



Dragon Threat Labs

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Insight in to a strategic web compromise and attack campaign against Hong Kong infrastructure. Revealing an attacker's persistence, sophistication and aggression.

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Introduction

Over the past several months an increasing number of strategic web compromises (“wateringholes”) have been discovered on websites in Hong Kong. This rise in activity coincides with the Occupy Central protests. In this post we will talk about a single attack, whilst not trying to distract attention from the vast number of attacks and subsequent compromises that remain persistent in Hong Kong.

Whilst going about our daily business we were alerted to a website that began serving a malicious payload alongside its usual web page. The initial investigation revealed that the attack and associated payloads are part of an ongoing attack campaign by an Advanced Persistent Threat group that is known to target various sectors of industry and Government in Hong Kong.

In this instance we have chosen to obfuscate details of the compromised website to protect the identity of the victim. This website belongs to a private educational institution who, since being notified about the compromise, has removed the malicious executable and remediated the compromised of their server, thus breaking a crucial link in the chain of this attack.

Strategic web compromise

The website in question was implanted with some HTML code that simply reaches out to a secondary website and downloads malware. Whilst this in itself is not interesting the methodology used to obfuscate code and evade detection are. The underlying code in the first page that loads exploits a vulnerability in Internet Explorer ([CVE-2014-6332](#)) and runs several scripts, each with support for different operating systems and methods of downloading and executing a file from a website. The first script is obfuscated Visual Basic Script (“VBS”)

```
1 <html>
2 <meta http-equiv="X-UA-Compatible" content="IE=EmulateIE8" >
3 <head>
4 <title>Count</title>
5 </head>
6 <!-- saved from url=(0014)about:internet -->
7 <script language=vbscript>
8 function MoSaklgEs7(k)
9 s=Split(k,"")
10 t=""
11 For i = 0 To UBound(s)
12 t=t+Chr(eval(s(i)))
13 Next
14 MoSaklgEs7=t
15 End Function
16 t=
"60,33,100,111,99,116,121,112,101,32,104,116,109,108,62,13,10,60,104,116,109,108,
,108,101,34,32,99,111,110,116,101,110,116,61,34,73,69,61,69,109,117,108,97,116,10
,83,67,82,73,80,84,32,76,65,78,71,85,65,71,69,61,34,86,66,83,99,114,105,112,116,3
,82,101,115,117,109,101,32,78,101,120,116,13,10,13,10,83,101,116,32,111,83,104,10
41 13 10 115 116 114 72 111 109 101 70 111 108 100 101 114 32 61 32 111 83 104 10
```

By decoding this we can see the true intentions of the script – which opens a whole new can of worms.

First stage malware (Dropper component – “Swisyn” part 1)

The first binary payload that lands on the system is relatively simple and serves as a method of yet again detecting the operating system version and where to drop a secondary payload file. Whilst this binary is not complicated by nature it has been designed to masquerade as a legitimate application and contains functionality to evade anti-virus. This malware implant is commonly detected by anti-virus as [Swisyn](#).

Upon running this malware determines the operating system version, but only delineating between Windows XP, Vista and above.

```
.text:004015A3 GetOS      proc near          ; CODE XREF: sub_4016C9+C94p
.text:004015A3
.text:004015A3 var_A0      = dword ptr -0A0h
.text:004015A3 VersionInformation= _OSVERSIONINFOA ptr -9Ch
.text:004015A3
.text:004015A3      push    ebp
.text:004015A4      mov     ebp, esp
.text:004015A6      sub     esp, 0A0h
.text:004015A8      push    edi
.text:004015AA      mov     [ebp+VersionInformation.dwOSVersionInfoSize], 0
.text:004015AB      mov     ecx, 26h
.text:004015AD      xor     eax, eax
.text:004015AE      lea     edi, [ebp+VersionInformation.dwMajorVersion]
.text:004015AF      rep stosl
.text:004015B0      mov     [ebp+VersionInformation.dwOSVersionInfoSize], 9Ch
.text:004015B1      lea     eax, [ebp+VersionInformation]
.text:004015B2      push    eax
.text:004015B3      call   ds:GetVersionExA
.text:004015B4      test    eax, eax
.text:004015B5      jnz     short loc_4015E8
.text:004015B6      xor     eax, eax
.text:004015B7      jmp     loc_4016C4
.text:004015E8      ;-----
.text:004015E8 loc_4015E8:      ; CODE XREF: GetOS+3C7j
.text:004015E8      cmp     [ebp+VersionInformation.dwMajorVersion], 5
.text:004015E9      jnz     short loc_401609
.text:004015EA      cmp     [ebp+VersionInformation.dwMinorVersion], 0
.text:004015EB      jnz     short loc_401609
.text:004015EC      mov     [ebp+var_A0], 0
.text:004015ED      jmp     loc_4016BE
.text:00401609      ;-----
.text:00401609 loc_401609:      ; CODE XREF: GetOS+4C7j
.text:00401609      ; GetOS+551j
.text:00401609      cmp     [ebp+VersionInformation.dwMajorVersion], 5
.text:0040160A      jnz     short loc_40162A
.text:0040160B      cmp     [ebp+VersionInformation.dwMinorVersion], 1
.text:0040160C      jnz     short loc_40162A
000009E8 004015E8: GetOS:loc_4015E8
```

