FoSAP Lecture Notes

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1 Introduction

Look at the Following problem:

```
   Input:  a string ‘w‘ consisting of 0s and 1s
Examples:  0101, 1001, 00110, 0101010
Question:  Are these requirements fullfilled?
```

→ 11 is not a sub-word

→ `w` is divisible by 3 in binary

1.1 A Possible Solution (featuring cryptic C code)

```c
int F[] = { 1, 0, 0, 0, 1, 0, 0 };
int delta[][2] = {{ 0, 1 },{ 3, 6 },{ 3, 4 },{ 2, 5 },{ 0, 6 },{ 2, 6 },{ 6, 6 }};

int three_not_11(char *w)
{
    int q = 0;
    while( *w ) q = delta[ q ][ *w++ - '0' ];
    return F[q];
}

```

1.2 Finite Automation

The above C code essentially simulates the following finite automaton:

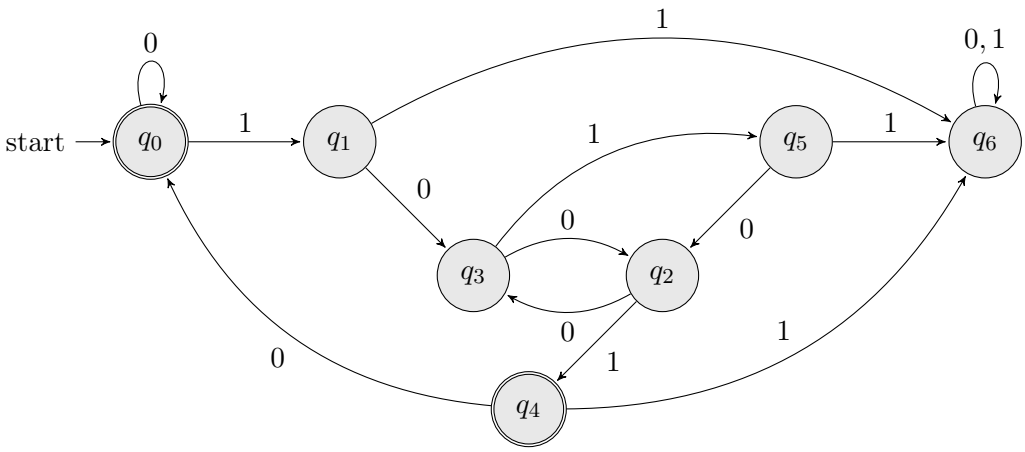

Figure 1: The Finite Automaton

This automaton works in the following way:

q_0 - the inital state, as well as an accepting state, we return true, if we end here

q_4 - an accepting state, we return true, if we land here

q_6 - the default false state, if we have `11` as subword

other - if we end on any other state, we also return false, because it isn't divisible by 3

2 Words and Languages

2.1 What is a word and what is a language?

An informal answer to the question could be, that a word w is a concatenation of symbols s in a specified alphabet Σ. A language L would then be a set of words defined by a pattern, like $\{a^{n^2} | n \geq 0\} = \{\epsilon, a, aaaa, aaaaaaaa, \dots\}$.
Words have a natural and intuitive mathematical structure. We can concatenate words, split a word, and parse sentences naturally. We can also get a section of the word and remove parts to get a new word.

$$\text{Flughafen} \rightarrow \text{Flug \& Hafen}$$

$$\text{Baumhaus} \rightarrow \text{Baum \& Haus}$$

In the case of natural language, if it uses latin characters, the alphabet and language is easily defined:

$$L = \{s^* | s \in \Sigma\} = \{a, a, they, them, you, ich, er, sie, \dots\} \tag{1}$$
$$\Sigma = \{a, b, c, \dots, x, y, z, A, B, C, \dots, X, Y, Z\} \tag{2}$$

As we can see here, the alphabet Σ is just every uppercase and lowercase latin character and our language L is no more than a set of words that have one or more character $s \in \Sigma$.

2.2 The Formal Definition

- A semigroup $(H, \circ)$ consists of a set H and an associative relation $\circ : H \times H \rightarrow H$
- A *monoid* is a semigroup with a neutral element
- Let $(M, \circ)$ be a monoid and $E \subseteq M$
  E is a generating system of $(M, \circ)$, if every $m \in M$ can be represented as $m = e_1 \circ \dots e_n$ with $e_i \in E$. A neutral element e is left and right neutral. $\forall x : e \circ x = x \circ e = x$.

A great example, using what we already established, would be to define the Alphabet Σ as the generating system of our monoid $(L, \circ)$ where $\circ$ is sefined as concatenation.