```
Graph Theory
Bridge:
vector<int>vc[100001];
int vis[100001];
int cnt,in[100001],val[100001];
set< pair < int, int > > bridge;
void dfs(int v,int p) {
 vis[v] = 1;
 in[v] = val[v] = cnt++;
 REP(i, vc[v].size()) {
   if(vc[v][i] == p)continue;
   if(vis[vc[v][i]])val[v] = min(val[v], in[vc[v][i]]);
   else{
     dfs(vc[v][x], v);
     val[v] = min(val[v], val[vc[v][i]]);
     if (val[vc[v][x]] > in[v]){
       pair<int,int> pp = make_pair(v,vc[v][x]);
       if(pp.xx > pp.yy)swap(pp.xx, pp.yy);
                              bridge.insert(pp);
       }
    }
  }
Articulation Point:
vector<int>vc[100001];
int vis[100001];
int cnt = 0,in[100001],val[100001];
set< int > apNodes;
void dfs(int v,int p){
vis[v] = 1;
in[v] = val[v] = cnt++;
int c=0;
REP(i, vc[v].size()){
 if(vc[v][i] == p)continue;
 if(vis[vc[v][i]])val[v] = min(val[v], in[vc[v][i]]);
 else{
   dfs(vc[v][x],v);
   val[v]=min(val[v], val[vc[v][i]]);
   if (val[vc[v][i]] >= in[v] && p!=-1)
       apNodes.insert(v);
   C++;
if(p==-1 \&\& c>1) apNodes.insert(v);
Euler's theorem. For any planar graph, V - E + F =
1 + C, where V is the number of graph's
vertices, E is the number of edges, F is the number of // havel-hakimi ends here. f = 1 shows invalid graph
```

faces in graph's planar drawing, and C is the number of connected components. Corollary: V – E + F = 2 for a 3D polyhedron.

Vertex covers and independent sets. Let M, C, I be a max matching, a min vertex cover, and a max independent set. Then |M| |C| = N - |I|, with equality for bipartite graphs. Complement of an MVC is always a MIS, and vice versa. Given a bipartite graph with partitions (A, B), build a network: connect source to A, and B to sink with edges of capacities, equal to the corresponding nodes' weights, or 1 in the unweighted case. Set capacities of the original graph's edges to the infinity. Let (S, T) be a minimum s-t cut. Then a maximum(-weighted) independent set is I = $(A \cap S) \cup (B \cap T)$, and a minimum(-weighted) vertex cover is $C = (A \cap T) \cup (B \cap S)$.

Matrix-tree theorem. Let matrix $T = [t \ ij]$, where t ij is the number of multiedges between i and j, for i = j, and t ii = -deg i. Number of spanning trees of a graph is equal to the determinant of a matrix obtained by deleting any k-th row and k-th column from T.

Euler tours. Euler tour in an undirected graph exists iff the graph is connected and each vertex has an even degree. Euler tour in a directed graph exists iff in-degree of each vertex equals its out-degree, and underlying undirected graph is connected. Construction:

doit(u):

for each edge e = (u, v) in E, do: erase e, doit(v) prepend u to the list of vertices in the tour.

Havel Hakimi:

```
f = 0;
REP(i, b)
 sort(tmp + j, tmp + b);
 reverse(tmp + j, tmp + b);
 if(tmp[j] < 0){
    f = 1;
    break;
 for(int k = 1; k \le tmp[i]; k++){
     if(j + k >= b){
        f = 1;
        break;
     tmp[i + k]--; } }
```

```
Stable Marriage:
                                                           memset (d, -1, size of d); d[s] = 0;
scanf("%d", &n);
                                                        while (qh < qt && d[t] == -1)
                                                          int v = q [qh ++];
for( int i = 0; i < n; i++ ) {
       for( int j = 0; j < n; j++ ) {
                                                          for ( size_t i = 0; i < g [ v ] . size ( ); ++ i ) {
              scanf("%d", &mPref[i][j]);
                                                            int id = g [ v ] [ i ],
                                                           to = e [ id ] . b;
                                                      if (d [to] == -1 \&\& e [id]. flow < e [id]. cap) {
for( int i = 0; i < n; i++) {
                                                        q [qt ++] = to;
       for( int j = 0; j < n; j++ ) {
                                                        d[to] = d[v] + 1;  }
              int k:
                                                        return d [ t ] != - 1; }
              scanf("%d", &k);
                                                        int dfs (int v, int flow) {
              wPref[i][k-1] = n - j;
                                                           if (! flow ) return 0;
                                                          if (v == t) return flow;
                                                      for (; ptr [v] < (int) g [v]. size (); ++ ptr [v]) {
memset(R, -1, sizeof(R));
                                                             int id = g[v][ptr[v]], to = e[id].b;
memset( P, 0, sizeof(P) );
                                                             if (d[to]!=d[v]+1) continue;
for( int i = 0; i < n; i++ ) {
                                                      int pushed = dfs(to,min(flow,e[id].cap-e[id].flow ) );
       int man = i;
                                                             if (pushed) {
       while (man \ge 0)
                                                                e [ id ] . flow += pushed;
          int woman = mPref[man][ P[man]++ ];
                                                                e [ id \land 1 ] . flow -= pushed;
    if( R[woman] == -1 || wPref[woman]
                                                                return pushed; } }
[ R[woman] ] < wPref[woman][ man ] ) {
                                                           return 0;
       L[man] = woman;
                                                        }
       swap( man, R[woman] );
                                                      int dinic () {
                                                          int flow = 0;
                                                           for (;;) {
                                                      if (!bfs ()) break;
printf("Case %d:", ++caseno);
                                                             memset (ptr, 0, size of ptr);
for(int i=0;i<n; i++) printf("%d %d", i+1, L[i]+n+1);
                                                             while (int pushed = dfs (s, INF)) {
                                                                flow += pushed;
             FLOW
Dinic:
                                                                if(pushed == 0)break;
                                                             }
struct edge {
int a, b, cap, flow;
edge(int _a,int _b,int _c,int _d) {
                                                           return flow;
  a=_a,b=_b,cap=_c,flow=_d; } };
                                                        }
int INF = 1500000001;
                                                      MCMF:
int n,s,t,d[30001], ptr [30001], q [30001 * 10];
                                                      const int MAXN = 7210;
vector < edge > e,g [ 30001 ];
                                                      const int MAXM = 100010;
void add_edge ( int a, int b, int cap ) {
                                                      const int INF = 1000000000;
edge e1 =edge(a, b, cap,0); edge e2=edge(b,a,0, 0);
                                                      struct Edge
g[a] . push_back ( ( int ) e. size ( ) );
e. push_back (e1);
                                                        int u, v, cap, cost;
  g[b].push_back((int)e.size());
                                                        int next;
    e. push_back (e2);
                                                             Edge(){}
                                                             Edge(int u,int v,int cap,int cost,int
bool bfs () {
                                                      next):u(u),v(v),cap(cap),cost(cost),next(next){}
    int qh = 0, qt = 0;
                                                      }edge[MAXM*3];
    q [qt ++] = s;
```

```
int NE;
int head[MAXN], dist[MAXN], pp[MAXN];
                                                              }
bool vis[MAXN];
void init(int n)
                                                           if(dist[t] == INF) return false;
                                                           return true;
  NE = 0;
  fill(head,head+n+1,-1);
                                                        pair<int,int> MCMF(int s, int t, int n) //
                                                        minCostMaxFlow
void addedge(int u, int v, int cap, int cost)
                                                           int flow = 0; // 总流量
  edge[NE] = Edge(u,v,cap,cost,head[u]);
                                                           int i, minflow, mincost;
       head[u] = NE++;
                                                           mincost = 0;
  edge[NE] = Edge(v,u,0,-cost,head[v]);
                                                          for(;SPFA(s, t, n);)
  head[v] = NE++;
                                                             minflow = INF + 1;
int qu[300010];
                                                              for(i = pp[t]; \sim i; i = pp[edge[i].u])
inline bool Min(int &x,int y)
                                                                  Min( minflow, edge[i].cap);
                                                              flow += minflow;
  if(y < x) {x = y;return true;}
                                                              for(i = pp[t]; \sim i; i = pp[edge[i].u])
  return false;
                                                                edge[i].cap -= minflow;
bool SPFA(int s, int t, int n)
                                                                edge[i^1].cap += minflow;
