

1. Drew Sadler
2. 262144 page frames
3.  $2.81475^{14}$  bytes or 256 terabytes  
68719476736 Pages
4. Page 3  
With a offset of 12288-10500: 1788 bytes
5. Physical Address: 0x274432
6. 3,932,160 Page faults
7. 491,520 Page faults
8. Having a smaller page size allows the file to load faster and therefore could go through multiple small pages while still waiting on 1 large page size, data is able to be transferred quicker and if faults happen they can be recovered or resent faster
9. Since the cpu is holding onto 1 bytes but memory must address it into separate bytes give space to put the program. We it would probably transfer 2 bytes as it need to send and store the actual information, while also maybe clearing other programs or spreading the bits out into a addressable program in memory