

# CS4500

# Operating Systems

Week 3



University of Colorado  
Colorado Springs

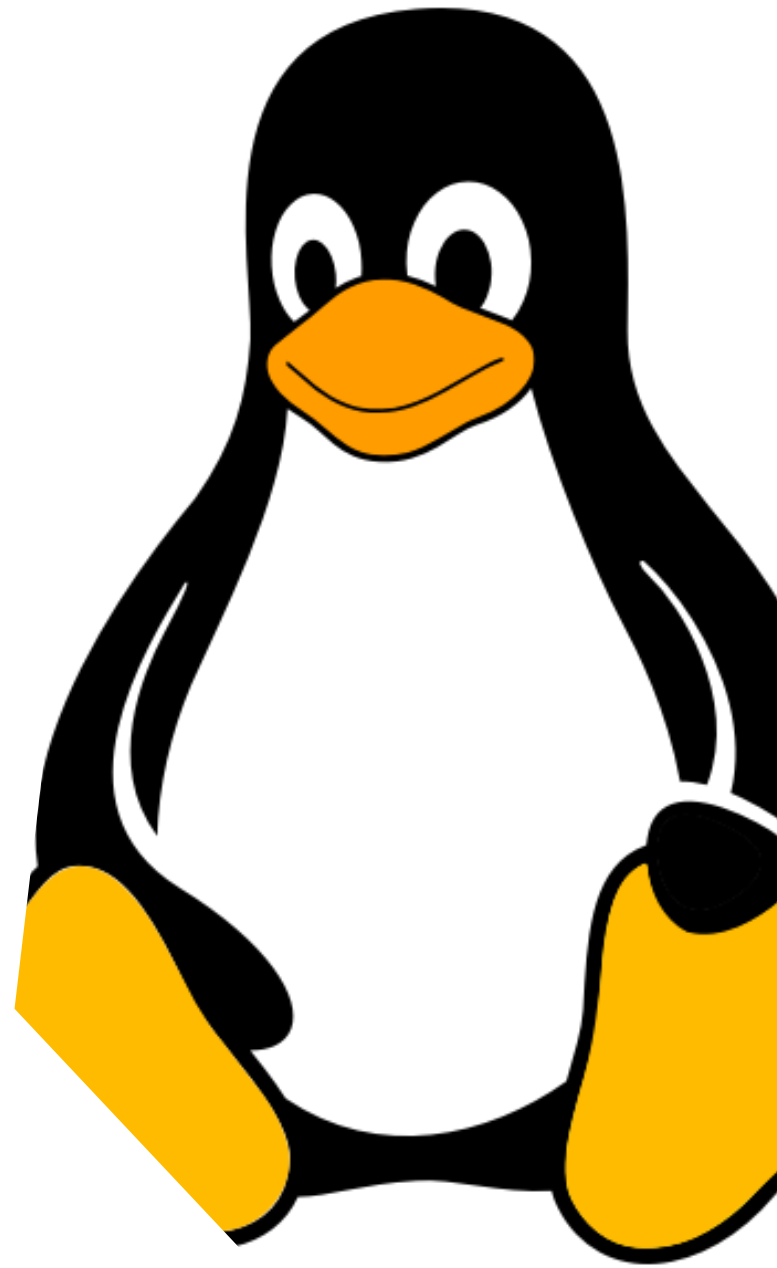


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# Introduction to Operating Systems: Processes



# What to expect today:

An overview of processes

Assigned reading

- Andrew S. Tanenbaum, Modern Operating Systems, 4thd edition, 2014, Prentice Hall Chapter 2

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# Operating Systems

Three Easy Pieces

Remzi Arpaci-Dusseau  
Andrea Arpaci-Dusseau

## Additional Resource

Remzi H. Arpaci-Dusseau, R. H., Arpaci-Dusseau, A.C. *Operating systems: Three easy pieces: Chapter 10 and 28*. (2018). Arpaci-Dusseau Books.

