Graph algorithms in their many shapes and sizes

Daniel Epstein, 5/8/14

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https://github.com/depstein/programming-competitions

Housekeeping

Amazon Programming Contest

Thanks for your interest in the programming contest on May 13th. This page has all the details you need to make the competition fun.

Prizes

First place gets a Kindle Fire HDX!

Second and third place get Kindle Paperwhites!

All competitors get awesome Amazon swag!

Rules

- 1. Any current UW student is eligible to participate. Registration is limited to the first 120 participants.
- 2. Individual competitors only.
- 3. Previous winners can compete but are ineligible for prizes.
- 4. Contest runs 2 hours and includes 4 different short problems.
- Programming in Java 7.
- The only reference material allowed is this website and Oracle's online Java API: http://docs.oracle.com/javase/7/docs/api/. No
 other materials, online or printed, are allowed.

The number of players is limited, so it is first come first served! You will receive an email confirmation of your spot within 72 hours of signing up.

Dates and Locations

We will have food, drinks, and the opportunity to meet some of our developers at the end of the competition.

Date: Tuesday, May 13, 2014

Time: 6:00 PM

Location: Paul Allen Center for CSE ATRIUM

https://github.com/depstein/programming-competitions

\$\mathcal{P} branch: master ▼ programming-com	petitions / problems / dynamic programming / 🕒	• History
Added solutions to Problem 2728 and 4123		
ping128 authored 8 days ago		latest commit 41b420e535 🚉
-		
2728 (A Spy on the Metro)	Added solutions to Problem 2728 and 4123	8 days ago
4123 (Glenbow Museum)	Added solutions to Problem 2728 and 4123	8 days ago
4131 (Currency Shopping)	Solution to Problem 4131	10 days ago
a 4213 (DNA Sequences)	Solutions to Problem 4213, 4877, 4905, 5945, 6088	10 days ago
4280 (Pencils from the Ninet	Solution to Problem 4280	13 days ago
4794 (Sharing Chocolate)	Yeah, solved Problem 4794	9 days ago
4877 (Non-Decreasing Digits)	Solutions to Problem 4213, 4877, 4905, 5945, 6088	10 days ago
4905 (Pro-Test Voting)	Solutions to Problem 4213, 4877, 4905, 5945, 6088	10 days ago
5945 (Raggedy, Raggedy)	Solutions to Problem 4213, 4877, 4905, 5945, 6088	10 days ago
6088 (Approximate Sorting)	Solutions to Problem 4213, 4877, 4905, 5945, 6088	10 days ago

Practice Contest

- 5 hours (+setup time)
- Real problem set!
- Real submission system
- Teams of 3
- Food?

Practice Contest

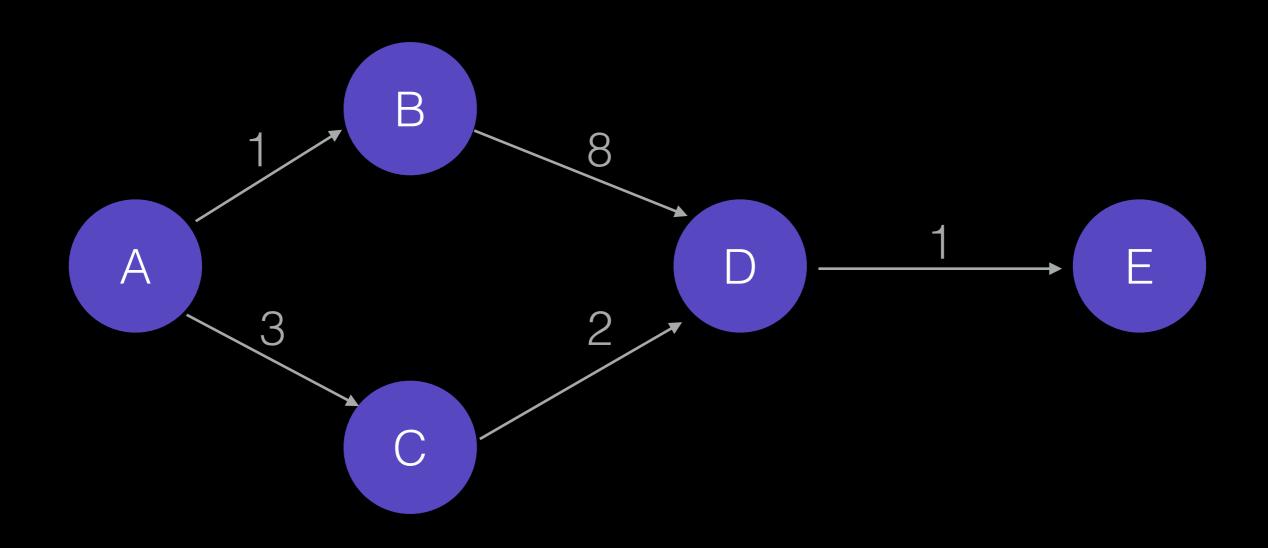
- 5/24, 5/25 (Memorial Day weekend)
- Monday 5/26 (Memorial Day)
- 5/31, 6/1
- 6/7, 6/8 (Weekend before finals)
- 6/9-6/13 (Finals week)

In mathematics and computer science, **graph theory** is the study of *graphs*, which are mathematical structures used to model pairwise relations between objects. A "graph" in this context is made up of "vertices" or "nodes" and lines called *edges* that connect them. A graph may be *undirected*,

In mathematics and computer science, **graph theory** is the study of *graphs*, which are mathematical structures used to model pairwise relations between objects. A "graph" in this context is made up of "vertices" or "nodes" and lines called *edges* that connect them. A graph may be *undirected*,

Nodes and Edges!

Graph Representation



Arrays

```
import java.io.*;
import java.util.*;
public class arrays {
  public static int[][] dist = new int[5][5];
  public static void main(String[] args) {
    for (int i=0;i<5;i++)</pre>
      for(int j=0;j<5;j++) {</pre>
        if(i != j)
          dist[i][j] = 1000; // Not using Integer.MAX VALUE to avoid integer overflowing
    //Initialize graph as described
    dist[0][1] = 1;
    dist[0][2] = 3;
    dist[1][3] = 8;
    dist[2][3] = 2;
    dist[3][4] = 1;
                                                          В
                                                                         D
```

Arrays

```
import java.io.*;
import java.util.*;
* Know the max number of nodes
 public static int[][] dist = new int[5][5];
 public static void main(String[] args)
   Adjaceney Matrix
       dist[i][j] = 1000; // Not using Integer.MAX VALUE to avoid integer overflowing
        Easy to look up edge weight/edge existence
   Slow to get all edges out of a node
       Uses N^2 memory
```

I tend to only use them for all-pairs, shortest path

Dedicated classes

```
class Node {
   public ArrayList<Edge> neighbors = new ArrayList<Edge>();
class Edge {
   public Node dest;
   public int distance;
   public Edge(Node dest, Node source, int distance) {
      this.dest = dest;
      this.distance = distance;
      source.neighbors.add(this);
                                                           В
```

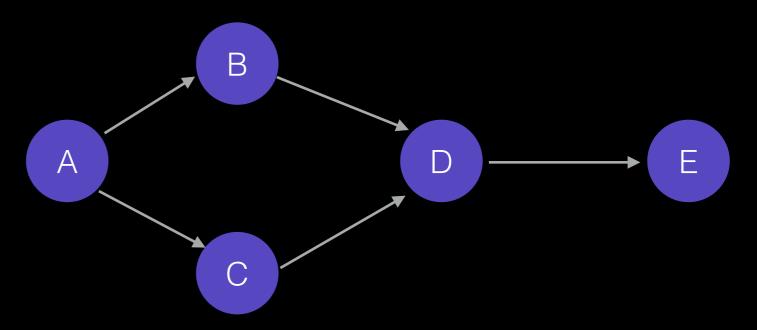
Dedicated classes

Well-organized

```
class Nede Adjacency Lists
   public ArrayList<Edge> neighbors = new ArrayList<Edge>();
   Edge of Easy to get all edges out of a node public Node dest;
   public int distance;
   public edge Difficult touting a particular edge
      this.distance = distance;
      • This is what I use for weighted graphs
```

Unweighted Graphs

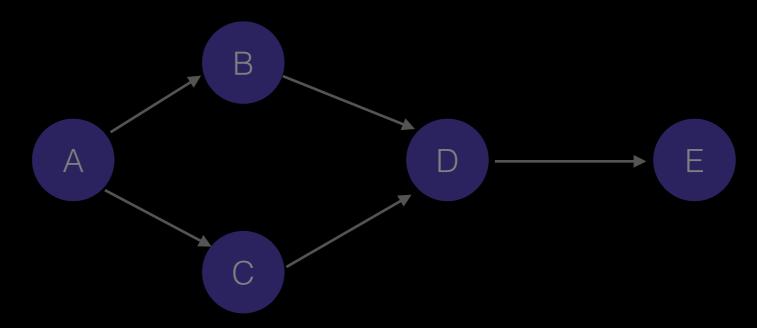
```
class Node {
   public ArrayList<Node> neighbors = new ArrayList<Node>();
}
```



Unweighted Graphs

Easy. You don't need an edge class!

```
class Node {
   public ArrayList<Node> neighbors = new ArrayList<Node>();
}
```



HashMaps

//For a weighted graph

```
class Node {
   public HashMap<Node, ArrayList<Node, Integer>> neighbors = new HashMap<Node,</pre>
ArrayList<Node, Integer>>();
//For a weighted graph, naming the nodes with Strings
public HashMap<String, HashMap<String, Integer>> neighbors = new HashMap<String,</pre>
HashMap<String, Integer>>();
//For an unweighted graph, naming the nodes with Strings
public HashMap<String, ArrayList<String>> neighbors = new HashMap<String,</pre>
ArrayList<String>>();
                                                             В
```

HashMaps

A little harder to organize, debug

//For a weighted graph

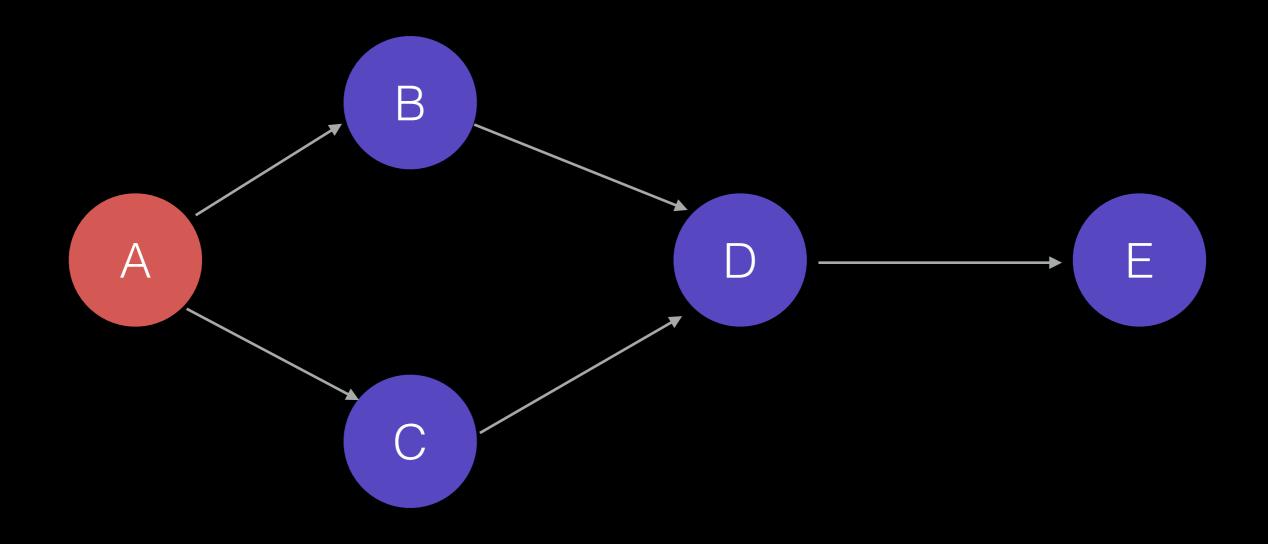
```
class Node {
Public But If you get used to it, can write shorter code
//For a weighted graph, naming the nodes with Strings
public HashMap<String, HashMap<String, Integer>> neighbors = new HashMap<String,</pre>
HashMap<String, Integer>>();
//For an unweighted graph, naming the nodes with Strings
public HashMap<String, ArrayList<String>> neighbors = new HashMap<String,</pre>
ArrayList<String>>();
                                                        В
```