It appears that this functionality is included because the secondary payload comes in both 32bit and 64bit versions.

Both of these second stage payloads are obfuscated but decoded with a simple bitwise operation as per below

```

.text:00401566 loc_401566:                ; CODE XREF: sub_401536+67↓j
.text:00401566      mov     edx, [ebp+var_8]
.text:00401569      add     edx, 1
.text:0040156C      mov     [ebp+var_8], edx
.text:0040156F loc_40156F:                ; CODE XREF: sub_401536+2E↑j
.text:0040156F      mov     eax, [ebp+arg_4]
.text:00401572      sub     eax, 1
.text:00401575      cmp     [ebp+var_8], eax
.text:00401578      jnb     short loc_40159F
.text:0040157A      mov     ecx, [ebp+arg_0]
.text:0040157D      mov     dl, [ecx]
.text:0040157F      sub     dl, [ebp+var_4]
.text:00401582      mov     eax, [ebp+arg_0]
.text:00401585      mov     [eax], dl
.text:00401587      mov     ecx, [ebp+arg_0]
.text:0040158A      mov     dl, [ecx]
.text:0040158C      xor     dl, [ebp+var_4]
.text:0040158F      mov     eax, [ebp+arg_0]
.text:00401592      mov     [eax], dl
.text:00401594      mov     ecx, [ebp+arg_0]
.text:00401597      add     ecx, 1
.text:0040159A      mov     [ebp+arg_0], ecx
.text:0040159D      jmp     short loc_401566
.text:0040159F      ; -----
.text:0040159F loc_40159F:                ; CODE XREF: sub_401536+42↑j
0000099F  0040159F: sub_401536:loc_40159F

```

In this scenario, the secondary payloads can be decoded using a simple subtraction of 3 followed by an XOR of 3 from each byte. This file is then written to %User%\Application Data\Microsoft in a newly created folder name 'wuauctl'. The filename depends on the operating system version, for Windows XP it is "clbcatq.dll", for Windows Vista and above it is "profapi.dll". Once this file has been written to disk a file from the Windows System32 folder is copied into the directory. This file, named 'wuauctl.exe', is the Windows update client interface and it a standard windows file. By executing this file in a specific manner it will load the freshly dropped DLL file – affectively this is known as DLL hijacking.

Following this action another file, named 'wuauctl.dat', is written on to the disk under the same directory. This file is encoded and not decoded at this stage of the attack. To complete the file drop wuauctl.exe is executed.

