

Debit

A \rightarrow (-60)

 \times
A (-10)

Credit

B (50)

C (10)

\times \times \times

B (0)

C (10)

P1 → P2
 Payer: Pulkit, Sanket
 Payee: Deepak, Deepak, Riya, Riya
 Amount: 100, 200, 300, 200
 Add: 1
 Payer: Pulkit, Sanket
 Payee: Deepak, Deepak, Riya, Riya
 Amount: 100, 200, 300, 200
 + ADD

BUILD GRAPH SIMPLIFY SETTLEMENTS

Simplified Settlement

Payer: Riya, Riya
 Payee: Sanket, Pulkit
 Amount: 400, 100
 1

Deepak → Pulkit (100)

Deepak → Sanket (200)

Riya → Deepak (300)

Riya → Sanket (200)

key → value → Hash Map
 Deepak → 0

Sanket → ~~400~~ 0

Riya → ~~300~~ ~~600~~ 0

Pulkit → ~~100~~ 0

① Calculate the total debit or credit for a person

A B C D \rightarrow again

Cost of 1 ticket \rightarrow 30

A paid for ticket \rightarrow 200 \rightarrow $\begin{matrix} A \\ B \\ C \\ D \end{matrix}$

$\left\{ \begin{array}{l} A \rightarrow B \rightarrow 50 \\ A \rightarrow C \rightarrow 50 \\ A \rightarrow D \rightarrow 50 \end{array} \right\}$

A \rightarrow A 50

How to segregate who is under credit &

who is under debit.

$p \rightarrow a$

$a < 0$ ($p \rightarrow \text{debit}$)

$a > 0$ ($p \rightarrow \text{credit}$)

$a = 0$ (p is settled)

① from all the person who are under debit,
get the one with max debit.

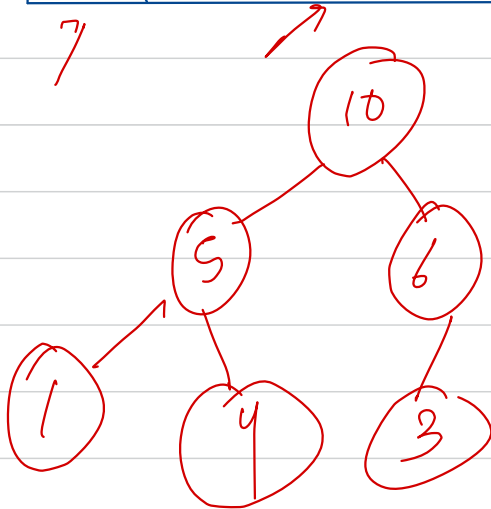
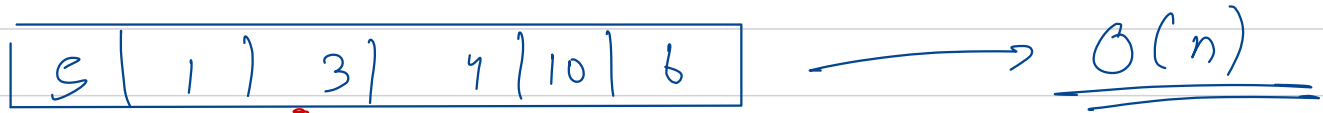
from all the persons who are under credit
get the one with max credit.

