-Lab 1-

Warm-up: the target machine: RISCV

Objective

- Be familiar with the RISCV instruction set.
- Understand how it executes on the RISCV processor with the help of a simulator.
- Write simple programs, assemble, execute.

1.1 The RISCV processor, instruction set, simulator

EXERCISE #1 ► Lab preparation

Clone the Git repository for this year's labs:

```
git clone https://github.com/Drup/cap-labs21.git
```

If you haven't done so already, follow the instructions to compile riscv-xxx-gcc and spike on your machine (see INSTALL.md file). On the ENS machines, all installations have been done for you. However, you still have to add the following lines to your .bashrc:

```
RISCV="/usr/local/riscv"
export PATH="$RISCV/bin:$PATH"
```

Then, run the command source .bashrc or restart your terminal to update your environment.

EXERCISE #2 \triangleright RISCV C-compiler and simulator, first test

In the directory TP01/riscv/:

- Compile the provided file ex1.c with:
 - riscv64-unknown-elf-gcc ex1.c -o ex1.riscv

It produces a RISCV binary.

- Execute the binary with the RISCV simulator:
 - spike pk ex1.riscv
 - This should print 42. If you get a runtime exception, try running spike -m100 pk ex1.riscv instead: this limits the RAM usage of spike to 100 MB (the default is 2 GB).
- The corresponding RISCV can be obtained in a more readable format by:
 - riscv64-unknown-elf-gcc ex1.c -S -o ex1.s -fverbose-asm

(have a look at the generated .s file!)

The objective of this sequence of labs is to design our own (subset of) C compiler for RISCV.

EXERCISE #3 **▶ Documents**

Some documentation can be found in the RISCV ISA on the course webpage and in Appendix A.

```
https://compil-lyon.gitlabpages.inria.fr/
```

In the architecture course, you already saw a version of the target machine RISCV. The instruction set is depicted in Appendix A.

1.1.1 Hand exercises

EXERCISE #4 ► TD

On paper, write (in RISCV assembly language) a program which initializes the t_0 register to 1 and increments it until it becomes equal to 8.

EXERCISE #5 ► TD: sum

Write a program in RISCV assembly that computes the sum of the 10 first positive integers (excluded 10).

1.1.2 Assembling, disassembling

EXERCISE #6 \triangleright Hand assembling, simulation of the hex code

Assemble by hand (on paper) the instructions:

You will need the set of instructions of the RISCV machine and their associated opcode. All the info is in the ISA documentation.

To check your solution (after you did the job manually), you can redo the assembly using the toolchain:

```
riscv64-unknown-elf-as -march=rv64g asshand.s -o asshand.o
```

asshand.o is an ELF file which contains both the compiled code and some metadata (you can try hexdump asshand.o to view its content, but it's rather large and unreadable). The tool objdump allows extracting the code section from the executable, and show the binary code next to its disassembled version:

```
riscv64-unknown-elf-objdump -d asshand.o
```

Check that the output is consistent with what you found manually.

EXERCISE #7 \triangleright Hand disassembling

Guess a RISCV program that assembles itself into:

Listing 1.1: disass.lst

```
disass.o: format de fichier elf64-littleriscv
```

Déassemblage de la section .text:

0000000000000000 <main>:

```
0: 00128313 xx
4: ffdff06f yy
8: 00008067 zz
```

From now on, we are going to write programs using an easier approach. We are going to write instructions using the RISCV assembly.

1.2 RISCV Simulator

EXERCISE #8 ► **Execution and debugging**

See https://www.lowrisc.org/docs/tagged-memory-v0.1/spike/for details on the Spike simulator.

test_print.s is a small but complete example using Risc-V assembly. It uses the println_string, print_int, print_char and newline functions provided to you in libprint.s. Each function can be called with call print_... and prints the content of register a0 (call newline takes no input and prints a newline character).

- First test assembling and simulation on the file test_print.s: riscv64-unknown-elf-as -march=rv64g test_print.s -o test_print.o
- 2. Optionally, run riscv64-unknown-elf-objdump -D as in previous exercise. The -D option shows all sections, including .rodata.
- 3. The libprint.s library must be assembled too:

riscv64-unknown-elf-as -march=rv64g libprint.s -o libprint.o

4. We now link these files together to get an executable:

riscv64-unknown-elf-gcc test_print.o libprint.o -o test_print

The generated test_print file should be executable, but since it uses the Risc-V ISA, we can't execute it natively (try./test_print, you'll get an error like Exec format error).

