



# How to design TinyURL?

# What is the problem?

<http://www.hugeurl.com/?>

NmU1OGQ3ZTY5ZmQ1NmYzNTc3ODE2MDc3ZDU1  
YTc4YTQmMTImVm0wd2QyUX1VWGxXYTJoV1YwZG9NV1l3Wkc5a1JsWjBUVlpP  
VOZac2JET1hhMUpUVmpGYWMvSkVUbGhoTWsweFZqQmFTMk15U2tWVWJHaG9U  
VmhDVVZadGVGwmxSbGw1Vkd0c2FsSnRhRzlVVjNOM1pVWmFkR05GZEZT1ZU  
VkpWbTEwYTFk5FNrZGpTRUpYVFVad1NGU1VSbUZgVmtaMFVteFNUbUY2U1RG  
V1ZFBd3dWak2hV020cmJGSmlSMmh2V1d4b2IwMHhXbGRYY1VaclVsUkdXbGt3  
WkRSVklrcElaSHBHViJFeVVY2FpWRVpvVTBaT2NscEhiR1jTV1hCW1ZrWldh  
MVV5VW50a1JtU11ZbF2hY1ZscldtRmxWbVJ5VjI1a1YwMUVSa1pWYk2KRFza  
QXhkV1Z1V2xaaGEexcFlXa1ZhVJDJodFnRZFRIv3hYVWpOb1dGWnRNSGrsU1Bs  
NFUvdGthVkJGV2xSWmJHaFRWMVpXY1ZKcmRGU1diRm93V2xWb2ExWXdNVVZT  
YTFwWF1rZG9jbFpxU2tabF2sW1pXa1prYUdFeGNGaFhiRnBoVkrkt2RGSnJh  
R2hTYXpWe1dXeG91mWRHV25ST1NhaFBVbTE0V1FSVmFHOvhSMHBJV1d4c1dt  
SkhhR1JXTU2wVFZgRmtkRkp0ZU2kaWEwcE1WbXBK2UUxR1dsafRhM1JxVWtW  
YVYxWnFUbT1sYkZweFUvdGthbUpWVmpaW1ZWcGhZVWRGZUdOSE9WaGhNvnBv  
VmtSS1RtVkdjRWxVY1doVFRXNW9WVmRXVWs5Uk1rbDRWMWhvWVZKR1NtR1dh  
a1pIVGtaYVdHUkhkRmRpV1hCN1ZUSTFUMVp0Um5KT1ZsS1hUVVp3VkZacVJu  
ZFNWa1pSVDFkclUwMH1hRmxXY1hCTFpXczFXRkpyWkZoaWF6VnhWVEJvUTFs  
V1VsWlhibVVjWWtad2VGvnRkREJWIWtwSVZXCENXbFpXYOROWmEvUkdaVWRP  
U0dGR2FHbFNiaOp2Vm10U1MxUnRwbGRUYmtwb1VgTm9WRmxZY02ka01WcFla  
VWM1VWsxWFVraFdNa1ZUVkd4S1Jz2HVTbF2XYkhCN1ZHeGF2VmR1VmtoalIv  
aHBVbGhDTmxkVVFtRmpNV1IwVWVxcoc2FGS1dTbUZhViNSaF1VWnJ1RmRyZEd0  
U2EzQ1ZwbGQ0YTJGV1Ns2GhNM1JYWx0Q1MxcF2Xa3BsUm1Se11VW1NhVkp1  
UWxwV2JYU1haREZrUjJKSVRtaFNgelZQVkZaYWQvVkdWWGxrUjBacFVteHdl  
bF152UhkWF1wV1RZMFJPV21FeVWrZGFVW1JQVTBVNVYxcEhiRmhTVlhCS1Zq  
SiBVMU14VFhsVVdHeF2ZVEZ3YU2wCvNtOVdSbEp2VGxiNvdGwnNjREJVV1ZK  
SFZXC3hXR1ZvYU2kTmFsW1VWa2Q0YTFOR1ZuT1hiR1pYWWtad1dWWkh1R0Za  
V1FKR1RsWmFVR1p0VW5CV2JHaERVMVprV1ZGdFJsWk5WbXcxV1d4b2MxWnNX  
a1pUYkdoWF1XczFkbGxWV21GalZrcHpXa1pvViJKc1NrbFdNbVewV1ZaWnVG  
TnJXbE5XU1ZVNQ==

Such a long url is common!



