# Linked Lists: Locking, Lock-Free, and Beyond ...



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#### Linked Lists

- We can make effective spin locks
  - Correct
  - Work well under contention
- Are we done with concurrent data structures?



#### Not Over Yet

- Contention
  - Solved by MCS or CLH locks
- Sequential Bottleneck
  - No "bag of tricks"
- Linked Lists
  - Simple data structure
  - Good testbed



#### Set Interface

- Unordered collection of objects
- No duplicates
- Methods
  - Add a new object
  - Remove an object
  - Test if object is present



#### List-Based Sets

```
public interface Set {
  public boolean add(Object x);
  public boolean remove(Object x);
  public boolean contains(Object x);
}
```



```
public class Entry {
  public Object object;
  public int key;
  public Entry next;
}
```



```
public class Entry
public Object object;
public Entry ne
               Object of interest
```



```
public class Entry {
  public Object object;
  public int key;
  public Entry Next;
}
```

Sort by key value (usually hash code)



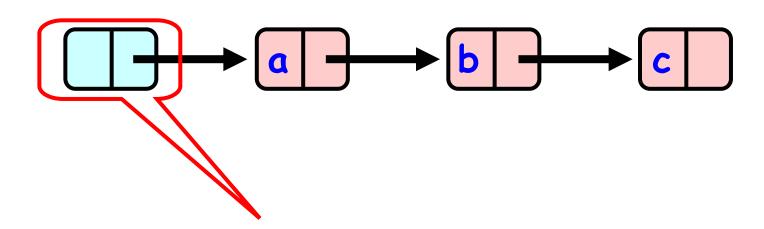
```
public class Entry {
 public Object;
public int key;
     Sorting makes it
 easy to detect absence
```



```
public class Entry {
  public Object object;
  public int key;
  public Entry next;
}
Reference to next entry
```



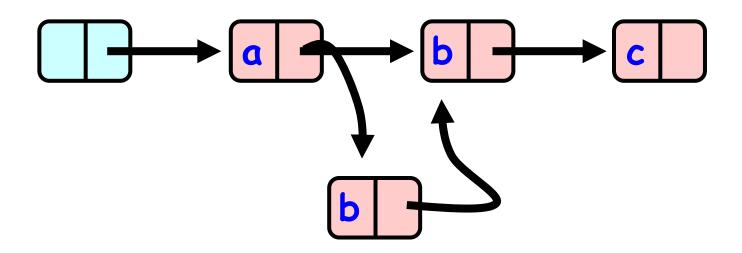
#### List-Based Set



# Sentinel node never deleted (minimum possible key)

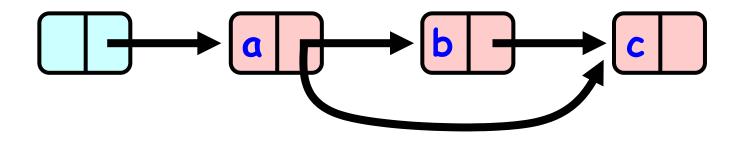


# Adding an Entry





### Removing an Entry





#### Nu, What About Concurrency?

- Our bag of tricks
  - Coarse-grained locks
  - Fine-grained locks
  - Optimistic synchronization
  - Lock-free synchronization



### Coarse-Grained Locking

- · Easy, same as synchronized methods
  - "One lock to rule them all ..."
- · Simple, clearly correct
  - Deserves respect!
- · Works poorly with contention
  - Queue locks help
  - But bottleneck still an issue



### Fine-grained Locking

- · Requires careful thought
  - "Do not meddle in the affairs of wizards, for they are subtle and quick to anger"
- Split object into pieces
  - Each piece has own lock
  - Methods that work on disjoint pieces need not exclude each other



### Optimistic Synchronization

- Requires very careful thought
  - "Do not meddle in the affairs of dragons, for you are crunchy and taste good with ketchup."
- Try it without synchronization
  - If you win, you win
  - If not, try it again with synchronization

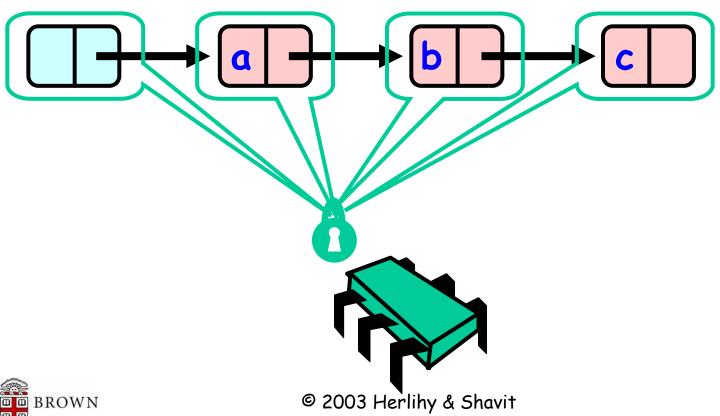


### Lock-Free Synchronization

- · Dump locking altogether ...
  - "You take the red pill and you stay in Wonderland and I show you how deep the rabbit-hole goes"
- · No locks, just native atomic methods
  - Usually compareAndSet

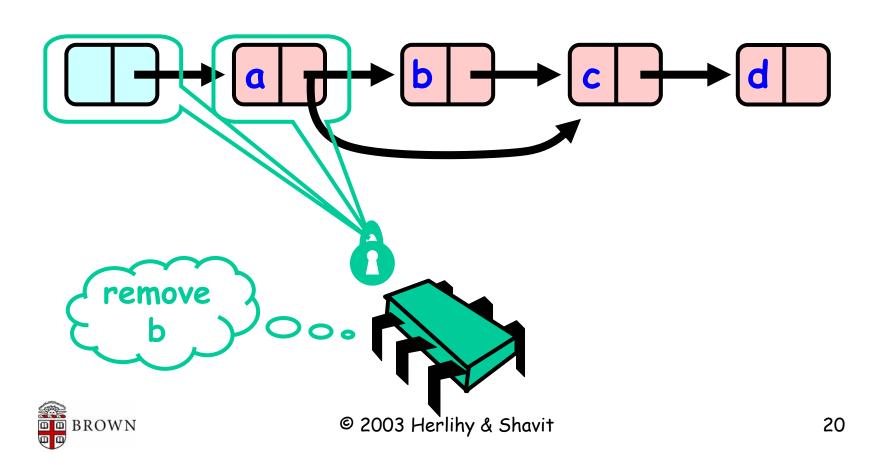


### Hand-over-Hand locking

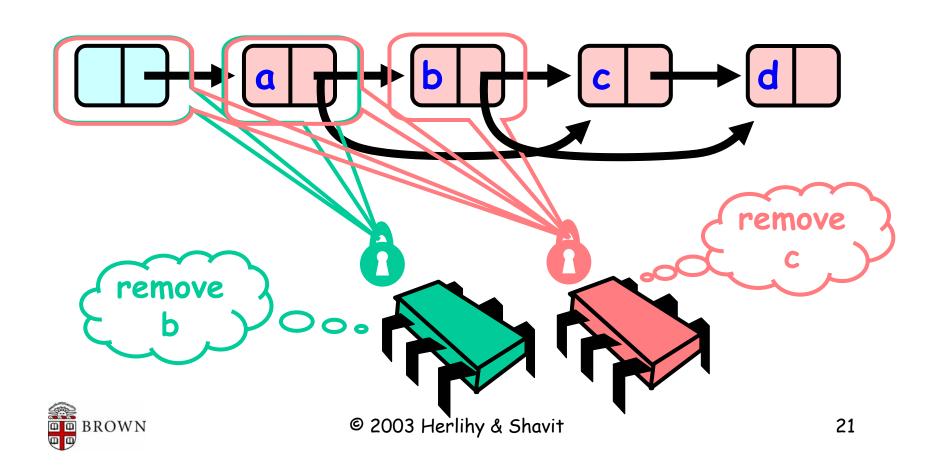


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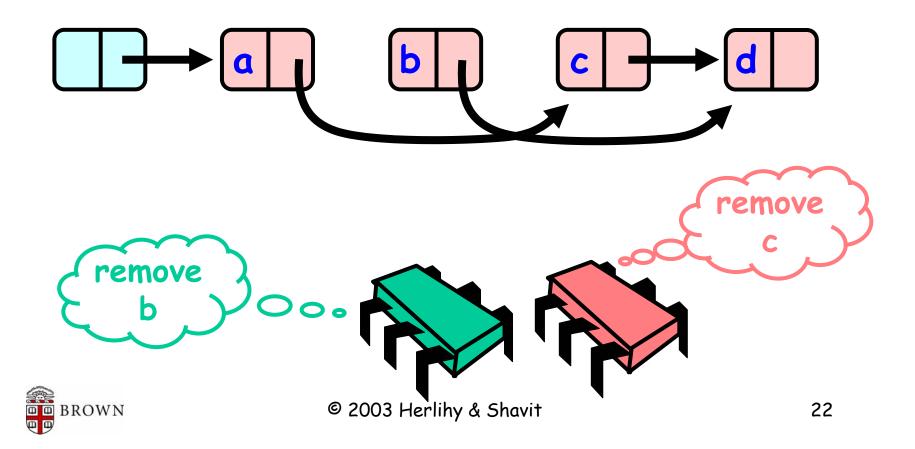
### Removing an Entry



