Virtual destructor example

```
virtual void * thiscall CMruLongList::`scalar deleting
destructor'(unsigned int) proc near
  push
      ebp
  mov ebp, esp
  push esi
  mov esi, ecx
  test [ebp+arg 0], 1
  jz short loc_42D93260
  push esi
                             ; lpMem
  call operator delete(void *)
  pop
       ecx
loc 42D93260:
  mov eax, esi
  pop esi
  pop ebp
  retn 4
virtual void * __thiscall CMruLongList::`scalar deleting
destructor'(unsigned int) endp
```

RTTI: Run-time type information

- Necessary for dynamic_cast<> and typeid() operators
- Only required for polymorphic classes (with virtual methods)
- Because of this, usually attached to the virtual table
- MSVC uses a complex set of structures, see my OpenRCE article¹
- GCC puts a pointer to typeinfo class instance just before the method addresses
- First data member of that instance (after vfptr) is a pointer to the mangled name of the class.

¹ http://www.openrce.org/articles/full_view/23

RTTI alternatives: MFC

- MFC does not use standard RTTI
- All MFC classes inherit from CObject
- First virtual method of CObject is GetRuntimeClass()
- Returns a pointer to a of CRuntimeClass instance
- The object contains the MFC class name, instance size and functions for dynamic creation

RTTI: MFC example

```
const CConfirmDriverListPage::`vftable'
  dd offset CConfirmDriverListPage::GetRuntimeClass(void)
  dd offset CConfirmDriverListPage::`vector deleting
destructor'(uint)
  dd offset CObject::Dump(CDumpContext &)
CConfirmDriverListPage::classCConfirmDriverListPage
  dd offset aCconfirmdriver ; m lpszClassName
                            ; m_nObjectSize
  dd 164h
                            ; m wSchema
  dd 0FFFFh
  dd offset CConfirmDriverListPage::CreateObject(void)
  dd offset CTypAdvStatPage:: GetBaseClass(void)
  dd 0
```

RTTI alternatives: Qt

- Qt uses a completely custom OOP model
- Slots and signals used instead of virtual methods
- Leaves a lot of meta information, including slot method names
- The whole implementation was described by Daniel Pistelli
- See http://www.ntcore.com/files/qtrev.htm for details and some IDC scripts

RTTI alternatives: Apple IOKit

- IOKit is the base framework for implementing drivers for Apple OS X and iOS
- Uses a subset of C++: no exceptions, templates, multiple inheritance or standard RTTI
- Uses its own implementation with support for dynamic object creation
- All classes inherit from OSObject
- One of its methods is getMetaClass()
- Metaclass instance contains instance size and class name
- A static instance of metaclass is created for each class

Names and hierarchy can be tracked from metaclasses

C++ and Hex-Rays

- IDA type system does not support C++ (yet)
- Hex-Rays is a C decompiler
- C++ constructs have to be emulated using C ones

C++	IDA/Hex-Rays
classes	structures
class inheritance	nested structures
virtual function table	function pointer table
implicit arguments	explicit arguments

C++ and Hex-Rays: classes and inheritance

C++ and Hex-Rays: function prototypes

- Some conversion is necessary if C++ prototypes are known (from headers or demangled symbol names)
- Non-static methods need a this pointer added
- Structure/class returns take an additional result pointer
- References to pointers (done by IDA automatically)

<pre>unsigned longthiscall CMachine::Initialize(void)</pre>	<pre>unsigned longthiscall CMachine::Initialize(CMachine* this)</pre>
<pre>myStruct MyClass::getStruct()</pre>	<pre>myStruct* MyClass::getStruct (myStruct* result, MyClass*this)</pre>
<pre>static unsigned long CFolderRedirector::GetDefaultA ttributes(struct _GUID const &)</pre>	<pre>static unsigned long CFolderRedirector::GetDefaultAttributes (struct _GUID *)</pre>