Graph Algorithms

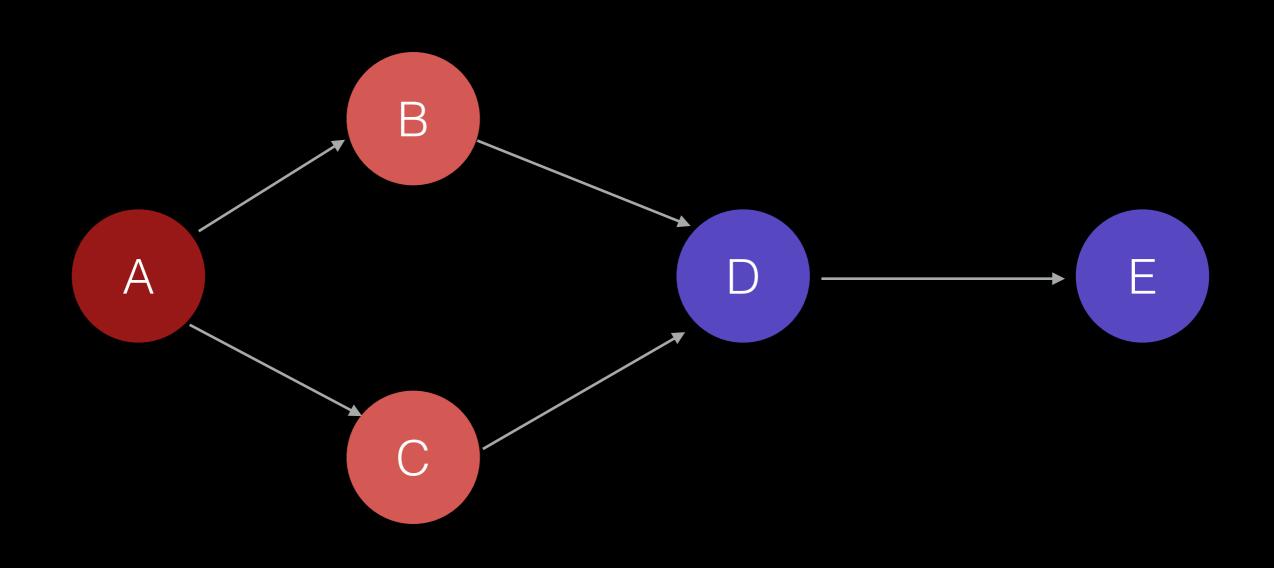
- Searches (Breadth-First, Depth-First)
- Shortest Path (Dijkstra's)
- Minimum Spanning Tree (Prim's, Kruskal's)
- Topological Sort
- Negative-Edge Shortest Path (Bellman-Ford)
- All-Pairs, Shorest Path (Floyd-Warshall)
- Max Flow (Ford-Fulkerson, Edmonds-Karp, Preflow Push)
- Minimum-Cost Flow

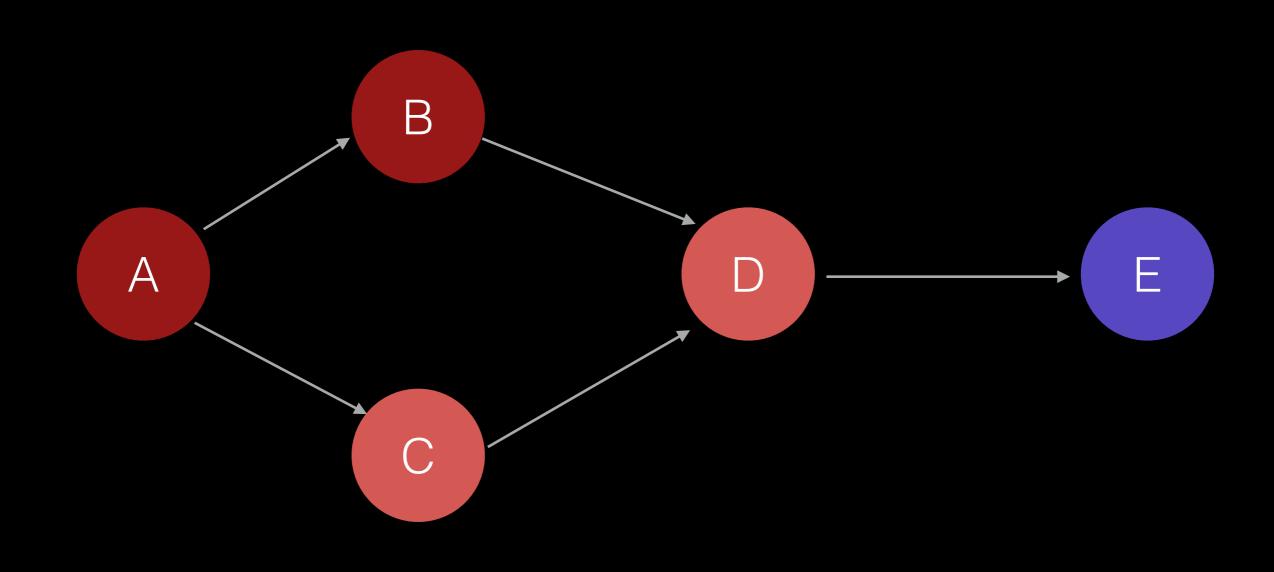
Graph Algorithms

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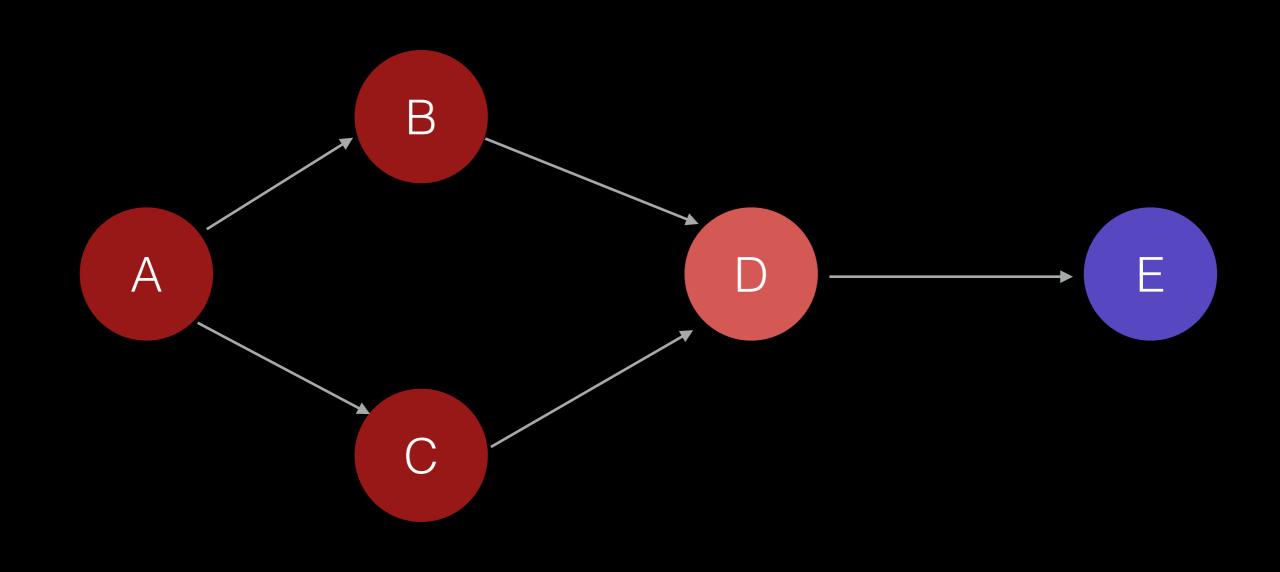
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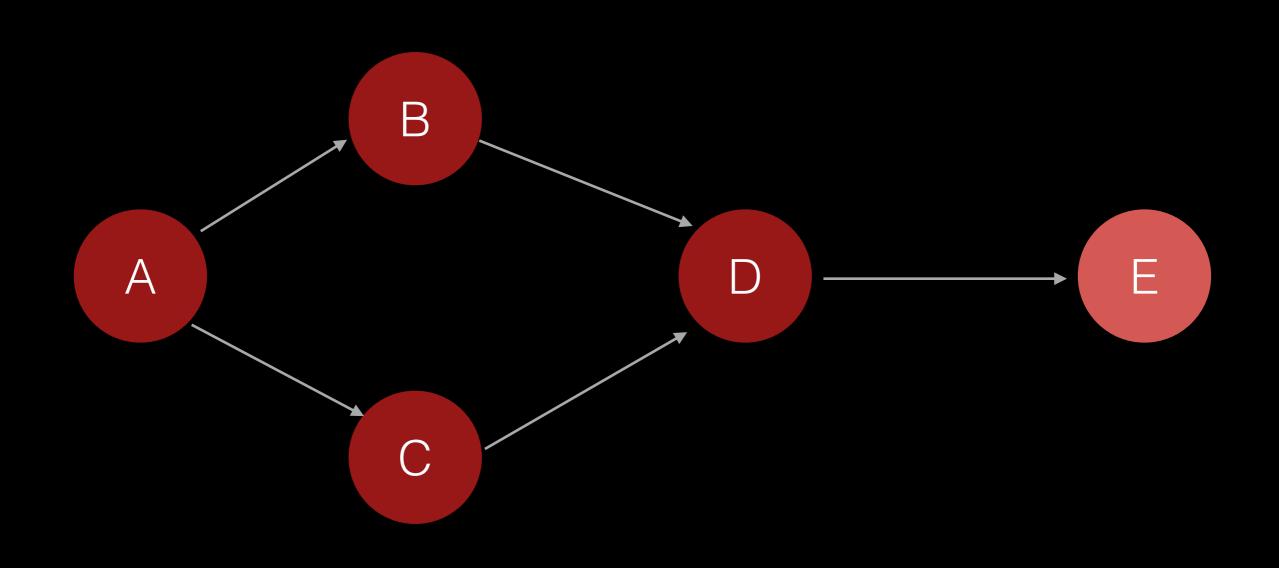
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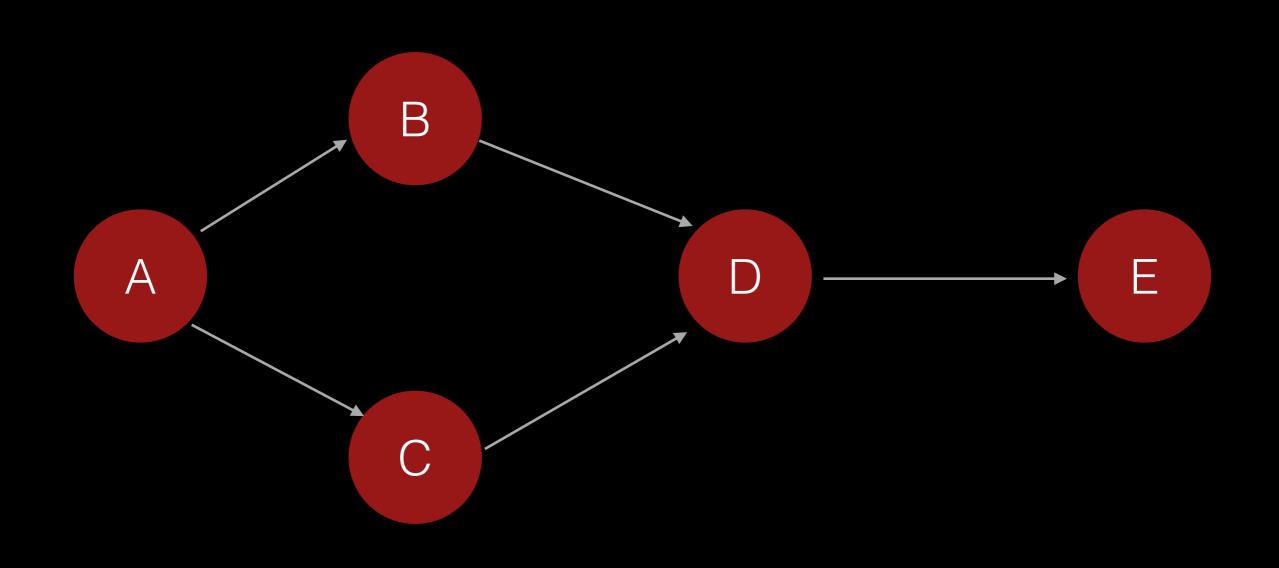