### Removing an Entry



#### Uh-oh



#### Problem

- To delete entry b
  - Swing entry a's next field to c
- · Problem is,
  - Someone could delete c concurrently

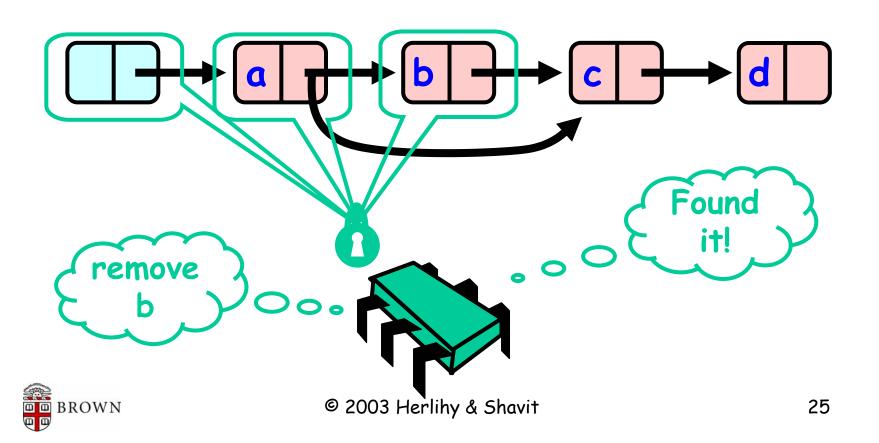


### Insight

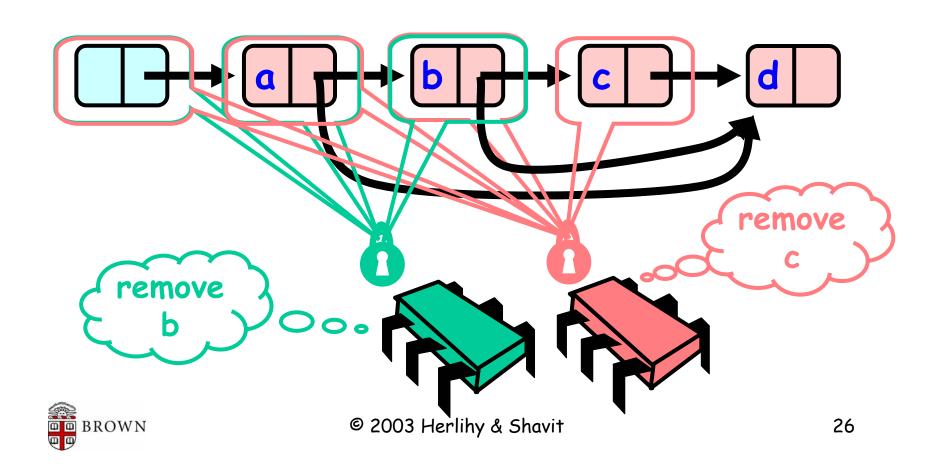
- · If an entry is locked
  - No one can delete entry's successor
- If a thread locks
  - Entry to be deleted
  - And its predecessor
  - Then it works



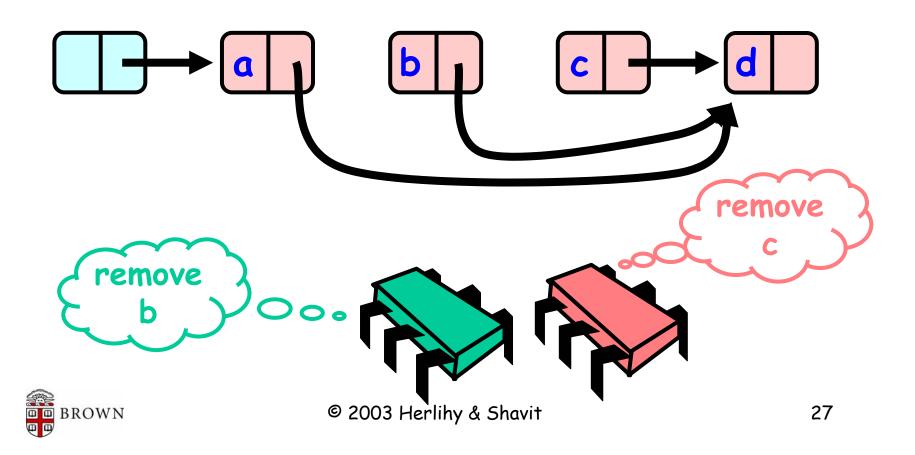
### Hand-Over-Hand Again



### Removing an Entry



### Removing an Entry



```
public boolean remove(Object object) {
  int key = object.hashCode();
  Entry predEntry, currEntry;
  try {
    ...
  } finally {
    currEntry.unlock();
    predEntry.unlock();
  }}
```



```
public boolean remove(Object object) {
int key = object.hashCode();
 Entry prevEntry, currEntry;
 try {
} finally {
  currEntry.unlock();
  prevEntry.unlock();
```

#### Key used to order entry



```
public boolean remove(Object object) {
int key = object.hashCode();
Entry predEntry, currEntry;
} finally {
  currEntry.unlock();
  prevEntry.unlock();
```



Predecessor and current entries



```
public boolean remove(Object object) {
int key = object.hashCode();
 Entry prevEntry, currEntry;
  currentry.unlock()
                       Everything else
  prevEntry.unlock();
```



```
try {
  predEntry = this.head;
  predEntry.lock();
  currEntry = predEntry.next;
  currEntry.lock();
....
} finally { ... }
```



```
lock previous
predEntry = this.head;
predEntry.lock();
 currentry = predentry.next;
 currEntry.lock();
} finally { ... }
```



```
try {
                      Lock current
 prevEntry = this.head:
currEntry = predEntry.next;
currEntry.lock();
} finally { ... }
```



```
try {
 prevEntry = this.head;
 prevEntry.lock(): Traversing list
 currentry = preventry.next;
 currentry.lock();
  finally { ... }
```



```
while (currEntry.key <= key) {</pre>
  if (object == currEntry.object) {
   predEntry.next = currEntry.next;
   return true;
  predEntry.unlock();
  predEntry = currEntry;
  currentry = currentry.nex
  currEntry.lock();
 return false;
```



```
while (currEntry.key <= key) {</pre>
  if (object == currEntry.object) {
   prevEntry.next = currEntry.next;
   return true;
                   Search key range
  prevEntry.unlock();
  prevEntry = currEntry;
  currentry = currentry.nex
  currEntry.lock();
 return false;
```



```
while (currEntry.key <= key) {</pre>
  if (object == currEntry.object) {
   prevEntry.next = currEntry.next;
   return true;
  preventry At start of each loop:
  prevEntry
            currentry and predentry
  currEntry
                             locked
  return false;
              © 2003 Herlihy & Shavit
BROWN
```

```
(currentry key <-
 if (object == currEntry.object) {
  predEntry.next = currEntry.next;
   return true;
 prevEntry.unlock();
 prevEntry = currEntry;
 currentry = currentry.nex
 currentry.lock():
If entry found, remove it
                © 2003 Herlihy & Shavit
BROWN
                                             <del>4</del>0
```

```
(currentry key
 if (object == currEntry.object) {
  predEntry.next = currEntry.next;
  return true;
 prevEntry.unlock();
 prevEntry = currEntry;
 currentry = currentry.nex
 currentry.lock();
If entry found, remove it
               © 2003 Herlihy & Shavit
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```

Unlock predecessor while (currEntry.key <= key) if (object == currEntry.object) { carrEntry.next; prevEntry.next return true; predEntry.unlock(); currentry = currentry.ne currEntry.lock(); return false;



Only one entry locked!