The 64bit version of this dropper is vastly similar in functionality although it offers slightly more efficiency in the code. The decoding routine is more simplified than its 32bit counterpart and the decoding key is hardcoded

```

.text:00000000100015B0 DecodeLoop:                                ; CODE XREF: StartAddress+182↓j
.text:00000000100015B0      movzx  ecx, byte ptr [rdx] ; ecx = DecodeBuffer[i]
.text:00000000100015B3      inc    rdx                ; Increase buffer index
.text:00000000100015B6      sub    cl, 3              ; DecodeBuffer[i] -= 3
.text:00000000100015B9      xor    cl, 3              ; DecodeBuffer[i] ^= 3
.text:00000000100015BC      dec    r8
.text:00000000100015BF      mov    [rdx-1], cl
.text:00000000100015C2      jnz    short DecodeLoop ; ecx = DecodeBuffer[i]
.text:00000000100015C4      loc_100015C4:            ; CODE XREF: StartAddress+16A↑j
.text:00000000100015C4      lea    rcx, unk_10004030
.text:00000000100015C8      mov    r8d, eax
.text:00000000100015CE      mov    rdx, rbx
.text:00000000100015D1      call   MemorySetup
.text:00000000100015D6      test   eax, eax
.text:00000000100015D8      jnz    short loc_100015E2
.text:00000000100015DA      mov    rcx, rbx          ; void *
.text:00000000100015DD      loc_100015DD:            ; CODE XREF: StartAddress+14C↑j

```

The following pseudo-code can decode both 32bit and 64bit versions of the DLL stored in 'wuauclt.dat'

```

void decode(BYTE *data, int len, BYTE xorkey)
{
    for (int i = 0; i < len; i++)
    {
        data[i] -= xorkey;
        data[i] ^= xorkey;
    }
    printf("%s", data);
}

```

Not to dwell on a dropper, let's move on to the second stage malware.

Second stage (Decoder and loader – “Swisyn” part 2)

The malware second stage is now loaded and running. Interestingly, this payload is also detected by anti-virus as [Swisyn](#). This DLL is again fairly simple and acts as a secondary dropper. It primarily serves as a method of decoding one of the files dropped by the previous malware stage and creating a method to start the malware on system boot-up or user login. In order to do this the malware firstly decodes a file that was dropped by the previous stage – in this case it is “wuauclt.dat”. The decoding routine is again overly complex but ultimately amounts to a simple subtraction and XOR, again both of these operations are performed by the number 3, thus each byte is subtracted by 3 and then XOR'd with 3

```

.text:100018A0 DecodeConfig      proc near                                ; CODE XREF: sub_100018E0+34↓p
.text:100018A0
.text:100018A0     arg_0          = dword ptr 4
.text:100018A0     arg_4          = dword ptr 8
.text:100018A0     arg_8          = dword ptr 0Ch
.text:100018A0
.text:100018A0     mov     eax, [esp+arg_8]
.text:100018A4     mov     ecx, 1Ah
.text:100018A9     and     eax, 0FFh
.text:100018AE     cdq
.text:100018AF     idiv    ecx
.text:100018B1     mov     eax, [esp+arg_4]
.text:100018B5     inc     dl
.text:100018B7     dec     eax
.text:100018B8     test    eax, eax
.text:100018BA     jbe     short locret_100018D0
.text:100018BC     push    esi
.text:100018BD     mov     esi, eax
.text:100018BF     mov     eax, [esp+4+arg_0]
.text:100018C3
.text:100018C3 loc_100018C3:                ; CODE XREF: DecodeConfig+2D↓j
.text:100018C3     mov     cl, [eax]
.text:100018C5     sub     cl, dl
.text:100018C7     xor     cl, dl
.text:100018C9     mov     [eax], cl
.text:100018CB     inc     eax
.text:100018CC     dec     esi
.text:100018CD     jnz     short loc_100018C3
.text:100018CF     pop     esi
.text:100018D0
.text:100018D0 locret_100018D0:            ; CODE XREF: DecodeConfig+1A↑j
.text:100018D0     retn
.text:100018D0 DecodeConfig      endp

```

Once this file has been decoded it is loaded into memory and executed. This file, once decoded is the third and final payload. The method of leaving the encoded file on disk and only decoding it in memory is to thwart poorly configured anti-virus or disk surface heuristic scanners.

