COMSM1302
Overview of Computer Architecture

Lecture 4

Simple devices

In this lecture

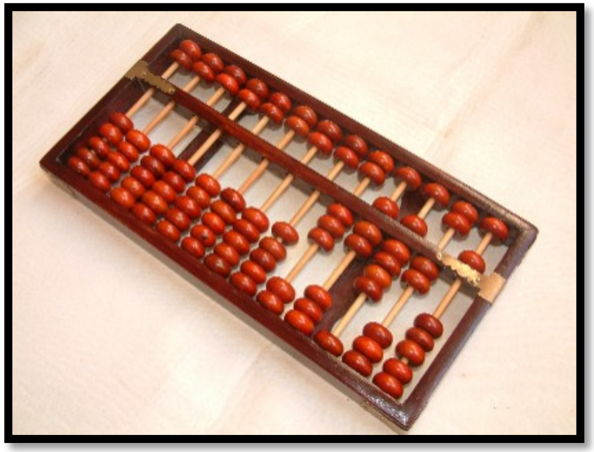
Foundations	• Data representation, logic, Boolean algebra.
Building blocks	• Transistors, transistor based logic, **simple devices**, storage.
Modules	• Memory, simple controllers, FSMs, processors and execution.
Programming	• Machine code, assembly, high-level languages, compilers.
Wrap-up	• Operating systems, energy aware computing.

Today, we learn to add!

- And also…
  - Subtract
  - Select 1 signal from many
  - Distribute 1 signal to many
- The circuits shown hereafter will be drawn in Logisim.
  - You can download it from:

  http://sourceforge.net/projects/circuit/

  - Install Logisim on your own computer and practice.



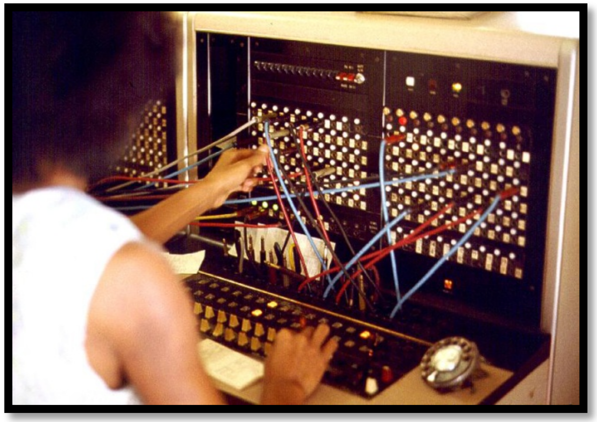


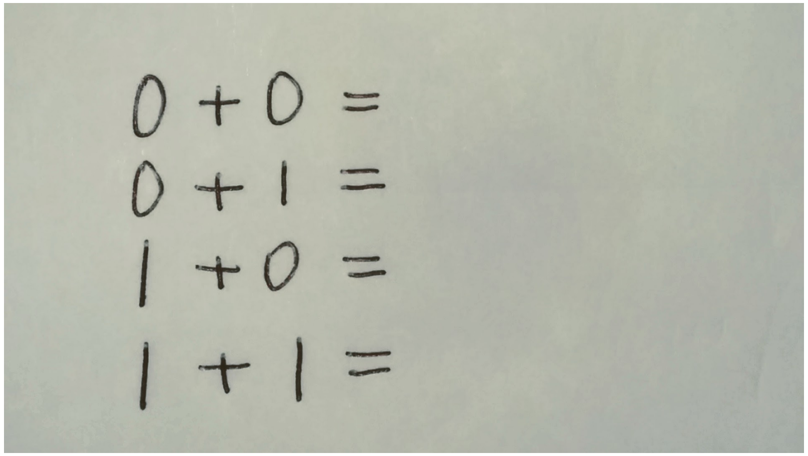
Photo by Joseph A. Carr, 1975

COMSM1302 NAND board kit

- Your task in the NAND labs will involve building some of the circuits that will be introduced today.
- **You should always design with Logisim BEFORE you start building with the NAND board kit.**

The simplest of binary addition

- How to add two, single-digit binary numbers?
  - Each digit is either 0 or 1
  - What are the possible results?