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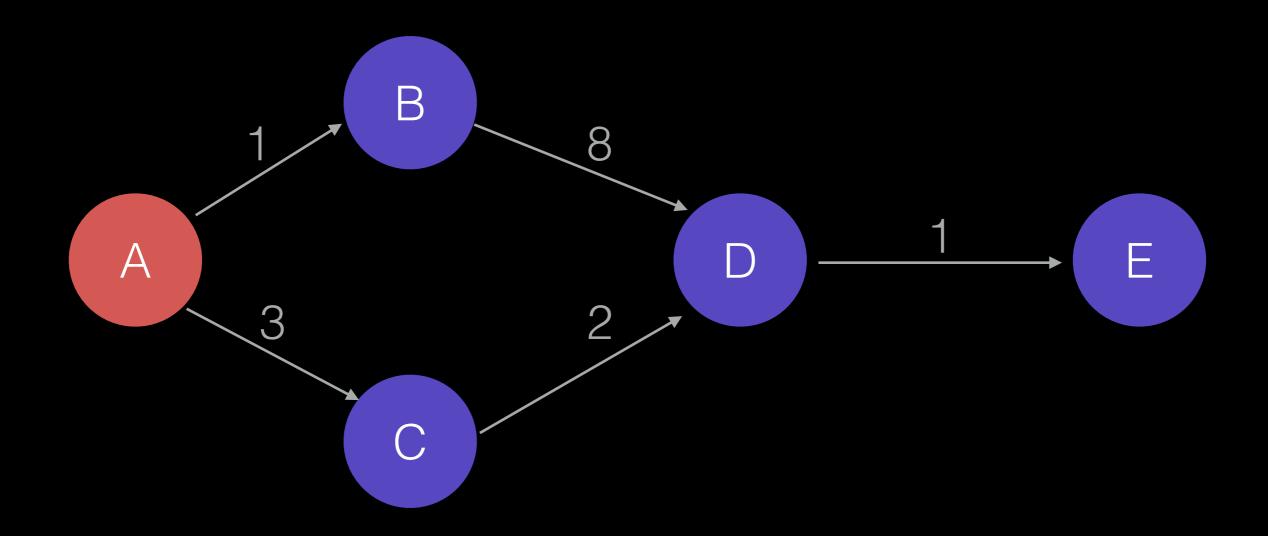
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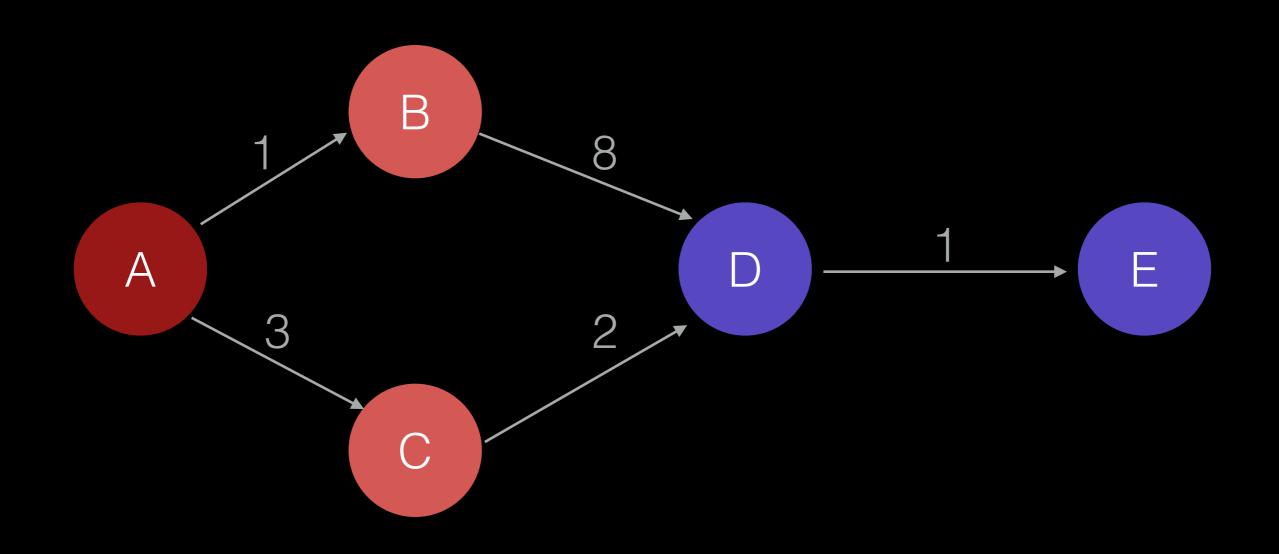
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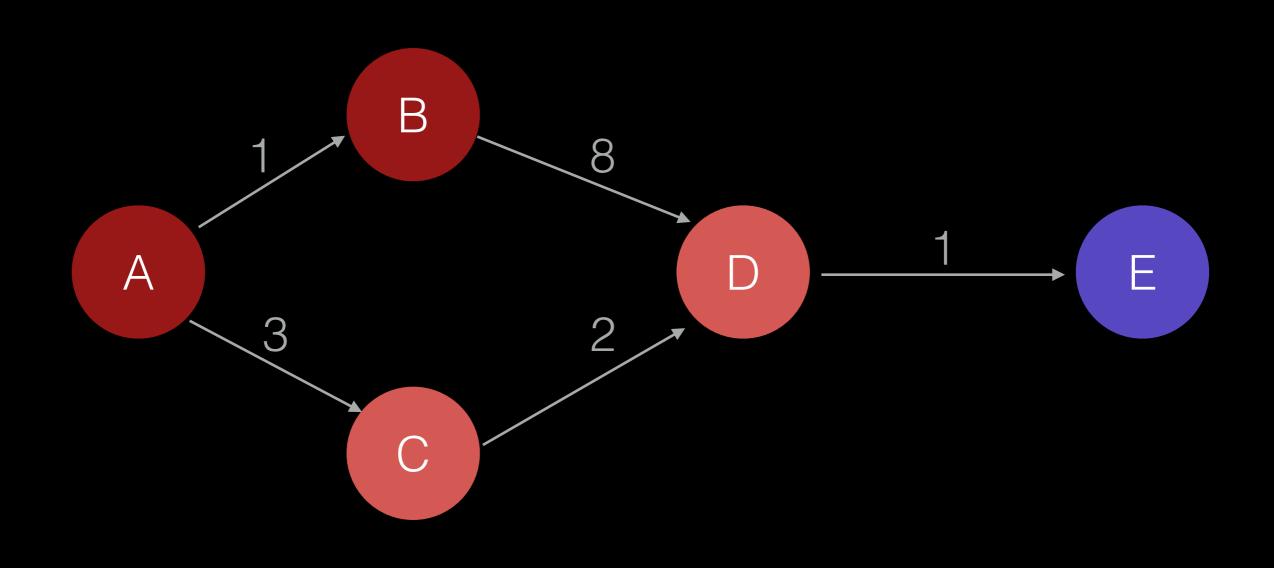
В

```
import java.io.*;
import java.util.*;
public class bfs {
  public static void main(String[] args) {
    Node a = new Node();
   Node b = new Node();
   Node c = new Node();
   Node d = new Node();
   Node e = new Node();
    a.edges.add(b);
    a.edges.add(c);
   b.edges.add(d);
    c.edges.add(d);
    d.edges.add(e);
   ArrayList<Node> allNodes = new ArrayList<Node>(Arrays.asList(new Node[]{e, d, c, b, a}));
   bfs(a, allNodes);
    System.out.printf("Distance to Node e is: %d\n", e.distance);
  public static void bfs(Node root, ArrayList<Node> allNodes) {
    Queue<Node> q = new LinkedList<Node>();
    root.distance = 0;
    q.add(root);
    while(q.size() > 0) {
     Node u = q.poll();
      for (Node n : u.edges) {
        if (n.distance == Integer.MAX VALUE) { // Has not been visited yet
          n.distance = u.distance + 1;
          q.add(n);
class Node implements Comparable<Node> {
  public ArrayList<Node> edges = new ArrayList<Node>();
 public int distance = Integer.MAX VALUE;
  public int compareTo(Node o) {
    return (distance < o.distance) ? -1 : ((distance == o.distance) ? 0 : 1);
```