  int i, u, v, tmp;
                                                             mincost += dist[t] * minflow;
  int he = 0, ta = 0;
  fill(vis,vis+n+1,false);
                                                           return make_pair(flow,mincost);
  fill(pp,pp+n+1,-1);
  fill(dist,dist+n+1,INF);
                                                        BPM:
                                                        int pre[201], vis[201];
  vis[s] = true; dist[s] = 0;
  qu[++ta] = s;
                                                        vector<int>vc[201];
  while(he!=ta)
                                                        int find(int id) {
                                                           int x,y,z;
    u = qu[++he]; vis[u] = false;
                                                           for(x=0;x<vc[id].size();x++) {
    if(he \geq 250000) he = 0;
                                                             if(vis[vc[id][x]]==0) {
     for(i = head[u]; \sim i; i = edge[i].next)
                                                                vis[vc[id][x]]=1;
                                                                if(pre[vc[id][x]] = -1 \parallel find(pre[vc[id][x]])) 
       v = edge[i].v;
                                                                   pre[vc[id][x]]=id;
                      tmp = dist[u] + edge[i].cost;
                                                                   return 1; } }
       if(edge[i].cap && dist[v] > tmp)
                                                           return 0;
          dist[v] = tmp;
                                                        int bpm(int n) {
          pp[v] = i;
                                                           int a=0,x;
          if(!vis[v])
                                                           memset(pre,-1,sizeof pre);
                                                           for(x=1;x\leq n;x++) {
                                                             memset(vis,0,sizeof vis);
            qu[++ta] = v;
            if(ta \geq 250000) ta = 0;
                                                             a+=find(x);
            vis[v] = true;
                                                        } return a;
          }
```

```
STRING
                                                        Manacher:
KMP:
                                                        void pre_process(){
                                                           s = "%";
int k,l,c; // 1 indexed
char txt[1000001],pat[10000001];
                                                           REP(i,tmp.size()){
int fnd[1000001],fn[1000001];
                                                             s += '\#':
void comp()
                                                             s += tmp[i];
  fn[1]=0;
                                                           s += "#\&";
  int m=0;
  for(int x=2;x\leq l;x++)
                                                        LL P[300001];
                                                        void manacher(){ // manacher's algo - collected
                                                        from internet
    while(m>0 && pat[m+1]!=pat[x])m=fn[m];
                                                           LL c = 0, r = 0;
    if(pat[m+1]==pat[x])m++;
    fn[x]=m;
                                                           FOR(i, 1, s.size() - 1){
  }
                                                             LL i mirror = 2 * c - i;
                                                             P[i] = (r > i) ? min(r-i, P[i\_mirror]) : 0;
void job()
                                                             while (s[i + 1 + P[i]] == s[i - 1 - P[i]])P[i]++;
                                                             if (i + P[i] > r) {
  int b,d,e,f,x,y,z;
                                                                c = i;
  b=0:
                                                                r = i + P[i];
  for(x=1;x<=k;x++)
                                                           }
     while(b>0 && pat[b+1]!=txt[x])b=fn[b];
                                                        }
                                                        Z-Algo:
    if(pat[b+1]==txt[x])b++;
    if(b==1)
                                                        char s[1000003];
                                                        int zz[1000003];
       C++;
       b=fn[b];
                                                        void Z_algo(int n) {
                                                           int L = 0, R = 0;
  }
                                                           for (int i = 1; i < n; i++) {
                                                             if (i > R) {
Finite Automata:
                                                                L = R = i;
int TF[1001][27];
                                                                while (R < n \&\& s[R-L] == s[R]) R++;
char pat[1001];
                                                                zz[i] = R-L;
void computeFA(int M){
                                                                R--;
  int lps = 0;
                                                             } else {
                                                                int k = i-L;
  for(int i = 0; i < 26; i++) TF[0][i] = 0;
                                                                if (zz[k] < R-i+1) zz[i] = zz[k];
  TF[0][pat[0] - 'a'] = 1;
                                                                else {
                                                                  L = i:
  for(int i = 1; i < M; i++){
                                                                  while (R < n \&\& s[R-L] == s[R]) R++;
    for(int j = 0; j < 26; j++) TF[i][j] = TF[lps][j];
                                                                  zz[i] = R-L;
    TF[i][pat[i] - 'a'] = i + 1;
                                                                  R--;
    if(i < M) lps = TF[lps][pat[i] - 'a'];
                                                                }
  }
                                                             }
}
                                                           }
```

```
Suffix Array:
                                                                                                                         }
int P[20][100001], N, stp, cnt;
char A[100001];
                                                                                                                   int LCA(int p, int q){
                                                                                                                        if(lvl[p] < lvl[q])swap(p,q);
struct info {
     int nr[2], p;
                                                                                                                         int log;
     bool operator < ( const info &b ) const {
                                                                                                                         for(log = 1; (1 << log) <= lvl[p]; log++);
          return nr[0] == b.nr[0] ? (nr[1] < b.nr[1] ) :
(nr[0] < b.nr[0]);
                                                                                                                         for(int i = log; i >= 0; i--){
                                                                                                                              if(lvl[p] - (1 << i) >= lvl[q])p = P[p][i];
} L[100001];
                                                                                                                        if(p == q)return p;
void suffixArray() {
                                                                                                                         for(int i = log; i >= 0; i--){
     for( int i = 0; i < N; i++) P[0][i] = A[i] - 'a';
                                                                                                                              if(P[p][i] != -1 && P[p][i] != P[q][i]){
     for( stp = 1, cnt = 1; cnt >> 1 < N; ++stp, cnt <<=
                                                                                                                                   p = P[p][i];
1){
                                                                                                                                   q = P[q][i];
          for( int i = 0; i < N; i++) {
               L[i].nr[0] = P[stp - 1][i];
               L[i].nr[1] = i + cnt < N ? P[stp - 1][i + cnt] :
                                                                                                                        return pre[p];
-1;
               L[i].p = i;
                                                                                                                   DIVIDE AND CONQUER:
                                                                                                                   CONVEX HULL DP
                                                                                                                   int M[501][20001];
          sort(L, L + N);
                                                                                                                   int B[501][20001];
          for ( int i = 0; i < N; i++)
               P[stp][L[i].p] = i > 0 \&\& L[i].nr[0] == L[i -
                                                                                                                   int bad(int l1,int l2,int l3,int l){