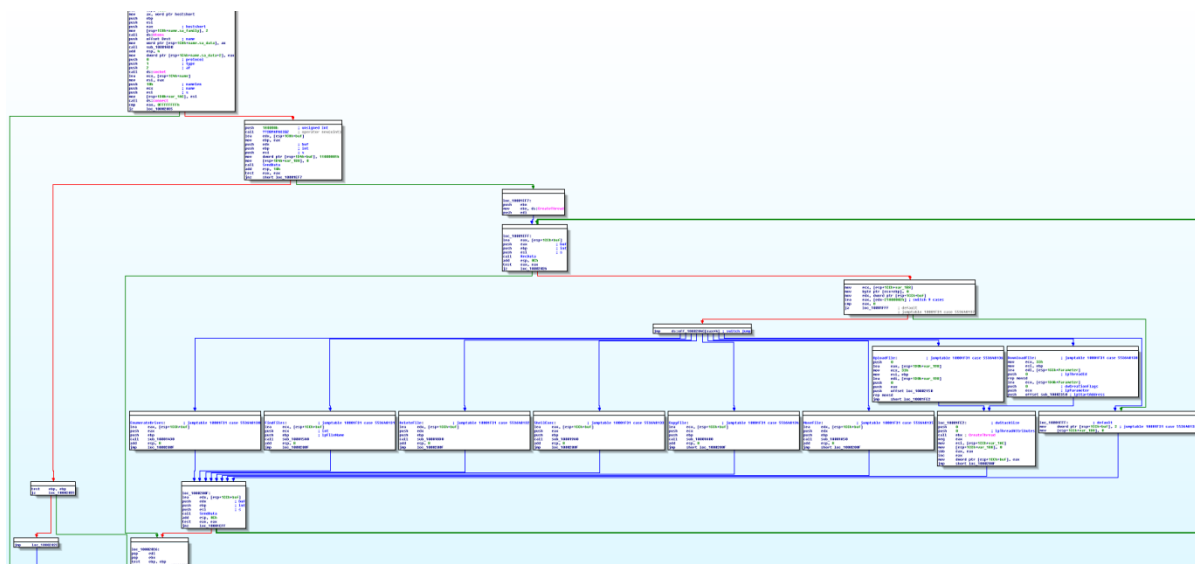
Finally to wrap up, an entry is created in the registry named 'wuauclt' is created under 'HKCU\Software\Microsoft\Windows\Current Version\Run' to ensure that this file is executed upon user-login.

Final stage (RAT – “PCClient”)

Finally we are left with a full payload. Unsurprisingly the 3rd and final stage of this part of the attack is a fully fledged RAT (Remote Administration Tool), which is detected by anti-virus as [PCClient](#). This RAT allows the attacker to control the infected workstation and perform a vast array of administrative functions such as:

- Downloading files to the infected workstation
- Uploading from the infected workstation files to the attackers
- Enumerate/list all connected drives such as network shares or external devices
- Search the infected workstations hard drive for files
- Deleting, copying and moving files on the infected workstation
- Executing commands on the infected workstation

A high-level view of the command structure gives us an idea as to how simple this functionality can seem, but does not turn away from how damaging the affects can be:



Once the RAT has been loaded on the infected machine it begins calling out to the command and control server (“phoning home”) and waits for the attackers to issue one of the above commands to the victim. As we usually see with APT attacks the malware controllers use a specific ID to code their attack campaign, which in this case is ‘C00BBB’. Information about the victim system is collected and posted off to the command and control server. This information gives the attacker a brief description about the machine. The information consists of:

- Machine hostname
- Total amount of RAM memory
- Operating system and service pack level
- Attack campaign code

This information is encoded using a simple bitwise operation and then sent to the command and control server. For example:

Unencoded data	Encoded data
44 45 4C 4C 2D 31 37 38 DELL-178	BA B9 B2 B2 51 4D 47 46 9122QMGF
44 33 43 00 00 00 00 00 D3C.....	BA 4B BB 7E 7E 7E 7E 7E 7E 9K»~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~
35 31 32 4D 42 00 00 00 512MB...	49 4D 4C B1 BC 7E 7E 7E 7E IML±¼~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~
57 69 6E 20 58 50 20 53 Win XP S	A7 95 90 5E A6 AE 5E AB §••^!®^«
50 33 20 28 42 75 69 6C P3 (Buil	AE 4B 5E 56 BC 89 95 92 ®K^V¼%•’