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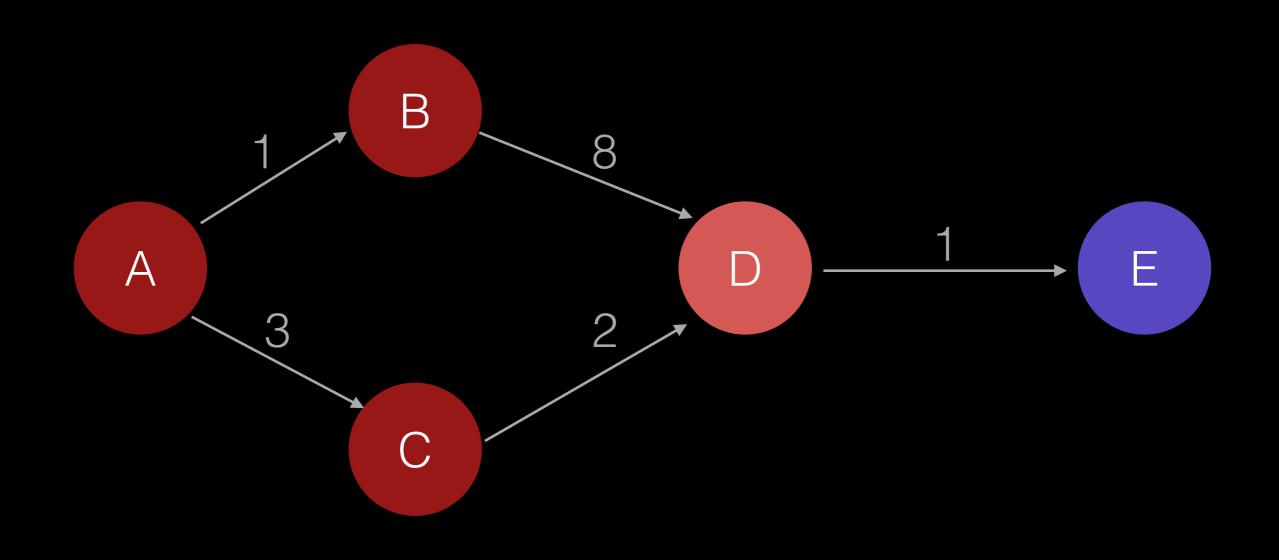




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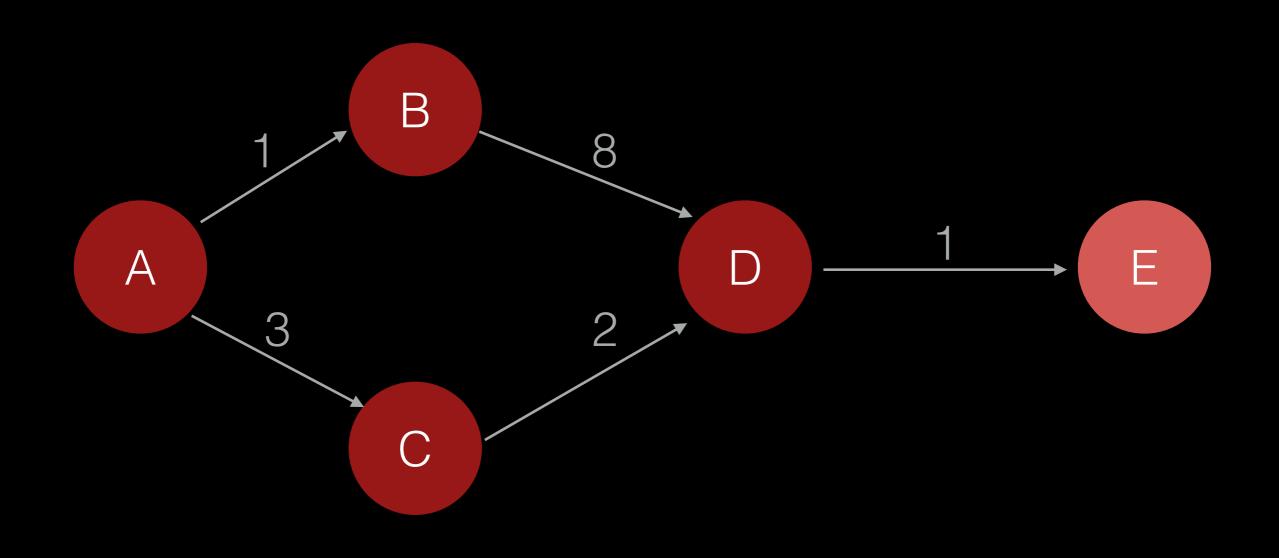
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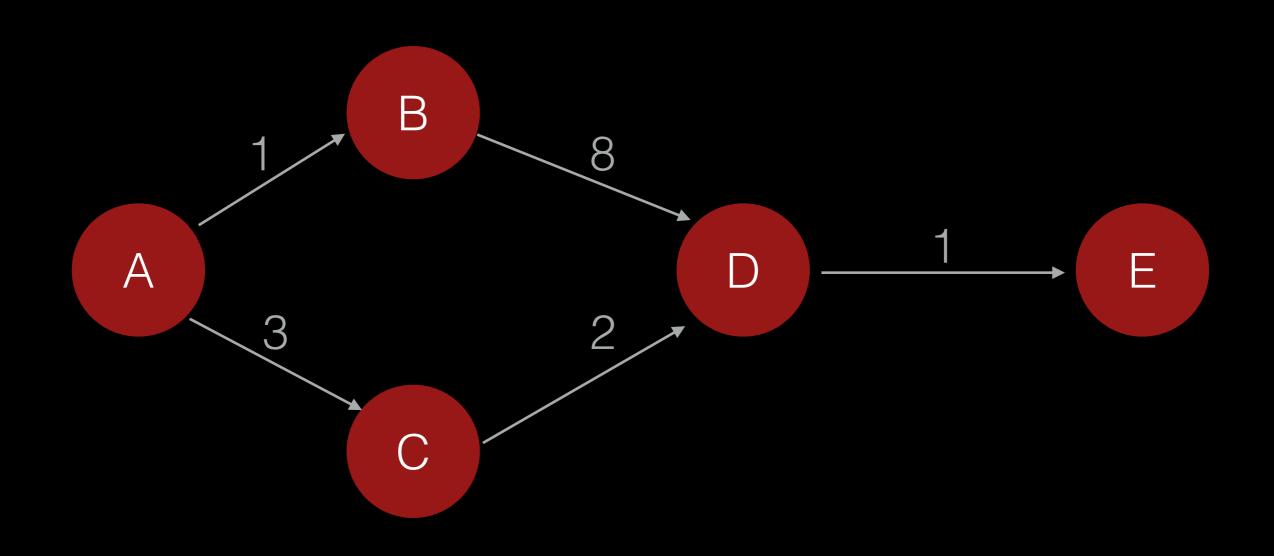
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```
import java.io.*;
import java.util.*;
public class dijkstra {
  public static void main(String[] args) {
   Node a = new Node();
   Node b = new Node();
   Node c = new Node();
   Node d = new Node();
   Node e = new Node();
   a.edges.put(b, 1);
   a.edges.put(c, 3);
   b.edges.put(d, 8);
   c.edges.put(d, 2);
    d.edges.put(e, 1);
   ArrayList<Node> allNodes = new ArrayList<Node>(Arrays.asList(new Node[]{e, d, c, b, a}));
    dijkstra(a, allNodes);
   System.out.printf("Distance to Node e is: %d\n", e.distance);
  public static void dijkstra(Node root, ArrayList<Node> allNodes) {
    PriorityQueue<Node> q = new PriorityQueue<Node>();
    root.distance = 0;
    q.add(root);
    while (q.size() > 0) {
     Node u = q.poll();
      for (Node n : u.edges.keySet()) {
       if(n.distance == Integer.MAX VALUE) { // Update the distance to node n
          q.remove(n);
       n.distance = Math.min(n.distance, u.distance + u.edges.get(n));
        q.add(n);
class Node implements Comparable<Node> {
  public HashMap<Node, Integer> edges = new HashMap<Node, Integer>();
  public int distance = Integer.MAX VALUE;
  public int compareTo(Node o) {
    return (distance < o.distance) ? -1 : ((distance == o.distance) ? 0 : 1);
```

```
public static void bfs(Node root, ArrayList<Node> allNodes) {
    Queue < Node > q = new LinkedList < Node > ();
    root.distance = 0;
    q.add(root);
    while (q.size() > 0) {
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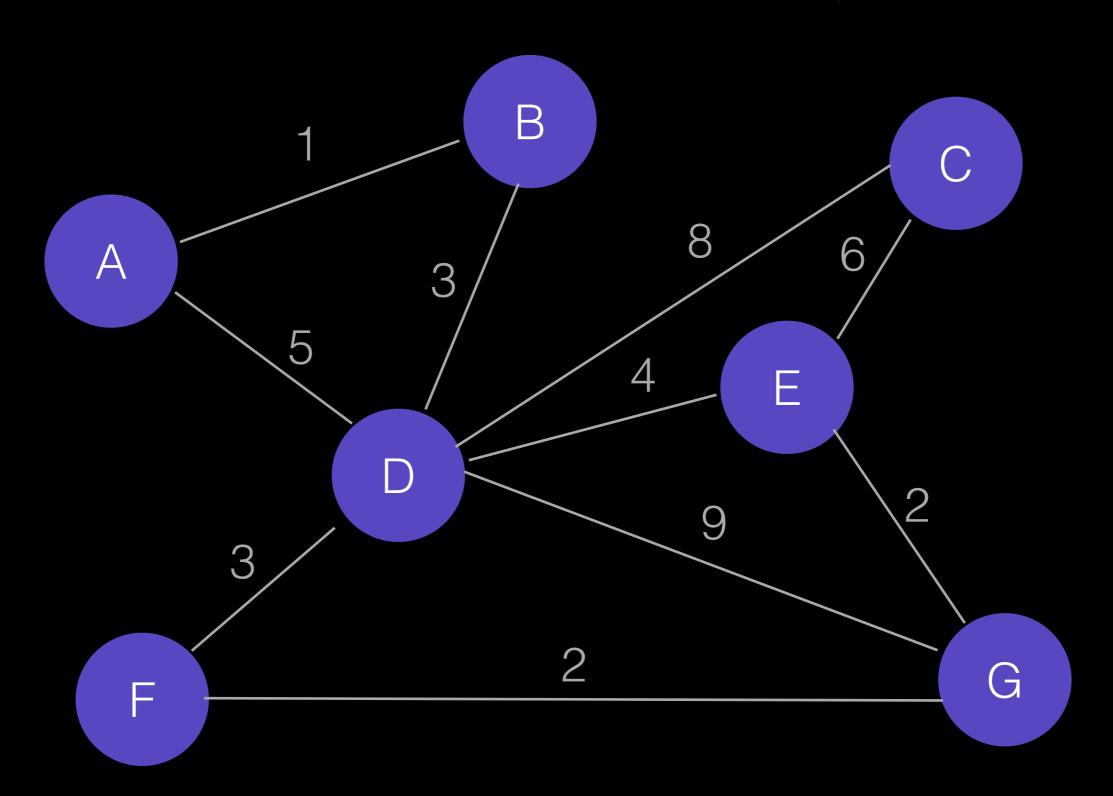
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```