And settle then

Max heap

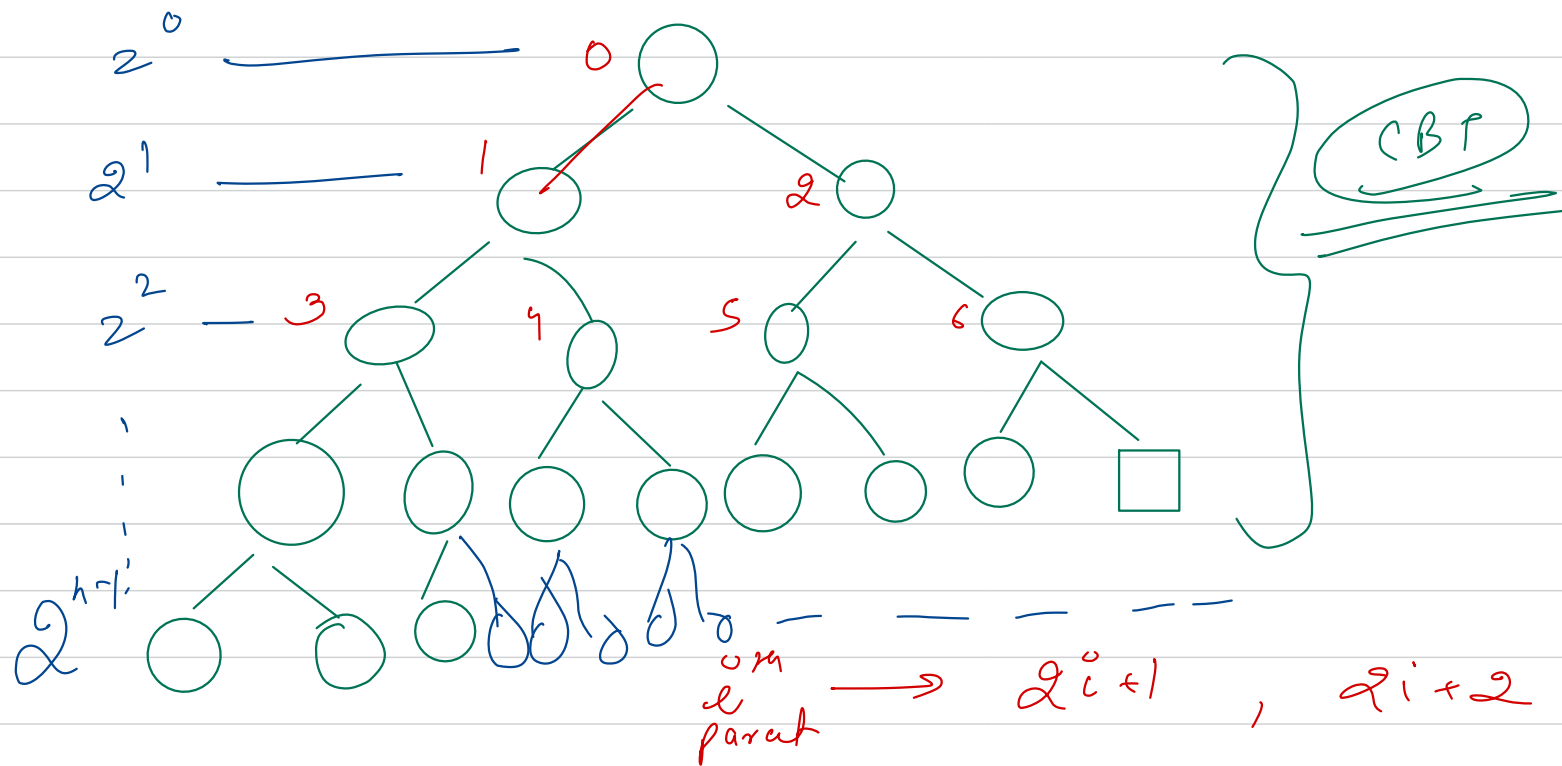
① heap is a complete B.T. ←

② Priority of a parent is higher than child.



upheapify

$O(n \log n)$



$$j^{\text{th}} \text{ child} \rightarrow \frac{j-1}{2} \rightarrow \text{parent}$$

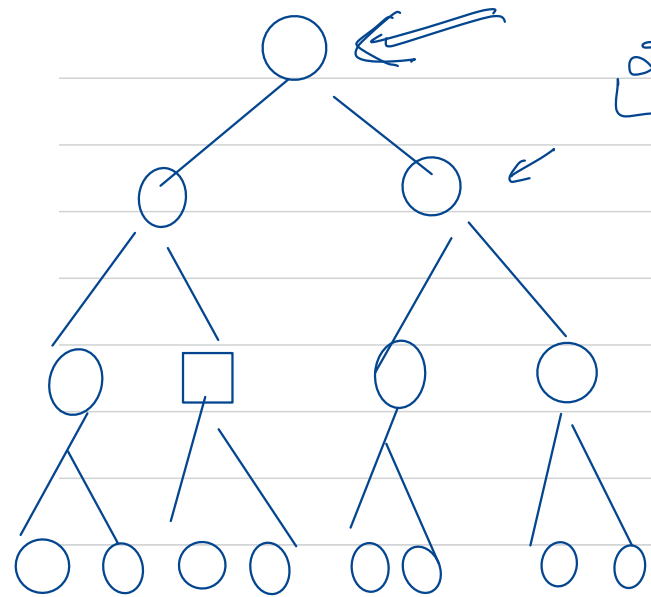
$$\underbrace{2^0 + 2^1 + 2^2 + \dots + 2^{h-1}} = n$$