<https://pages.cs.wisc.edu/~remzi/OSTEP/cpu-sched-multi.pdf>

Copy of .pdf available in Canvas



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# Recap of the Last Class

## Computer hardware

- Time-multiplexed – resource not divided into units
- Space-multiplexed – resource split in pieces, then assigned a process

## OS components

- Process management
- Memory management
- File and storage management

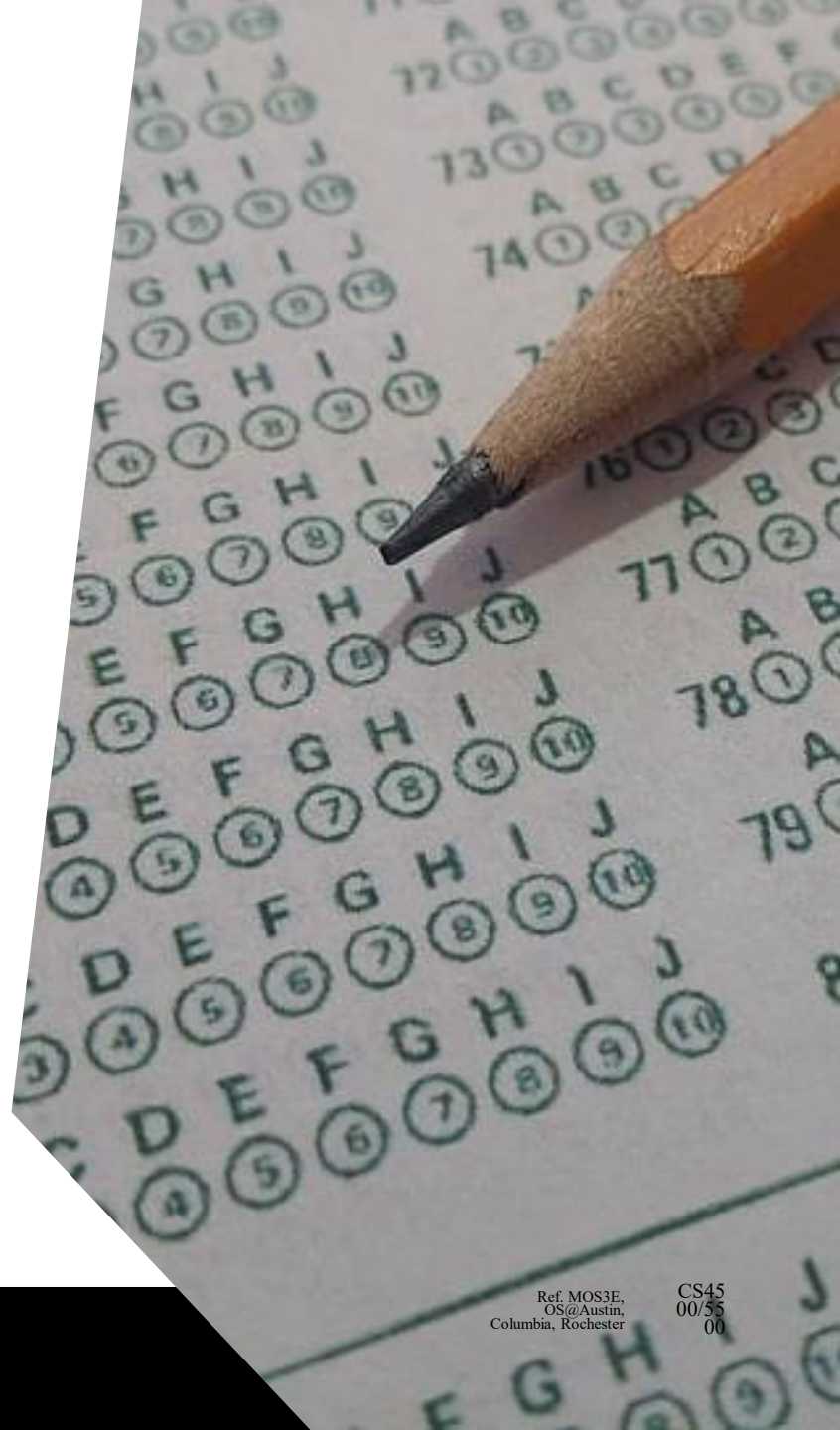
# Which of the following are likely components of an operating system?