```
while (currentry.key <= key) {
  if (object == currEntry.object) {
   preventry.next = currEntry.next;
   return true;
  prevEntry.unlock();
  currentry = currentry.ne
  currEntry.lock();
 return false;
```



```
demote current
               (allence, object) {
  prevEntry.r
                 = currEntry.next;
  return true
 prevEntry.un ock
predEntry = currEntry;
currEntry.lock();
return false;
```



```
While (CURRENTRY KOY - KOY) 5
   Find and lock new current
   prevEntry next = currEntry.next;
   return true;
  prevEntry.unlock();
  prevEntry = currEntry;
 currEntry = currEntry.next
 currEntry.lock();
 return false;
```



```
Lock invariant restored y) {
 ir (object == currentry.object) {
         ry.next = currEntry.next;
  return
 prevEntry.unlock();
 preventry = currentry;
currEntry = currEntry.next;
currEntry.lock();
return false;
```



```
while (currEntry.key <= key) {
  if (object == currEntry.object) {
   prevEntry.next = currEntry.next;
   return true;
            Otherwise, error 404, dude!
  prevEntry.unlock();
  prevEntry = currEn
  currentry = currentry.next;
  currEntry.]
 return false:
```



#### Why does this work?

- To remove entry e
  - Must lock e
  - Must lock e's predecessor
- · Therefore, if you lock an entry
  - It can't be removed
  - And neither can its successor



#### Lock and Load

- To move to successor entry for e
  - Lock e

Until next entry

- Lock e.next

idenified and locked

- Unlock e
- While traversing

Don't release e

- e cannot be removed
- e.next cannot be removed



## Adding Entries

- To add entry e
  - Must lock predecessor
  - Must lock successor
- Neither can be deleted
  - (Is successor lock actually required?)



#### Drawbacks

- Better than coarse-grained lock
  - Threads can traverse in parallel
- Still not ideal
  - Long chain of acquire/release
  - Inefficient

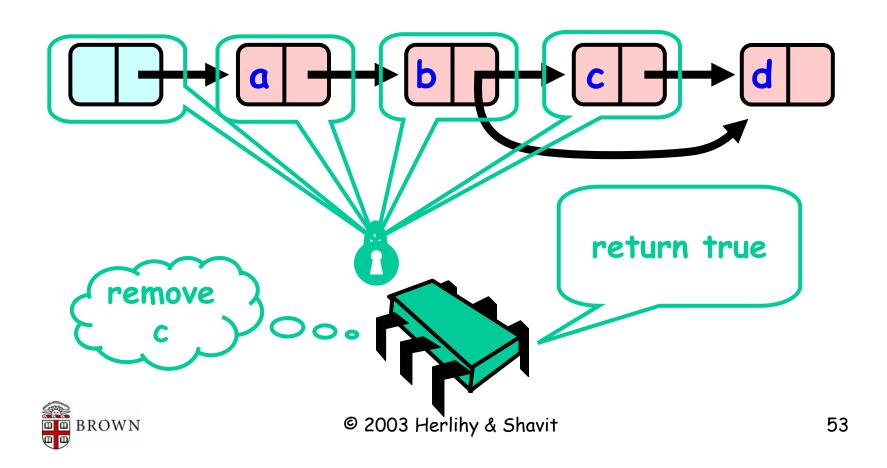


#### Optimistic Synchronization

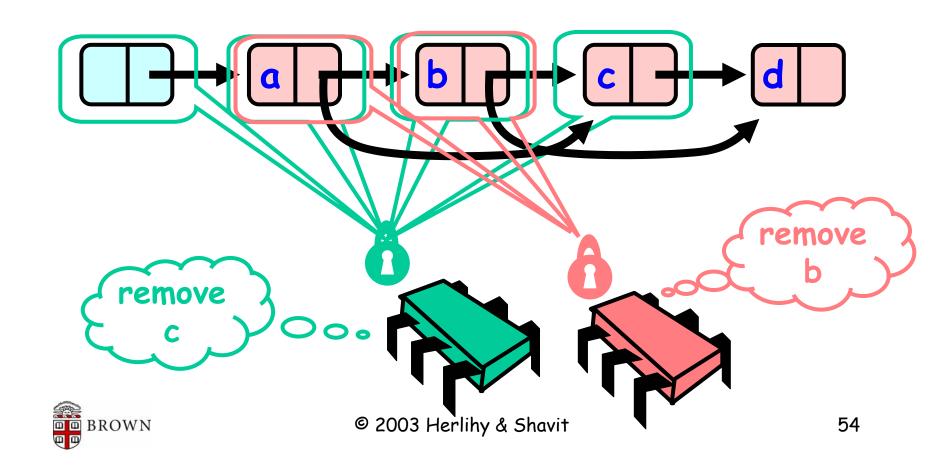
- Find entries without locking
- Lock entries
- Check that everything is OK



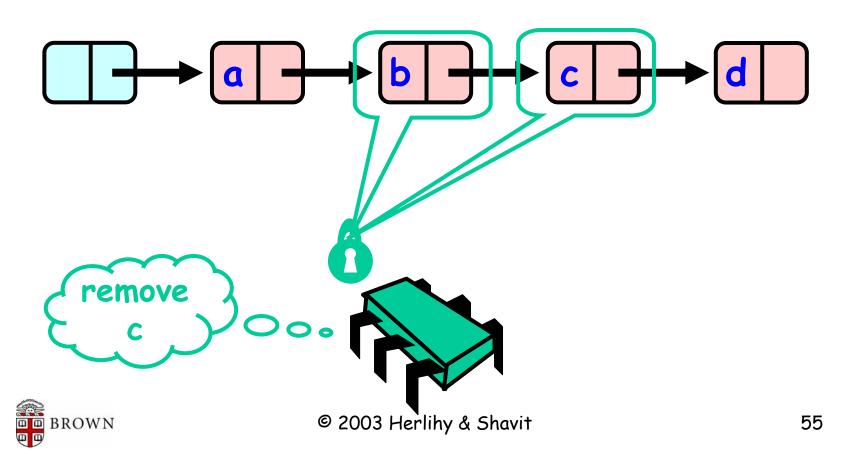
## Removing an Entry



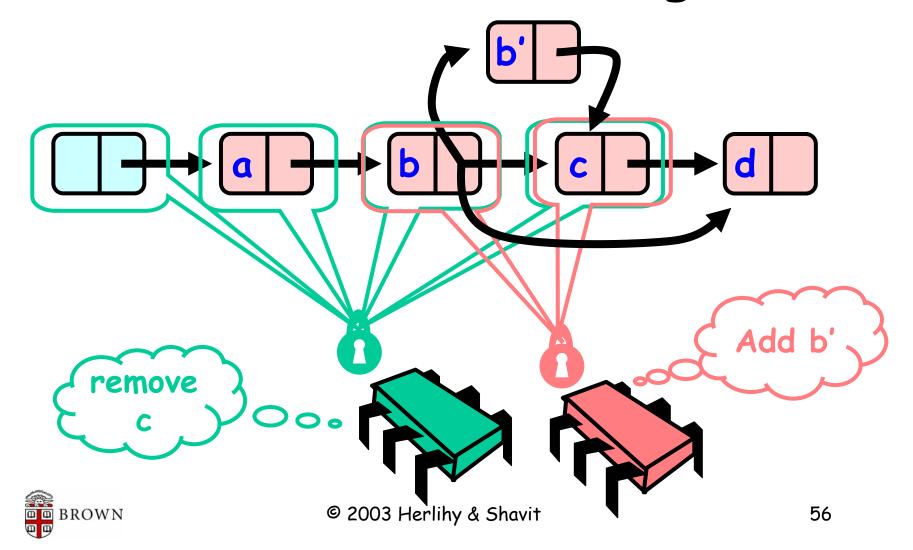
## What Can Go Wrong?



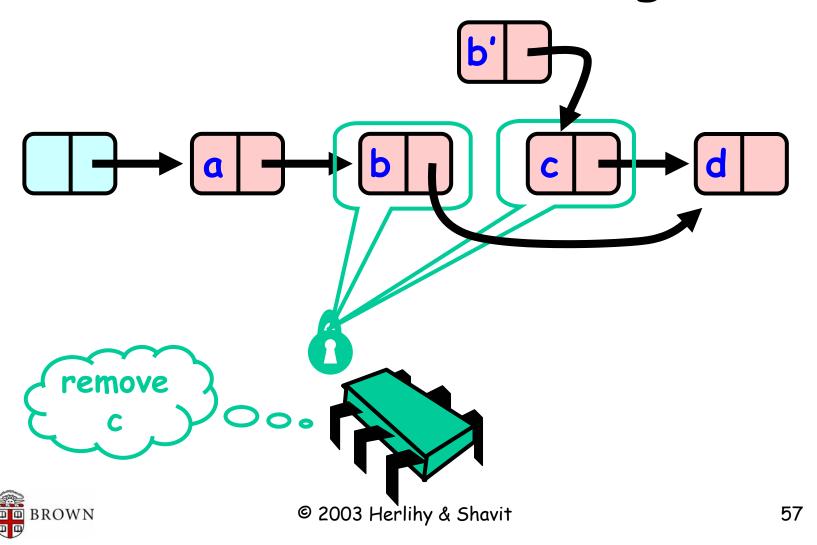
## Check that Entry is Still Accessible



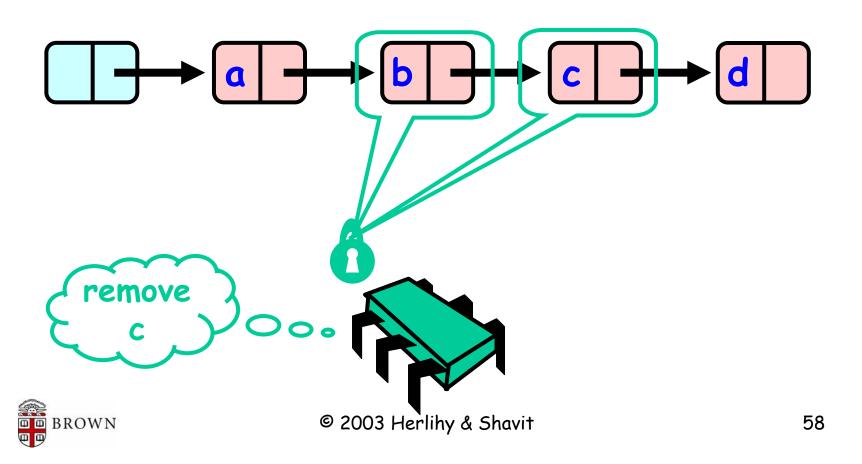
## What Can Go Wrong?



#### What Can Go Wrong?



# Check that Entries Still Adjacent

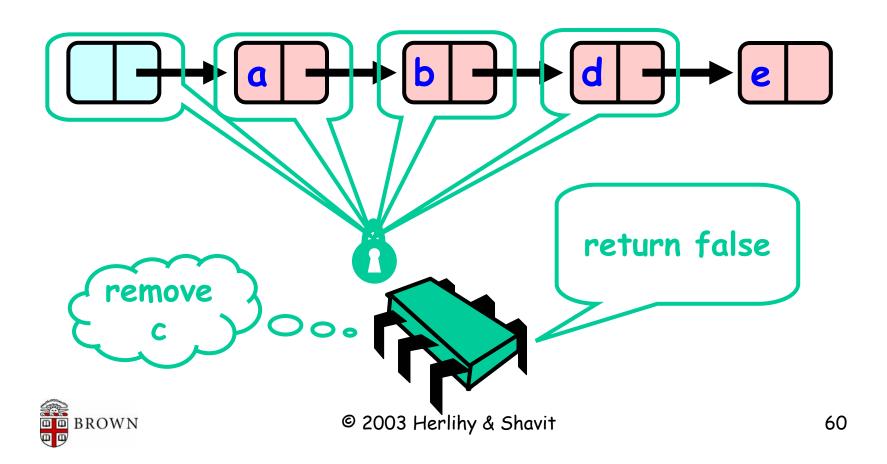


#### Correctness

- · If
  - Entry b and Entry c both locked
  - Entry b still accessible
  - Entry c still successor to b
- Then
  - Neither will be deleted
  - OK to delete and return true



## Removing an Absent Entry



#### Correctness

- · If
  - Entry b and Entry d both locked
  - Entry b still accessible
  - Entry d still successor to b
- Then
  - Neither will be deleted
  - No thread can add c after b
  - OK to return false