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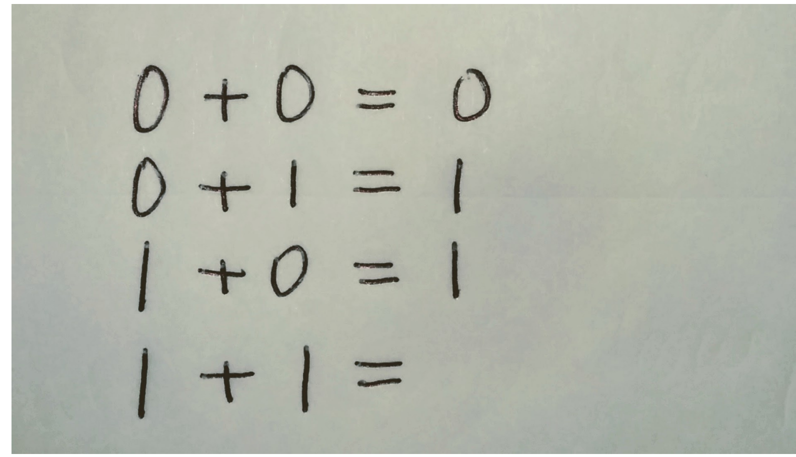

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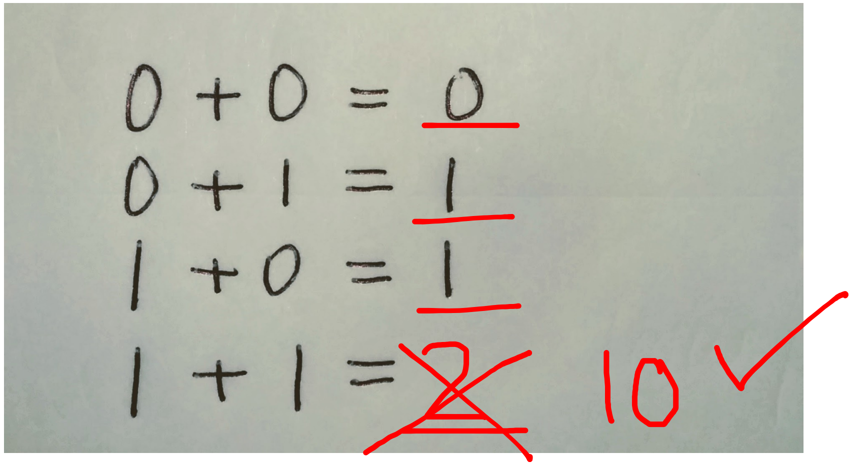

The simplest of binary addition

- How to add two, single-digit binary numbers?
  - Each digit is either 0 or 1
  - There are three possible results
    - 0, 1, 2
    - 0b00, 0b01, 0b10
  - Two inputs
    - A, B
  - Output
    - Sum (S)

A		B		S
0	+	0	=	0
0	+	1	=	1
1	+	0	=	1
1	+	1	=	10

The simplest of binary addition

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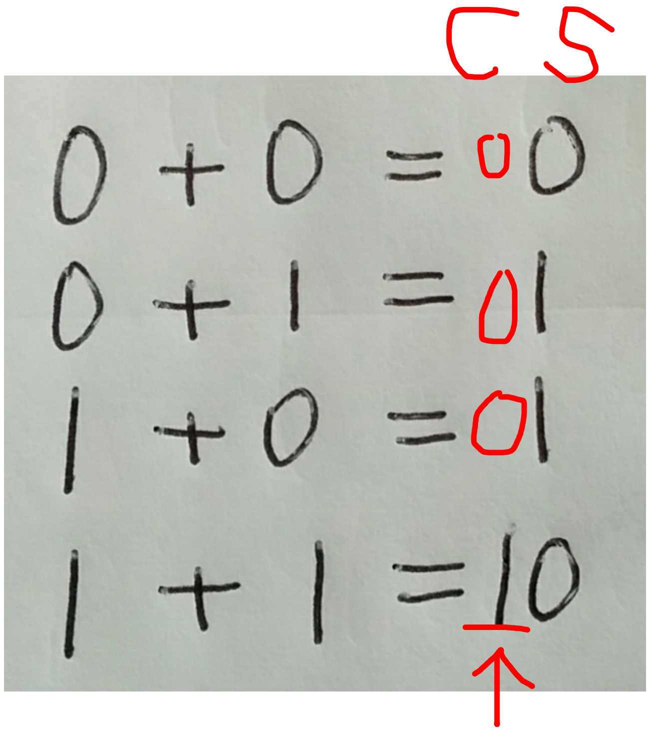
    - Carry (C)