5. Run the simulator:

```
spike pk ./test_print
The output should look like:
bbl loader
HI MIF08!
42
a
```

The first line comes from the simulator itself, the next two come from the println_string, print_int and print_char calls in the assembly code.

6. We can also view the instructions while they are executed:

```
spike -l pk ./test_print
```

Unfortunately, this shows all the instructions in pk (Proxy Kernel, a kind of mini operating system), and is mostly unusable. Alternatively, we can run a step-by-step simulation starting from a given symbol. To run the instructions in main, we first get the address of main in the executable:

```
$ riscv64-unknown-elf-nm test_print | grep main
00000000001014c T main
```

This means: main is a symbol defined in the .text section (T in the middle column), it is global (capital T), and its address is 1014c. Now, run spike in debug mode (-d) and execute code up to this address (until pc 0 1014c, i.e. "Until the program counter of core 0 reaches 1014c"). Press Return to move to the next instruction and q to quit:

```
$ spike -d pk ./test_print
: until pc 0 1014c
bbl loader
core
       0: 0x000000000001014c (0xff010113) addi
                                                   sp, sp, -16
       0: 0x0000000000010150 (0x00113423) sd
core
                                                   ra, 8(sp)
core
       0: 0x0000000000010154 (0x0001d7b7) lui
                                                   a5. 0x1d
       0: 0x0000000000010158 (0x02078513) addi
                                                   a0, a5, 32
core
: q
$
```

Remark: For your labs, you may want to assemble and link with a single command (which can also do the compilation if you provide . c files on the command-line):

```
riscv64-unknown-elf-gcc -march=rv64g libprint.s test_print.s -o main
```

In real-life, people run compilation+assembly and link as two different commands, but use a build system like a Makefile to re-run only the right commands.

EXERCISE #9 \triangleright Algo in RISCV assembly

Write (in minmax.s) a program in RISCV assembly that computes the min of two integers, and stores the result in a precise location of the memory that has the label min. Try with different values. We use 64 bits of memory to store ints, i.e., use .dword directive and ld and sd instructions.

EXERCISE #10 ► (Advanced) Algo in RISCV assembly

Write and execute the following programs in assembly:

• Count the number of non-nul bits of a given integer, print the result.

• Draw squares and triangles of stars (character '*') of size *n*, *n* being stored somewhere in memory. Examples:

```
n=3 square:
***

***
n=3 triangle:

* * *
```

1.3 Finished?

If you're done with the lab, do the python tutorial at the following address:

https://docs.python.org/fr/3.5/tutorial/

Appendix A

RISCV Assembly Documentation (ISA), rv64g

About

- RISCV is an open instruction set initially developed by Berkeley University, used among others by Western Digital, Alibaba and Nvidia.
- We are using the rv64g instruction set: Risc-V, 64 bits, General purpose (base instruction set, and extensions for floating point, atomic and multiplications), without compressed instructions. In practice, we will use only 32 bits instructions (and very few of floating point instructions).
- Document: Laure Gonnord and Matthieu Moy, for CAP and MIF08.

This is a simplified version of the machine, which is (hopefully) conform to the chosen simulator.

A.1 Installing the simulator and getting started

To get the RISCV assembler and simulator, follow instructions of the first lab (git pull on the course lab repository).

A.2 The RISCV architecture

Here is an example of RISCV assembly code snippet (a proper main function would be needed to execute it, cf. course and lab):

```
addi a0, zero, 17 # initialisation of a register to 17
loop:
  addi a0, a0, -1 # subtraction of an immediate
  j loop # equivalent to jump xx
```

The rest of the documentation is adapted from https://github.com/riscv/riscv-asm-manual/blob/master/riscv-asm.md and https://github.com/jameslzhu/riscv-card/blob/master/riscv-card.pdf

A.3 RISC-V Assembly Programmer's Manual - adapted for CAP and MIF08

A.3.1 Copyright and License Information - Documents

The RISC-V Assembly Programmer's Manual is

© 2017 Palmer Dabbelt palmer@dabbelt.com © 2017 Michael Clark michaeljclark@mac.com © 2017 Alex Bradbury asb@lowrisc.org

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- Official Specifications webpage: https://riscv.org/specifications/
- Latest Specifications draft repository: https://github.com/riscv/riscv-isa-manual

This document has been modified by Laure Gonnord & Matthieu Moy, in 2019.

