### **Operations ... SQL Specifics**



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### Certain SQL specific features

- Functions for computed attributes in projection
- Certain issues related to NULLs,
  - UNKNOWN is third TRUTH value while evaluating SQL expressions
  - IS NULL as additional predicate
- UNION/INTERSECT/EXCEPT ALL in SQL
- Correlated Sub-Queries
- EXISTS
- Some (relation as BAG) comparison operators



### **Functions and Operators in SQL**

### Functions and Operators

 SQL provides various functions and operators that can be used to create a new attribute in resultant relations

 There are typically, type conversion, arithmetic operators, mathematical, and string manipulation operators and functions. For example: substring, upper, lower, sqrt, ln, etc.

 Details for PostgreSQL functions can be seen at: <a href="http://intranet.daiict.ac.in/~pm\_jat/postgres/html/functions.html">http://intranet.daiict.ac.in/~pm\_jat/postgres/html/functions.html</a>.



```
SELECT ssn,
  fname || ' ' || minit || '. ' || lname AS name,
  current_date-bdate AS age FROM employee;

SELECT essn, hours*50 AS amount FROM works_on;

SELECT upper(fname) AS name, ln(salary) AS x FROM employee;
```



```
SELECT * FROM employee
     WHERE upper(fname) = 'FRANKLIN';

SELECT essn FROM dependent WHERE
    age(d.bdate) > interval '18 years');
```

- **BETWEEN**, **LIKE** are used in predicate:
  - SELECT .... WHERE A BETWEEN 10 TO 20;
  - SELECT ... WHERE A1 LIKE '%IX%' OR A2 LIKE 'ABC%' OR A3 LIKE '%XYZ';
  - SELECT ... WHERE A1 LIKE \\_X\_%\;

• Also: NOT BETWEEN and NOT LIKE.

# Regular Expression Matching in PostgreSQL

 PostgreSQL also allows regular expression matching in string match using IS SIMILAR TO <reg-ex>



#### **Issues with Null Values**

## Issues with Null Values

- An attribute having NULL could mean either of following-
  - Value is unknown or not available right now
  - Value is not application for the tuple: a employee not having supervisor will have null in this attribute
- Arithmetic expressions (+,-,\*,/) involving null values result null value for result

### Null Values and Comparisons

- There are more issues with Null value when it appears in a attribute used in WHERE clause of SQL-SELECT
- What should (r.a < 10) where r.a happens to be null</li>
- When we compare a NULL value with another value including NULL, result is UNKNOWN.
- In relational operations UNKNOWN is another truth-value like TRUE and FALSE

# Truth values for UNKOWN

- NOT
  - NOT UNKOWN -> UNKWON
- AND
  - TRUE AND UNKOWN —> UNKOWN
  - FALSE AND UNKOWN -> FALSE
  - UNKWON AND UNKOWNUNKOWN
- OR
  - TRUE OR UNKOWN -> TRUE
  - FALSE OR UNKOWN -> UNKOWN
  - UNKWON OR UNKOWNUNKOWN

### **Null Values and Comparisons**

- While evaluating WHERE clause tuples with UNKOWN or FALSE truth values are not included in result
- Following query will not include any tuple where either of value in NULL irrespective value in other attribute

```
SELECT * FROM EMPLOYEE WHERE
bdate < DATE '2001-01-01' AND salary > 30000
```

 Following query will not include a tuple only when both are NULL, if one of attribute meets the condition then it will get included in result

```
SELECT * FROM EMPLOYEE WHERE
bdate < DATE '2001-01-01' OR salary > 30000
```

# Null Values and Comparisons - IS NULL

Following will not give desired result. Why? -

```
SELECT * FROM employee

WHERE superssn = NULL;
```

- This is so because Null = Null is also UNKOWN. For checking an attribute for having NULL value, SQL provides IS NULL (and IS NOT NULL)
- We write as following for such situations –

```
SELECT * FROM employee
WHERE superssn IS NULL;
```



#### **Bags and Relational Operations**

# Relational Operations and multiset (or bag)

- By Definition, relations are set; but implementations may permit duplicate tuples and such relations are called bags
- Normally stored relations (base) relations should still be sets, because most relations have PK
- However SQL SELECT results are often bags, possibly because duplicate removal is expensive.
- To get set you use DISTINCT keyword