C++ and Hex-Rays: virtual tables

- Table structure can be made manually
- You can use "Create struct from data" to generate initial structure
- Then set types of each member to be a function pointer
- Not very hard to create a script which analyzes vtable and creates a structure
- Hint: add a repeatable comment with the target address and number of purged bytes
- One table per class or share among many classes
- First approach is more universal but recovered prototypes have to be copied to other tables

Second one doesn't work for tree inheritance

```
CMachine::`vftable'
  dd offset CMachine::`scalar deleting destructor'(uint)
  dd offset CMachine::Initialize(void)
00 CMachine vtable struc
00 __delDtor dd ? ; 0101A9C3
04 Initialize dd ? ; 0101A6C0
08 CMachine vtable ends
struct CMachine vtable
  int (__thiscall *__delDtor)(CMachine *, int);
  int ( thiscall *Initialize)(CMachine *);
};
```

```
public: ___thiscall CMachine::CMachine(void) proc near
 push
      esi
 mov esi, ecx
 and dword ptr [esi+24h], 0
 or dword ptr [esi+28h], 0FFFFFFFh
 and dword ptr [esi+2Ch], 0
 or dword ptr [esi+30h], 0FFFFFFFh
 mov dword ptr [esi], offset const
CMachine::`vftable'
 mov eax, esi
 pop esi
 retn
public: thiscall CMachine::CMachine(void) endp
```

```
00 CMachine struc
00 _ CDataStoreObject ?
24 dword24 dd ?
28 dword28 dd ?
20 dword2C dd ?
30 dword30 dd ?
34 CMachine ends

struct CMachine
{
    CDataStoreObject _;
    _ DWORD dword24;
    _ DWORD dword28;
    _ DWORD dword2C;
    _ DWORD dword30;
};
```

```
public: ___thiscall CMachine::CMachine(void) proc near
 push esi
 mov esi, ecx
 call CDataStoreObject::CDataStoreObject(void)
 and [esi+CMachine.dword24], 0
 or [esi+CMachine.dword28], 0FFFFFFFh
 and [esi+CMachine.dword2C], 0
 or [esi+CMachine.dword30], 0FFFFFFFFh
 mov [esi+CMachine. . vtable], offset const
CMachine::`vftable'
 mov eax, esi
 pop esi
 retn
public: thiscall CMachine::CMachine(void) endp
```

```
CMachine * thiscall CMachine::CMachine(CMachine)
*this)
  CDataStoreObject::CDataStoreObject(&this-> );
  this->dword24 = 0;
  this->dword28 = -1;
  this->dword2C = 0;
  this->dword30 = -1;
  this-> . vtable = (CMachine vtable
*)&CMachine:: vftable;
  return this;
```

```
ecx, [esi+4]
                                  ; CMachine *
mov
cmp ecx, edi
jz
     short loc 100C57C
mov eax, [ecx+CMachine. . vtable]
push
       1
                                  ; int
call [eax+CMachine vtable. delDtor]; 0101A9C3
       [esi+4], edi
mov
_machine = context->m_machine;
if ( machine )
  _machine->_._vtable->__delDtor(_machine, 1);
  context->m machine = 0;
}
```

C++ and Hex-Rays: vtables redux

- Nesting of complete class structures works for simple cases, but not good for complex inheritance
- We cannot set different types for the vtable pointer shared between classes
- Nesting breaks down in case of virtual inheritance because virtual bases are shuffled around
- A different approach is to use two structures: one for just the fields, and one for the complete class
- The complete class structure includes the fields structure and adds virtual table pointers

C++ and Hex-Rays: vtable redux example

```
00 CDataStoreObject fields struc
                                        00 CDataStoreObject struc
00 dword4 dd?
                                        00 vtable dd ?
04 m cs RTL CRITICAL SECTION ?
                                        04 f CDataStoreObject fields ?
1C dword20 dd?
                                        24 CDataStoreObject ends
20 CDataStoreObject fields ends
00 CMachine fields struc
                                        00 (Machine struc
00 dword24 dd ?
                                        00 vtable dd?
04 dword28 dd?
                                        04 __f1 CDataStoreObject_fields ?
08 dword2C dd ?
                                        24 f2 CMachine fields ?
OC dword30 dd?
                                        34 CMachine ends
10 CMachine fields ends
                                        00 CSession struc
                                        00 vtable dd?
                                        04 _f1 CDataStoreObject_fields ?
                                        24 _f2 CSession_fields ?
                                        84 CSession ends
```