1].nr[0] && L[i].nr[1] == L[i-1].nr[1] ? P[stp][L[i-1].nr[1] ? P[stp][L[i-1].nr[1] ? P[stp][L[i-1].nr[1] ? P[stp][L[i-1].nr[1]] ? P[stp
                                                                                                                        return (B[l][l3]-B[l][l1])*(M[l][l1]-M[l][l2])<(B[l]
1].p]:i;
                                                                                                                   [l2]-B[l][l1])*(M[l][l1]-M[l][l3]);
                                                                                                                   }
     }
                                                                                                                   int cnt;
                                                                                                                   void insert_line(int m,int b,int l){
int lcp(int x, int y) {
                                                                                                                         M[l][cnt]=m;
     int k, ret = 0;
                                                                                                                         B[l][cnt++]=b;
                                                                                                                         while (cnt>=3 && bad(cnt-3,cnt-2,cnt-1,l)>0){
     if(x == y) return N - x;
                                                                                                                              swap(M[l][cnt-1],M[l][cnt-2]);
     for(k = stp - 1; k \ge 0 && x \le N && y \le N; k--)
if(P[k][x] == P[k][y])
                                                                                                                              swap(B[l][cnt-1],B[l][cnt-2]);
          x += 1 << k, y += 1 << k, ret += 1 << k;
                                                                                                                              cnt--;
                                                                                                                         }
     return ret;
                                                                                                                   }
                              DATA STRUCTURE
                                                                                                                   int p;
LCA:
                                                                                                                   int query(int x,int l)
int pre[100001],lvl[100001];
int P[100001][22];
                                                                                                                        if(p>=cnt)p=cnt-1;
void pre_process(int N){
                                                                                                                         while(p<cnt-1 && M[l][p+1]*x+B[l][p+1]<M[l]
     memset(P,-1,sizeof P);
                                                                                                                   [p]*x+B[l][p])p++;
     for(int i = 1; i \le N; i++)P[i][0] = pre[i];
                                                                                                                        return M[l][p]*x+B[l][p];
     for(int j = 1; (1 << j) <= N; j++){
          for(int i = 1; i \le N; i++){
          if(P[i][j-1]!=-1) P[i][j]=P[P[i][j-1]][j-1];
```

```
SHANKS:
LL solve (LL a, LL b, LL m) {
                                                         i64 myfloor( i64 a, i64 b ) {
       LL n = (LL) sqrt (m + .0) + 1;
                                                                 i64 c = a / b:
                                                                 if( (a \% b) \&\& a < 0 ) c--;
       LL an = 1:
                                                                 return c:
       for ( int i = 0; i < n; ++ i)
                                                         }
               an = (an * a) \% m;
                                                         i64 myceil( i64 a, i64 b ) {
       map < LL, LL > vals;
                                                                 i64 c = a / b;
       for ( int i = 1 , cur = an; i \le n; ++ i ) {
                                                                 if( (a \% b) \&\& a > 0 ) c++;
               if (! vals. count (cur))
                                                                 return c:
                                                         }
                       vals [ cur ] = i;
               cur = (cur * an) % m:
                                                         i64 solve() {
       LL ret = m;
                                                                 i64 a, b, c, x1, x2, y1, y2, X, Y, n1, n2, m1,
       for ( int i = 0 , cur = b ; i \le n ; ++ i ) {
                                                         m2;
               if (vals. count (cur)) {
                                                                 scanf("%lld %lld %lld %lld %lld %lld %lld",
                       int ans = vals [ cur ] * n - i ;
                                                         &a, &b, &c, &x1, &x2, &y1, &y2);
                       if (ans < ret)
                              ret = ans;
                                                                 c *= -1:
                                                                 if( a < 0 ) a *= -1, x1 *= -1, x2 *= -1,
               cur = (cur * a) % m;
                                                         swap(x1, x2);
                                                                 if(b < 0) b *= -1, v1 *= -1, v2 *= -1,
       if(ret < m)return ret;
                                                         swap(y1, y2);
       return - 1;
                                                                 if(!a &&!b) return!c?(x2 - x1 + 1) * (y2 -
MATH:
                                                         y1 + 1) : 0;
Extended Euclid:
                                                                 if(b == 0) {
typedef long long i64;
                                                                        if( c % a ) return 0;
int cases, caseno;
                                                                        i64 x = c / a;
                                                                        return (x \ge x1 & x \le x2) * (y2 - x2)
struct Euclid {
                                                         y1 + 1);
       i64 x, y, d;
                                                                 if( a == 0 ) {
       Euclid() {}
       Euclid( i64 xx, i64 yy, i64 dd ) { x = xx, y =
                                                                        if( c % b ) return 0;
yy, d = dd; 
                                                                        i64 v = c / b;
};
                                                                        return ( y \ge y1 & y \le y2 ) * (x2 - y2)
                                                         x1 + 1);
Euclid egcd( i64 a, i64 b ) {
       if(!b) return Euclid(1, 0, a);
       Euclid r = \text{egcd} (b, a \% b);
                                                                 Euclid s;
       return Euclid( r.y, r.x - a / b * r.y, r.d );
}
                                                                 s = egcd(a, b);
i64 abs64( i64 a ) {
                                                                 if( c % s.d ) return 0;
       return a > 0? a : -a;
                                                                 a = s.d;
}
```

```
b = s.d;
                                                   res += bigmod( rest % P[i], P[i] - 2, P[i] ) * R[i]
       c = s.d;
                                                   * rest;
       s.d = 1;
                                                                 res %= lcm;
      X = s.x * c;
                                                                 printf("Case %d: %lld\n", +
       Y = s.y * c;
                                                   +caseno, res);
      n2 = min(myfloor(X - x1, b), myfloor(y2 -
                                                   Matrix Expo:
Y, a));
                                                   LL mul_res[MX][MX];
      n1 = -min(myfloor(Y - y1, a), myfloor(x2 -
                                                   void multiply(LL a[][MX], LL b[][MX]){
X, b));
                                                      set0(mul_res);
                                                      for(int i = 0; i < MX; i++){
      return (n2 < n1)? 0: n2 - n1 + 1;
                                                        for(int j = 0; j < MX; j++){
Chinese Remainder Theorem:
                                                           for(int k = 0; k < MX; k++)
i64 P[NN], R[NN];
                                                             mul_res[i][j] = (mul_res[i][j] + a[i][k]
int cases, caseno, n;
                                                   * b[k][j]) % mod;
                                                        }
i64 bigmod( i64 a, i64 p, i64 MOD ) {
                                                      }
      i64 res = 1;
                                                   }
       while(p) {
             if( p \& 1 ) res = (res * a) % MOD;
                                                   LL big_res[MX][MX];
             a = (a * a) \% MOD;
                                                   LL init[MX][MX];
             p >>= 1;