- A. Cache memory
- B. Scheduler
- C. File editor
- D. File system
- E. Device driver
- F. Web browser



Which of the following are likely components of an operating system?

- A. Cache memory
- B. **Scheduler**
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# Process

Process == Program??

# Process

## Definitions

- An instance of a program running on a computer
- An abstraction that supports running programs
- An *execution stream* in the context of a particular *process state*
- A *sequential* stream of execution in its *own address space*

Ref. CS4  
MOS3E, 500/  
OS@Austin, 550  
Columbia, 0  
Rochester



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# Process

Two parts of a process

- Sequential execution of instructions
- Process state
  - Registers: PC, SP,...
  - Memory: address space, code, stack, heap ...
  - I/O status: opened files ...

# Program vs. Process

Program != Process

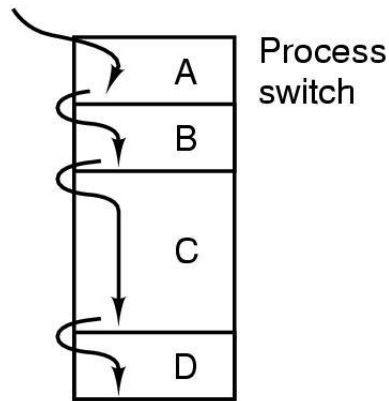
- Program = static code + data
- Process = dynamic instantiation of code + data + files
- ...

No 1:1 mapping

- A program can invoke many processes
  - Running the same program twice
  - A program contains `fork()`

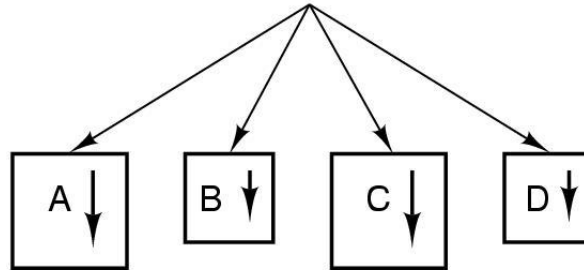
# The Process Model

One program counter

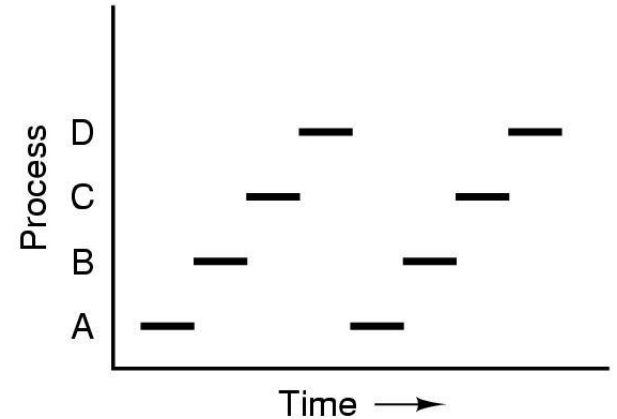


(a)

Four program counters



(b)



(c)

- (a) Multiprogramming of four programs
- (b) Conceptual model of 4 *independent, sequential processes*
  - Sequential process mode: hiding the effects of interrupts, and support blocking system calls
- (c) Only one program active at any instant

# Process Creation

## Principal events that cause process creation

- System initialization; foreground and background
- Execution of a process creation system call
- User request to create a new process; interactive systems
- Initiation of a batch job

# Process Creation

## UNIX example

- ***fork*** system call creates an **exact copy** of calling process
  - Same memory image, environment settings, and opened files
  - After fork, caller is parent, newly-created process is child



# After fork()

The child process returns executing at the exact same point after its parent called fork()

- fork() returns **twice**: the new PID to the parent, and 0 to the child

```
pid= fork();
if (pid == 0) {
    /* I am the child (0: invalid PID)
    */
} else {
    /* I am the parent */
}
```

**Two processes  
execute the code!  
(parent/child share  
same text)**

- All memory contents of parent/child are identical
- Both have the same files open at the same position (point to the same file objects)