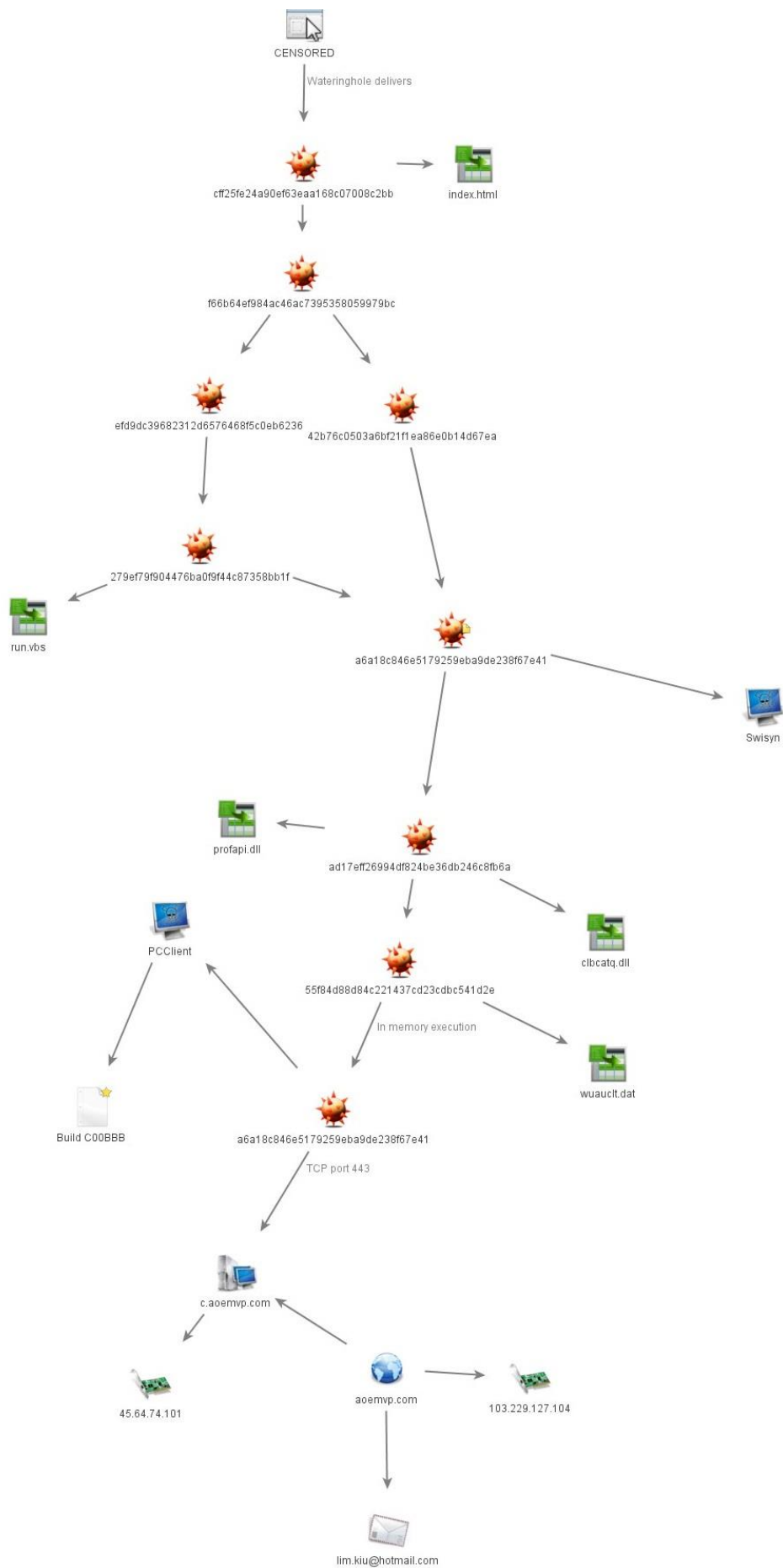
64 20 32 36 30 30 29 00 d 2600).	9A 5E 4C 48 4E 4E 55 7E Š^LHNNU~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~~
00 00 00 00 00 00 00 00	7E 7E 7E 7E 7E 7E 7E 7E ~~~~~~
43 30 30 42 42 42 00 00 CO0BBB..	BB 4E 4E BC BC BC 7E 7E »NN¼¼¼~