## 崩溃的瞬间：

- 在文章中引用链接.....10行结束了
- 在报纸上发现这个链接想浏览一下
- 复制粘贴到浏览器.....落下几行
- 在Twitter中share该链接.....140个字符根本不够用
- .....

# How to solve it?

<http://tinyURL/a1d3isf>



- TinyURL is a URL shortening web service, which provides short aliases for redirection of long URLs.
- First launched by Gilbertson in January 2002

# Review: SNAKE

- Scenario: case/interface
- Necessary: constrain/hypothesis
- Application: service/algorithm
- Kilobit: data
- Evolve

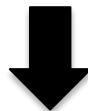
# Scenario

Question: How many services should TinyURL provide?

1. One
2. Two
3. Three
4. Four

It requires at least **TWO** basic services:

- Given a long url, it generates a **unique** short url
- Given a short url, it is able to retrieve the corresponding long url, thus,sends http request to the correct address



**short\_url insert(long\_url)**  
**long\_url lookup(short url)**

# Review: Performance v.s. Scalability

Question: Which describes performance? Which describes scalability?

1. How long does it take to respond to a request?

2. How many simultaneous users can you support?

3. Measured by the simultaneous users to support

4. Measured by delay from sending requests to displaying response

5. Influenced by code quality, which database to use, database schema, web framework, etc.

6. Influenced by performance per machine, how many machines to use, and how are they coordinated to work, etc.

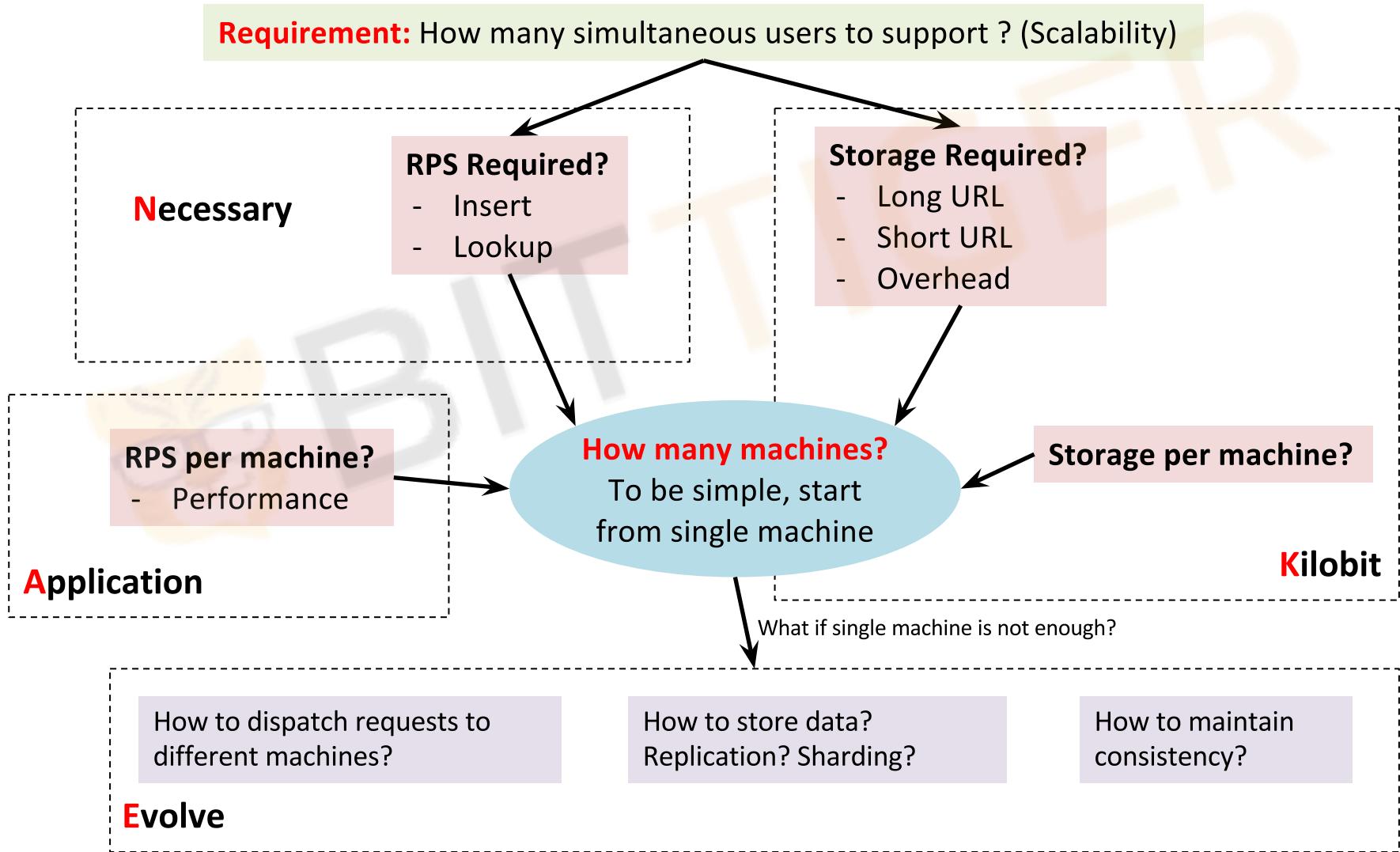
**Performance:**

?      ?      ?  
1      4      5

**Scalability:**

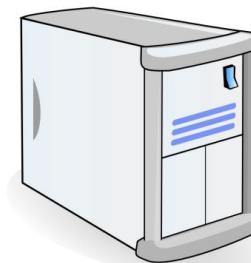
?      ?      ?  
2      3      6

# Decision Process



# Necessary

- Requirement: 1 Million Daily Active Users
- Insert Operations
  - $1,000,000 * 1\% \text{ (function usage)} * 10 \text{ (function frequency)} = 100,000$
  - RPS:  $100,000 / 86400 = 1.2$