### SQL and Multiset (or Bag)

- SET operations, that are UNION, INTESECT, and EXCEPT in SQL yield their result as SET, that means duplicates are removed
- SQL however provides options by which you can have bag results by adding ALL keyword to operation name, i.e. UNION ALL, EXCEPT ALL or so.
- Let us see an example-

### UNION/INTERSECT/EXCEPT ALL in SQL

Compare result of following queries:

```
SELECT superssn FROM employee; --Q1
SELECT mgrssn FROM department; --Q2
SELECT superssn FROM employee
UNION
SELECT mgrssn FROM department; --Q3
SELECT superssn FROM employee
UNION ALL
```

SELECT mgrssn FROM department; --Q4



R	superssn
	102
	101
	101
	102
	101
	(Null)
	108
	108

S	mgrssn	
		101
		102
		108

R UNION S	superssn
	101
	102
	108
	(Null)

R UNION ALL S	superssn
	102
	101
	101
	102
	101
	(Null)
	108
	108
	101
	102
pecifics	108



R	superssn
	102
	101
	101
	102
	101
	(Null)
	108
	108

S	mgrssn	
		101
		102
		108

R EXCEPT S	superssn
	(Null)

R EXCEPT ALL S	superssn	
		101
		101
		102
		108
	(	Null)

# UNION/INTESECT/EXCEPT ALL in SQL

- UNION ALL
  - count of an element e in result is sum of count in R and S
- INTERSECT ALL
  - min(count-r, count-s) of an element in R and S, is taken as result
- EXCEPT ALL:
  - Every occurrence of an element e in S decreases its count in R by one.



### **Sub-queries in SQL**

# Subquery in SQL

 A Query that is part of another query is subquery. A subquery may also have subquery, and so forth upto any level

 A subquery in SQL is written as a query expression enclosed in parentheses, and is in following form-

"(SELECT ... FROM ...)" as a part of some existing query

Result of sub-query is again a relation;



### **Subquery in FROM clause**

 FROM clause of SQL SELECT can have a sub-query, as following-SELECT e.ssn, fname, dno, dname FROM employee AS e NATURAL JOIN (SELECT mgrssn AS ssn, dno, dname FROM department) AS dept;

 In queries like above naming to relation returned by subquery is required, even if the relation name is not used; and that is the only relation as in query below-SELECT \* FROM (SELECT mgrssn AS ssn, dno, dname FROM department) AS dept;

```
    We have seen sub-query in IN, as
    WHERE ... IN (SELECT ...)
```

For example,
 SELECT \* FROM employee WHERE ssn IN
 ( SELECT essn FROM works\_on );



### **Subquery in WHERE clause**

- When used in where clause and a sub-query returns a single column, single tuple relation, it can be interpreted as single value as following
- SELECT pname FROM project
  WHERE dno = (SELECT dno FROM department
  WHERE dname = 'Research');
- Note: resultant relation of sub-query is getting compared with a attribute value – appropriate conversion takes place;
  - Has a underlying assumption that sub-query return a single tuple

## **Execution of Subquery**

- SUB-Query may not execut for every tuple of outer query
- Consider the query below-

```
SELECT pname FROM project
WHERE dno = (SELECT dno FROM department
WHERE dname = 'Research');
```

It is to typically executed as following: Execute inner query, and let us say returns 5, and then places it in outer query, and the query to be executed becomes following SELECT pname FROM project
 WHERE dno = 5;

## **Execution of Subquery**

- SUB-Query may not execute for every tuple of outer query
- Consider another query-

 Typically, after execution of inner query, outer query may be translated to:

```
SELECT * FROM student WHERE progid IN (BEC, BEE);
```



#### **Correlated Sub-Queries**

### Correlated Sub-Queries

 When inner query makes a reference to tuple of outer query then it is correlated sub-query. Consider following query -

 List employees, whose salary is more than department average:

```
SELECT ssn, fname FROM employee as e WHERE salary > (SELECT AVG(salary) FROM employee WHERE dno = e.dno)
```



#### **Execution of Correlated Sub-Queries**

Consider same query

```
SELECT ssn, fname FROM employee as e WHERE salary > (SELECT AVG(salary) FROM employee WHERE dno = e.dno)
```

- Logically, it is as following: For each tuple of outer query, execute inner query.
- Note that it can not be executed once for all tuples of outer query, as the case be with un-related inner query, and we have to execute SUB-Query for every tuple of outer query
- This is identified problem with correlated sub-queries.