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    - A, B
  - Two outputs
    - Carry (C)
    - Sum (S)

A	B	C	S
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

The simplest of binary addition

- How to add two, single-digit binary numbers?
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  - Two inputs
    - A, B
  - Two outputs
    - Carry (C)
    - Sum (S)

C S

A	B	2^1 2	2^0 1
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

With some Boolean algebra

- Remember our truth tables for the connectives in Boolean algebra, e.g. AND, OR, XOR, etc?
- Which operations can be used to help us generate S and C?

A	B	C	S
0	0	0	0
0	1	0	1
1	0	0	1
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With some Boolean algebra

- Remember our truth tables for the connectives in Boolean algebra, e.g. AND, OR, XOR, etc?
- Which operations can be used to help us generate S and C?

A	B	C	S
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

XOR

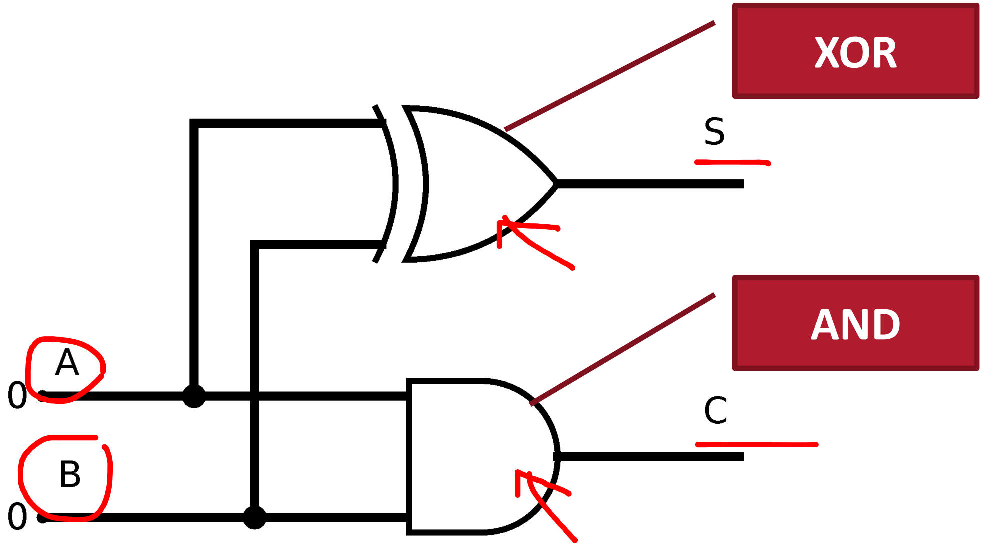
A	B	¬A	A ∧ B	A ∨ B	A ⊕ B	A → B	A≡B
0	0	1	0	0	0	1	1
0	1	1	0	1	1	1	0
1	0	0	0	1	1	0	0
1	1	0	1	1	0	1	1