# Putting it Together

```
/* now create new process */
pid = fork();
char *const parmList[] = {"/Helloworld", NULL};
if (pid == 0) /* fork() returns 0 to the child process */
{
    sleep(1);
    printf("CHILD: My parent's PID: %d\n", getppid());
    execve("./Helloworld", parmList);
    printf("retval of Helloworld: %d\n", retval);
    exit(retval);
}
else /* fork() returns new pid to the parent process */
{
    printf("PARENT: my child PID: %d\n", pid);
    wait(&status);
    printf("PARENT: Child's exit code is: %d\n", WEXITSTATUS(status));
    exit(0);
}
```

# Process Termination

## Conditions which terminate processes

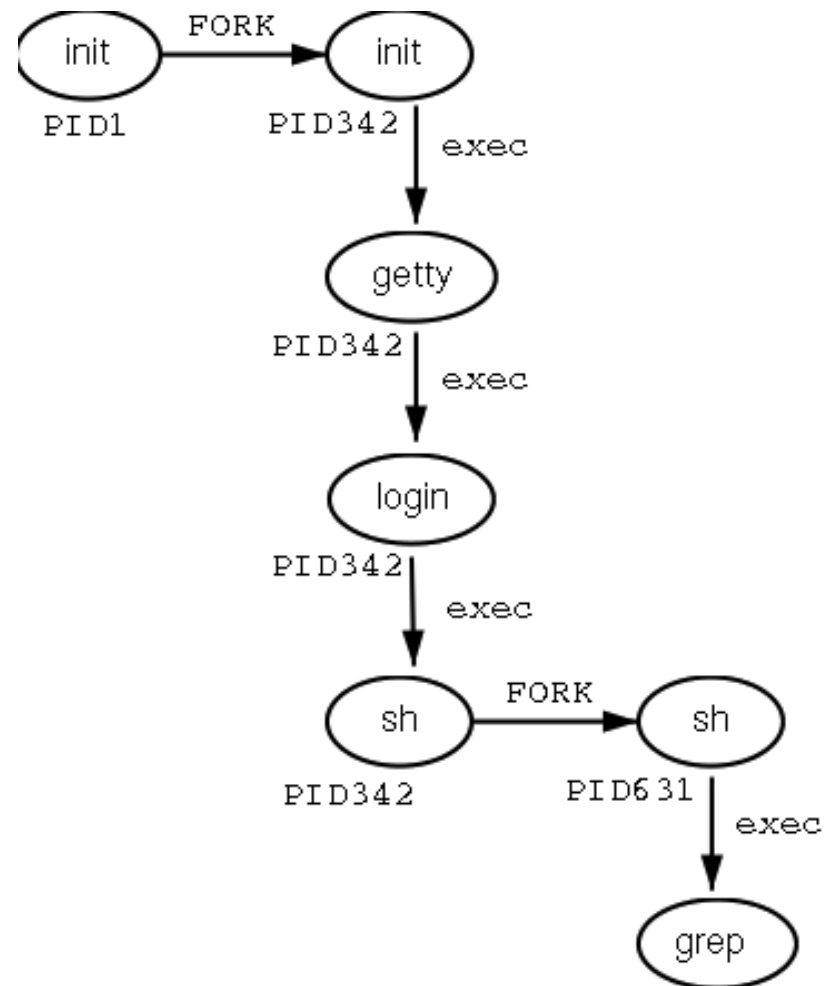
- Normal exit (voluntary)
- Error exit (voluntary)
- Killed by another process (involuntary)

