

## 6.824 2018 Lecture 20: Bitcoin

Bitcoin: A Peer-to-Peer Electronic Cash System, by Satoshi Nakamoto, 2008

bitcoin:

- a digital currency
- a public ledger to prevent double-spending
- no centralized trust or mechanism <-- this is hard!
- malicious users ("Byzantine faults")

why might people want a digital currency?

- might make online payments easier
- credit cards have worked well but aren't perfect
  - insecure -> fraud -> fees, restrictions, reversals
- record of all your purchases

what's hard technically?

- forgery
- double spending
- theft

what's hard socially/economically?

- why do Bitcoins have value?
- how to pay for infrastructure?
- monetary policy (intentional inflation &c)
- laws (taxes, laundering, drugs, terrorists)

idea: signed sequence of transactions

- (this is the straightforward part of Bitcoin)
- there are a bunch of coins, each owned by someone
- every coin has a sequence of transaction records
  - one for each time this coin was transferred as payment
- a coin's latest transaction indicates who owns it now

what's in a transaction record?

- pub(user1): public key of new owner
- hash(prev): hash of this coin's previous transaction record
- sig(user2): signature over transaction by previous owner's private key
- (Bitcoin is much more complex: amount (fractional), multiple in/out, ...)

transaction example:

- Y owns a coin, previously given to it by X:
  - T7: pub(Y), hash(T6), sig(X)
- Y buys a hamburger from Z and pays with this coin
  - Z sends public key to Y
  - Y creates a new transaction and signs it
    - T8: pub(Z), hash(T7), sig(Y)
  - Y sends transaction record to Z
  - Z verifies:
    - T8's sig() corresponds to T7's pub()
  - Z gives hamburger to Y

Z's "balance" is set of unspent transactions for which Z knows private key

- the "identity" of a coin is the (hash of) its most recent xaction
- Z "owns" a coin = Z knows private key for "new owner" public key in latest xaction

can anyone other than the owner spend a coin?

- current owner's private key needed to sign next transaction
- danger: attacker can steal Z's private key, e.g. from PC or smartphone

can a coin's owner spend it twice in this scheme?

- Y creates two transactions for same coin: Y->Z, Y->Q
  - both with hash(T7)
- Y shows different transactions to Z and Q
- both transactions look good, including signatures and hash
- now both Z and Q will give hamburgers to Y

why was double-spending possible?

- b/c Z and Q didn't know complete set of transactions

what do we need?

- publish log of all transactions to everyone, in same order
  - so Q knows about Y->Z, and will reject Y->Q
- a "public ledger"
- ensure Y can't un-publish a transaction

why not rely on CitiBank, or Federal Reserve, to publish transactions?  
not everyone trusts them  
they might be tempted to reverse or restrict

why not publish transactions like this:  
1000s of peers, run by anybody, no trust required in any one peer  
peers flood new transactions over "overlay"  
transaction Y->Z only acceptable if majority of peers think it is valid  
i.e. they don't know of any Y->Q  
hopefully majority overlap ensures double-spend is detected  
how to count votes?  
how to even count peers so you know what a majority is?  
perhaps distinct IP addresses?  
problem: "sybil attack"  
IP addresses are not secure -- easy to forge  
attacker pretends to have 10,000 computers -- majority  
when Z asks, attacker's majority says "we only know of Y->Z"  
when Q asks, attacker's majority says "we only know of Y->Q"  
voting is hard in "open" p2p schemes

the BitCoin block chain  
the goal: agreement on transaction log to prevent double-spending  
the block chain contains transactions on all coins  
many peers  
each with a complete copy of the whole chain  
proposed transactions flooded to all peers  
new blocks flooded to all peers  
each block:  
hash(prevblock)  
set of transactions  
"nonce" (not quite a nonce in the usual cryptographic sense)  
current time (wall clock timestamp)  
new block every 10 minutes containing xactions since prev block  
payee doesn't believe transaction until it's in the block chain

who creates each new block?  
this is "mining"  
all peers try  
requirement: hash(block) has N leading zeros  
each peer tries nonce values until this works out  
trying one nonce is fast, but most nonces won't work  
it's like flipping a zillion-sided coin until it comes up heads  
each flip has an independent (small) chance of success  
mining a block \*not\* a specific fixed amount of work  
it would likely take one CPU months to create one block  
but thousands of peers are working on it  
such that expected time to first to find is about 10 minutes  
the winner floods the new block to all peers

how does a Y->Z transaction work w/ block chain?  
start: all peers know ...<-B5  
and are working on B6 (trying different nonces)  
Y sends Y->Z transaction to peers, which flood it  
peers buffer the transaction until B6 computed  
peers that heard Y->Z include it in next block  
so eventually ...<-B5<-B6<-B7, where B7 includes Y->Z