The half-adder

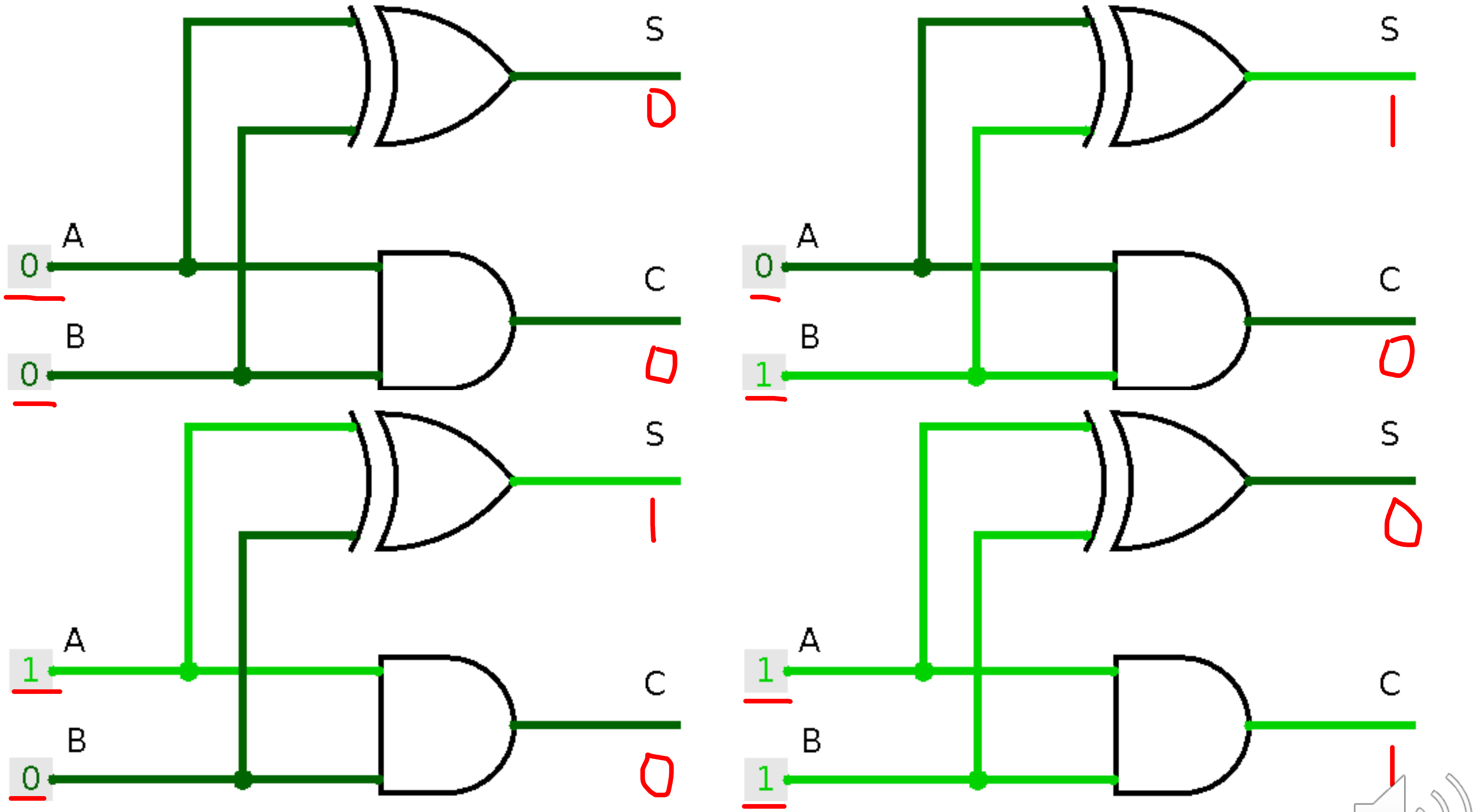
- $S = A \oplus B$
- $C = A \wedge B$

add two bits together, no carry in

- Now let's build it with logic gates.



The half-adder in action



Now what?

- We can add two bits together.
  - By generating a sum and a carry bit.
- How do we add multiple bits together?