Sum of
gp

$$2^h - 1 = n$$

$$2^h = n + 1$$

$$h = \log_2(n+1)$$



$$2^0 + 2^1 + 2^2 + \dots + 2^{h-2} + \boxed{2^{h-1}} = n$$

$$2^{h-1} - 1 + 2^{h-1} = n$$

$$2(2^{h-1}) = n + 1$$

$$\boxed{2^{h-1}} = \frac{(n+1)}{2}$$

nodes
on
last level
legs

→ How many elements on last level

$$2^{h-1}(h-1) + 2^{h-2}(h-2) + \dots + 2(1) + 1(0)$$

$$S = \underbrace{2^{h-1}(h-1)} + 2^{h-2}(h-2) \dots - 2^2(2) + 2^1(1) + \cancel{1(0)}$$

$$2S = 2^h(h-1) + \underbrace{2^{h-1}(h-2)} \dots - 2^3(2) + 2^2(1) \cdot$$

$$2S - S = 2^h(h-1) + \underbrace{2^{h-1}(-1) + 2^{h-2}(-1) \dots - 2^2(-1)}_{+2}$$

$$S = 2^h(h-1) + (-1)2^2(2^{h-2} - 1) + 2$$

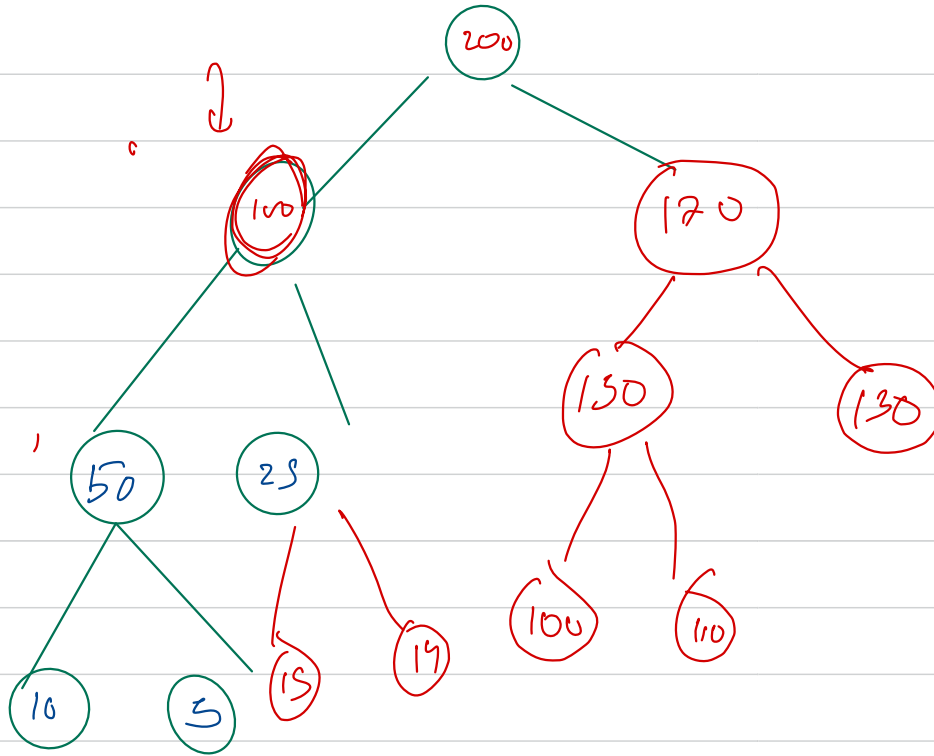
$$S = 2^h h - 2^h + (-1)(2^h - 2) + 2$$