Whilst this may seem to make the data harder to recover it actually makes detection of the traffic easier. To decode the traffic a simple calculation can be performed by reversing the encoding operations. In this case the malware simply increases the initial encoding key by 1, then adds this value to each byte in the buffer and finally XOR's each byte. Once again, the following pseudo-code can decode this data

```
void decode(BYTE *data, int len, BYTE xorkey)
{
    for (int i = 0; i < len; i++)
    {
        data[i] -= xorkey;
        data[i] ^= xorkey;
    }
    printf("%s", data);
}
```

Infrastructure & associations

During in the investigation we performed analysis of the infrastructure that this malware communicates with. On this occasion we have not been able to gain physical access to the command and control server as it is legitimate, but compromised production infrastructure. The graph below shows the flow in which various parts of the attack are loaded and how they chain together.



Detection & mitigation

This attack can be detected and/or mitigated at each stage. In order to help organisations protect themselves we have created a number of network IDS rules and disk-scan rules that can be used with Snort and Yara. Rules are provided in a best-effort basis and we cannot vouch for their efficiency in your environment.

Wateringhole code

```
rule apt_win_wateringhole {
meta:
    author = "@dragonthreatlab "
    description = "Detects code from APT wateringhole"
strings:
    $str1 = "function runmumaa()"
    $str2 = "Invoke-Expression $(New-Object IO.StreamReader ($(New-Object IO.Compression.DeflateStream ($(New-Object IO.MemoryStream ([Convert]::FromBase64String("
    $str3 = "function MoSaklgEs7(k)"
condition:
    any of ($str*)
}
```

Swisyn

```
rule apt_win_swisyn {
meta:
    author = "@dragonthreatlab"
    md5 = "a6a18c846e5179259eba9de238f67e41"
    description = "File matching the md5 above tends to only live in memory, hence the lack of MZ header check."
strings:
    $mz = {4D 5A}
    $str1 = "/ShowWU" ascii
    $str2 = "IsWow64Process"
    $str3 = "regsvr32 "
    $str4 = {8A 11 2A 55 FC 8B 45 08 88 10 8B 4D 08 8A 11 32 55 FC 8B 45 08 88 10}
condition:
    $mz at 0 and all of ($str*)
}
```

Malware dropper 32bit

```
rule apt_win32_dropper {
meta:
    author = "@dragonthreatlab"
    md5 = "ad17eff26994df824be36db246c8fb6a"
    description = "APT malware used to drop PcClient RAT"
strings:
    $mz = {4D 5A}
    $str1 = "clbcaiq.dll" ascii
    $str2 = "profapi_104" ascii
    $str3 = "/ShowWU" ascii
    $str4 = "Software\\Microsoft\\Windows\\CurrentVersion\\" ascii
```

```
$str5 = {8A 08 2A CA 32 CA 88 08 40 4E 75 F4 5E}
condition:
    $mz at 0 and all of ($str*)
}
```

Malware dropper 64bit

```
rule apt_win64_dropper {
meta:
    author = "@dragonthreatlab"
    md5 = "ad17eff26994df824be36db246c8fb6a"
    description = "APT malware used to drop PcClient RAT"
strings:
    $mz = {4D 5A}
    $str1 = "clbcaiq.dll" ascii
    $str2 = "profapi_104" ascii
    $str3 = "\\Microsoft\\wuauc\\wuauc\\wuauc.dat" ascii
    $str4 = {0F B6 0A 48 FF C2 80 E9 03 80 F1 03 49 FF C8 88 4A FF
75 EC}
condition:
    $mz at 0 and all of ($str*)
}
```

