

Informatics II, Spring 2023, Solution 12

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Task 1

1. A. Note C is no cycle since cycle must not visit an vertex twice
2. B
3. C, D
4. A, D

Graphs

5. FALSE. G_2 contains 2 cycles and is thus not acyclic. The first cycle is formed by the vertices 3 - 5 - 3 (Backedge). The second cycle is 3 - 4 - 5 - 3.
6. TRUE. This statement can easily be verified, since the adjacency matrix is not symmetric which must be the case for undirected Graphs. $G_2(1,2)$ shows there is an edge from 1 to 2 but $G_2(2,1)$ shows there is no edge back from 2 to 1.
7. TRUE. The Dijkstra algorithm can only be used for Graphs with non-negative edge weights. As soon as one edge is negative only one iteration of the Dijkstra Algorithm may return a suboptimal result. Therefore the Bellman-Ford must be used which basically runs the Dijkstra algorithm multiple times to adjust for the negative edge weights.

MST and SSSP

8. C
9. A
10. B, D

Task 2: Finding cycle in undirected Graph

```
algo: DFS-Acyclic(G, s) isAcyclic = TRUE;  
  
s.col = G;  
foreach  $u \in s.adj$  do  
    if  $u.col == W$  and  $isAcyclic == TRUE$  then  
         $u.pred = s$ ;  
         $isAcyclic = isAcyclic$  and DFS-Acyclic(G,u);  
    end  
    else if  $u.col == G$  and  $u \neq s.pred$  then  
        return FALSE;  
    end  
end  
s.col = B;  
return isAcyclic;
```

Algorithm 1: DFS-Acyclic

Task 3: Count all possible Paths between two Vertices

```
algo: Count-Paths(M, visited, s, t) PathCount = 0;  
  
visited[s] = true;  
for  $i = 0$  to  $M.length$  do  
    if  $M[s][i] == 1$  then  
        if  $i == t$  then  
            PathCount = PathCount + 1;  
        end  
        else if  $visited[i] == false$  then  
            PathCount = PathCount + CountPaths(M, visited, i, t);  
        end  
    end  
end  
visited[s] = false;  
return PathCount;
```

Algorithm 2: Count-Paths

Task 4: k-hop [20 FS Final Exam]

1.

Adjacency Matrix:

	1	2	3	4	5	6
1	0	1	0	1	1	0
2	0	0	1	0	1	0
3	1	0	0	0	0	0
4	0	0	0	0	0	1
5	0	0	0	0	0	0
6	0	0	0	0	0	0

2.

- The 2-hop neighbors of node 1 are nodes 3 and 6.
- Run a BFS search until first node in the queue has a distance of k . Dequeue and report nodes as long as they have a distance equal to k .

3. **Pseudocode:**

Algorithm: khop(G, v, k)

```
for ( $i = 1; i \leq n; i++$ ) do  $dist[i] = -1$  ;  
;  
 $dist[v] = 0$ ;  
InitQueue(Q);  
while ( $v \neq -1 \ \&\& \ dist[v] < k$ ) do  
    for ( $i = 1; i \leq n; i++$ ) do  
        if ( $a[v, i] == 1 \ \&\& \ dist[i] == -1$ ) then  
             $dist[i] = dist[v] + 1$ ;  
            Enqueue(Q, i);  
     $v = Dequeue(Q)$ ;  
while  $v \neq -1 \ \&\& \ dist[v] == k$  do  
    print(v);  
     $v = Dequeue(Q)$ ;
```

C code:

```
1 void neighbors(int edges[7][2], int m, int n, int v, int k){
2
3     int graph[n + 1][n + 1];
4
5     for(int i = 0; i < m; i++) {
6         graph[edges[i][0]][edges[i][1]] += 1;
7     }
8
9     int distance[n + 1]; // visited[i] means minimum distance from v to i .
10    node_t *head = NULL;
11
12    enqueue(&head, v);
13    distance[v] = -1;
14
15    int steps = 1;
16    // k times BFS
17    while(steps <= k) {
18        node_t *new_head = NULL;
19        int node;
20        while ((node=dequeue(&head)) > 0) {
21            for(int i = 1; i <= n; i++){
22                if(graph[node][i] > 0 && distance[i] == 0) {
23                    distance[i] = steps; // update minimum distance.
24                    enqueue(&new_head, i);
25                }
26            }
27        }
28        head = new_head;
29        steps += 1;
30    }
31    printf("%d-hop_neighbors_of_node_%d is:\n", k, v);
32    for(int i = 1; i <= n ; i++) {
33        if(distance[i] == k) {
34            printf("node_%d\n", i);
35        }
36    }
37 }
```