$$S = \boxed{2^h h} - 2^h - 2^h + 2 + 2$$

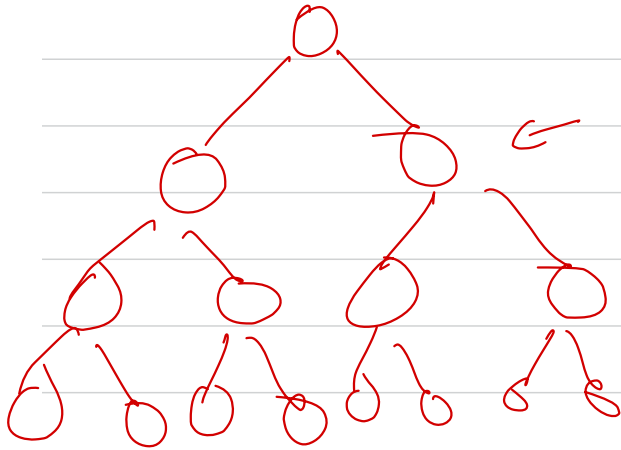
$2^h \rightarrow 2^{10} 2^n \rightarrow \underline{n}$

$$\underline{\underline{O(n \log n)}}$$

MAX
heap



Down heapify
Remain $\rightarrow O(\log n)$



$$S = 2^{h-1} \times 0 + (2^{h-2} \times 1) + (2^{h-3} \times 2) - - - - 2^{(h-2)} + 1^{(h-1)}$$

$$S = \underbrace{(2^{h-2} \times 1) + (2^{h-3} \times 2) + \dots + 2^2 \times (h-3) + 2^{(h-2)} + 1}_{(h-1)}$$

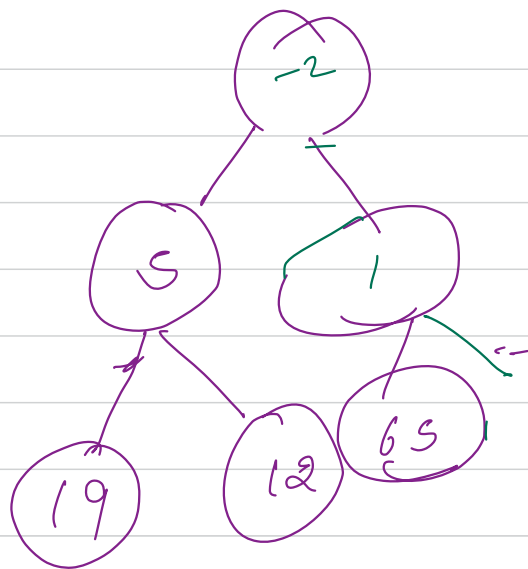
$$2S = \underbrace{2^{h-1} \times 1 + 2^{h-2} \times 2 + 2^{h-3} \times 3 + \dots + 2^2 \times (h-2) + 2^{(h-1)}}_{(h-1)}$$

$$2S - S = 2^{h-1} \times 1 + \underbrace{2^{h-2} + 2^{h-3} + \dots + 2^2 + 2}_{\text{GP}} + (h-1)$$