                                                   void matrix_expo(LL n){
                                                      if(n == 0){
      return res;
                                                        set0(big_res);
                                                        REP(i, MX)big\_res[i][i] = 1;
                                                        return;
int main() {
      freopen("a.in", "r", stdin);
                                                      if(n \% 2 == 0){
      freopen("a.ans", "w", stdout);
                                                        matrix_expo(n / 2);
                                                        multiply(big_res, big_res);
      double cl = clock();
                                                        REP(i, MX){
                                                           REP(j, MX)big\_res[i][j] = mul\_res[i][j];
      scanf("%d", &cases);
       while( cases-- ) {
                                                      }
             scanf("%d", &n);
                                                      else {
             for( int i = 0; i < n; i++)
                                                       matrix expo(n - 1);
scanf("%lld %lld", &P[i], &R[i]);
                                                         multiply(big_res, init);
                                                         REP(i, MX){
             i64 lcm = 1, res = 0;
                                                         REP(j, MX)big_res[i][j] = mul_res[i][j];
             for( int i = 0; i < n; i++) {
lcm *= P[i];
                                                      }
                    i64 rest = 1:
                    for( int j = 0; j < n; j++ ) if( i
                                                          n <<= 1; /// not needed for this problem
!= j) rest = rest * P[j];
```

```
FFT:
typedef complex < double > base;
void fft ( vector < base > & a, bool invert ) {
      int n = (int) a. size();
      for (int i = 1, j = 0; i < n; ++ i) {
             int bit = n >> 1:
             for (; j \ge bit; bit \ge 1)
                    j -= bit ;
             i += bit;
             if (i < j)
                    swap(a[i],a[j]);
       }
      for ( int len = 2; len \leq n; len \leq 1) {
             double ang = 2 * PI / len *
(invert? - 1:1);
             base wlen (cos (ang), sin (ang))
             for ( int i = 0; i < n; i += len ) {
                    base w (1);
                    for ( int j = 0; j < len / 2; +
+ j ) {
                           base u = a[i+j], v
= a [i + j + len / 2] * w;
                           a[i+j] = u + v;
                           a[i+j+len/2] = u
- v ;
                           w = wlen:
                    }
              }
      if (invert)
             for ( int i = 0; i < n; ++ i)
                    a[i]/=n;
int multiply ( const vector < int > & a, const
vector < int > & b) {
      vector < base > fa ( a. begin ( ) , a. end
()), fb (b. begin (), b. end ());
      size t n = 1;
      while ( n < max ( a. size ( ) , b. size ( ) )
            n <<= 1:
```

```
fa. resize (n), fb. resize (n);
      fft (fa, false), fft (fb, false);
      for ( size_t i = 0; i < n; ++ i)
             fa[i]*=fb[i];
      fft (fa, true);
  int ret=0;
      for ( size t i = 0; i < n; ++ i)
    int x=(int) ( fa [ i ] . real ( ) + 0.5 );
    if(x>0) ret++;
  return ret;
CLOSEST PAIR:
const int N = int(1e5) + 10; //sorted by X axis
const LL INF = 1LL << 60;
struct Point {
    LL x, y;
} point[N];
int n;
int tmpt[N];
bool cmpxy(const Point& a, const Point& b) {
         if (a.x != b.x)
                  return a.x < b.x;
         return a.y < b.y;
bool cmpy(const int& a, const int& b) {
         return point[a].y < point[b].y;</pre>
LL dis2(int i, int j) {
         return (point[i].x - point[j].x) *
(point[i].x - point[j].x)
                            + (point[i].y -
point[j].y) * (point[i].y - point[j].y);
LL sqr(LL x) {
         return x * x;
LL Closest_Pair(int left, int right) {
         LL d = INF:
         if (left == right)
                  return d;
         if (left + 1 == right)
                  return dis2(left, right);
         int mid = (left + right) >> 1;
         LL d1 = Closest_Pair(left, mid);
         LL d2 = Closest_Pair(mid + 1, right);
         d = min(d1, d2);
```

```
int i, j, k = 0;
                                                               ll b=a[i][y],c=a[x][y]; ///for mod k
          for (i = left; i <= right; i++) {</pre>
                    if (sqr(point[mid].x -
                                                               rep(j,y,col) a[i][j] = ((c*a[i][j] - b*a[x][j])
point[i].x) \ll d
                                                       %k+k)%k; ///for mod k
                              tmpt[k++] = i;
          sort(tmpt, tmpt + k, cmpy);
          for (i = 0; i < k; i++) {
                                                               double tmp=a[i][y]/a[x][y]; ///for real values
                    for (j = i + 1; j < k &&
                                                               rep(j,y,col) a[i][j] = a[x][j]*tmp; ///for real
sqr(point[tmpt[j]].y - point[tmpt[i]].y) < d;</pre>
                                                       values
                              LL d3 =
dis2(tmpt[i], tmpt[j]);
                              if (d > d3)
                                                            }
                                        d = d3;
                                                            x++:
                    }
          return d;
                                                          /// for GF(2) and mod k
GAUSSIAN ELIMINATION:
                                                          int ans=1;
int gauss()
                                                          rep(i,1,row-x) ans=(2LL*ans)%MOD;
  int x,y;
  for(x=0,y=0;x<row && y<col;y++)
                                                          ans--;
    int p=x;
                                                          return (ans%MOD+MOD)%MOD;
    rep(i,x,row-1) if(a[i][y]) {p=i;break;} ///for
                                                          ///for real values only
GF(2)
                                                          rep(j,0,col-1)
    rep(i,x+1,row-1) if(abs(a[i][y])>abs(a[p][y]))
p=i; ///for mod k
                                                            if(loc[j]==-1) continue;
    rep(i,x+1,row-1) if(fabs(a[i][y])>fabs(a[p][y])
                                                            a[loc[i]][col]/=a[loc[i]][i];
+eps) p=i; ///for real values
                                                            a[loc[j]][j]/=a[loc[j]][j];
                                                          }
    if(!a[p][y]) continue; ///for GF(2) and mod k
    if(fabs(a[p][y])<eps) continue; ///for real values
                                                       void init()
    if(p!=x) rep(j,y,col) swap(a[p][j],a[x][j]);
                                                          num=row=col=100;
    loc[y]=x; ///only for real value
                                                          CLR(a);
                                                          SET(s);
    rep(i,0,row-1)
                                                          SET(loc);
       if(i==x \parallel !a[i][y]) continue; ///for GF(2) and
                                                       int main() {
mod k
                                                          init();
       if(i==x \parallel fabs(a[i][v]) < eps) continue; ///for
                                                          gauss();
real values
                                                          printf("%.10lf\n",a[loc[pos]][col]);
       rep(j,y,col) a[i][j]^=a[x][j]; ///for GF(2)
```

```
GEOMETRY:
int dblcmp(double d)
  if (fabs(d)<eps)return 0;
  return d>eps?1:-1;
inline double sqr(double x)\{return x*x;\}
point()
                                             - Empty constructor
point(double x, double y)
                              - constructor
input()
                                             - double input
                                             - .2lf output
output()
operator ==
                                             - compares x and y
                                             - compares first by x, then by y
operator <
                                             - gives length from origin
len()
                                             - gives square of length from origin
len2()
                                     - gives distance from p
distance(point p)
                                     - returns new point after adding curresponging x and y
add(point p)
sub(point p)
                                     - returns new point after subtracting curresponging x and y
                                     - returns new point after multiplieing x and y by b
mul(double b)
                                     - returns new point after divideing x and y by b
div(double b)
dot(point p)
                                     - dot product
det(point p)
                                     - cross product of 2d points
                              - Probably radius of circumcircle of the triangle
rad(point a, point b)
trunc(double r)
                                             - return point that is truncated the distance from center to r
rotleft()
                                             - returns 90 degree ccw rotated point
rotright()