- Lookup Operations
  - $1,000,000 * 100\% \text{ (function usage)} * 3 \text{ (function frequency)} = 3,000,000$
  - RPS:  $3,000,000 / 86400 = 35$
- Peak Estimation
  - Assume Peak Traffic = 5 \* Average Daily Traffic
  - Insertion Peak RPS =  $1.2 * 5 = 6$
  - Lookup Peak RPS =  $35 * 5 = 175$



# Application – Architecture

Naïve Thought:

short\_url insert(long\_url)  
long\_url lookup(short url)



HashTable/HashMap

To be simple, first, assume In-memory storage:

```
Class TinyURL {  
    map <longURL, shortURL> LongToShortMap;  
    map<shortURL, longURL> ShortToLongMap;  
  
    shortURL insert(longURL) {  
        if LongToShortMap not containsKey longURL:  
            generate shortURL;  
            put <longURL, shortURL> into LongToShortMap;  
            put <shortURL, longURL> into ShortToLongMap;  
            Return LongToShortMap.get(longURL);  
    }  
  
    longURL lookup(shortURL) {  
        return ShortToLongMap.get(shortURL);  
    }  
}
```

How to generate the  
shortURL?

How to define the hash  
function from longURL  
to shortURL?

# Application – Traditional Hash

Question: How are you familiar with Hash Function?

- 1. Know a lot
- 2. Learn a little
- 3. Have no idea

[http://www.construx.com/Blogs/10x\\_Software\\_Development/?id=15082](http://www.construx.com/Blogs/10x_Software_Development/?id=15082)

Hash Method	Hash Result	Length
Adler32	399014e3	8 (32bit)
CRC32	78aa9d1a	8 (32bit)
MD2	286c50c2db4fcad77adb4edeb3a937b2	32 (128bit)
MD4	387ac3f6aae7956c4fab176271bb4518	32 (128bit)
MD5	f061a171dfc30635462850684f98b886	32 (128bit)
SHA-1	3c93b6d332091b2970fb660d644d0ba3d756e322	40 (160bit)

Each Hexadecimal digit represents four binary digits

# Application – Traditional Hash

Assume we use **CRC32**

```
shortURL GenerateShortURL(longURL) {  
    return CRC32(longURL)  
}
```