# Process Hierarchies (Trees)

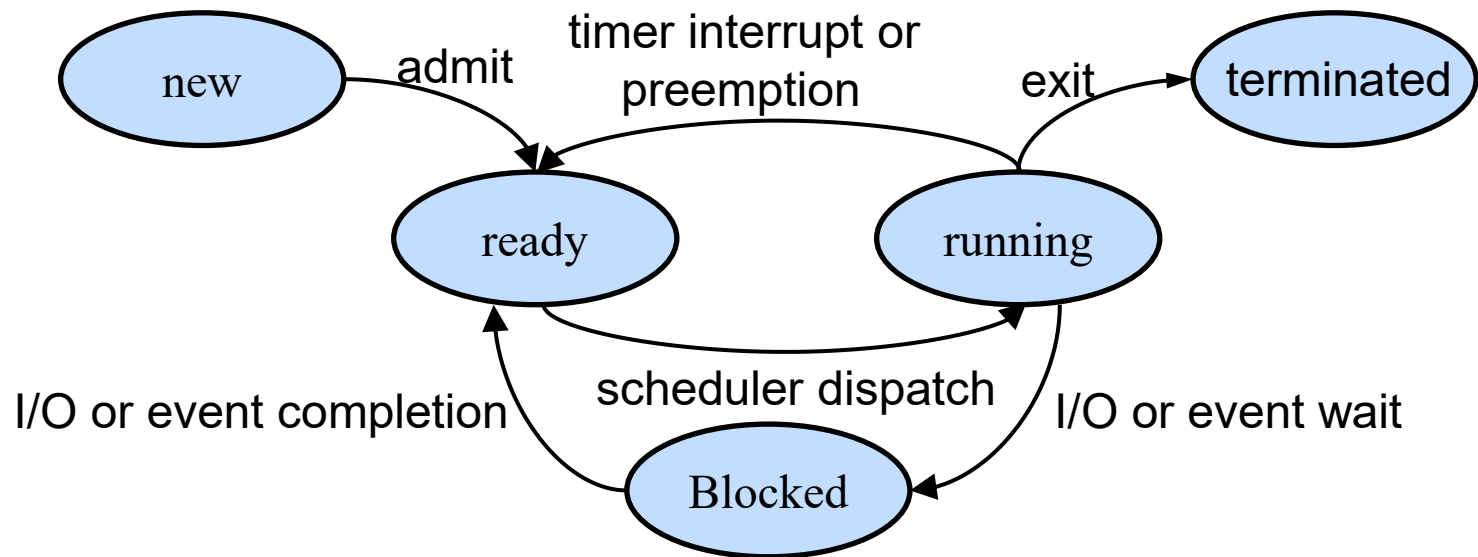
Parent creates a child process,  
child processes can create its  
own process

Forms a hierarchy

- UNIX: a process and all its children and further descendants form a "process group"
- *init*, a special process present in the boot image
- Try: `ps tree -h`



# Process Life Cycle



# Process Life Cycle Stages

- **New** - The process is in the stage of being created.
- **Ready** - The process has all the resources available that it needs to run, but the CPU is not currently working on this process's instructions.
- **Running** - The CPU is working on this process's instructions.
- **Waiting** - The process cannot run because it is waiting for some resource to become available or for some event to occur. **Terminated** - The process has completed.

# Implementation of Processes

## Process table

- One entry per process
- Each entry is called a process control block (PCB)

## Process control block (PCB)

- OS data structure containing data associated with processes

# Process Elements

- Identifier
- State
- Priority
- Program counter
- Memory pointers
- Context data
- I/O status
- Accounting information



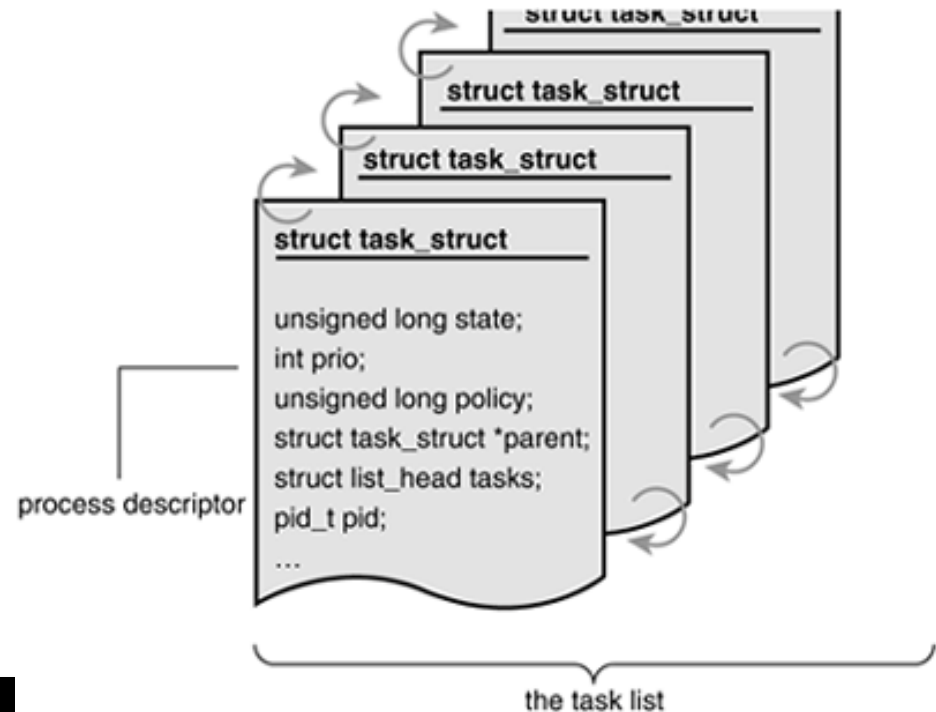
# Process Control Block (PCB)