                                             - returns 90 degree cw rotated point
rotate(point p, double angle) - returns point after rotateing the point centering at p by angle radian ccw
*/
struct point
  double x,y;
  point()
  point(double \_x,double \_y){ x = \_x; y = \_y;
  void input()
                   { scanf("%lf%lf",&x,&y);
                                                       }
  void output()
                    { printf("%.2f %.2f\n",x,y);
  bool operator==(point a)const{
     return dblcmp(a.x - x) == 0 \&\& dblcmp(a.y - y) == 0;
  bool operator<(point a)const{</pre>
     return dblcmp(a.x - x) == 0 ? dblcmp(y - a.y) < 0 : x < a.x;
  point operator-(point a)const{
               return point(x-a.x, y-a.y);
  double len()
                    { return hypot(x, y);
                    { return x * x + y * y;
  double len2()
  double distance(point p){return hypot(x - p.x, y - p.y); }
  point add(point p) { return point(x + p.x, y + p.y); }
```

```
point sub(point p) { return point(x - p.x, y - p.y); }
  point mul(double b) { return point(x * b, y * b);
  point div(double b) { return point(x / b, y / b);
                                                       }
  double dot(point p) { return x*p.x+y*p.y;
  double det(point p) { return x*p.y-y*p.x;
                                                       }
  double rad(point a,point b){
     point p=*this;
     return fabs(atan2(fabs(a.sub(p).det(b.sub(p))),a.sub(p).dot(b.sub(p))));
  point trunc(double r){
     double l=len();
     if (!dblcmp(l))return *this;
     return point(x*r,y*r);
  point rotleft() { return point(-y,x);
  point rotright() { return point(y,-x);
  point rotate(point p,double angle){
     point v=this->sub(p);
     double c=cos(angle),s=sin(angle);
     return point(p.x+v.x*c-v.y*s,p.y+v.x*s+v.y*c);
  }
};
Stores two points
line()
                                                             - Empty constructor
line(point a, point b)
                                             - line through a and b
operator ==
                                                             - checks if two points are same
line(point p, double angle)
                                             - one end p, another end at angle degree
line(double a, double b, double c) - line of equation ax + by + c = 0
                                                             - inputs a and b
input()
adjust()
                                                             - orders in such a way that a < b
length()
                                                            - distance of ab
                                                             - returns 0 <= angle < 180
angle()
relation()
                                                             - 0 if collinear
                                                                      1 if ccw
                                                                      2 if cw
                                                     - returns 1 if point is on segment
pointonseg(point p)
parallel(line v)
                                                     - returns 1 if they are parallel
segcrossseg(line v)
                                                     - returns 0 if does not intersect
                                                                     returns 1 if non-standard intersection
                                                                     returns 2 if intersects
                                             - returns 1 if intersects strictly inside
segcrossseg_inside(line v)
                                                                     returns 0 if not
linecrossseg(line v)
                                             - v is line
linecrossline(line v)
                                             - 0 if parallel
                                                                      1 if coincides
                                                                      2 if intersects
crosspoint(line v)
                                                     - returns intersection point
```

```
dispointtoline(point p)
                                                    - distance from point p to the line
                                            - distance from p to the segment
dispointtoseg(point p)
                                                    - returns projected point p on ab line
lineprog(point p)
                                                    - returns reflection point of p over ab
symmetrypoint(point p)
*/
struct line
  point a,b;
  line()
  line(point _a,point _b){ a=_a; b=_b;
  bool operator==(line v){ return (a==v.a)&&(b==v.b);
                                                            }
  line(point p,double angle){
     a=p;
     if (dblcmp(angle-pi/2)==0){
       b=a.add(point(0,1));
     }else{
       b=a.add(point(1,tan(angle)));
     }
  //ax+bv+c=0
  line(double _a,double _b,double _c){
     if (dblcmp(_a)==0){
       a=point(0,-_c/_b);
       b=point(1,-_c/_b);
     else if (dblcmp(_b)==0){
       a=point(-_c/_a,0);
       b=point(-_c/_a,1);
     }else{
       a=point(0,-_c/_b);
       b=point(1,(-_c-_a)/_b);
  }
  void input()
                   { a.input(); b.input();
                                                 }
                   { if(b<a)swap(a,b);
  void adjust()
  double length()
                   { return a.distance(b);
                                                   }
  double angle(){
     double k=atan2(b.y-a.y,b.x-a.x);
     if (dblcmp(k)<0)k+=pi;
     if (dblcmp(k-pi)==0)k-=pi;
     return k;
  }
  int relation(point p){
     int c=dblcmp(p.sub(a).det(b.sub(a)));
     if (c<0)return 1;
     if (c>0)return 2;
     return 3;
  bool pointonseg(point p){
     return dblcmp(p.sub(a).det(b.sub(a)))==0\&\&dblcmp(p.sub(a).dot(p.sub(b)))<=0;
```

```
bool parallel(line v){
  return dblcmp(b.sub(a).det(v.b.sub(v.a)))==0;
int segcrossseg(line v){
  int d1=dblcmp(b.sub(a).det(v.a.sub(a)));
  int d2=dblcmp(b.sub(a).det(v.b.sub(a)));
  int d3=dblcmp(v.b.sub(v.a).det(a.sub(v.a)));
  int d4=dblcmp(v.b.sub(v.a).det(b.sub(v.a)));
  if ((d1^d2)=-2&&(d3^d4)=-2) return 2;
  return (d1==0\&\&dblcmp(v.a.sub(a).dot(v.a.sub(b)))\leq=0||
                    d2==0\&\&dblcmp(v.b.sub(a).dot(v.b.sub(b)))\leq=0
                    d3==0\&\&dblcmp(a.sub(v.a).dot(a.sub(v.b)))\leq=0
                    d4==0\&\&dblcmp(b.sub(v.a).dot(b.sub(v.b)))\leq=0);
}
int segcrossseg_inside(line v){
            if(v.pointonseg(a) || v.pointonseg(b) || pointonseg(v.a) || pointonseg(v.b)) return 0;
  int d1=dblcmp(b.sub(a).det(v.a.sub(a)));
  int d2=dblcmp(b.sub(a).det(v.b.sub(a)));
  int d3=dblcmp(v.b.sub(v.a).det(a.sub(v.a)));
  int d4=dblcmp(v.b.sub(v.a).det(b.sub(v.a)));
  if ((d1^d2)==-2&&(d3^d4)==-2) return 1;
  return (d1==0\&\&dblcmp(v.a.sub(a).dot(v.a.sub(b)))\leq=0||
                    d2==0\&\&dblcmp(v.b.sub(a).dot(v.b.sub(b)))\leq=0
                    d3==0\&\&dblcmp(a.sub(v.a).dot(a.sub(v.b))) <=0
                    d4==0\&\&dblcmp(b.sub(v.a).dot(b.sub(v.b)))\leq=0);
int linecrossseg(line v){//*this seg v line
  int d1=dblcmp(b.sub(a).det(v.a.sub(a)));
  int d2=dblcmp(b.sub(a).det(v.b.sub(a)));
  if ((d1^d2)==-2) return 2;
  return (d1==0)|d2==0;
int linecrossline(line v){
  if ((*this).parallel(v)){
     return v.relation(a)==3;
  return 2;
point crosspoint(line v){
  double a1=v.b.sub(v.a).det(a.sub(v.a));
  double a2=v.b.sub(v.a).det(b.sub(v.a));
  return point((a.x*a2-b.x*a1)/(a2-a1),(a.y*a2-b.y*a1)/(a2-a1));
}
double dispointtoline(point p){
  return fabs(p.sub(a).det(b.sub(a)))/length();
}
double dispointtoseg(point p){
  if (dblcmp(p.sub(b).dot(a.sub(b)))<0||dblcmp(p.sub(a).dot(b.sub(a)))<0){
```

```
return min(p.distance(a),p.distance(b));
    return dispointtoline(p);
 point lineprog(point p){
    return a.add(b.sub(a).mul(b.sub(a).dot(p.sub(a))/b.sub(a).len2()));
 point symmetrypoint(point p){
    point q=lineprog(p);
    return point(2*q.x-p.x,2*q.y-p.y);
/* 2D Geometry: Circle tangents
                                                        derangements : dr[i] = (i - 1) * (dr[i - 1] + dr[i - 2]);
                                                        dr[0] = dr[2] = 1; dr[1] = 0;
                                                        number of ways to ditribute n chocolate among r
 Description: Given a circle (defined by a center
                                                        persons is (n+r-1)C(r-1)
point and radius)
          and a point strictly outside the circle,
                                                        picks theorem : A = i + (b/2) - 1, i = interrior points,
                                                        b = boundary points.
returns the
          two points of tangency.
                                                        number of lattice points in between a segment of (x1,
                                                        v1) and (x2,v2) is gcd(abs(x1-x2), abs(v1-v2)) + 1.
 Complexity: O(1)
                                                        Strongly-connected components. Kosaraju's
                                                        algorithm.
            Ashley Zinyk
 Author:
                                                        1. Let G T be a transpose G (graph with reversed
 Date:
            Nov 19, 2002
                                                        edges.)
                                                        1. Call DFS(G T ) to compute finishing times f [u]
 References:
                                                        for each vertex u.