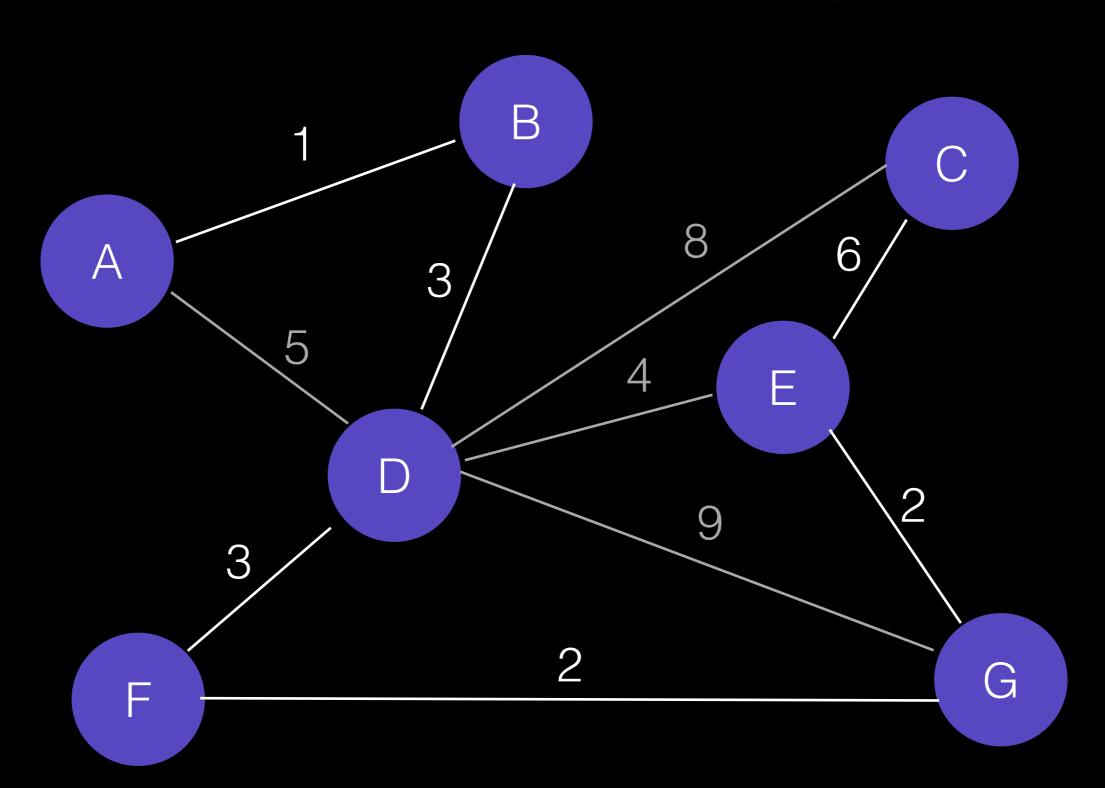
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        q.remove(n);
      n.distance = Math.min(n.distance, u.distance + u.edges.get(n));
      q.add(n);
```

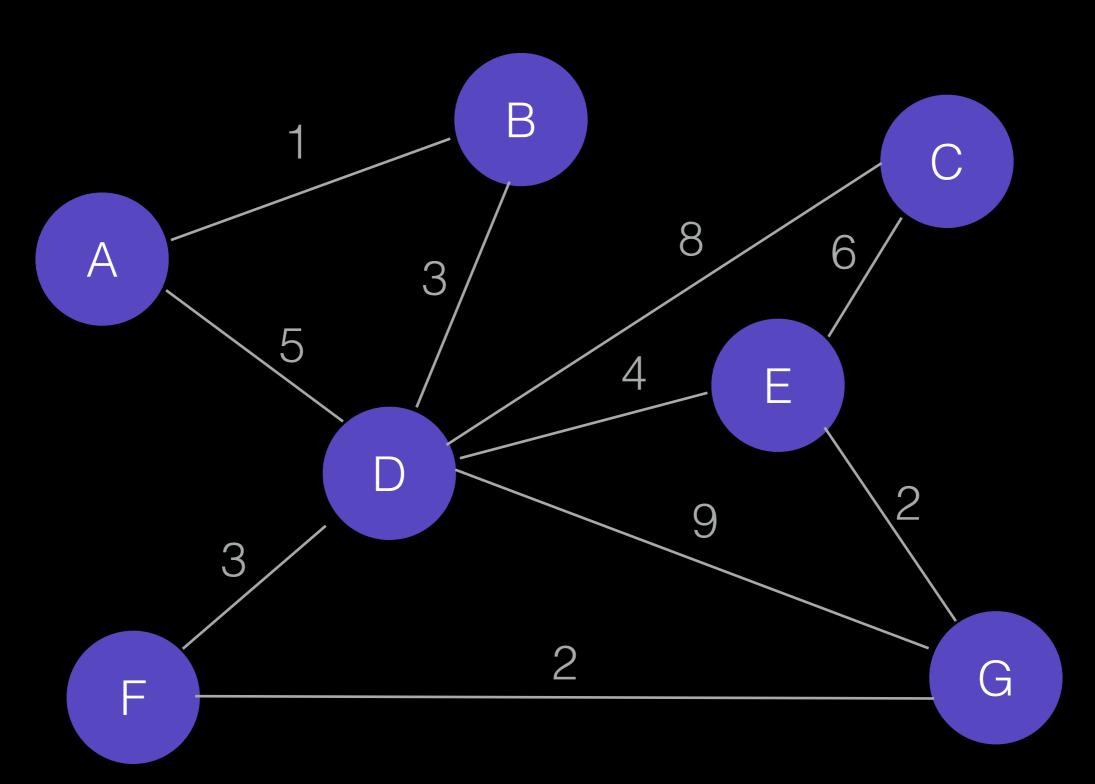
Minimum Spanning Tree

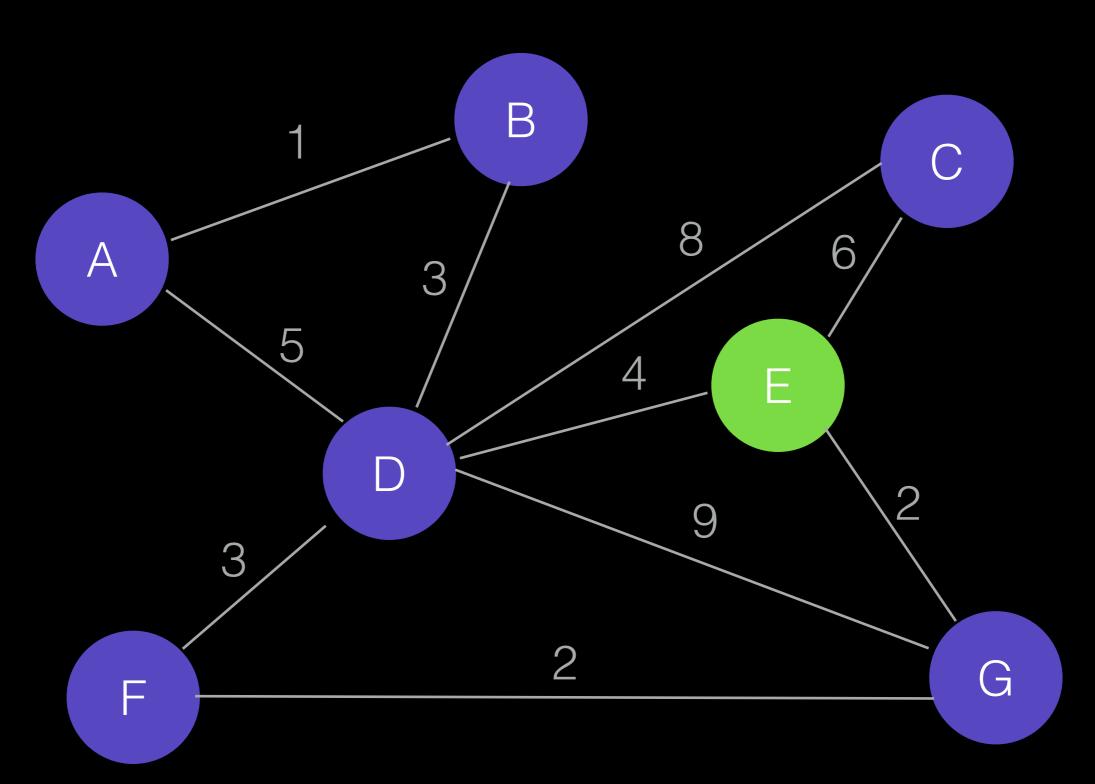


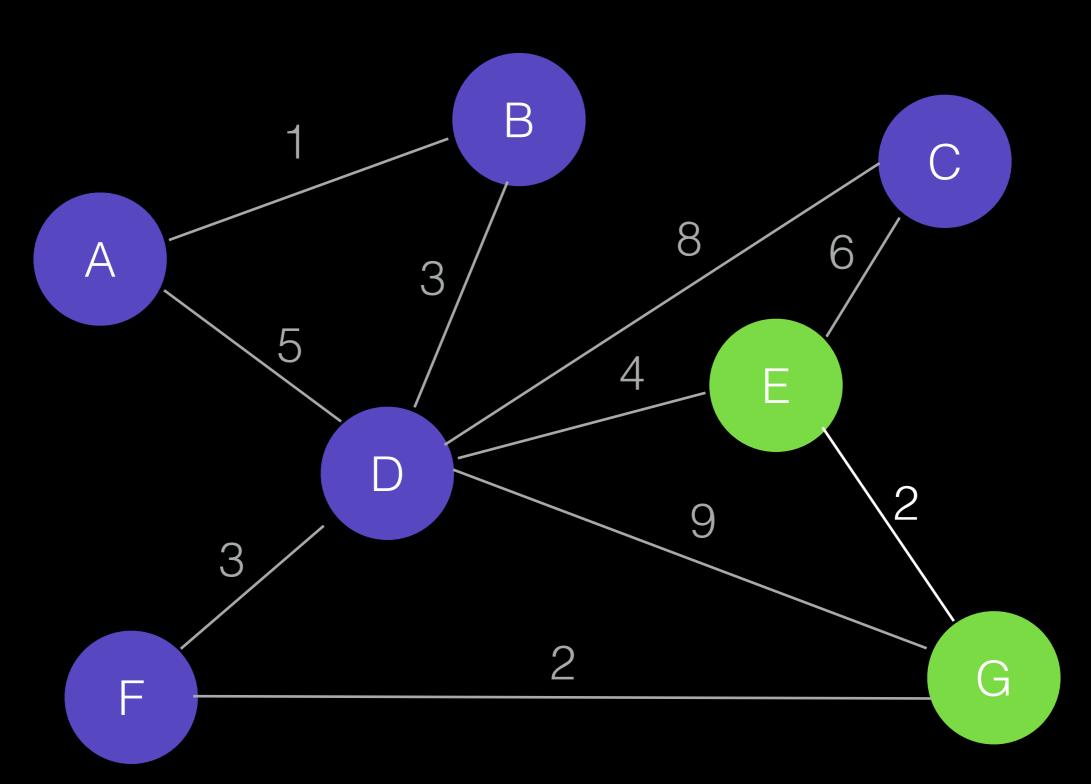
Minimum Spanning Tree

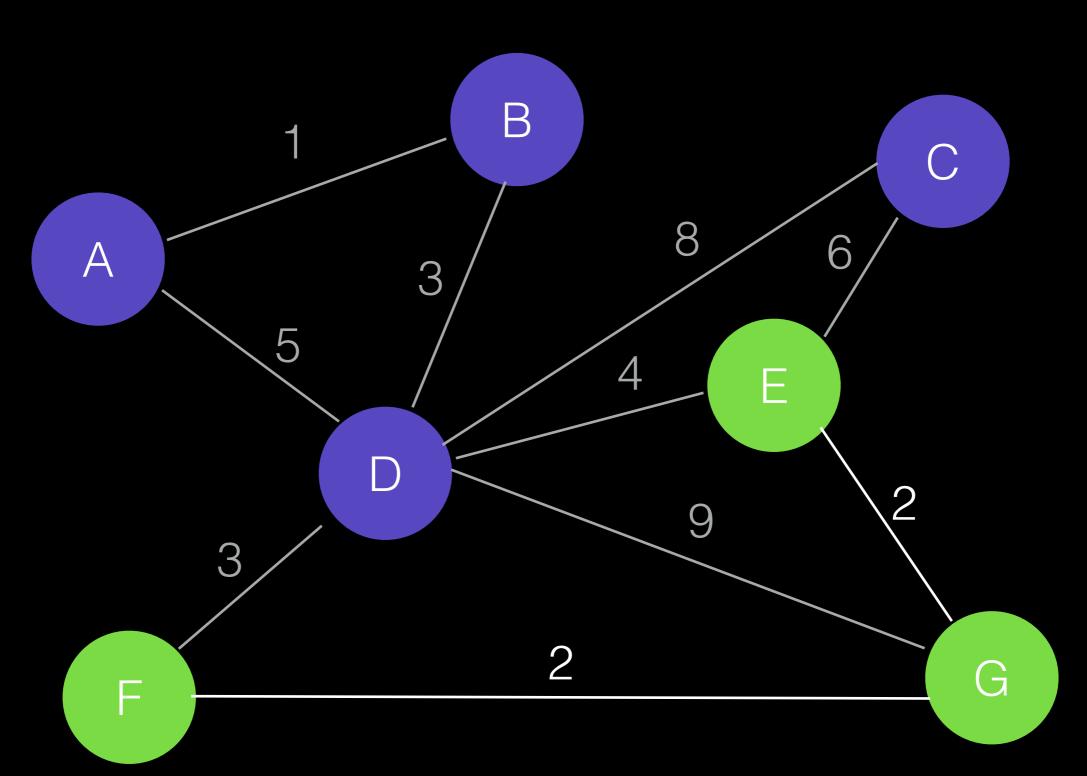


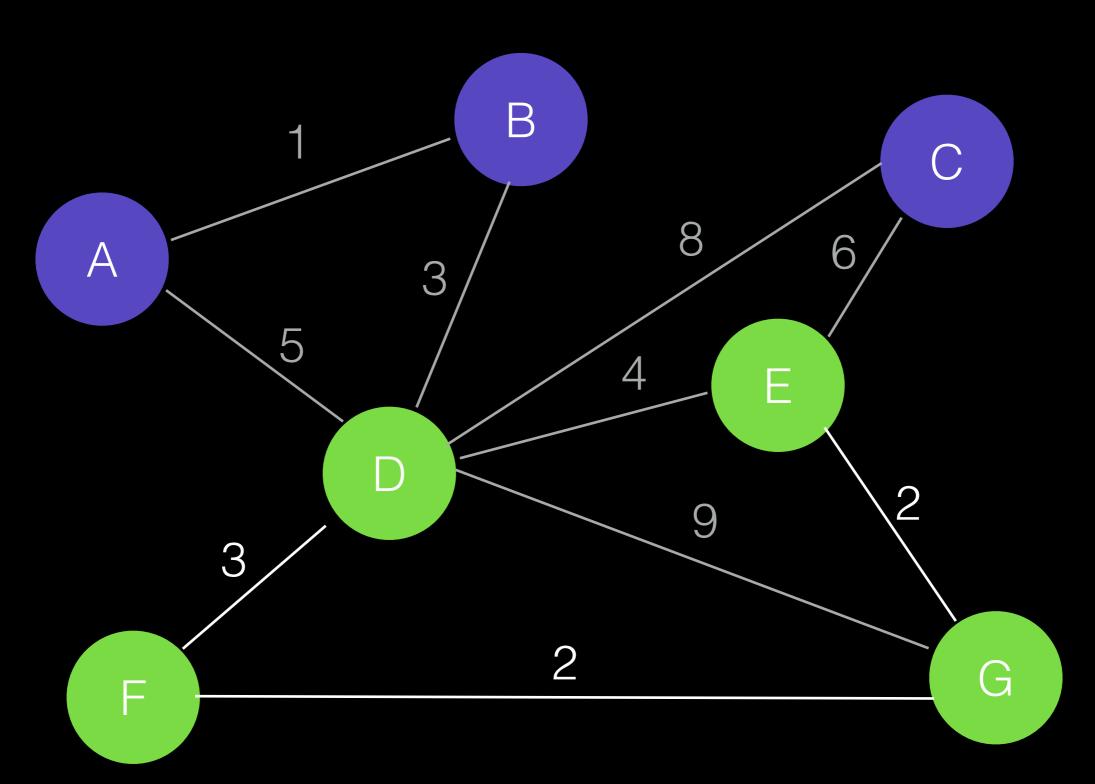
- Initialize a tree with an arbitrarily-chosen node
- Grow tree by one edge
 - Select edge of minimal weight that connects tree to a node not in the tree