- **How many different URL can it support?**

$2^{32} = 4,294,967,296 \sim 4 \text{ billion}$  regardless of storage capacity

- **How many attempts might cause a collision?**

According to [birthday paradox](#), P=50%, 77,000; P=75%, 110,000

*Basically you can expect to see the first collision after hashing  $2^{n/2}$  items, where n is number of bits. Say, 32bit,  $2^{16} = 65536$*

# Application – Traditional Hash

## What are the pros and cons?

### PROS

- 1. No need to write addition hash function, easy to implement.
- 2. Are randomly distributed
- 3. Support url clean

### CONS

Problem	Possible Solution
Not short enough (At least 4 bytes)	Only retrieve first 4-5 digit of the hash result
Collision cannot be avoided	Use (long_url + timestamp) as hash argument, if conflict happens, try again (timestamp changes) -> multiple attempts, highly possible conflicts when data is big
Hash computation costs a lot	

**Feasible, but not efficient enough!**

# Application – Base10

**Break Assumption:** Do we need to generate direct relationship between long\_url and short\_url, that is, String to String?

**What we want to do:** Given a long\_url, in any way, find a shorter sth, i.e.

- Keep only one-to-one relationship (no conflict)
- Easy to find the shorter sth when we want to insert another long\_url

**AUTO-INCREMENTAL INT → Index (Size)**

LongToShortMap

longURL	shortURL
Bittiger.io	0
www.inoreader.com/all_articles	1
http://www.construx.com/Blogs/10x_Software_Development/?id=15082	2
.....	3

ShortToLongMap

shortURL	longURL
0	Bittiger.io
1	www.inoreader.com/all_articles
2	http://www.construx.com/Blogs/10x_Software_Development/?id=15082
3	.....

# Application – Base10

Assume we use **Size of map**

```
shortURL GenerateShortURL() {  
    return ShortToLongMap.size();  
}
```

- **How many URL can it support?**

$2^{31} - 1 = 2,147,483,647 \sim 2$  billion in in-memory map  
basically as big as the capacity in database  
(if consider multiple machines)

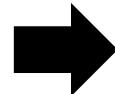
- **How many attempts might cause a collision?**

Within the capacity, no collision.

- **How long will be the url?**

$\text{NumOfURL} = 10^{\text{Length}}$

$\text{Length} = \log_{10}(\text{NumOfURL})$



**BASE What** is important!

# Application – Base62

Use more representations in one digit: a-z, A-Z, 0-9

```
shortURL GenerateShortURL() {  
    return ConvertTo62(ShortToLong.size());  
}  
  
shortURL ConvertTo62(mapSize) {  
    char Encode[62] = {'0', ..., '9', 'a', ..., 'z', 'A', ..., 'Z'};  
    shortURL = "";  
    while (mapSize > 0) {  
        shortURL = Encode[mapSize % 62] + shortURL;  
        mapSize /= 62;  
    }  
    return shortURL;  
}
```

- **How many URL can it support?**  
Same as Base10
- **How many attempts might cause a collision?**  
Within the capacity, no collision.
- **How long will be the url?**  
 $\text{NumOfURL} = 62^{\text{Length}}$   
 $\text{Length} = \log_{62}(\text{NumOfURL})$

# Let's see the Difference

Assume average 100,000 url are inserted each day



	Base10	Base62
Year URL	36,500,000	36,500,000
Usable characters	[0-9] = 10	[0-9a-zA-Z] = 62
Encoding length	$\log_{10}(36,500,000) \sim 8$	$\log_{62}(36,500,000) \sim 5$
Example	Goo.gl/36500000	Goo.gl/2t9jG

# Application – Base62

**What are the pros and cons?**

## PROS

- 1. Short url -> save space
- 2. Simple computation
- 3. No conflicts

## CONS

- 1. No support for url clean

# Application – In Database

We need a two-way lookup:

long\_url -> short\_url (index)  
short\_url (index) -> long\_url

**A table is enough** (auto-increment id, long\_url):

- Auto-increment id is index as default
- Establish another index on long\_url

## INSERT

LongURL  →

Long_url	id
http://bittiger.io	1
http://google.com	2
...	3

If contains such long\_url

return convertBase10toBase62(id)