                                                        3. For each vertex u, in the order of decreasing f [u],
                                                        perform DFS(G, u).
 Reliability: 0
                                                        4. Each tree in the 3rd step's DFS forest is a separate
 Notes:
            Assumes a non-zero distance between p
                                                        SCC.
and c.
                                                        2-SAT. Build an implication graph with 2 vertices for
*/
                                                        each variable – for the variable and its inverse;
                                                        for each clause x \ V \ y add edges (-x, y) and (-y, x).
                                                        The formula is satisfiable iff x and x are in distinct
#define SQR(x)((x)*(x))
                                                        SCCs, for all x. To find a satisfiable assignment,
                                                        consider the graph's SCCs in topological order from
typedef struct {
double x, y;
                                                        sinks to sources (i.e. Kosaraju's last step), assigning
} Point;
                                                        'true' to all variables of the current SCC (if it
                                                        hasn't been previously assigned 'false'), and 'false'
                                                        to all inverses.
double dist2(Point a, Point b) {
return SQR(a.x-b.x) + SQR(a.y-b.y);
Point a, b;
```

```
b.x = a.x;
void circ_tangents(Point c, double r, Point p) {
                                                        b.y = a.y;
 double perp, para, tmp = dist2(p,c);
                                                        dp = (b.x*(c.x-a.x) + b.y*(c.y-a.y)) / (SQR(b.x))
                                                       +SQR(b.y));
 para = r*r/tmp;
                                                        p.x = b.x*dp + a.x;
                                                        p.y = b.y*dp + a.y;
 perp = r*sqrt(tmp-r*r)/tmp;
                                                        return p;
 a.x = c.x + (p.x-c.x)*para - (p.y-c.y)*perp;
 a.y = c.y + (p.y-c.y)*para + (p.x-c.x)*perp;
                                                       Point closest_pt_lineseg(Point a, Point b, Point c) {
b.x = c.x + (p.x-c.x)*para + (p.y-c.y)*perp;
                                                        Point p;
b.y = c.y + (p.y-c.y)*para - (p.x-c.x)*perp;
                                                        double dp;
                                                        b.x = a.x:
                                                        b.y = a.y;
/* 2D-Geometry: Closest Point on a Line/Segment
                                                        if (fabs(b.x) \le EPS \&\& fabs(b.y) \le EPS) return a;
                                                        dp = (b.x*(c.x-a.x) + b.y*(c.y-a.y))/(SQR(b.x)
                                                       +SQR(b.v));
                                                        if (dp > 1) dp = 1;
 Description: Given the end points of a line
                                                        if (dp < 0) dp = 0;
segment, A and B, and
                                                        p.x = b.x*dp + a.x;
         another point C, returns the point on the
                                                        p.y = b.y*dp + a.y;
segment
                                                        return p;
              closest to C.
              If a line perpendicular to A,B
intersects A, B then
                                                       /* 2D-Geometry: Circle-Line Intersection
              intersection will be returned.
Otherwise the closer
              endpoint will be returned.
              If the segment has a length of zero, an
                                                         Description: Given either a line or a line segment,
                                                       and a circle
endpoint
              will be returned.
                                                                 these routines calculate the number and
                                                       coordinates
 Complexity: O(1)
                                                                      of the intersection points.
                                                         Complexity: O(1)
            Ashley Zinyk
 Author:
            Nov 9, 2002
 Date:
 References: 0
                                                                    Gilbert Lee
                                                         Author:
                                                         Date:
                                                                    Sept 8, 2002
                                                         References: mathworld.wolfram.com/Circle-
                                                       LineIntersection.html
 Reliability: 1 (Spain 10263 - Railway)
 Notes:
Point closest_pt_iline(Point a, Point b, Point c) {
                                                         Reliability: 0
 Point p;
 double dp;
```

```
Notes:
int sgn(double x){
                                                       Description: Given a set of points, this returns the
return x < 0? -1:1;
                                                      circle with
                                                                the minimum area which completely
                                                      contains all those
double dist_2d(Point a, Point b){
                                                                    points
return sqrt(SQR(a.x-b.x)+SQR(a.y-b.y));
                                                        Complexity: O(n^3) worst case, where n is the
                                                      number of points
int circ_iline_isect(Circle c, Point a, Point b,
                                                                but on average O(n)
                 Point *r1, Point *r2){
 double dx = b.x-a.x, dy = b.y-a.y;
 double sdr = SQR(dx) + SQR(dy), dr = sqrt(sdr);
 double D,disc,x,y;
                                                       Author:
                                                                   Gilbert Lee
                                                                  Jan 24, 2003
                                                       Date:
 a.x = c.o.x; a.y = c.o.y;
                                                        References:
 b.x = c.o.x; b.y = c.o.y;
                                                      http://www.cs.unc.edu/~eberly/gr_cont.htm
 D = a.x*b.y - b.x*a.y;
 disc = SQR(c.r*dr)-SQR(D);
if(disc < 0) return 0;
                                                        Reliability: 2 (Spain 10005 - Packing polygons)
 x = sgn(dy)*dx*sqrt(disc);
                                                                  Problem C Aliens Jan 21, 2003
 y = fabs(dy)*sqrt(disc);
                                                       Notes:
                                                                  This is a simplification of the old
r1->x = (D*dy + x)/sdr + c.o.x;
                                                      min_circle code
 r2->x = (D*dy - x)/sdr + c.o.x;
                                                                by Scott Crosswhite (/Old/min circle.c)
 r1->y = (-D*dx + y)/sdr + c.o.y;
                                                                    The input array is sorted to increase
r2->y = (-D*dx - y)/sdr + c.o.y;
                                                      stability of
return disc == 0 ? 1 : 2;
                                                                    an answer. This may be removed to
                                                      increased speed
                                                      */
int circ_lineseg_isect(Circle c, Point a, Point b,
                                                      int inside(Point p, Circle c){
                  Point *r1, Point *r2){
                                                       return SQR(p.x-c.x)+SQR(p.y-c.y) \le SQR(c.r);
 double d = dist 2d(a,b);
int res = circ_iline_isect(c,a,b,r1,r2);
                                                      Circle Circle1(Point p){
if(res == 2 \&\& dist_2d(a,*r2)+dist_2d(*r2,b) != d)
                                                       Circle c;
                                                       c.x = p.x; c.y = p.y; c.r = 0;
if(res \ge 1 \&\& dist_2d(a,*r1)+dist_2d(*r1,b) != d){