```
private boolean
validate(Entry predEntry,
          Entry currEntry) {
Entry entry = head;
while (entry.key <= predEntry.key) {</pre>
  if (entry == predEntry)
   return predEntry.next == currEntry;
  entry = entry.next;
 return false;
```



```
private boolean
validate Entry predEntry,
         Entry currEntry {
Entry entry = head;
while (entry key <= predEntr
 if (entry == predEntry)
   return/redEntry.next ==
         entry.next;
  Predecessor &
 current entries
```



```
private boolean
validate(Entry predEntry,
          Entry currEntry) {
Entry entry = head;
while (entry key <= predEntry key) {
 if (entry == predIntry)
   return predEntry next == currEntry;
  entry = entry.next;
                           Start at the
 return false;
                             beginning
```



```
private boolean
validate(Entry predEntry,
          Entry currentry) {
 Entry entry = head;
while (entry.key <= predEntry.key) {</pre>
 if (entry == predEntry)
   return predEntry.next == turrEntry;
  entry = entry.next;
                   Search range of keys
 return false;
```



```
private boolean
validate(Entry predEntry,
          Entry currentry) {
Entry entry = head;
while (entry.key <= predEntry.key) {</pre>
if (entry == predEntry)
   return predEntry.next == currEntry;
   entry = entry.next
 return false;
                   Predecessor reachable
```



```
private boolean
validate(Entry predEntry,
          Entry currentry) {
Entry entry = head;
while (entry.key <= predEntry.key) {</pre>
 if (entry == predEntry)
  return predEntry.next == currEntry;
  entry = entry.next;
 return false;
                 Is current entry next?
```



```
private boolean
                   Otherwise move on
validate(Entry predEntry,
          Entry currentry
Entry entry = head;
while (entry.key <= predEntr
 if (entry == predEntry)
   return predEntry hext ==
 entry = entry.next;
 return false;
```



```
private boolean Predecessor not reachable
 validate(Entry predEntry
          Entry currentry
 Entry entry = head;
 while (entry.key / predEntr
  if (entry == predEntry)
   return predentry.next ==
 entry = entry hext;
return false;
```



```
public boolean remove(Object object) {
 int key = object.hashCode();
 retry: while (true) {
 Entry predEntry = this.head;
 Entry currEntry = predEntry.next;
while (currEntry.key <= key) {</pre>
  if (object == currEntry.object)
   break:
  predEntry = currEntry;
  currentry = currentry.next;
```



```
public boolean remove(Object object) {
int key = object.hashCode();
retry: while (true) {
Entry prevEntry = this.head;
Entry currentry = prevEntry
while (currEntry.key <= k
 if (object == currEntry
  break;
 prevEntry = currEntry;
 currentry = currentry.next;
                      Search key
```



```
public boolean remove(Object object) {
int key = object.hashCode();
retry: while (true) {
 Entry prevEntry = this.head;
 Entry currentry = prevEntry.
while (currEntry.key <= key)
 if (object == currEntry.obj
   break;
  prevEntry = currEntry;
  currentry = currentry.next;
       Retry on synchronization conflict
```



```
public boolean remove(Object object) {
int key = object.hashCode();
 retry: while (true)
 Entry predEntry = this.head;
 Entry currEntry = predEntry.rext;
while (currentry key <= key)
 if (object == currEntry.obj
   break:
  prevEntry = currEntry;
  CURRENTRY - CUR Entry novt
 Examine predecessor and current entries
```



```
public boolean remove(Object object) {
int key = object.hashCode();
 retry: while (true) {
Entry prevEntry = this.head;
    rv currentry = preventry.next;
while (currEntry.key <= key)
 if (object == currEntry.obj
  break;
  prevEntry = currEntry;
                   try.next;
   Search by key
```



```
public boolean remove(Object object) {
int key = object.hashCode();
 retry: while (true) {
 Entry prevEntry = this.head;
 Entry currentry = prevEntry.next;
while (currEntry.key <= key) {
 if (object == currEntry.object)
   break:
  preventry currentry;
  currentrv = currentrv.ne
 Stop if we find object
```



```
retry: while (true) {
Entry prevEntry = this.head;
Entry currentry = preventry.
while (currentry.key <= key)
 if (object == currEntry.ob)
  break;
 predEntry = currEntry;
 currentry = currentry.next;
```



## On Exit from Loop

- If object is present
  - currEntry holds object
  - predEntry just before currEntry
- · If object is absent
  - currEntry has first higher key
  - predEntry just before currEntry
- Assuming no synchronization problems



```
try {
  predEntry.lock(); currEntry.lock();
  if (validate(predEntry, currEntry) {
   if (currEntry.object == object) {
    predEntry.next = currEntry.next;
    return true;
   } else {
    return false;
   }}} finally {
     predEntry.unlock();
     currEntry.unlock();
   }}}
```



```
dEntry.lock(); currEntry.lock();
if (validate(predEntry, currEntry) {
if (currentry.object == object) {
  predEntry.next = currEntry.next;
  return true;
} else {
  return false:
                       Always unlock
finally {
   predEntry.unlock();
  currEntry.unlock();
```



```
predEntry.lock(); currEntry.lock();
if (validate(predEntry,currEntr
 if (currEntry.object == object) {
  predEntry.next = currEntry.next;
  return true;
 } else {
  return false;
                    Lock both entries
 }}} finally {
   predEntry.unlock();
   currEntry.unlock();
 }}}
```



```
predEntry.lock(); currEntry.lock();
if (validate(predEntry,currEntry) {
 if (currentry.object == object)
  predEntry.next = currEntry.next;
  return true;
 } else { Check for synchronization
  return false; conflicts
 }}} finally {
   predEntry.unlock();
   currEntry.unlock();
 }}}
```



```
try {
  predEntry.lock(); currEntry.lock();
  if (validate(predEntry.currEntry)
  if (currEntry.object == object) {
    predEntry.next = currEntry.next;
    return true:
   } else {
    return false;
                          Object found,
   }}} finally {
                          remove entry
     predEntry.unlock();
     currEntry.unlock();
   }}}
```