If not contains such long\_url

Add the long\_url to the end  
return convertBase10toBase62(id)

## LOOKUP

ShortURL → convertBase62toBase10

id in database

→ Search id

id	long_url
1	http://bittiger.io
2	http://google.com
3	...

return long\_url

# Kilobit

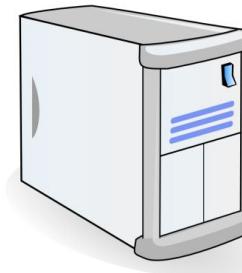
Average size of longURL = 100 bytes

Average size of id = 4 bytes (int)

Status = 4 bytes

Daily new URL =  $100,000 * 108 = 10.8\text{MB}$

Yearly URL =  $10.8 \text{ MB} * 365 \sim 4\text{GB}$



Single  
Machine



# Evolve – Single Machine

- How to implement time-limited service?
  - Set a expiration period (in status)
    - \* delete url in 3 years after its creation
    - \* delete url if it is not used during the past 3 years
- How to cache?
  - Pre-load: load data into memory when starting server
  - Replacement: LRU Cache

# Evolve – CAP

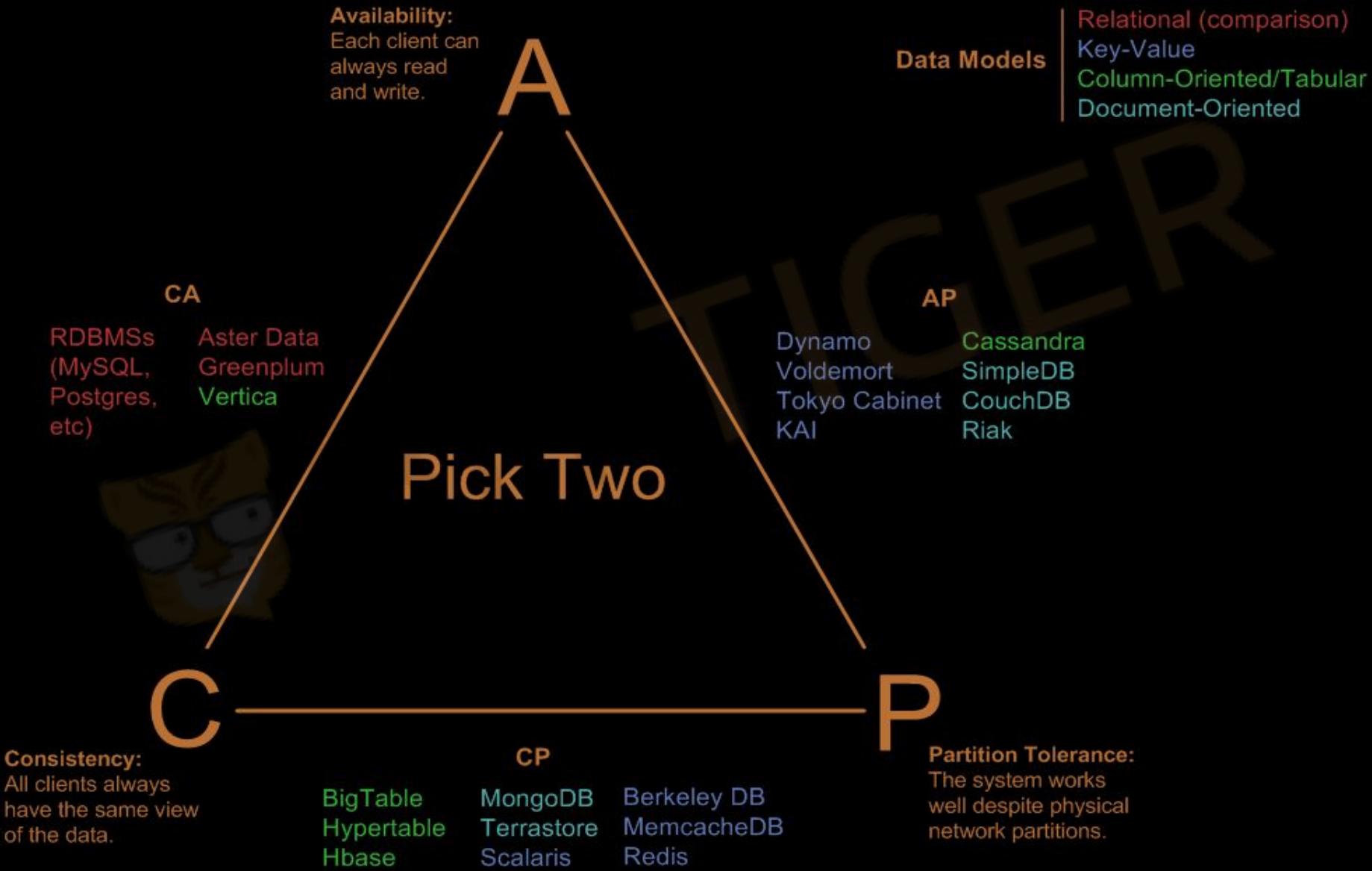
Guarantees	Definition
Consistency	Every read receives the most recent write or an error
Availability	Every request receives a response, without guarantee that it contains the most recent version of the information
Partition Tolerance	The system continues to operate despite an arbitrary number of messages being dropped by the network between nodes