                                                       return c;
  r1 = r2;
                                                      }
  res--;
                                                      Circle Circle2(Point p1, Point p2){
 }
return res;
                                                       Circle c;
                                                       c.x = 0.5*(p1.x + p2.x);
                                                       c.y = 0.5*(p1.y + p2.y);
/* 2D Geometry: Minimum bounding circle
                                                       c.r = 0.5*sqrt(SQR(p1.x-p2.x)+SQR(p1.y-p2.y));
                                                       return c;
______
```

```
Circle res; double a,b,c,d,e,f,g;
                                                         Circle Circle3(Point p1, Point p2, Point p3){
 a = p2.x - p1.x; b = p2.y - p1.y;
                                                         if( fabs( p.x ) <= EPS ) return ( p.y > EPS ? 1.0 :
 c = p3.x - p1.x; d = p3.y - p1.y;
                                                         3.0) * acos(0);
 e = (p2.x + p1.x)*a + (p2.y + p1.y)*b;
                                                           double theta = atan( 1.0 * p.y / p.x );
                                                           if( p.x > EPS ) return( p.y >= -EPS ? theta : ( 4 *
 f = (p3.x + p1.x)*c + (p3.y + p1.y)*d;
 g = 2.0*(a*(p3.y - p2.y) - b*(p3.x - p2.x));
                                                         acos(0) + theta);
 if (fabs(g) < EPS){
                                                           return(2 * acos(0) + theta);
  res.x = res.y = res.r = DBL\_MAX;
  return res;
                                                         bool pointInPoly( P p, vector< P > &poly )
res.x = (d*e - b*f) / g;
                                                           int n = poly.size();
                                                           double ang = 0.0;
res.y = (a*f - c*e) / g;
                                                           for( int i = n - 1, j = 0; j < n; i = j++)
 res.r = sqrt(SQR((p1.x-res.x))+SQR((p1.y-res.y)));
return res;
                                                              P v(poly[i].x - p.x, poly[i].y - p.y);
                                                              P w(poly[j].x - p.x, poly[j].y - p.y);
                                                              double va = polarAngle(v);
Circle min_circle(Point *p, int n){
                                                              double wa = polarAngle( w );
int i, j, k; Point t; Circle c = Circle1(p[0]);
                                                              double xx = wa - va;
                                                              if( va < -0.5 \parallel wa < -0.5 \parallel fabs( fabs( xx ) - 2 *
 /* Randomize point array to avoid doctored input -
                                                        acos(0) > EPS
may modify the
  limit on the for loop to increase/decrease
                                                              {
randomness */
                                                                // POINT IS ON THE EDGE
 for(i = 0; i < n; i++){
                                                                assert( false );
  i = rand() \% n;
                                                                ang += 2 * acos(0);
  k = rand() \% n;
                                                                continue:
  t = p[j]; p[j] = p[k]; p[k] = t;
                                                             if( xx < -2 * acos(0) ) ang += xx + 4 *
                                                         acos(0):
 for(i = 1; i < n; i++) if(!inside(p[i], c)){ c =
                                                              else if( xx > 2 * acos(0) ) ang += xx - 4 * acos(
                                                         0);
Circle1(p[i]);
for(j = 0; j < i; j++) if(!inside(p[j], c)){ c =
                                                              else ang += xx;
Circle2(p[i],p[j]);
for(k = 0; k < j; k++) if(!inside(p[k], c)) c =
                                                           return( ang * ang > 1.0 );
Circle3(p[i],p[j],p[k]);}}
return c;
                                                         bool lineIntersect( Pa, Pb, Pc, Pd, P&r)
bool pointInsideTriangle( double x[], double y[],
                                                           P n; n.x = d.y - c.y; n.y = c.x - d.x;
double xx, double yy)
                                                           double denom = n.x * (b.x - a.x) + n.y * (b.y -
                                                         a.y );
  return leftTurn( x[0], y[0], x[1], y[1], xx, yy)
                                                           if( fabs( denom ) < EPS ) return false;</pre>
                                                           double num = n.x * (a.x - c.x) + n.y * (a.y - c.y);
     && leftTurn( x[1], y[1], x[2], y[2], xx, yy )
                                                           double t = -num / denom;
     && leftTurn( x[2], y[2], x[0], y[0], xx, yy );
                                                           r.x = a.x + t * (b.x - a.x);
double polarAngle( P p )
                                                           r.y = a.y + t * (b.y - a.y);
                                                           return true;
  if( fabs( p.x ) <= EPS && fabs( p.y ) <= EPS )
return -1.0;
                                                         bool lineSegIntersect( vector< T > &x, vector< T >
```

```
&y)
  double ucrossv1 = (x[1] - x[0]) * (y[2] - y[0]) - (
y[1] - y[0]) * (x[2] - x[0]);
  double ucrossv2 = (x[1] - x[0]) * (y[3] - y[0]) - (
                                                        * Computes the center of a circle containing the 2
y[1] - y[0]) * (x[3] - x[0]);
  if( ucrossv1 * ucrossv2 > 0 ) return false;
                                                        * points. The circle has the given radius. The
  double vcrossu1 = (x[3] - x[2]) * (y[0] - y[2]) - (
                                                       returned
y[3] - y[2]) * (x[0] - x[2]);
                                                        * center is never to the right of the vector
  double vcrossu2 = (x[3] - x[2]) * (y[1] - y[2]) - (
                                                        * (x1, y1) --> (x2, y2).
y[3] - y[2]) * (x[1] - x[2]);
                                                        * If this is possible, returns true and passes the
  return( vcrossu1 * vcrossu2 <= 0 );</pre>
                                                        * through the ctr array. Otherwise, returns false.
                                                        * #include <math.h>
/**********
                                                        * FIELD TESTING:
                                                             - Valladolid 10136: Chocolate Chip Cookies
* Circle through 3 points *
* Computes the circle containing the 3 given points.
                                                       bool circle2ptsRad( double x1, double y1, double x2,
* The 3 points are
                                                       double v2, double r, double ctr[2])
     (x[0], y[0]), (x[1], y[1]) and (x[2], y[2]).
* The centre of the circle is returned as (r[0], r[1]).
                                                          double d2 = (x1 - x2) * (x1 - x2) + (y1 - y2) * (
* The radius is returned normally. If the circle is
                                                       y1 - y2);
* undefined (the points are collinear), -1.0 is
                                                          double det = r * r / d2 - 0.25;
returned.
                                                          if( det < 0.0 ) return false;
* #include <math.h>
                                                          double h = sqrt( det );
* REQUIRES: lineIntersect
                                                          ctr[0] = (x1 + x2) * 0.5 + (y1 - y2) * h;
* FIELD TESTING: Passed UVA 190
                                                          ctr[1] = (y1 + y2) * 0.5 + (x2 - x1) * h;
                                                          return true:
double circle3pts( double x[], double y[], double r[] )
  double lix[4], liy[4];
                                                        * Great Circle *
                                                        *****
  lix[0] = 0.5 * (x[0] + x[1]); liy[0] = 0.5 * (y[0] +
y[1]);
                                                        * Given two pairs of (latitude, longitude), returns the
  lix[1] = lix[0] + y[1] - y[0]; liy[1] = liy[0] + x[0] -
                                                        * great circle distance between them.