- Contains the elements of the process
- Created and maintained by OS
- Supports multiple processes

# Linux Processes

Process table: implemented as a linked list/  
hashtable

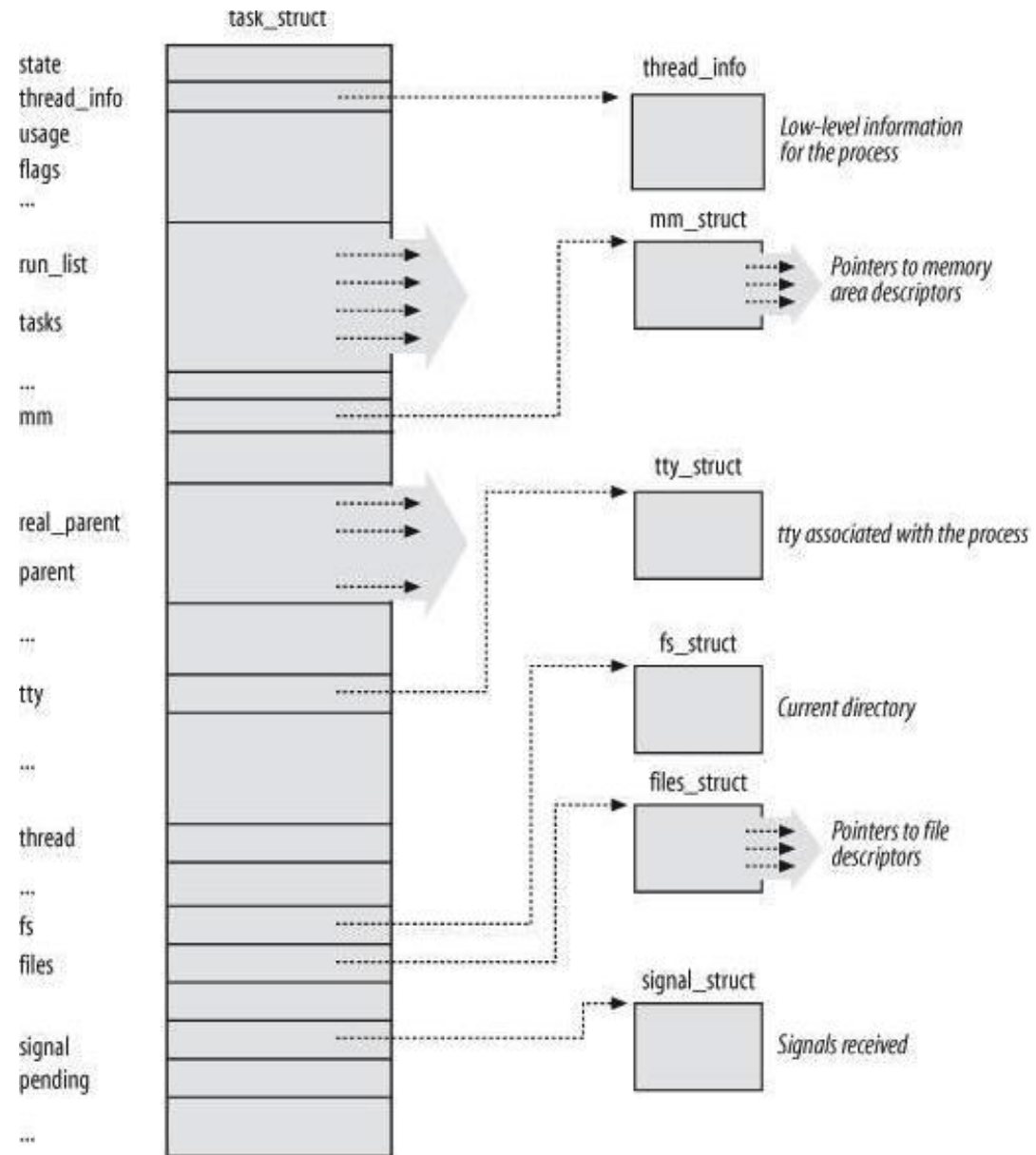
- Each element: a process descriptor of type **task\_struct**
- Dynamically allocated for each process
- Process descriptor
  - Contains all info about a specific process