**ONLY PICK TWO!**

# Distributed Systems

- Distributed System Features:
  - No distributed system is safe from network failure
  - Networking partition has to be tolerated
- Choose between Consistency and Availability in the presence of a partition
- In the absence of network failure, both availability and consistency can be satisfied

# Visual Guide to NoSQL Systems

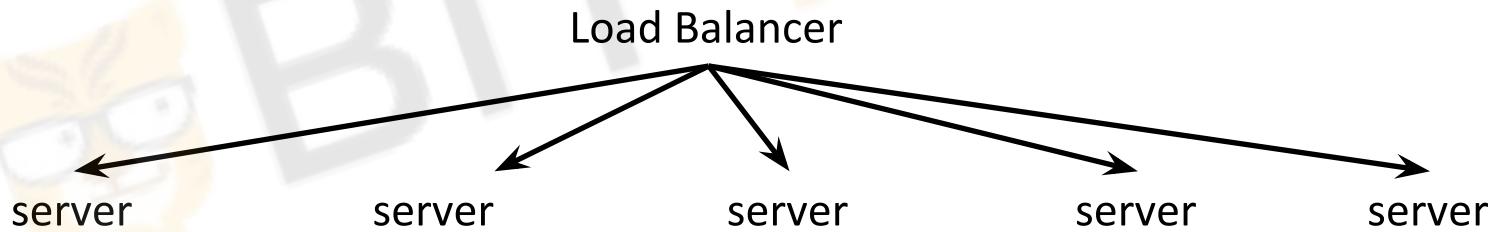


# Evolve – Multiple Machine

Problem: What if a large amount of users frequently visit limited url?

Demand: High RPS for lookup, limited amount of frequent url

**REPLICATION – store full data on all machines**

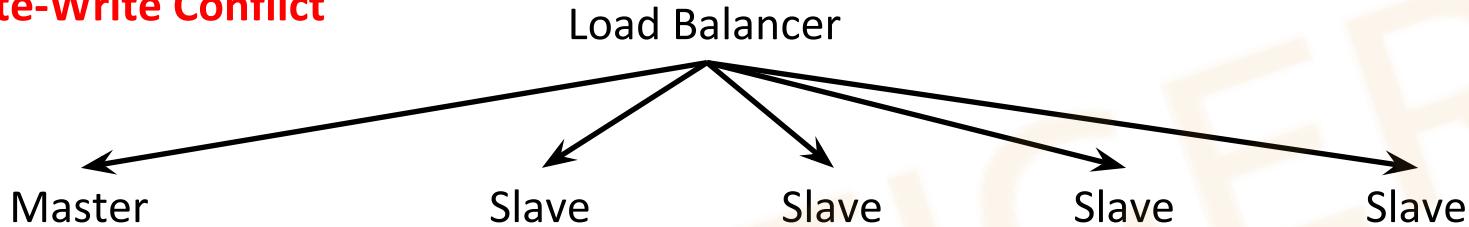


Problem:

- Write-Write conflict:  
What if A writes google.com on server1, while B writes google.com on server2?
- Read-Write conflict:  
What if A writes google.com on server1, while B reads google.com on server2?