                                                        * FIELD TESTING
x[1];
  lix[2] = 0.5 * (x[1] + x[2]); liy[2] = 0.5 * (y[1] +
                                                             - Valladolid 535: Globetrotter
y[2]);
  lix[3] = lix[2] + y[2] - y[1]; liy[3] = liy[2] + x[1] -
                                                       double greatCircle( double laa, double loa, double
                                                       lab, double lob)
  if(!lineIntersect(lix, liy, r)) return -1.0;
  return sqrt(
                                                          double PI = acos(-1.0), R = 6378.0;
    (r[0] - x[0]) * (r[0] - x[0]) +
                                                          double u[3] = {\cos(laa) * \sin(loa), \cos(laa) *}
    (r[1] - y[0]) * (r[1] - y[0]);
                                                       cos(loa), sin(laa)};
                                                          double v[3] = \{ \cos( lab ) * \sin( lob ), \cos( lab ) * \}
}
                                                       cos( lob ), sin( lab ) };
                                                          double dot = u[0]*v[0] + u[1]*v[1] + u[2]*v[2];
                                                          bool flip = false;
                                                          if (dot < 0.0)
                                                          {
* Circle of a given radius through 2 points *
```

```
flip = true;
     for( int i = 0; i < 3; i++) v[i] = -v[i];
                                                          PT c(0,0):
                                                           double scale = 6.0 * ComputeSignedArea(p);
  double cr[3] = \{ u[1]*v[2] - u[2]*v[1], u[2]*v[0] - u[2]*v[1], u[2]*v[0] - u[2]*v[1], u[2]*v[1] \}
                                                           for (int i = 0; i < p.size(); i++){
                                                             int j = (i+1) \% p.size();
u[0]*v[2], u[0]*v[1] - u[1]*v[0] ;
  double theta = asin( sqrt( cr[0]*cr[0] + cr[1]*cr[1]
                                                             c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
+ cr[2]*cr[2]));
  double len = theta * R;
                                                           return c / scale;
  if(flip) len = PI * R - len;
  return len;
                                                          // tests whether or not a given polygon (in CW or
                                                          CCW order) is simple
                                                          bool IsSimple(const vector<PT> &p) {
// compute intersection of circle centered at a with
                                                           for (int i = 0; i < p.size(); i++) {
radius r
                                                             for (int k = i+1; k < p.size(); k++) {
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b,
                                                              int j = (i+1) \% p.size();
double r, double R) {
                                                              int l = (k+1) \% p.size();
                                                              if (i == l \parallel j == k) continue;
 vector<PT> ret;
                                                              if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
 double d = sqrt(dist2(a, b));
 if (d > r+R \parallel d+min(r, R) < max(r, R)) return ret;
                                                               return false:
 double x = (d*d-R*R+r*r)/(2*d);
 double y = sqrt(r*r-x*x);
 PT v = (b-a)/d;
                                                           return true;
 ret.push_back(a+v*x + RotateCCW90(v)*y);
 if (y > 0)
                                                          bool cw(const point &a, const point &b, const point
  ret.push back(a+v*x - RotateCCW90(v)*v);
 return ret;
                                                             return (b.first - a.first) * (c.second - a.second) -
                                                          (b.second - a.second) * (c.first - a.first) < 0;
// This code computes the area or centroid of a
(possibly nonconvex)
                                                          vector<point> convexHull(vector<point> p) {
// polygon, assuming that the coordinates are listed in
                                                             int n = p.size();
a clockwise or
                                                             if (n \le 1)
// counterclockwise fashion. Note that the centroid is
                                                               return p;
                                                             sort(p.begin(), p.end());
often known as
// the "center of gravity" or "center of mass".
                                                             int cnt = 0;
double ComputeSignedArea(const vector<PT> &p) {
                                                             vector<point> q(n * 2);
 double area = 0;
                                                             for (int i = 0; i < n; q[cnt++] = p[i++])
 for(int i = 0; i < p.size(); i++) {
                                                               for (; cnt \geq 2 && !cw(q[cnt - 2], q[cnt - 1],
  int j = (i+1) \% p.size();
                                                          p[i]); --cnt)
  area += p[i].x*p[j].y - p[j].x*p[i].y;
                                                             for (int i = n - 2, t = cnt; i \ge 0; q[cnt++] = p[i--])
return area / 2.0;
                                                               for (; cnt > t && !cw(q[cnt - 2], q[cnt - 1], p[i]);
                                                          --cnt);
                                                             q.resize(cnt - 1 - (q[0] == q[1]));
double ComputeArea(const vector<PT> &p) {
                                                             return q;
 return fabs(ComputeSignedArea(p));
                                                          }
PT ComputeCentroid(const vector<PT> &p) {
                                                               int u,v;
```

```
struct trie
                                                             u=q.front();
                                                             q.pop();
  int next[26],fail,score;
  void clear(){ rep(i,0,CHZ-1) next[i]=0;
                                                             rep(i,0,CHZ-1)
fail=score=0: }
};
                                                                if(node[u].next[i])
trie node[MX];
                                                        node[node[u].next[i]].fail=node[node[u].fail].next[i];
vi vv[MX];
                                                        node[node[u].next[i]].score+=node[node[node[u].nex
                                                        t[i]].fail].score;
int n,m,b,mp[300],val[30],sc[110];
                                                                  q.push(node[u].next[i]);
char s[110][110];
                                                                else {
void init()
                                                                  node[u].next[i]=node[node[u].fail].next[i];
  n=root+1;
  rep(i,0,1) node[i].clear();
  rep(i,0,CHZ-1) node[0].next[i]=root;
                                                          return;
  rep(i,0,MX-1) vv[i].clear();
                                                        int dp[MX][205],vis[MX][205];
  return;
                                                        int go(int p,int c) {
                                                          if(c > b) return -inf;
void insert(int x)
                                                          int &ret=dp[p][c];
  int pos=root;
                                                          if(vis[p][c] == kk) return ret;
  for(int i=0;s[x][i];i++)
                                                          vis[p][c]=kk;
                                                          ret=0;
    if(!node[pos].next[mp[s[x][i]]])
                                                          rep(i,0,CHZ-1) {
       node[n].clear();
                                                             if(c + val[i] \le b) {
       node[pos].next[mp[s[x][i]]]=n++;
                                                                ret = max(ret, node[node[p].next[i]].score +
                                                        go(node[p].next[i],c+val[i]));
    pos=node[pos].next[mp[s[x][i]]];
                                                             }
                                                           }
  node[pos].score+=sc[x];
  return;
                                                          return ret;
void get_failure()
  queue<int>q;
  q.push(root);
  while(!q.empty())
```

```
TEMPLATE:
#define PCASE printf("Case %d: ",kk++)
#define PCASENL printf("Case %d:\n",kk++)
#define NL puts("")
#define sz(a) ((int)a.size())
#define repv(i,a) for(int i=0;i < sz(a);i++)
#define revv(i,a) for(int i=sz(a)-1;i>=0;i--)
#define rep(i,a,b) for(int i=a;i \le b;i++)
#define rev(i,a,b) for(int i=a;i>=b;i--)
#define FOR(I,A,B) for(int I = (A); I < (B); ++I)
#define REP(I,N) FOR(I,0,N)
#define all(a) a.begin(),a.end()
#define pb(a) push_back(a)
#define mp(a,b) make_pair(a,b)
#define pi (2.0*acos(0.0))
#define PI (2.0*acos(0.0))
#define SET(a) memset(a,-1,sizeof a)
#define CLR(a) memset(a,0,sizeof a)
#define set0(ar) memset(ar,0,sizeof ar)
#define setinf(ar) memset(ar,126,sizeof ar)
#define in(a,x,y) (a>=x && a<=y)
#define out(a,x,y) (!in(a,x,y))
#define xx first
#define yy second
using namespace std;
typedef long long ll;
typedef long long LL;
typedef unsigned long long ull;
typedef vector<int> vi;
typedef vector<vi>vvi;
typedef vector<ll> vll;
typedef pair<int,int> pii;
typedef pair<ll,ll> pll;
typedef vector<pii> vii;
template < class T > inline T _sq(T a) \{ return a * a; \}
template< class T, class X > inline T _{pow}(T a, X y) \{T z=1; rep(i,1,y)\{z*=a;\} return z; \}
template < class T > inline T extended < tag < and < tag < abs < abs < abs < by < abs < by < abs < abs < by < abs < abs < by < abs < abs < abs < by < abs < abs < abs < by < abs > abs < abs > abs < abs > abs 
\{x=1;y=0;g=a; return g;\}\ g=\_extended(b,a\%b,x1,y1);\ x=y1;\ y=x1-(a/b)*y1;\ return g;\}
template < class T, class X > inline bool getbit(T a, X i) { T t=1; return ((a&(t < i)) > 0); }
template < class T, class X > inline T setbit(T a, X i) { T t=1; return (a|(t<<i)); }
template < class T, class X > inline T resetbit(T a, X i) { T t=1; return (a&(\sim(t<<i))); }
template< class T, class X > inline T togglebit(T a, X i) { T t=1;return (a\land(t<<i)); }
template< class T,class X, class Y > inline T _bigmod(T n,X m,Y mod){ull ret=1, a = n%mod; while(m)
{ if(m&1)ret=(ret*a)%mod; m>>=1; a=(a*a)%mod; }ret%=mod;return (T)ret;}
template < class Y > inline T modiny (T n, Y mod) {return bigmod(n,mod-2,mod);}
```