# Linux Processes

## Process descriptor (PCB)

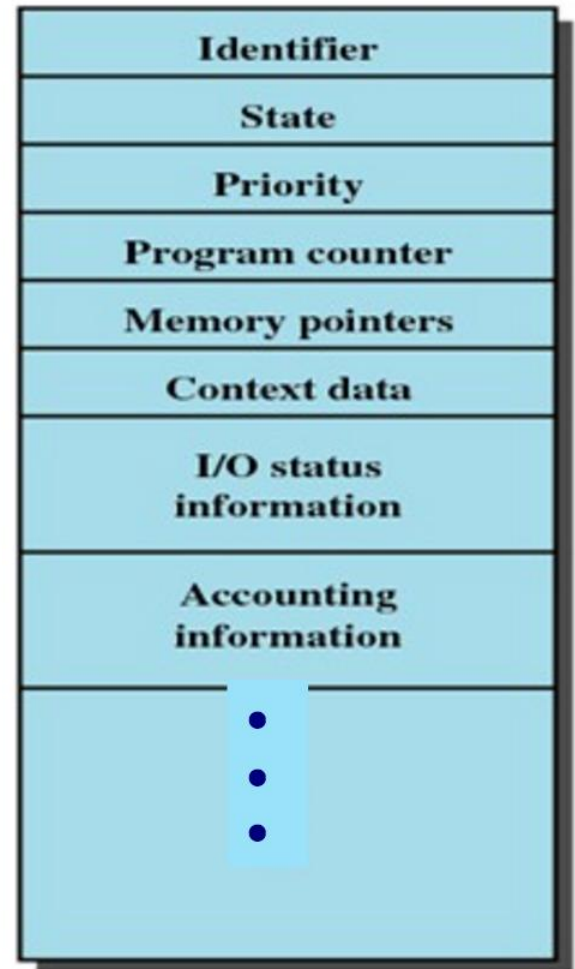
- State
- Identifiers
- Scheduling info
- File system
- Virtual memory
- ...



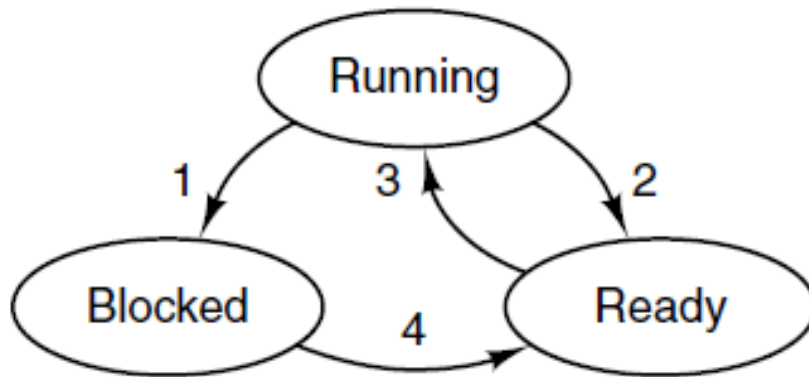
# Linux Process Descriptor

## State

- TASK\_RUNNING
  - Running
- TASK\_INTERRUPTIBLE
  - Blocked
- EXIT\_ZOMBIE
  - Terminated by not deallocated
- EXIT\_DEAD
  - Completely terminated



# Process States



1. Process blocks for input
2. Scheduler picks another process
3. Scheduler picks this process
4. Input becomes available

Figure 2-2. A process can be in running, blocked, or ready state. Transitions between these states are as shown.