$1+3=$ $3+3=$ $15+7=$

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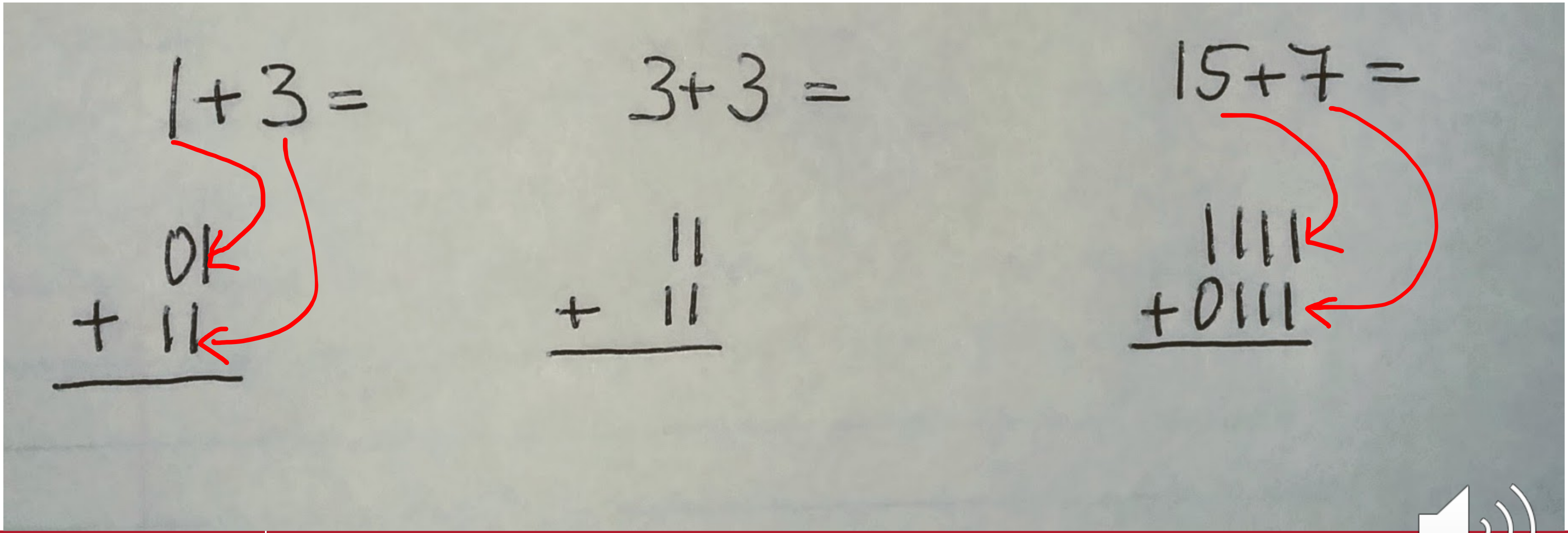

Now what?