A.3.2 Registers

Registers are the most important part of any processor. RISC-V defines various types, depending on which extensions are included: The general registers (with the program counter), control registers, floating point registers (F extension), and vector registers (V extension). We won't use control nor F or V registers.

General registers

The RV32I base integer ISA includes 32 registers, named x0 to x31. The program counter PC is separate from these registers, in contrast to other processors such as the ARM-32. The first register, x0, has a special function: Reading it always returns 0 and writes to it are ignored.

In practice, the programmer doesn't use this notation for the registers. Though x1 to x31 are all equally general-use registers as far as the processor is concerned, by convention certain registers are used for special tasks. In assembler, they are given standardized names as part of the RISC-V **application binary interface** (ABI). This is what you will usually see in code listings. If you really want to see the numeric register names, the -M argument to objdump will provide them.

Register	ABI	Use by convention	Preserved?
x0	zero	hardwired to 0, ignores writes	n/a
x1	ra	return address for jumps	no
x2	sp	stack pointer	yes
x3	gp	global pointer	n/a
x4	tp	thread pointer	n/a
x5	t0	temporary register 0	no
x6	t1	temporary register 1	no
x7	t2	temporary register 2	no
x8	s0 or fp	saved register 0 or frame pointer	yes
x9	s1	saved register 1	yes
x10	a0	return value or function argument 0	no
x11	al	return value <i>or</i> function argument 1	no
x12	a2	function argument 2	no
x13	a3	function argument 3	no
x14	a4	function argument 4	no
x15	a5	function argument 5	no
x16	a6	function argument 6	no
x17	a7	function argument 7	no
x18	s2	saved register 2	yes
x19	s3	saved register 3	yes
x20	s4	saved register 4	yes
x21	s5	saved register 5	yes
x22	s6	saved register 6	yes
x23	s7	saved register 6	yes
x24	s8	saved register 8	yes
x25	s9	saved register 9	yes
x26	s10	saved register 10	yes
x27	s11	saved register 11	yes
x28	t3	temporary register 3	no
x29	t4	temporary register 4	no
x30	t5	temporary register 5	no
x31	t6	temporary register 6	no
pc	(none)	program counter	n/a

Registers of the RV32I. Based on RISC-V documentation and Patterson and Waterman "The RISC-V Reader" (2017)

As a general rule, the saved registers s0 to s11 are preserved across function calls, while the argument

registers a0 to a7 and the **temporary registers** t0 to t6 are not. The use of the various specialized registers such as sp by convention will be discussed later in more detail.

A.3.3 Instructions

Arithmetic

```
add, addi, sub, classically.

addi a0, zero, 42

initialises a0 to 42.
```

Labels

Text labels are used as branch, unconditional jump targets and symbol offsets. Text labels are added to the symbol table of the compiled module.

```
loop:
    j loop
```

Jumps and branches target is encoded with a relative offset. It is relative to the beginning of the current instruction. For example, the self-loop above corresponds to an offset of 0.

Branching

Test and jump, within the same instruction:

```
beq a0, a1, end
```

tests whether a0=a1, and jumps to 'end' if its the case.

Absolute addressing

The following example shows how to load an absolute address:

which generates the following assembler output and relocations as seen by objdump:

```
0000000000000000 <_start>:
```

```
0: 000005b7 lui a1,0x0
0: R_RISCV_HI20 msg
4: 00858593 addi a1,a1,8 # 8 <.L21>
4: R_RISCV_L012_I msg
```

Relative addressing

The following example shows how to load a PC-relative address:

which generates the following assembler output and relocations as seen by objdump:

```
0000000000000000 <_start>:
```

Load Immediate

The following example shows the li pseudo instruction which is used to load immediate values:

```
li a0, 0x76543210
```

which generates the following assembler output as seen by objdump (generated code will be different depending on the constant):

```
0: 76543537 lui a0,0x76543
4: 2105051b addiw a0,a0,528
```

Load Address

The following example shows the la pseudo instruction which is used to load symbol addresses:

A.3.4 Assembler directives for CAP and MIF08

Both the RISC-V-specific and GNU .-prefixed options. The following table lists assembler directives:

Directive	Arguments	Description			
.align	integer	align to power of 2 (alias for .p2align)			

Directive	Arguments	Description
.file	"filename"	emit filename FILE LOCAL symbol table
.globl	symbol_name	emit symbol_name to symbol table (scope GLOBAL)
.local	symbol_name	emit symbol_name to symbol table (scope LOCAL)
.section	[{.text,.data,.rodata,.bss}]	emit section (if not present, default .text) and make current
.size	symbol, symbol	accepted for source compatibility
.text		emit .text section (if not present) and make current
.data		emit .data section (if not present) and make current
.rodata		emit .rodata section (if not present) and make current
.string	"string"	emit string
.equ	name, value	constant definition
.word	expression [, expression]*	32-bit comma separated words
.balign	b,[pad_val=0]	byte align
.zero	integer	zero bytes

A.3.5 Assembler Relocation Functions

The following table lists assembler relocation expansions:

Assembler Notation	Description	Instruction / Macro
%hi(symbol)	Absolute (HI20)	lui
%lo(symbol)	Absolute (LO12)	load, store, add
%pcrel_hi(symbol)	PC-relative (HI20)	auipc
%pcrel_lo(label)	PC-relative (LO12)	load, store, add

A.3.6 Instruction encoding

 $\label{lem:composition} \textbf{Credit} \quad \text{This is a subset of the RISC-V greencard, by James Izhu, licence CC by SA, $$https://github.com/jameslzhu/riscv-card$$

Core Instruction Formats

31	27	26	25	24	20	19		15	14	12	11	7	6		0	
	func	ct7		rs	32		rs1		fun	ct3		rd	op	code		R-type
	i	mm[11:0)]			rs1		fun	ct3		rd	op	code		I-type
i	mm[11:5]		rs	32		rs1		fun	ct3	imı	n[4:0]	op	code		S-type
im	m[12	2 10:5	5]	rs	32		rs1		fun	ct3	imm	[4:1 11]	op	code		B-type
imm[3				1:12]						rd	op	code		U-type		
imm[20 10:1					11 1	9:12]					rd	op	code		J-type	

RV32I Base Integer Instructions - CAP subset

Inst	Name	FMT	Opcode	funct3	funct7	Description (C)	Note
add	ADD	R	0110011	0x0	0x00	rd = rs1 + rs2	
sub	SUB	R	0110011	0x0	0x20	rd = rs1 - rs2	
xor	XOR	R	0110011	0x4	00x0	rd = rs1 ^ rs2	
or	OR	R	0110011	0x6	00x0	rd = rs1 rs2	
and	AND	R	0110011	0x7	0x00	rd = rs1 & rs2	
slt	Set Less Than	R	0110011	0x2	00x0	rd = (rs1 < rs2)?1:0	
sltu	Set Less Than (U)	R	0110011	0x3	00x0	rd = (rs1 < rs2)?1:0	zero-extends
addi	ADD Immediate	I	0010011	0x0		rd = rs1 + imm	
xori	XOR Immediate	I	0010011	0x4		rd = rs1 ^ imm	
ori	OR Immediate	I	0010011	0x6		rd = rs1 imm	
andi	AND Immediate	I	0010011	0x7		rd = rs1 & imm	
1b	Load Byte	I	0000011	0x0		rd = M[rs1+imm][0:7]	
lw	Load Word	I	0000011	0x2		rd = M[rs1+imm][0:31]	
lbu	Load Byte (U)	I	0000011	$1 0x4 \qquad \qquad rd = M[rs1+imm][0:7]$		zero-extends	
sb	Store Byte	S	0100011	0x0	M[rs1+imm][0:7] = rs2[0:7]		
SW	Store Word	S	0100011	0x2		M[rs1+imm][0:31] = rs2[0:31]	
beq	Branch ==	В	1100011	0x0		if(rs1 == rs2) PC += imm	
bne	Branch !=	В	1100011	0x1		if(rs1 != rs2) PC += imm	
blt	Branch <	В	1100011	0x4		if(rs1 < rs2) PC += imm	
bge	Branch ≥	В	1100011	0x5	if(rs1 >= rs2) PC += imm		
bltu	Branch < (U)	В	1100011	0x6	if(rs1 < rs2) PC += imm		zero-extends
bgeu	Branch ≥ (U)	В	1100011	1 $0x7$ if(rs1 >= rs2) PC += imm		zero-extends	
jal	Jump And Link	J	1101111			rd = PC+4; PC += imm	
jalr	Jump And Link Reg	I	1100111	0x0	rd = PC+4; PC = rs1 + imm		
lui	Load Upper Imm	U	0110111		rd = imm << 12		
auipc	Add Upper Imm to PC	U	0010111			rd = PC + (imm << 12)	