```
try {
  predEntry.lock(); currEntry.lock();
  if (validate(predEntry, currEntry) {
   if (currEntry.object == object) {
    predEntry.next = currEntry.next;
    return true;
                      Object not found
   } else {
   return false;
   ·}} finally {
     predEntry.unlock();
     currEntry.unlock();
   }}}
```



### So Far, So Good

- Much less lock acquisition/release
  - Performance
  - Concurrency
- · Problems
  - Need to traverse list twice
  - contains() method acquires locks
    - Most common method call



#### Marked Lists

- Remove Method
  - Scans list (as before)
  - Locks predecessor & current (as before)
  - Marks current entry as removed (new!)
  - Redirects predecessor's next (as before)

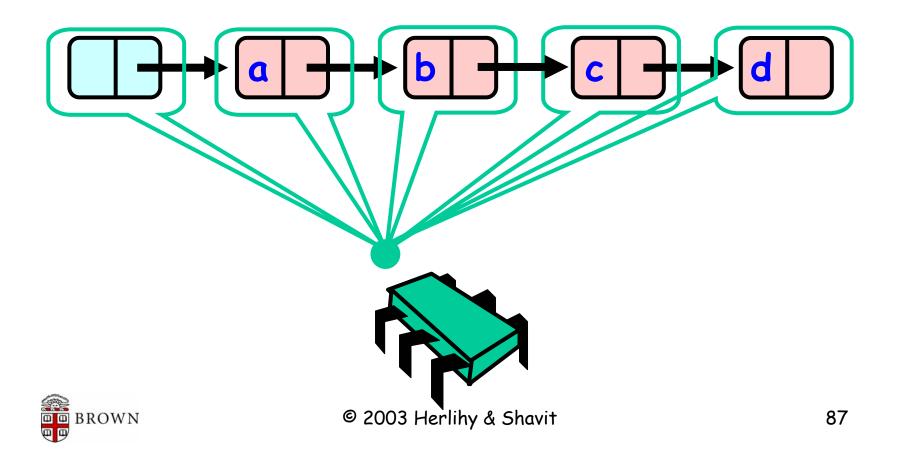


#### Marked Lists

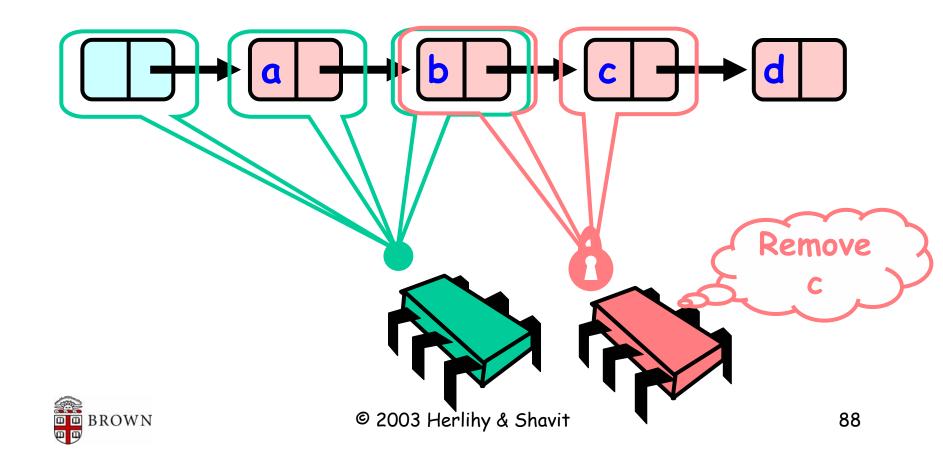
- All Methods
  - Scan list
  - Do not scan past marked entry
  - Instead, start over from list head



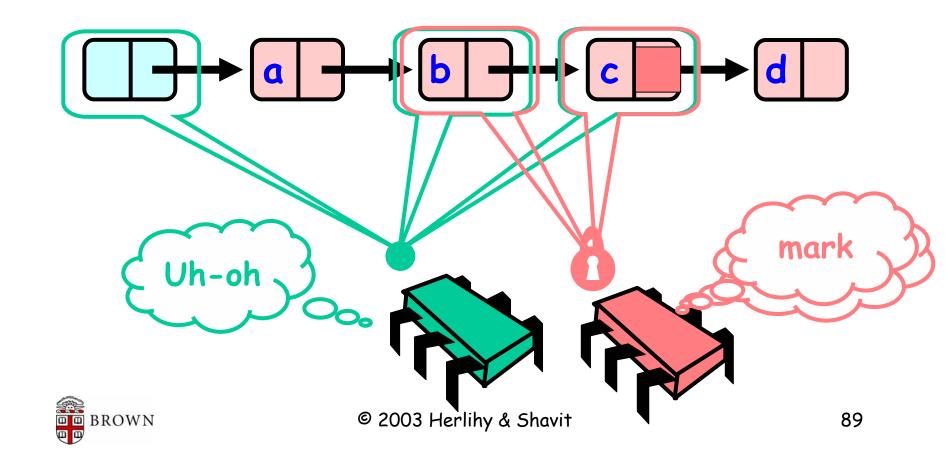
#### Business as Usual



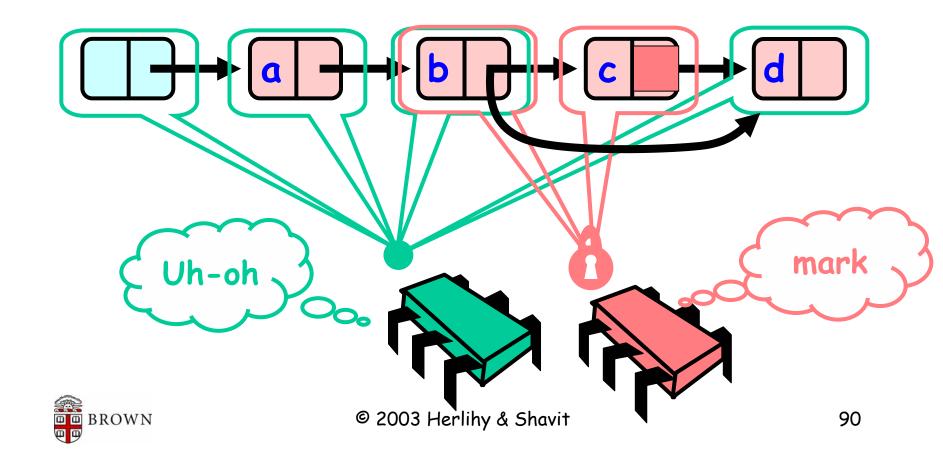
### Interference



### Interference



### Interference



# Marking a Node

- AtomicMarkableReference class
  - Java.util.concurrent.atomic package



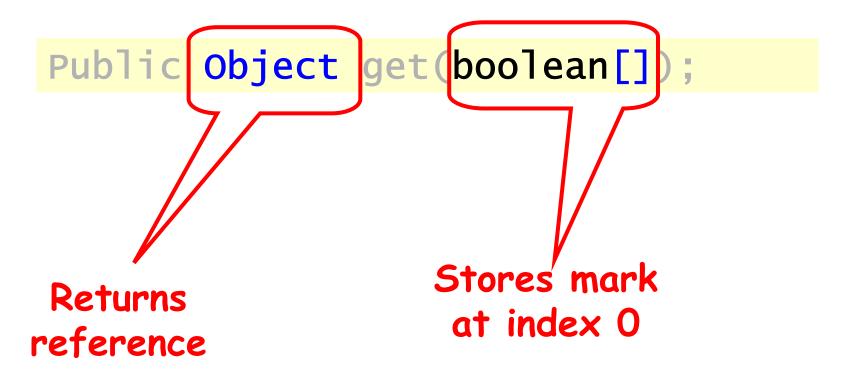


### Extracting Information

Public Object get(boolean[]);

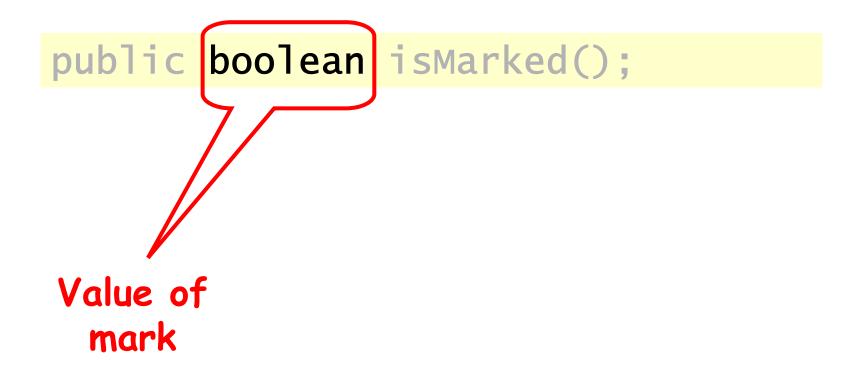


## Extracting Information





## Extracting Information





```
Public boolean compareAndSet(
   Object expectedRef,
   Object updateRef,
   boolean expectedMark,
   boolean updateMark);
```



If this is the current reference ...

```
Public boolean compareAndSet(
   Object expectedRef,
   Object updateRef,
   boolean expectedMark,
   boolean updateMark);
```

And this is the current mark ...