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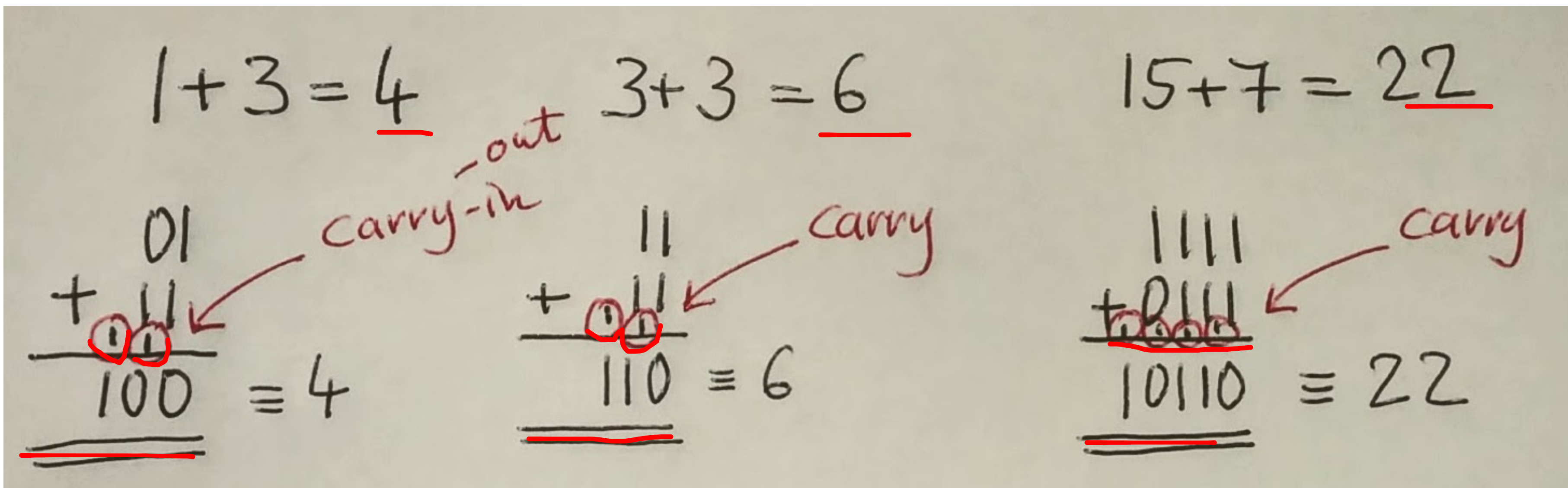
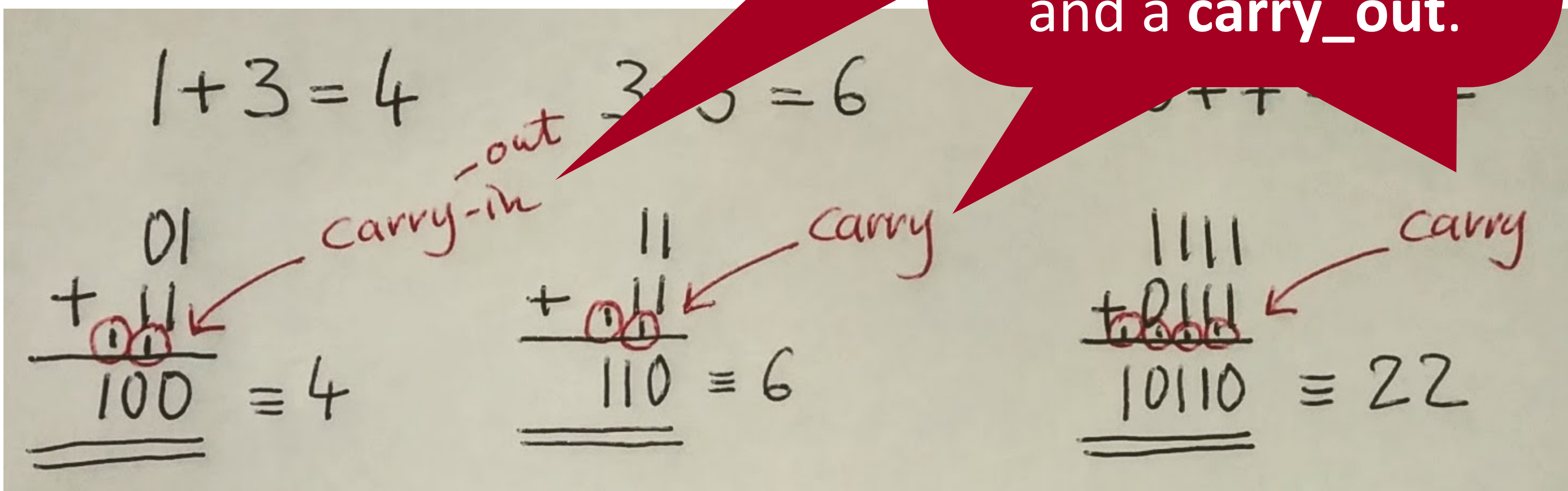


Now what?

- We can add two bits togeth
  - By generating a sum and a
- How do we add multip

For each addition (except for the first) we need to account for two input bits <u>and</u> a **carry_in**, and we produce a **sum** and a **carry_out**.



The full-adder

C_in	A	B	C_out	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

From the half-adder:

- $S = A \oplus B$
- $C\_out = A \land B$

(Note that the above covers only the top-half of this table!)

The full-adder

C_in	A	B	C_out	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

The full-adder

For multiple arguments, XOR is defined to be true iff an odd number of its arguments is true, and false otherwise.