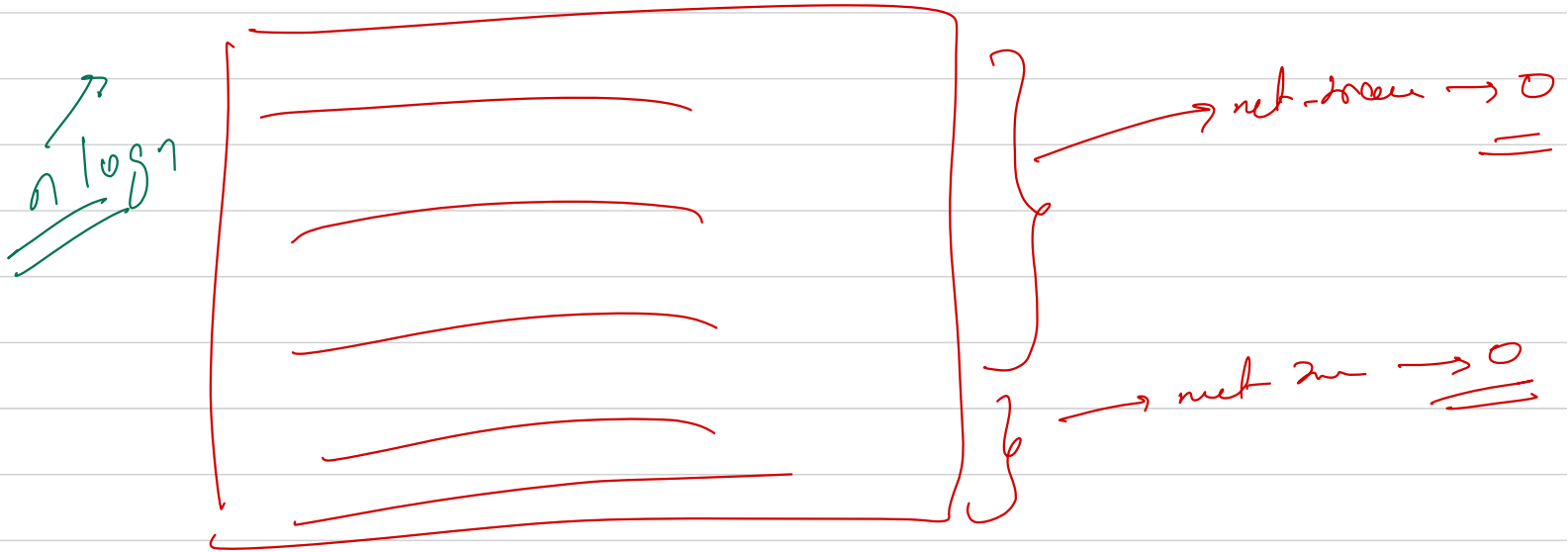
$$S = 2^{h-1} + 2(2^{h-2} - 1) + (h-1) \quad 2^{\log_2 n} = n$$

$$= 2^{h-1} + 2^{h-1} - 2 + h - 1$$

$$\Rightarrow 2 \times 2^{h-1} - 3 + h \Rightarrow \underline{2^h - 3 + h} \rightarrow \underline{O(n)}$$



Heap \rightarrow net transaction = 0



$$TC \rightarrow O(n + n \log n) \rightarrow \underline{\underline{O(n \log n)}}$$

not
solⁿ is not in
polynomial time, but verifiable in polynomial



no of candidates

find all subsets which sum to zero

Backtracking

2, 3, -5, 1, 2, -1, 3, -5

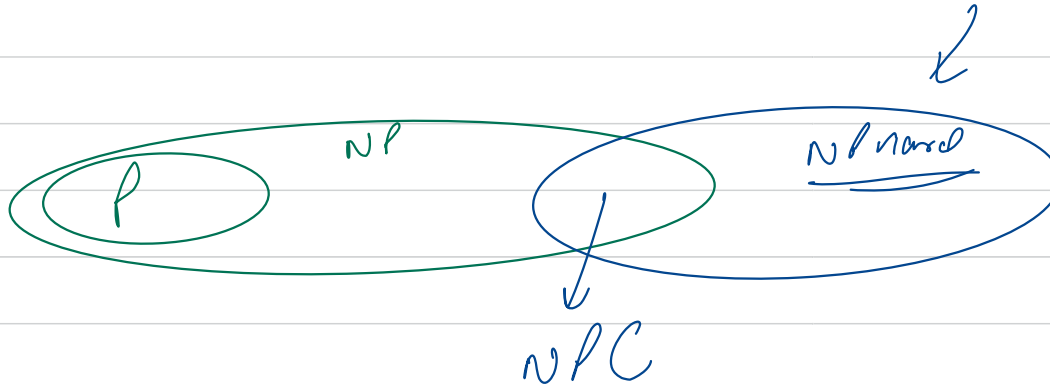
K goals

{2, 3, -5}

{1, 2, -1, 3, -5}

$P \rightarrow$
 $NP \rightarrow$

NP-Hard atleast as
 hard as
 the hardest prob
 Job



NP-C NP-C also can
 solve

So correctness is
 verifiable in polynomial
 time

install node

install yarn

yarn install

yarn start → local man

npm run build → build folder → git sub
relif