```
...then change to this
                   new reference ...
Public boolean compareAndSet(
  Object expectedRef,
 Object updateRef,
   oolean expectedMark,
 boolean updateMark);
                        and this new
                           mark
```



```
public boolean attemptMark(
   Object expectedRef,
   boolean updateMark);
```



```
public boolean attemptMark(
  Object expectedRef,
  bodleam updateMark);
If this is the current
    reference ...
```



```
public boolean attemptMark(
  Object expectedRef,
 boolean updateMark);
.. then change to
 this new mark.
```



#### Aside

- We will see that it is often useful to tag pointers with
  - Boolean values
  - Integer values
- Sometimes to mark, and sometimes to ensure pointers are unique



```
private boolean
 validate(Entry predEntry,
          Entry currentry) {
 return
  (!predEntry.next.isMarked()) &&
  (!currEntry.next.isMarked()) &&
  (predEntry.next.getReference()
   == currEntry);
```



```
private boolean Predecessor not removed
 validate(Entry predEntry
          Entry currentry
 return
 (!predEntry.next.isMarked()) &&
  !currEntry.next.isMarked(
  (predEntry.next.getReference()
    == currEntry);
```



```
private boolean Current not removed
 validate(Entry predEntry
          Entry currenty
 return
  (!predEntry.next.isMarked())
 (!currEntry.next.isMarked()) &&
    redEntry.next.getReference
    == currEntry);
```



```
private boolean Next field unchanged
 validate(Entry predEntry,
          Entry currentry
 return
  (!predEntry.next.isMarked
   !currEntry.next.isMarked(
 (predEntry.next.getReference()
    == currEntry);
```



```
public boolean remove(Object object) {
while (currEntry.key <= key) {</pre>
     Entry nextEntry =
       (Entry)currEntry.next.get(mark);
     if (mark[0])
       continue retry;
  if (object == currEntry.object)
   break:
  predEntry = currEntry;
  currEntry = currEntry.next;
```



```
public boolean remove(Object object) {
    Entry nextEntry =
       (Entry)currEntry.next.get(mark);
       (mark) U
      continue retry;
 if (object == currEntry object)
             Atomically read reference
  break;
 predEntry = C
                        & mark
  currentry = currentry.next;
```



```
ect object) {
 Panic if entry removed
while (currEntry
                 key <= key) {
    Entry pextEntry =
       Entry)currentry.next.get
    if (mark[0])
      continue retry;
 if (object == currEntry.object)
  break;
 predEntry = currEntry;
 currentry = currentry.next;
```



### Evaluation

#### · Good:

- Contains method doesn't need to lock
- Uncontended calls don't re-traverse

#### Bad

- Contended calls do re-traverse
- Traffic jam if one thread delays



# Traffic Jam

- Any concurrent data structure based on mutual exclusion has a weakness
- If one thread
  - Enters critical section
  - And "eats the big muffin" (stops running)
    - · Cache miss, page fault, descheduled ...
    - Software error, ...
  - Everyone else using that lock is stuck!



# Lock-Free Data Structures

- · No matter what ...
- 3
- Some thread will complete method call
- Even if others halt at malicious times
- Implies that
  - You can't use locks (why?)
  - Um, that's why they call it lock-free



### Lock-Free + Wait-Free

- · Wait-free synchronization
  - Every method call eventually finishes
  - What everyone really wants
- · Lock-free synchronization
  - Some method call eventually finishes
  - What we are usually willing to pay for
    - Starvation rare in practice ...

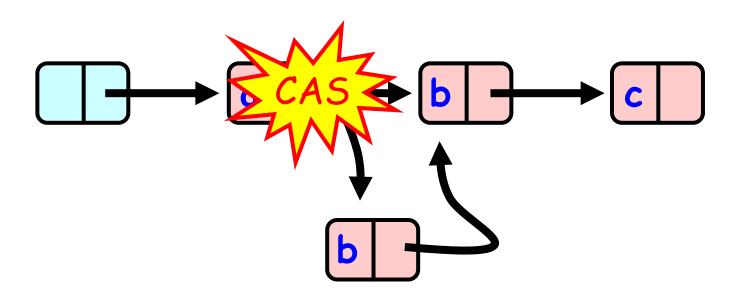


### Lock-Free Lists

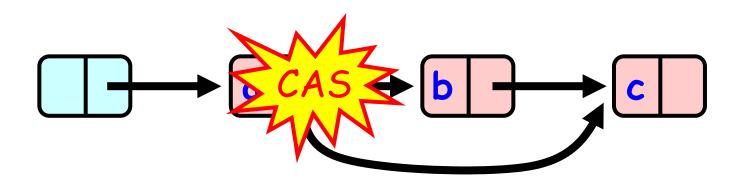
- Next logical step
- Eliminate locking entirely
- Use only compareAndSet()
- Invented by Maged Michael, 2003



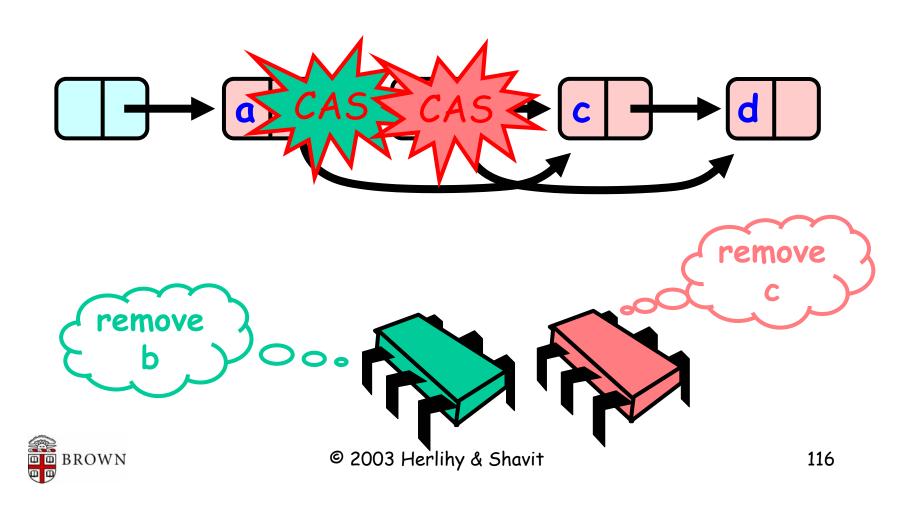
# Adding an Entry



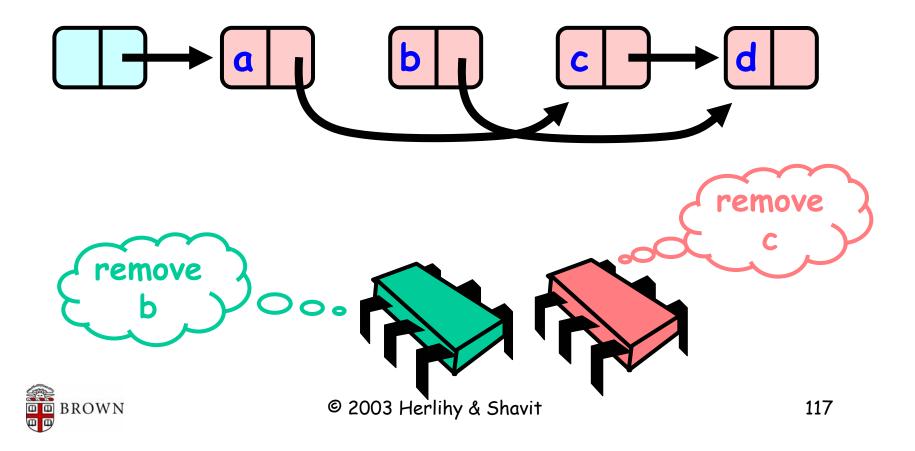








# Look Familiar?



#### Problem

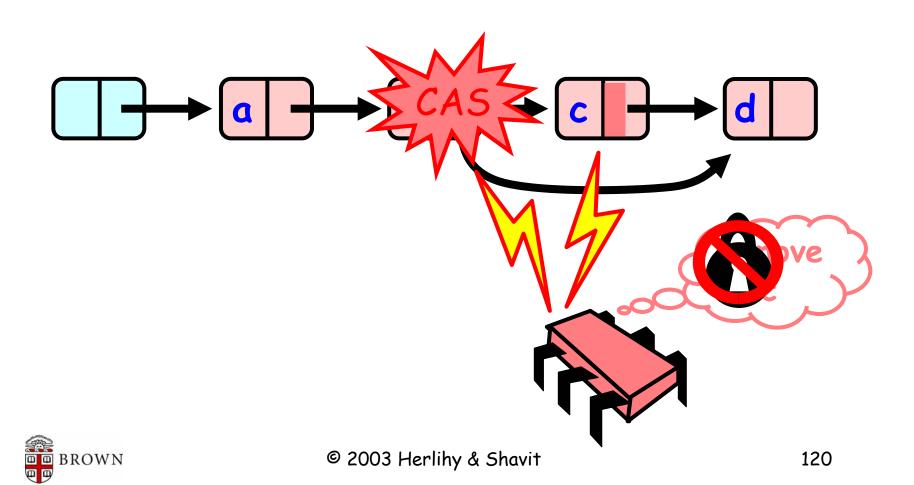
- · Method updates entry's next field
- · After entry has been removed

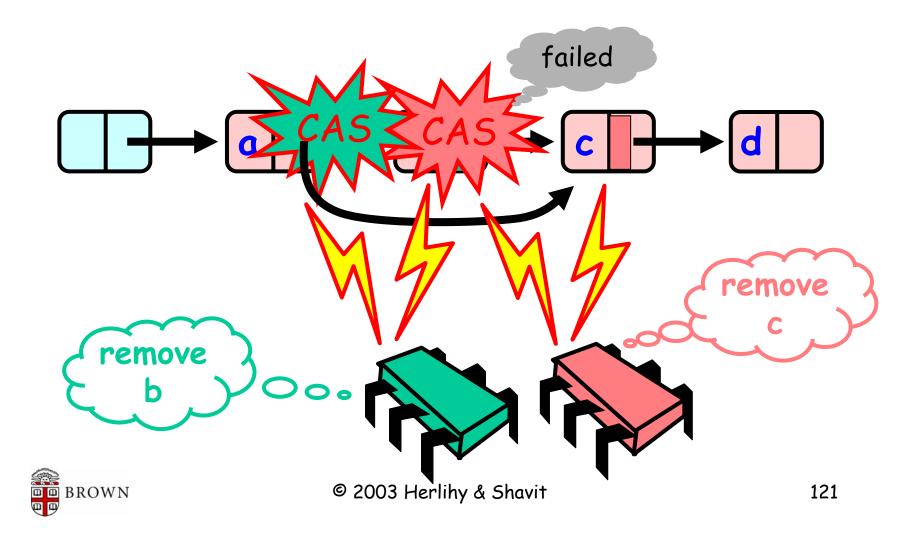


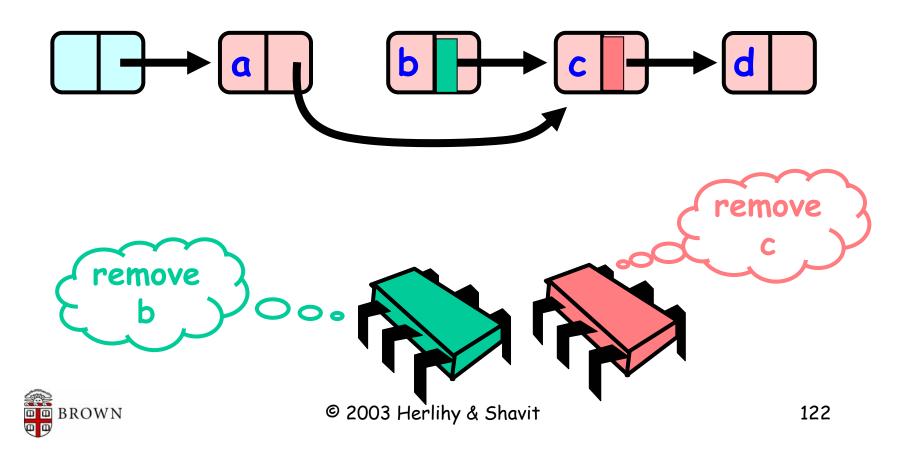
### Solution

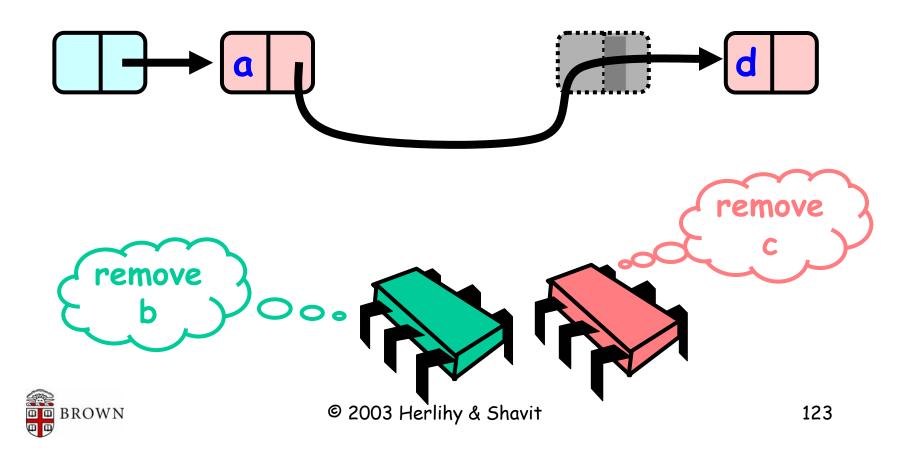
- Use AtomicMarkableReference
- Remove in two steps
  - Set mark bit in next field
  - Redirect predecessor's pointer
- · CAS
  - Fails if mark bit set (entry removed)









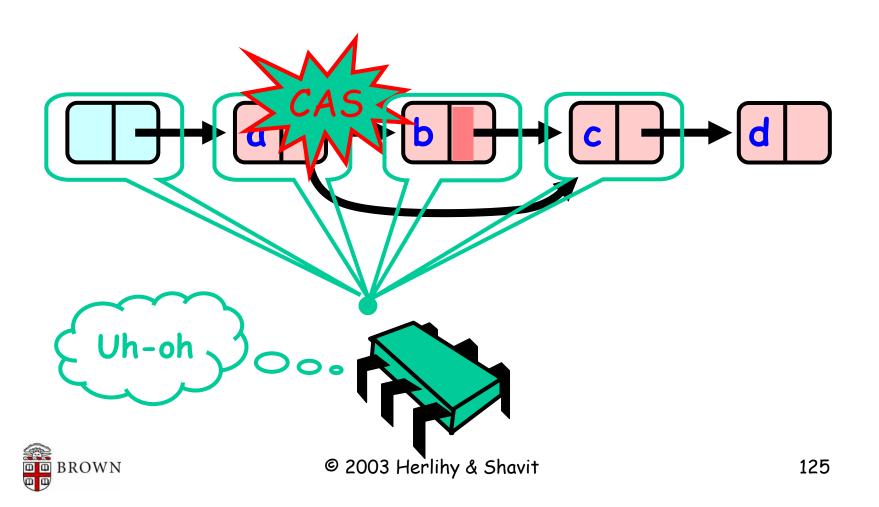


# Traversing the List

- Q: what do you do when you find a "logically" deleted entry in your path?
- A: finish the job.
  - CAS the predecessor's next field
  - Proceed (repeat as needed)



# Lock-Free Traversal

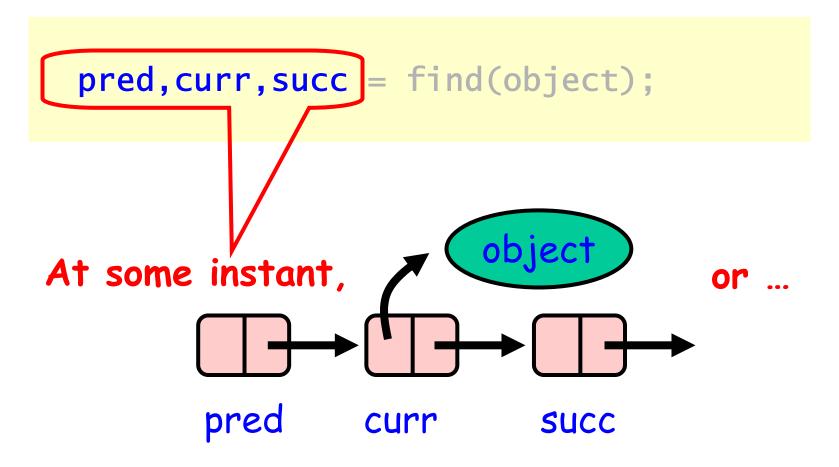


# The Find Method