- Rinse and repeat until all nodes are in the tree

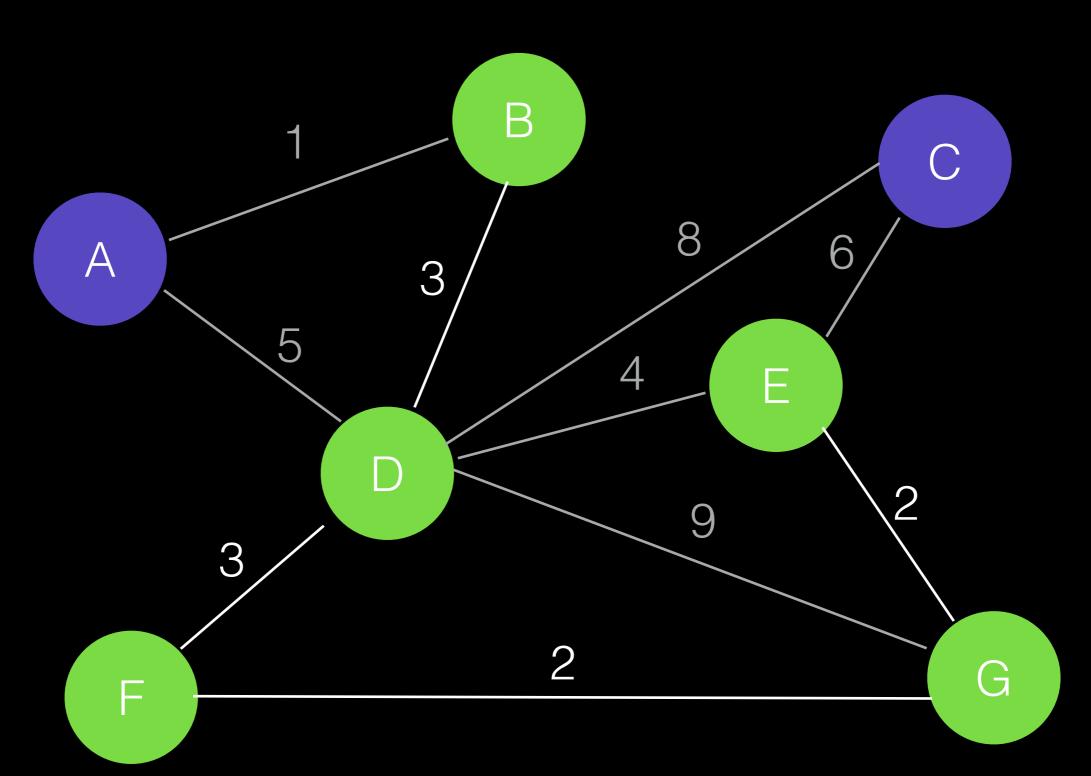


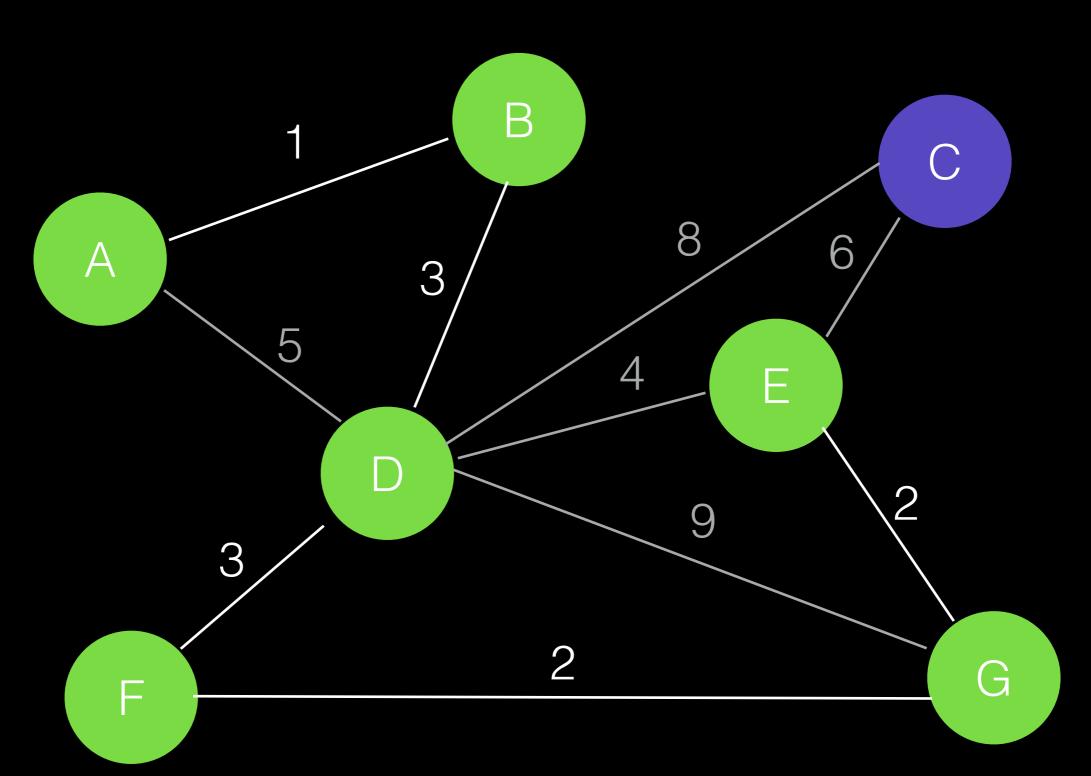


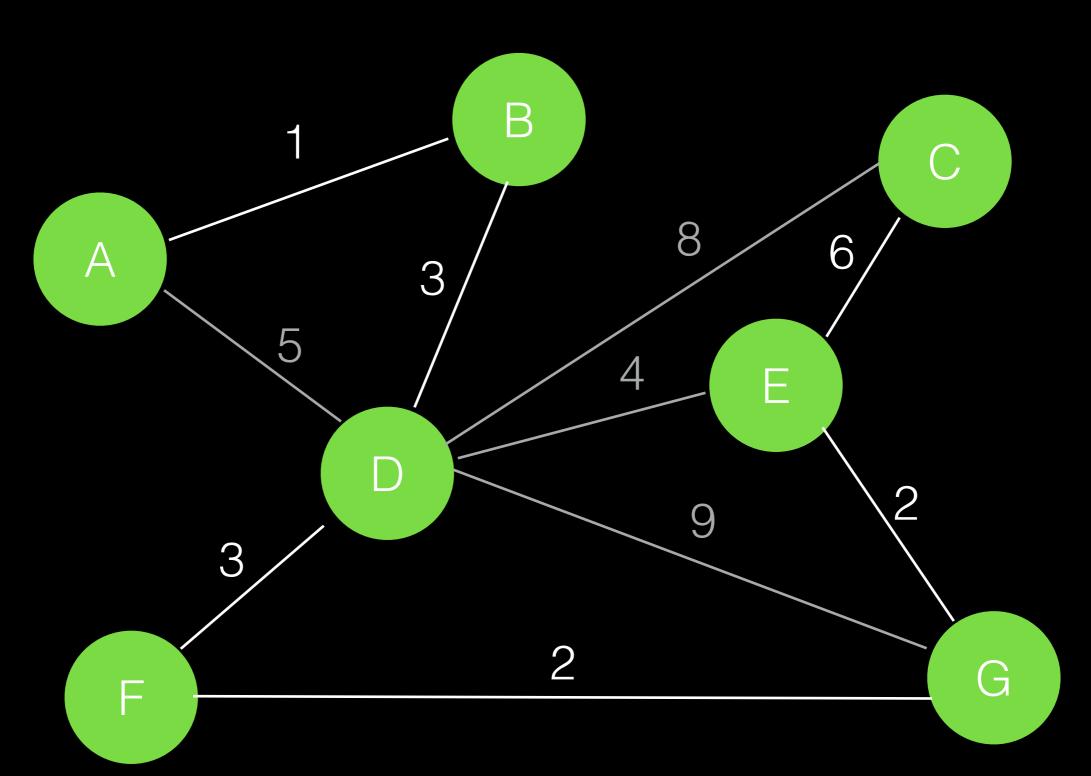


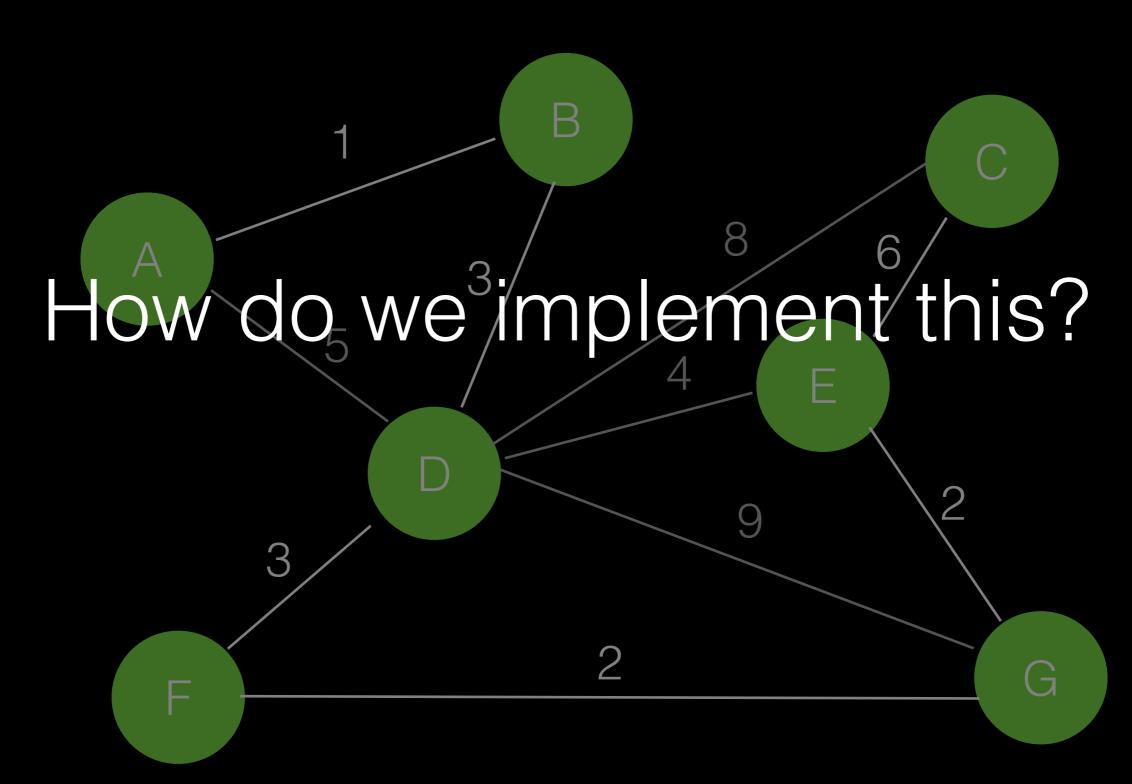












- Initialize a tree with an arbitrarily-chosen node
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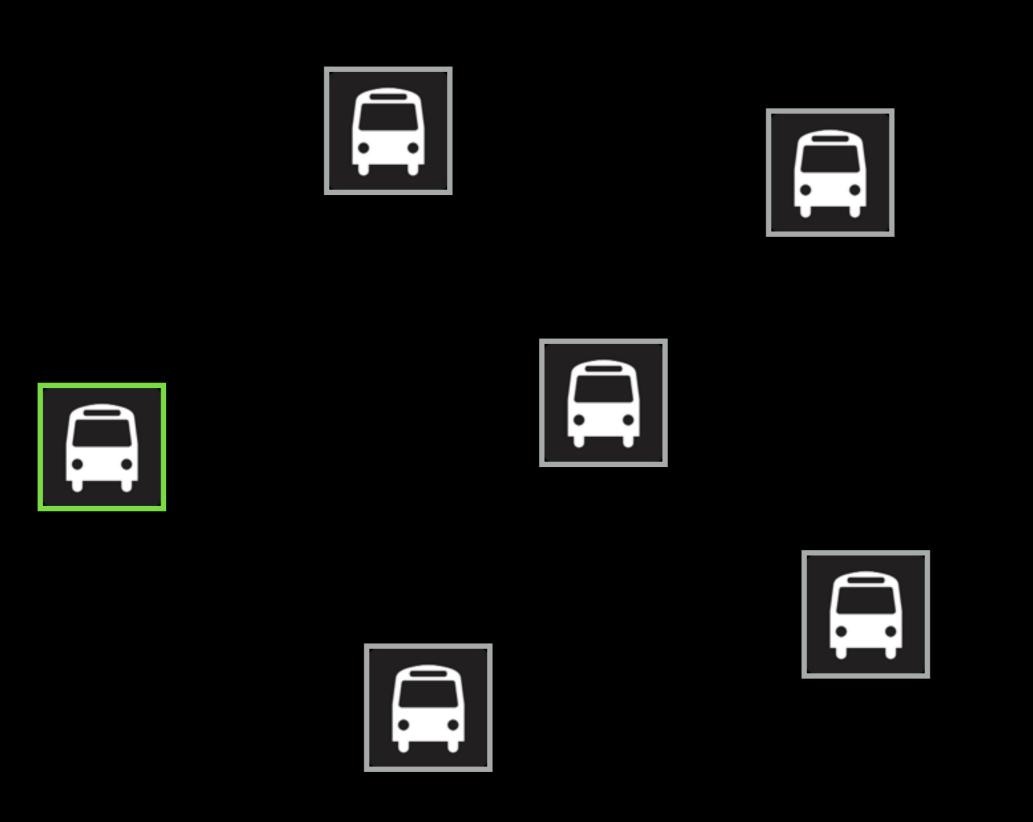
```
public static void prims(ArrayList<Node> allNodes) {
   PriorityQueue<Node> q = new PriorityQueue<Node>();
   Node root = allNodes.get(0);
   root.distance = 0;
   q.add(root);