Tanenbaum & Bo, Modern Operating Systems: 4th ed., (c) 2013 Prentice-Hall, Inc. All rights reserved.

# Process States

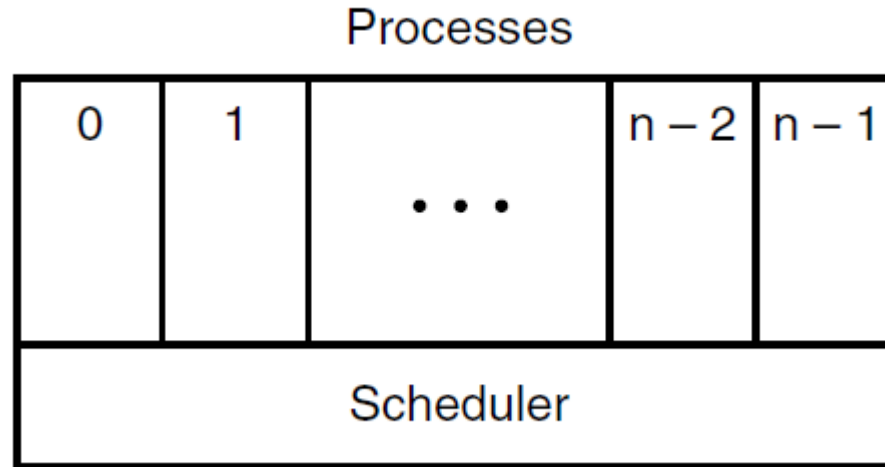


Figure 2-3. The lowest layer of a process-structured operating system handles interrupts and scheduling. Above that layer are sequential processes.

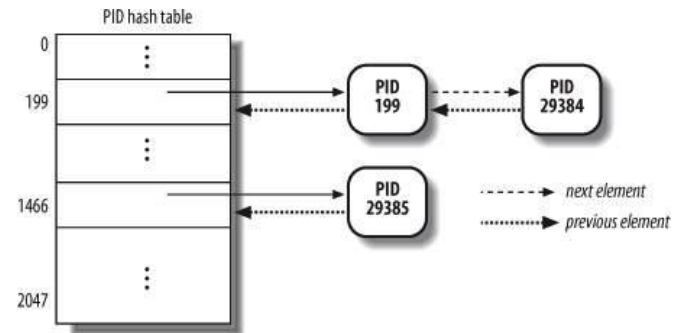
# Linux Process Descriptor (cont')

## Identifiers

- pid: PID of the process/thread
- pgrp: PID of the group leader

How to get the pointer to a specific process?

- The **current** macro
- The **init\_task** macro
- **find\_task\_by\_vpid(pid\_t pid)**



```
pid* pid_struct = find_get_pid(int pid);  
task_struct* task = pid_task(pid_struct, PIDTYPE_PID);
```

# Linux Process Descriptor (cont')

## Files

- fs\_struct
  - file system information: root directory, current directory
- files\_struct
  - Information on opened files



# Summary

What is a process?

- An instantiation of a program

Program life cycle

- Ready, running, blocked, new, terminated

Process implementation

- Process table, PCB

Additional practice

- Download Linux kernel source to your VM, find the following fields in structure `task_struct` (PCB) in `LINUX_SRC_FOLDER/include/linux/sched.h`
  - Program counter (try to google)
  - Stack pointer
  - Process ID
  - Opened file descriptors

# For Next Time

Read: Andrew S. Tanenbaum, Modern Operating Systems, 4thd edition, 2014, Prentice Hall Chapter 2



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