```
pred,curr,next = find(object);
```

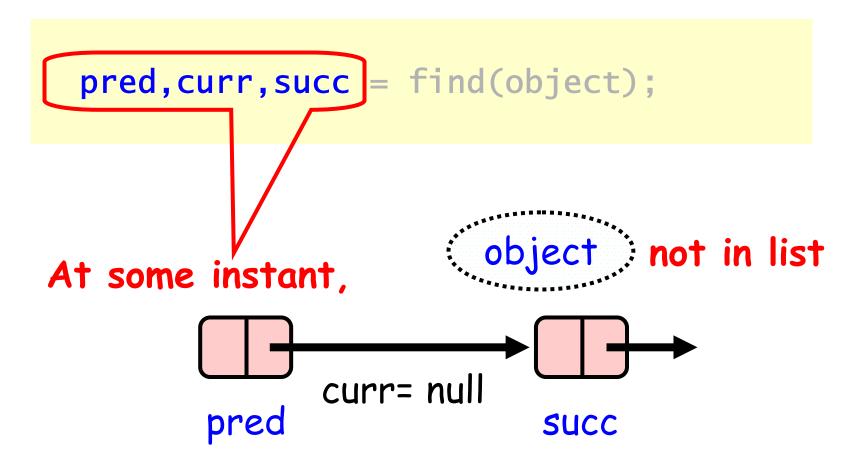


### The Find Method





### The Find Method





```
public boolean remove(Object object) {
while (true) {
  pred,curr,succ = find(object);
  if (curr == null)
   return false;
  if (!curr.next.attemptMark(succ,
                              true))
    continue;
  pred.next.compareAndSet(curr, succ,
                           false, false);
  return true;
  }}
```



```
public boolean remove(Object object) {
while (true) {
 pred,curr,succ = find(object);
 if (curr == null)
   return false
 if (!curr.next.
                 attemptMark(succ,
                              true))
    continue;
  pred.next.compareAndSet(curr, succ,
                           false, false);
  return true;
                         Keep trying
 }}
```



```
public boolean remove(Object object) {
while (true) {
 pred,curr,succ = find(object);
  1t (curr == n
   return false;
  if (!curr.next.attemptMark(succ,
                              true))
    continue;
  pred.next.compareAndSet(curr, succ,
                             lse, false);
  return true;
                 Find alleged neighbors
  }}
```