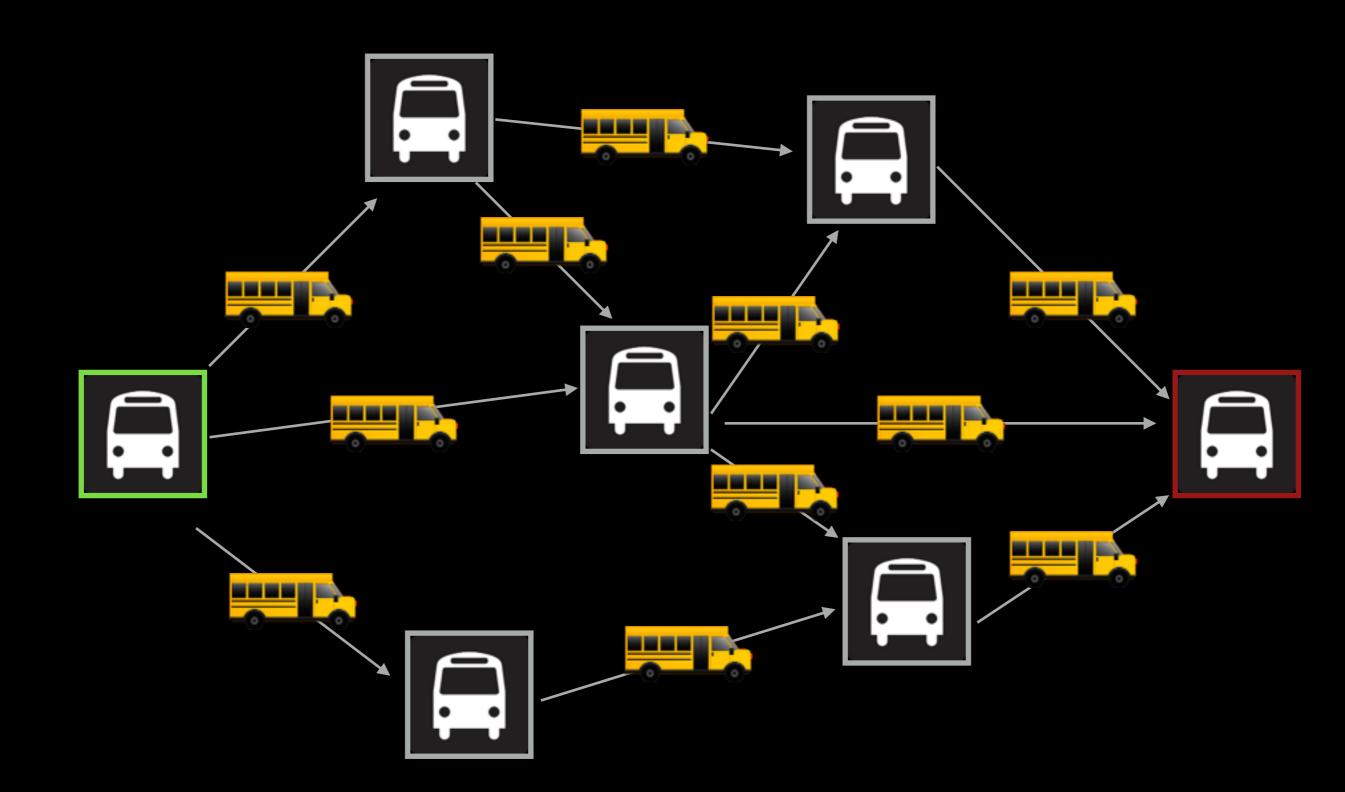
while(!q.isEmpty()) {
    Node cur = q.poll();
    cur.visited = true;
   for(Node n : cur.edges.keySet()) {
        if(!n.visited && n.distance > cur.edges.get(n)) {
            n.distance = cur.edges.get(n);
            q.remove(n);
            q.add(n);
        }
    }
}
```

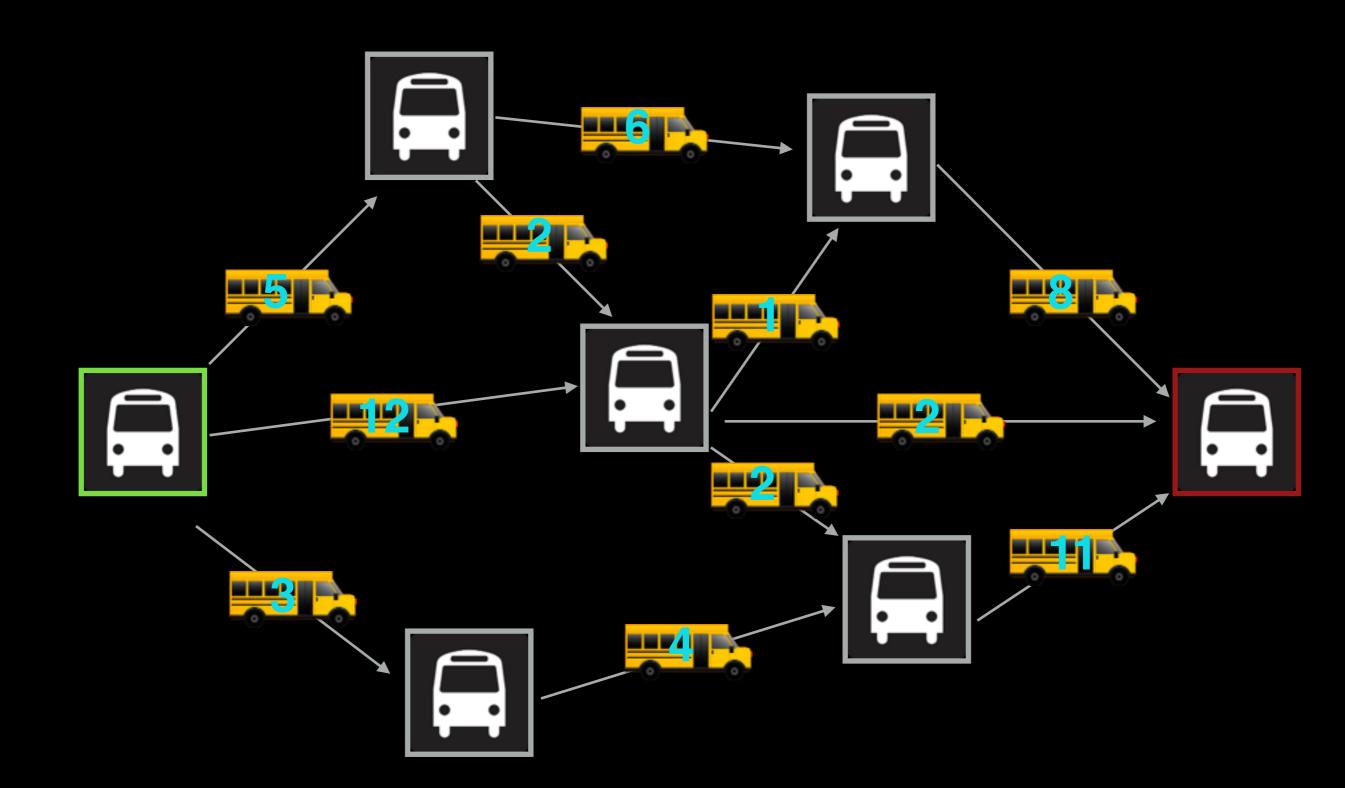
Maximum Flow

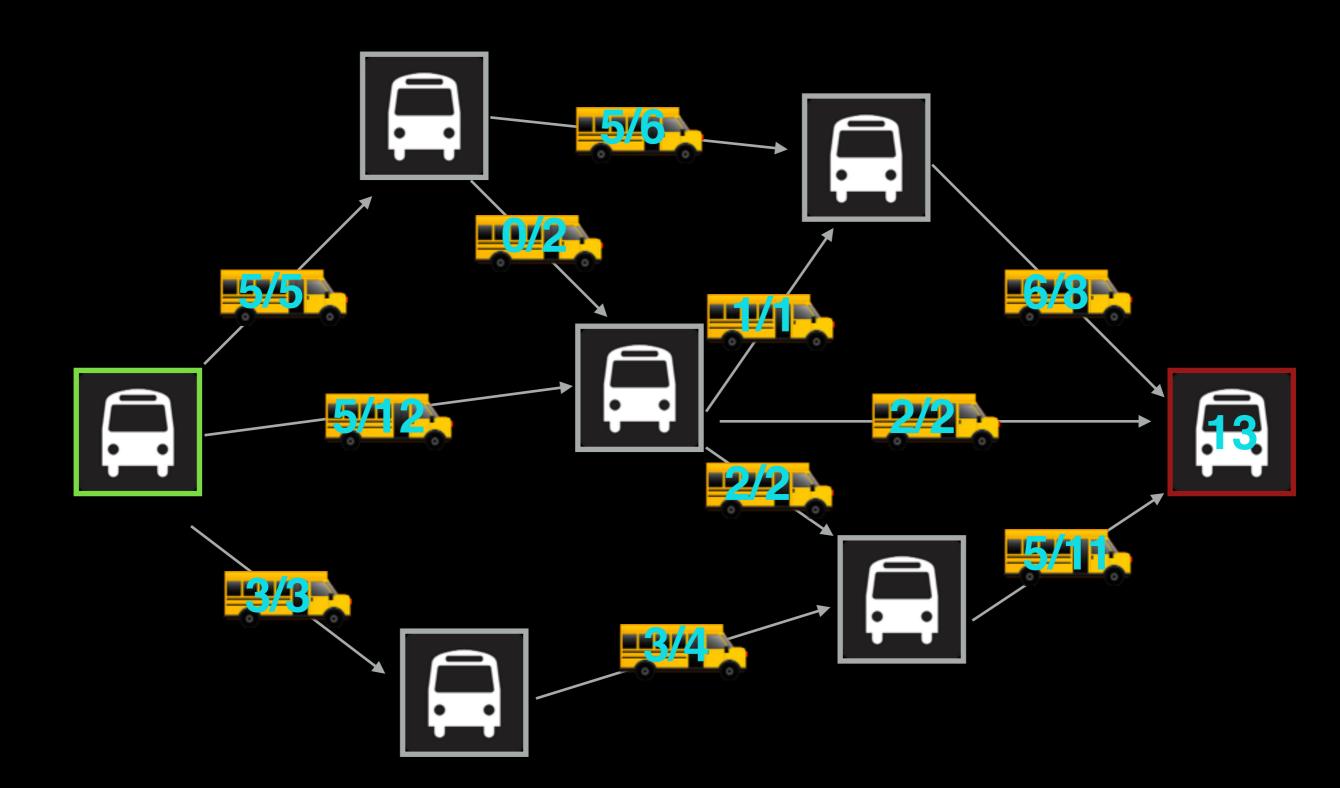










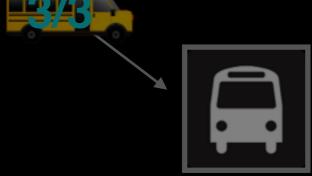




Attempt to push flow across an edge or set of edges/5

Provide a way to backtrack and undo that operation







Maximum Elow

Method	Complexity	Description
Linear programming		Constraints given by the definition of a legal flow. See the linear program here.
Ford-Fulkerson algorithm	O(E maxl fl)	As long as there is an open path through the residual graph, send the minimum of the residual capacities on the path. The algorithm works only if all weights are integers. Otherwise it is possible that the Ford–Fulkerson algorithm will not converge to the maximum value.
Edmonds-Karp algorithm	O(VE ²)	A specialization of Ford-Fulkerson, finding augmenting paths with breadth-first search.
Dinitz blocking flow algorithm	O(V ^E E)	In each phase the algorithms builds a layered graph with breadth-first search on the residual graph. The maximum flow in a layered graph can be calculated in $O(VE)$ time, and the maximum number of the phases is n -1. In networks with unit capacities, Dinic's algorithm terminates in $O(E\sqrt{V})$ time.
General push- relabel maximum flow algorithm	O(V ^E E)	The push relabel algorithm maintains a preflow, i.e. a flow function with the possibility of excess in the vertices. The algorithm runs while there is a vertex with positive excess, i.e. an active vertex in the graph. The push operation increases the flow on a residual edge, and a height function on the vertices controls which residual edges can be pushed. The height function is changed with a relabel operation. The proper definitions of these operations guarantee that the resulting flow function is a maximum flow.
Push-relabel algorithm with FIFO vertex selection rule	O(V ⁸)	Push-relabel algorithm variant which always selects the most recently active vertex, and performs push operations until the excess is positive or there are admissible residual edges from this vertex.



Maximum Flow