Pseudo Instructions

Pseudoinstruction	Base Instruction(s)	Meaning		
la rd, symbol	auipc rd, symbol[31:12]	Load address		
	addi rd, rd, symbol[11:0] auipc rd, symbol[31:12]			
{lb lh lw ld} rd, symbol	{lb lh lw ld} rd, symbol[11:0](rd)	Load global		
	auipc rt, symbol[31:12]			
{sb sh sw sd} rd, symbol, rt	s{b h w d} rd, symbol[11:0](rt)	Store global		
	auipc rt, symbol[31:12]			
{flw fld} rd, symbol, rt	fl{w d} rd, symbol[11:0](rt)	Floating-point load global		
66 16 12 1 1 1 1	auipc rt, symbol[31:12]	T1		
{fsw fsd} rd, symbol, rt	fs{w d} rd, symbol[11:0](rt)	Floating-point store global		
nop	addi x0, x0, 0	No operation		
li rd, immediate	Myriad sequences	Load immediate		
mv rd, rs	addi rd, rs, 0	Copy register		
not rd, rs	xori rd, rs, -1	One's complement		
neg rd, rs	sub rd, x0, rs	Two's complement		
negw rd, rs	subw rd, x0, rs	Two's complement word		
sext.w rd, rs	addiw rd, rs, 0	Sign extend word		
seqz rd, rs	sltiu rd, rs, 1	Set if = zero		
snez rd, rs	sltu rd, x0, rs	Set if ≠ zero		
sltz rd, rs	slt rd, rs, x0	Set if < zero		
sgtz rd, rs	slt rd, x0, rs	Set if > zero		
fmv.s rd, rs	fsgnj.s rd, rs, rs	Copy single-precision register		
fabs.s rd, rs	fsgnjx.s rd, rs, rs	Single-precision absolute value		
fneg.s rd, rs	fsgnjn.s rd, rs, rs	Single-precision negate		
fmv.d rd, rs	fsgnj.d rd, rs, rs	Copy double-precision register		
fabs.d rd, rs	fsgnjx.d rd, rs, rs	Double-precision absolute value		
fneg.d rd, rs	fsgnjn.d rd, rs, rs	Double-precision negate		
beqz rs, offset	beq rs, x0, offset	Branch if $=$ zero		
bnez rs, offset	bne rs, x0, offset	Branch if ≠ zero		
blez rs, offset	bge x0, rs, offset	Branch if ≤ zero		
bgez rs, offset	bge rs, x0, offset	Branch if ≥ zero		
bltz rs, offset	blt rs, x0, offset	Branch if < zero		
bgtz rs, offset	blt x0, rs, offset	Branch if > zero		
bgt rs, rt, offset	blt rt, rs, offset	Branch if >		
ble rs, rt, offset	bge rt, rs, offset	Branch if ≤		
bgtu rs, rt, offset	bltu rt, rs, offset	Branch if >, unsigned		
bleu rs, rt, offset	bgeu rt, rs, offset	Branch if ≤, unsigned		
j offset	jal x0, offset	Jump		
jal offset	jal x1, offset	Jump and link		
jr rs	jalr x0, rs, 0	Jump register		
jalr rs	jalr x1, rs, 0	Jump and link register		
ret	jalr x0, x1, 0	Return from subroutine		
call offset	<pre>auipc x1, offset[31:12]</pre>	Call far-away subroutine		
	<pre>jalr x1, x1, offset[11:0]</pre>	Can far away Sabioutific		
tail offset	<pre>auipc x6, offset[31:12]</pre>	Tail call far-away subroutine		
	jalr x0, x6, offset[11:0]			
fence	fence iorw, iorw	Fence on all memory and I/O		