Encoded version of PcClient

```
rule apt_win_disk_pcclient {
meta:
    author = "@dragonthreatlab "
    md5 = "55f84d88d84c221437cd23cdbbc541d2e"
    description = "Encoded version of pcclient found on disk"
strings:
    $header = {51 5C 96 06 03 06 06 06 0A 06 06 06 FF FF 06 06 BE
06 06 06 06 06 06 06 46 06 06 06 06 06 06 06 06 06 06 06 06 EE 06 06
06 06 06 06 06 06 06 06 06 06 06 06 06 06 06 06 06 06 EE 06 06
06 10 1F BC 10 06 BA 0D D1 25 BE 05 52 D1 25 5A 6E 6D 73 26 76 74 6F
67 74 65 71 26 63 65 70 70 6F 7A 26 64 69 26 74 79 70 26 6D 70 26 4A
4F 53 26 71 6F 6A 69 30 11 11 0C 2A 06 06 06 06 06 06 06 73 43 96 1B
37 24 00 4E 37 24 00 4E 37 24 00 4E BA 40 F6 4E 39 24 00 4E 5E 41 FA
4E 33 24 00 4E 5E 41 FC 4E 39 24 00 4E 37 24 FF 4E 0D 24 00 4E FA 31
A3 4E 40 24 00 4E DF 41 F9 4E 36 24 00 4E F6 2A FE 4E 38 24 00 4E DF
41 FC 4E 38 24 00 4E 54 6D 63 6E 37 24 00 4E 06 06 06 06 06 06 06 06
06 06 06 06 06 06 06 06 56 49 06 06 52 05 09 06 5D 87 8C 5A 06 06 06
06 06 06 06 06 E6 06 10 25 0B 05 08 06 06 1C 06 06 06 1A 06 06 06 06
06 06 E5 27 06 06 06 16 06 06 06 36 06 06 06 06 06 16 06 16 06 06 06
04 06 06 0A 06 06 06 06 06 06 06 0A 06 06 06 06 06 06 06 76 06 06
06 0A 06 06 06 06 06 06 04 06 06 06 06 16 06 06 16 06 06}
condition:
    $header at 0
}
```

In-memory version on PcClient

```
rule apt_win_memory_pcclient {
meta:
    author = "@dragonthreatlab "
    md5 = "ec532bbe9d0882d403473102e9724557"
    description = "File matching the md5 above tends to only live
in memory, hence the lack of MZ header check."
strings:
    $str1 = "Kill You" ascii
    $str2 = "%4d-%02d-%02d %02d:%02d:%02d" ascii
    $str3 = "%4.2f KB" ascii
    $encodefunc = {8A 08 32 CA 02 CA 88 08 40 4E 75 F4}
condition:
    all of them
}
```

PcClient malware beaconing

```
alert tcp $HOME_NET any -> $EXTERNAL_NET [80,443] (msg:"MALWARE -
DTL ID 12012015 - PcClient beacon"; flow:established,to_server;
content:"|BB 4E 4E BC BC BC 7E 7E|"; nocase; offset:160; depth:8;
classtype:trojan-activity; sid:YOUR_SID; rev:20122014;)
```

Malware domain

```
alert udp $HOME_NET any -> $EXTERNAL_NET 53 (msg:"MALWARE - DTL ID
12012015 - C2 Domain"; content:"|06|aoemvp|03|com";
classtype:trojan-activity; sid:YOUR_SID; rev: 20122014;)
```

C2 server IP #1

```
alert ip $HOME_NET any <> 45.64.74.101 any (msg:"MALWARE - DTL ID
12012015 - C2 IP Address"; classtype:trojan-activity; sid:YOUR_SID;
rev: 20122014;)
```

C2 server IP #2

```
alert ip $HOME_NET any <> 103.229.127.104 any (msg:"MALWARE - DTL ID
12012015 - C2 IP Address "; classtype:trojan-activity; sid:YOUR_SID;
rev: 20122014;)
```

Appendix

The following artefacts were found during the investigation

MD5s

```
a6a18c846e5179259eba9de238f67e41  
55f84d88d84c221437cd23cdb541d2e  
a6a18c846e5179259eba9de238f67e41  
279ef79f904476ba0f9f44c87358bb1f  
42b76c0503a6bf21f1ea86e0b14d67ea  
cfff25fe24a90ef63eaa168c07008c2bb  
ad17eff26994df824be36db246c8fb6a  
f66b64ef984ac46ac7395358059979bc  
efd9dc39682312d6576468f5c0eb6236
```

Network artefacts

```
c.aoemvp.com  
aoemvp.com  
lim.kiu@hotmail.com  
45.64.74.101  
103.229.127.104
```

Contact

For all questions relating to the publication or specifics in this document please contact us via one of the following methods:

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Kind regards,

Dragon Threat Labs