# Evolve – Consistency

## Write-Write Conflict



- Do lookup operations
  - Do insert operations, and populate to other slaves
  - Store Full data
- Only do lookup operations
  - Store Full data



## Read-Write Conflict

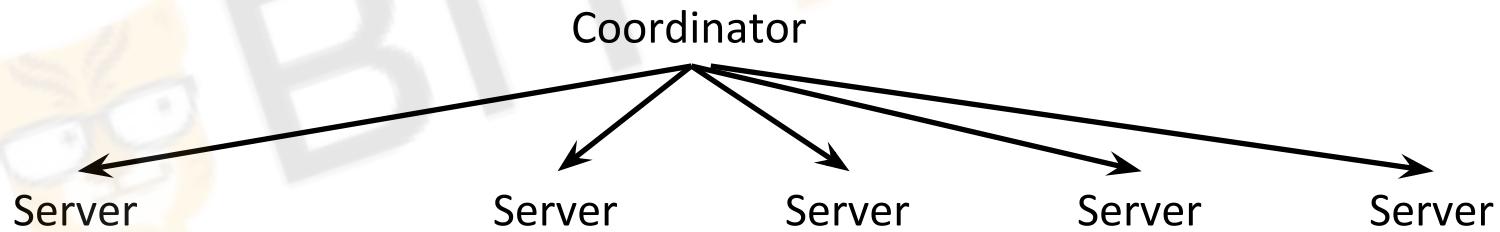
- Before finishing write on master, no response could be gotten from other servers.
- Append-only system, no need to worry update and choice between different versions.
- If allowing update, a choice between consistency and availability must be made.

# Evolve – Multiple Machine

Problem: What if a large amount of data to store?

Demand: A huge storage capacity for data, high RPS for both insert and lookup

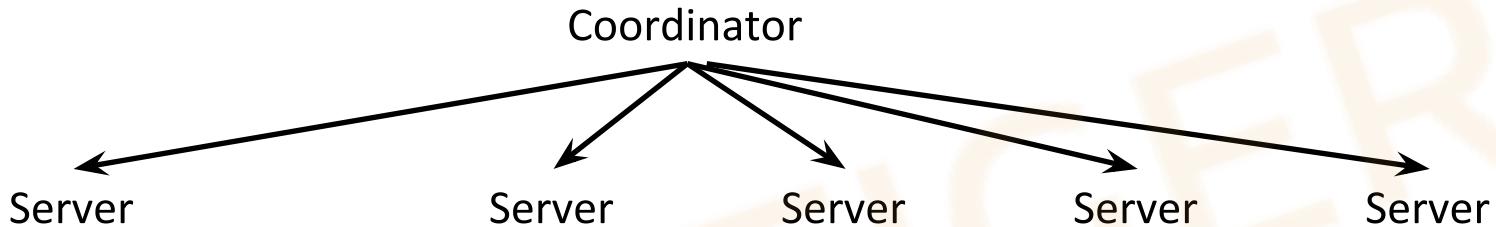
## SHARDING



- Do both lookup and insert operations
- Store part of the data in each server
- Coordinator decide which server to send (a hash function from id to server)

# Evolve – How to dispatch?

## SHARDING



### No id conflicts in all servers – mod

- Server #0 ~ 4, only save url that  $\text{url\_count} \% 5 == \#$
- For write request, dispatch it randomly. In server i,  $\text{actual\_size} = \text{size} * 5 + i$ , use  $\text{actual\_size}$  to calculate  $\text{short\_url}$
- For lookup request, compute  $\text{short\_url}$  to  $\text{actual\_size}$ , use  $\text{actual\_size} \% 5$  to decide which server to dispatch

### Con

Cannot eliminate duplicates for same long\_url

# Summary

- SNAKE Analysis
- Application:
  - traditional hash
  - base10
  - base62
- Evolve:
  - CAP Theorem in Distributed System
  - Replication & Sharding