C++ and Hex-Rays: vtable redux example

```
eax, [ecx+CMachine. vtable]
mov
push
                                    ; int
call
        [eax+CMachine vtable. delDtor] ; 0101A9C3
struct CMachine vtable {
  int ( thiscall * delDtor)(CMachine *, int);
  int ( thiscall *Initialize)(CMachine *);
};
result = machine-> vtable-> delDtor( machine, 1);
     eax, [ecx+CSession. vtable]
mov
push
call
     [eax+CSession vtable. delDtor]; 0101B90F
struct CSession vtable {
  int ( thiscall * delDtor)(CSession *, int);
  int ( thiscall *Initialize)(CSession *);
};
result = session-> vtable-> delDtor( session, 1);
```

C++ decompiling workflow

- Identify constructors and destructors from the global init tables or code patterns (stack construction, heap/new allocation, unwind funclets)
- Drill down to the most base constructors/destructors
- Make initial structures (e.g. using "Create new struct type")
- If has vtable, make a vtable stucture and set the vfptr type to be a pointer to it
- Follow cross-references to identify other methods of the class
- Fix up the structures and vtable function pointer types as necessary

Conclusion

- C++ decompilation is somewhat difficult but doable
- A lot of information can be extracted from RTTI and vtables/vbtables
- Many common tasks can be automated
- There is a lot of room for improvement

Links

MSVC: http://www.openrce.org/articles/full_view/23

GCC: http://www.codesourcery.com/public/cxx-abi/

Bonus matter: Hex-Rays 1.6 preview

Spoiler Alert

Hex-Rays 1.6: variable mapping

```
v1 = this;
 lck mtx lock(this-> b.field 228);
 if ( !OSIncrementAtomic(&v1->__b.field_19C) )
   v1-> vtbl->virt380(v1);
[map v1 to this]
 lck_mtx_lock(this->__b.field_228);
 if ( !OSIncrementAtomic(&this-> b.field 19C) )
   this->_vtbl->virt380(this);
```

Hex-Rays 1.6: support for unions

```
if ( StackLocation->Parameters.Read.ByteOffset.LowPart == 315396
    || StackLocation->Parameters.Read.ByteOffset.LowPart == 315412 )
{
    if ( StackLocation->Parameters.Create.Options >= 0x2C )

[choose correct union field]

if ( StackLocation->Parameters.DeviceIoControl.IoControlCode == 0x4D004
    || StackLocation->Parameters.DeviceIoControl.IoControlCode == 0x4D014 )
{
    if ( StackLocation->Parameters.DeviceIoControl.InputBufferLength >= 0x2C )
```

Hex-Rays 1.6: kernel idioms support

```
deviceInfo->ListEntry.Blink = &deviceInfo->ListEntry;
deviceInfo->ListEntry.Flink = deviceInfo->ListEntry.Blink;
InitializeListHead(&deviceInfo->ListEntry);
while ( (LIST ENTRY *)ListHead.Flink != &ListHead )
while ( !IsListEmpty(&ListHead) )
deviceInfo = (_DEVICE_INFO *)((char *)&thisEntry[-131] - 4);
if ( *p_serial == *((_DWORD *)thisEntry - 1) )
  break:
deviceInfo = CONTAINING_RECORD(thisEntry, DEVICE INFO, ListEntry);
if ( *p_serial == CONTAINING_RECORD(thisEntry, _DEVICE_INFO,
ListEntry)->SerialNo )
  break:
```

We're hiring!

We're looking for someone to help us improve the decompiler

If you liked this talk and would like to work on it, let us know

info@hex-rays.com

Thank you!

Questions?