```
public boolean remove(Object object) {
while (true) {
  pred,curr,succ = find(object);
 if (curr == null)
   return false;
 if (!curr.next.attemptMark(succ,
                              true))
    continue;
  pred.next.compareAndSet(curr, succ,
                           false, false);
  return true;
                      She's not there ...
  }}
```



```
Try to mark entry as deleted bject) {
 pred, curr, succ = find(object);
 if (curr \neq nu(1))
  return false;
 if (!curr.next.attemptMark(succ,)
                              true)
 pred.next.compareAndSet(curr, succ,
                           false, false);
 return true;
 }}
```



```
ect) {
If it doesn't work, just retry
pred, curr, succ = find(object);
if (curr == null)
 return false
if (!curr.next.attemptMark(succ,
                             true))
  continue;
pred.next.compareAn
                          false, false);
return true;
}}
```



```
If it works, our job is (essentially) done
 pred,curr,succ = find(object);
 if (curr == null)
   return false
 if (!curr.next.attemptMark(succ,
                              true))
    continue;
 pred.next.co
  return true;
```



```
Try to advance reference
(if we don't succeed, someone else did).
 if (curr == n
  return false;
 if (!curr next.attem
        Inue;
 pred.next.compareAndSet(curr, succ,
                           false, false);
 return true;
```



```
public boolean add(Object object) {
while (true) {
  pred,curr,succ= find(object);
  if (curr != null)
   return false;
  Entry entry = new Entry(object);
  entry.next = new AMR(succ,false);
  if (pred.next.CAS(succ, entry,
                    false, false))
        return true;
 }}
```



```
public boolean add(Object object) {
while (true) {
  pred,curr,succ= find(object);
 if (curr != null)
   return false:
  Entry entry = new Entry(object);
  entry.next = new AMR(succ, false);
 if (pred.next CAS(succ, entry,
          Object already there.
 }}
```



```
public boolean add(Obje
while (true) {
  pred,curr,succ= find(
  if (curr != null)
   return false:
 Entry entry = new Entry(object);
 entry.next = new AMR(succ,false)
       red.next.CAS(succ.entry,
                     alse, false))
           create new entry
```



```
public boolean add(Object object) {
                   find(object);
                      Install new entry
  Entry entry = new Entry(object):
  entry.next = new AMR(succ, false);
 if (pred.next.CAS(succ, entry,
                     false, false))
        return true;
}}
```



#### Contains

```
public boolean contains(Object obj){
  while (true) {
    prev,curr,succ = find(object);
    return (curr != null);
  }
}
```



#### Contains

```
public boolean contains(Object obj){
while (true) {
  prev, curr, succ = find(object);
  return (curr != null)
             Did we find anything?
```



### Find

```
private Entry,Entry
   find(Object object) {
   Entry pred, curr, succ;
   boolean[] pmark = new boolean[1];
   boolean[] cmark = new boolean[1];
   int key = object.hashCode();
   tryAgain: while (true) {
    ...
}}
```



# Find

```
private Entry, Entry, Entry
  find(Object object)
 Entry pred, curr, succ;
 boolean[] pmark = new boolean[1];
 boolean[] cmark = new boolean[1];
 int key = object.hashCode();
tryAgain: while (true) {
}}}
```



```
private Entry, Entry
   find(Object object) {
   Entry pred, curr, succ;
   boolean[] pmark = new boolean[1];
   boolean[] cmark = new boolean[1];
   int key = object.hashCode(),
   tryAgain: while (true) {
    ...
}}
```

# Deleted bits for pred and curr



```
private Entry,Entry
  find(Object object) {
  Entry pred, curr, succ;
  boolean[] pmark = new boolean[1];
  boolean[] cmark = new boolean[1];
  int key = object.hashCode();
  tryAgain: while (true) {
    ...
}}
```

If list changes while traversed, start over

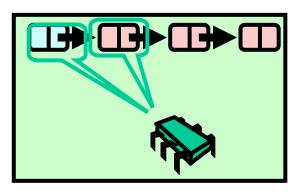


```
private Entry, Entry
   find(Object object) {
   Entry pred, curr, succ;
   boolean[] pmark = new boolean[1];
   boolean[] cmark = new boolean[1];
   int key = object.hashCode();
   tryAgain: while (true) {
    ...
}}
```

Lock-Free because we start over only if someone else makes progress



```
tryAgain: while (true) {
  pred = this.head.getReference();
  curr = pred.next.get(pmark);
  while (true) {
    ...
    Start with first two entries
```





```
tryAgain: while (true) {
  pred = this.head.getReference();
  curr = pred.next.get(pmark);
  while (true) {
    ...
}}
```

Move down the list



```
while (true) {
  if (curr == null)
    return pred, null, succ;
   succ = curr.next.get(cmark);
   int ckey = curr.key;
   if (isChanged(pred.next))
    continue tryAgain
}}}
```

Fell off the end of the list



```
while (true) {
  if (curr == null)
    return pred, null, succ;
  succ = curr.next.get(cmark);
  int ckey = curr.key;
  if (isChanged(pred.next))
    continue tryAgain;
}}
```

Get ref to successor and current deleted bit



```
Panic if predecessor's next
                    field changed
while (true) {
 if (curr == null)
  return pred, null,
 succ = curr.next.get(cmark);
 if (isChanged(pred.next))
  continue tryAgain;
```



```
while (true) {
if (!cmark[0]) {
     (curr.object == object)
    return pred
                 curr, succ;
   else if (ckey
    pred = curr;
  } else
   return prev, null, curk;
} else {
           If current node is not deleted
```



```
while (true) {
 if (!cmark[0]) {
  if (curr.object == object)
    return pred, curr, succ;
   else if (ckey <= key)
    pred = curr;
  } else
   return prev, null, cu
} else {
                    Object found
111
```

```
while (true) {
 if (!cmark[0]) {
   if (curr.object == object)
    return pred, curr, succ;
  else if (ckey <= key) {
    pred = curr;
  } else
   return prev, null,
} else {
                      Keep looking
111
```



```
while (true) {
                    Not there, give up
 if (!cmark[0]) {
   if (curr.object ==
    return pred, curr, succ;
   else if (ckey <= key)
    pred = curr
   else
   return prev, null, curr
} else {
BROWN
```

```
while (true) {
if (!cmark[0]) {
} else {
  f (pred_rext.compareAndSet()
  curr, succ, false, false))
   continue;
  else
                       Current entry is
   continue tryAgain;
                       logically deleted
```



```
Try to redirect predecessor's while true next reference
   if (!cmark[0]) {
     else {
    if (pred.next.compareAndSet()
     curr, succ, false, false))
       continue;
     else
       continue tryAgain;
```



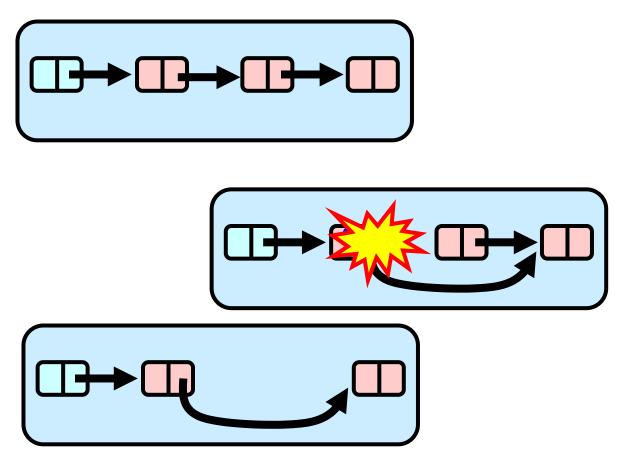
```
On success, keep going,
             on failure, start over
if (pred.next.compareAndSet()
curr, succ, false, false))
 continue;
else
 continue tryAgain;
```



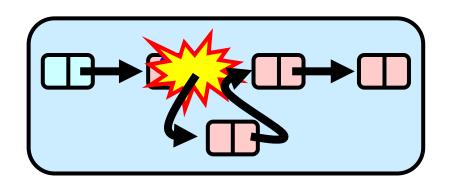
### Summary

- Coarse-grained locking
- Fine-grained locking
- · Optimistic syncrhronization
- · Lock-free synchronization

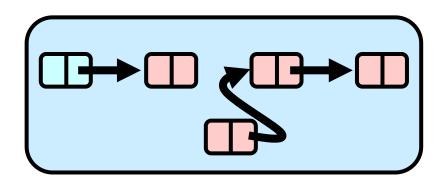




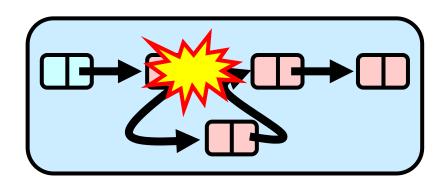




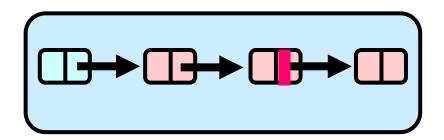




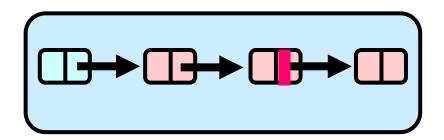














## Removing an Entry