### Correlated Sub-Queries could be expensive to execute – therefore should be avoided

- Correlated queries are expensive to execute, and can be avoided; for example the previous example
- SELECT ssn, fname FROM employee as e WHERE salary > (SELECT AVG(salary)
   FROM employee WHERE dno = e.dno)
- can be re-written as-SELECT ssn, fname, salary FROM employee as e NATURAL JOIN (SELECT dno, AVG(salary) as avg\_sal FROM employee GROUP BY dno) as av WHERE salary > av.avg\_sal;



### more Correlated Sub-queries

• List down employees having salary greater than their immediate supervisors.

```
select * from employee as e1 where e1.salary >
  (select salary from employee as e2 where e2.ssn =
  e1.superssn);
```

Select employees having dependents older than 18 years:
 SELECT \* FROM employee AS e WHERE ssn IN (SELECT essn FROM dependent AS d WHERE essn = e.ssn AND age(d.bdate) > interval '18 years');

Attempt re-writting them without correlated query.

### EXISTS and NOT EXISTS in SQL

- Checks for emptiness of a relation and returns true or false.
- EXISTS(r) can be interpreted as "is there some tuple exists in relation r"
- EXISTS(r) returning true says that argument relation r is not empty
- Similarly, NOT EXISTS(r) returning true says that argument relation r empty

### Example EXISTS

SELECT ssn FROM employee AS e WHERE EXISTS
 (SELECT \* FROM dependent AS d WHERE d.essn
 = e.ssn AND age(d.bdate) > interval '18
 years');

## SQL- EXISTS and IN

- While they might appear to be serving similar purposes, semantically are different.
- Both appear as part of predicate in WHERE clause of SELECT
- IN:
  - Syntax: x IN ( r )
  - Meaning: checks existence of tuple x in relation r, if found returns true, other wise false. Normally x is a scalar value and r is a single column relation.
- EXISTS:
  - Syntax: EXISTS ( r )
  - Meaning: checks if r is a non empty relation. Returns true if the relation has at least one tuple, otherwise false.
- In both above cases  $\mathbf{r}$  is a *relational expression* resulting a relation.

## DA-IICT DA-IICT DA TONES

### Compare a values with a bag of values (SQL)

 For example consider following two queries [Find out employee who have salary greater some or all employees of dno = 4]

SELECT ssn, fname FROM employee WHERE salary

> **SOME** (SELECT salary FROM employee WHERE dno = 4);

SELECT ssn, fname FROM employee WHERE salary

> ALL (SELECT salary FROM employee WHERE dno = 4);

### Compare a values with a bag of values (SQL)

- Note the equivalences:
  - SELECT ssn, fname FROM employee WHERE salary
  - > SOME (SELECT salary FROM employee WHERE dno = 4); and SELECT ssn, fname FROM employee WHERE salary
    - > (SELECT min(salary) FROM employee WHERE dno = 4);

SELECT ssn, fname FROM employee WHERE salary

- > ALL (SELECT salary FROM employee WHERE dno = 4); and SELECT ssn, fname FROM employee WHERE salary
  - > (SELECT max(salary) FROM employee WHERE dno = 4);

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### Compare a values with a bag of values (SQL)

Comparative operators could be, one of following-

```
>SOME, >=SOME, <=SOME, <SOME, =SOME, <>SOME
>ALL, >=ALL, <=ALL, =ALL, <>ALL
```

- Note: it can be easily proved that
  - **=SOME** is identical to **IN**, and
  - <>some is not identical to NOT IN
  - = ALL is not identical to IN
  - **ALL** (mean = NONE) and is same as **NOT** IN, and Earlier versions of SQL used **ANY** for **SOME**; today both keywords are used as synonymous.



### **Sub-queries in Update statements**

```
• UPDATE employee
   SET salary = salary * 1.1
   WHERE ssn IN ( ... );
```

```
• DELETE employee

WHERE ssn = ( ... );
```