```
public static void PreflowPush(Node source, Node sink)
    // source.height should be the number of nodes, sink.height should be 1
    LinkedList<Node> q = new LinkedList<Node>();
    q.add(source);
    while (!q.isEmpty()) {
       Node node = q.remove();
       int minHeight = Integer.MAX VALUE;
        for (int i = 0; i < node.neighbors.size(); i++) {</pre>
           if (node.neighbors.get(i).dest.height < minHeight</pre>
                   && node.neighbors.get(i).remaining() > 0)
               minHeight = node.neighbors.get(i).dest.height;
       if (minHeight != Integer.MAX VALUE && minHeight >= node.height)
            node.height = minHeight + 1;
        for (int i = 0; i < node.neighbors.size(); i++) {</pre>
           if (node.neighbors.get(i).dest.height < node.height) {</pre>
               int pushedFlow = node.neighbors.get(i).remaining();
               if (pushedFlow > node.excess)
                   pushedFlow = node.excess;
               if (pushedFlow > 0) {
                   node.neighbors.get(i).flow += pushedFlow;
                   node.neighbors.get(i).back.flow -= pushedFlow;
                   node.neighbors.get(i).dest.excess += pushedFlow;
                   node.excess -= pushedFlow;
                   if (!node.neighbors.get(i).dest.isSourceOrSink)
                       q.add(node.neighbors.get(i).dest);
                   if (node.excess <= 0)</pre>
                       break;
       if (node.excess > 0 && !node.isSourceOrSink)
           q.add(node);
```

Graph algorithms in their many shapes and sizes

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https://github.com/depstein/programming-competitions