Q: could there be \*two\* different successors to B6?  
A: yes, in (at least) two situations:  
1) two peers both get lucky (unlikely, given variance of block time)  
2) network partition  
in both cases, the blockchain temporarily forks  
peers work on whichever block they heard about before  
but switch to longer chain if they become aware of one

how is a forked chain resolved?  
each peer initially believes whichever of BZ/BQ it saw first  
tries to create a successor  
if many more saw BZ than BQ, more will mine for BZ,  
so BZ successor likely to be created first  
if exactly half-and-half, one fork likely to be extended first  
since significant variance in mining success time  
peers always switch to mining the longest fork, re-inforcing agreement

what if Y sends out Y->Z and Y->Q at the same time?  
i.e. Y attempts to double-spend  
no correct peer will accept both, so a block will have one but not both

what happens if Y tells some peers about Y->Z, others about Y->Q?  
perhaps use network DoS to prevent full flooding of either  
perhaps there will be a fork: B6<-BZ and B6<-BQ

thus:

- temporary double spending is possible, due to forks
- but one side or the other of the fork highly likely to disappear
- thus if Z sees Y->Z with a few blocks after it,
  - it's very unlikely that it could be overtaken by a different fork containing Y->Q
- if Z is selling a high-value item, Z should wait for a few blocks before shipping it
- if Z is selling something cheap, maybe OK to wait just for some peers to see Y->Z and validate it (but not in block)

can an attacker modify a block in the middle of the block chain?  
not directly, since subsequent block holds block's hash

could attacker start a fork from an old block, with Y->Q instead of Y->Z?  
yes -- but fork must be longer in order for peers to accept it  
so if attacker starts N blocks behind, it must generate N+M+1 blocks on its fork before main fork is extended by M  
i.e. attacker must mine blocks *\*faster\** than the other peers  
with just one CPU, will take months to create even a few blocks  
by that time the main chain will be much longer  
no peer will switch to the attacker's shorter chain  
if the attacker has 1000s of CPUs -- more than all the honest bitcoin peers -- then the attacker can create the longest fork, everyone will switch to it, allowing the attacker to double-spend

there's a majority voting system hiding here  
peers cast votes by mining to extend the longest chain

summary:

- if attacker controls majority of CPU power, can force honest peers to switch from real chain to one created by the attacker
- otherwise not

validation checks:

peer, new xaction:

- no other transaction spends the same previous transaction
- signature is by private key of pub key in previous transaction
- then will add transaction to txn list for next block to mine

peer, new block:

- hash value has enough leading zeroes (i.e. nonce is right, proves work)
- previous block hash exists
- all transactions in block are valid
- peer switches to new chain if longer than current longest

Z:

- (some clients rely on peers to do above checks, some don't)
- Y->Z is in a block
- Z's public key / address is in the transaction
- there's several more blocks in the chain
- (other stuff has to be checked as well, lots of details)

where does each bitcoin originally come from?

- each time a peer mines a block, it gets 12.5 bitcoins (currently)
- it puts its public key in a special transaction in the block
- this is incentive for people to operate bitcoin peers

Q: what if lots of miners join, so blocks are created faster?

Q: 10 minutes is annoying; could it be made much shorter?

Q: are transactions anonymous?

Q: if I steal bitcoins, is it safe to spend them?

Q: can bitcoins be forged, i.e. a totally fake coin created?

Q: what can adversary do with a majority of CPU power in the world?  
can double-spend and un-spend, by forking  
cannot steal others' bitcoins  
can prevent xaction from entering chain

Q: what if the block format needs to be changed?

esp if new format wouldn't be acceptable to previous s/w version?

Q: how do peers find each other?

Q: what if a peer has been tricked into only talking to corrupt peers?  
how about if it talks to one good peer and many colluding bad peers?

Q: could a brand-new peer be tricked into using the wrong chain entirely?  
what if a peer rejoins after a few years disconnection?  
a few days of disconnection?

Q: how rich are you likely to get with one machine mining?

Q: why does it make sense for the mining reward to decrease with time?

Q: is it a problem that there will be a fixed number of coins?  
what if the real economy grows (or shrinks)?

Q: why do bitcoins have value?  
e.g., 1 BTC appears to be around \$8700 on may 14 2018.

Q: will bitcoin scale well?  
as transaction rate increases?  
claim CPU limits to 4,000 tps (signature checks)  
more than Visa but less than cash  
as block chain length increases?  
do you ever need to look at very old blocks?  
do you ever need to xfer the whole block chain?  
merkle tree: block headers vs txn data.  
sadly, the maximum block size is limited to one megabyte

Q: could Bitcoin have been just a ledger w/o a new currency?  
e.g. have dollars be the currency?  
since the currency part is pretty awkward.  
(settlement... mining incentive...)

key idea: block chain  
public ledger is a great idea  
decentralization might be good  
mining is a clever way to avoid sybil attacks  
tying ledger to new currency seems awkward, maybe necessary