C_in	A	B	C_out	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

- $S = A \oplus B \oplus C\_in$
- $C\_out = (A \wedge B) \vee (C\_in \wedge (A \vee B))$
- Also valid:

$C\_out = (A \wedge B) \vee (C\_in \wedge (A \oplus B))$

  – **Why?**

The full-adder

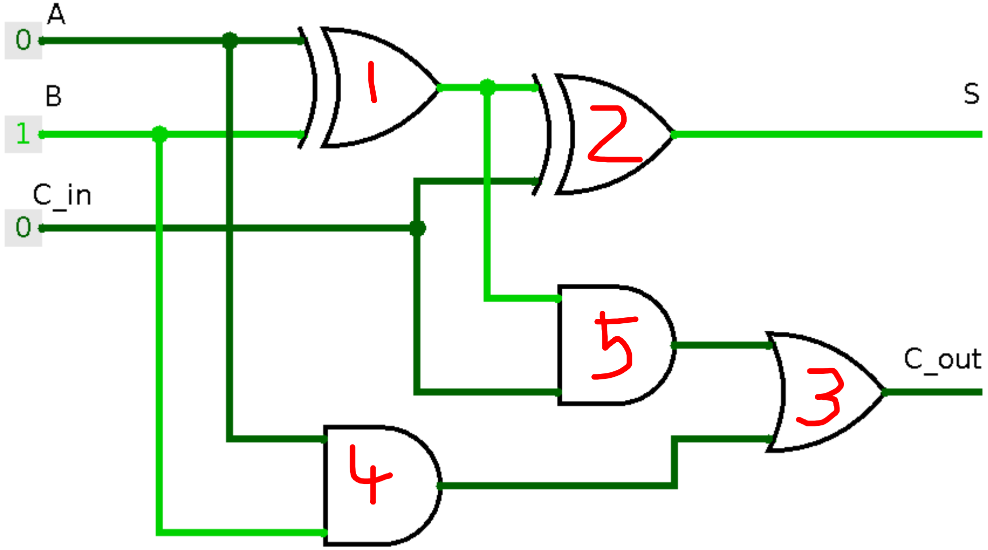
C_in	A	B	C_out	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

- $S = A \oplus B \oplus C\_in$
- $C\_out =$
  $(A \land B) \lor (C\_in \land (A \lor B))$

- Also valid:
  $C\_out = (A \land B) \lor (C\_in \land (A \oplus B))$
  - **Why?**

The full-adder

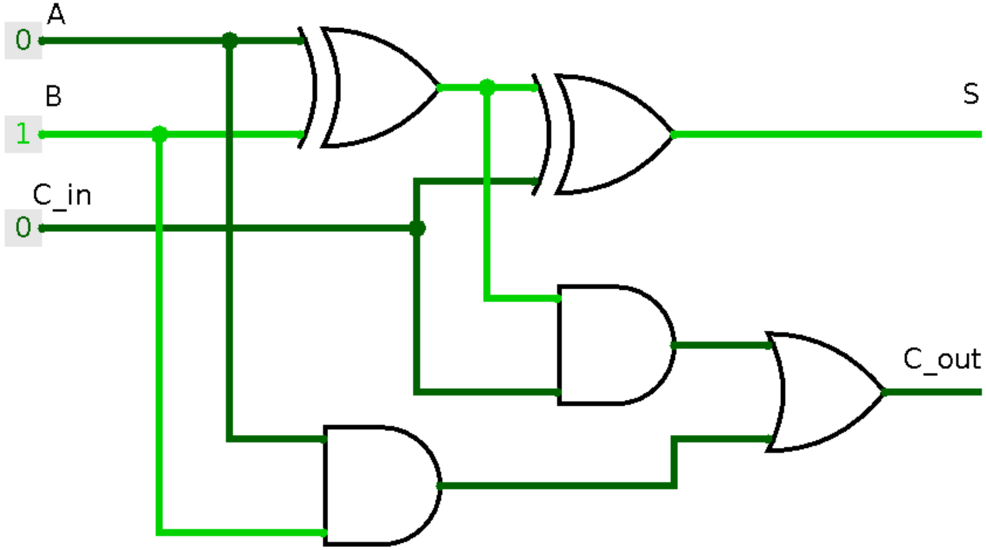
- 8 different combinations of input
- Try them yourself!



- S = A ⊕ B ⊕ C_in
- C_out =(A ∧ B) ∨ (C_in ∧ (A ⊕ B))

The full-adder

- 8 different combinations of input
- Try them yourself!



- To build this with the NAND kits you need to:
  - use Boolean algebra to obtain a design that is based purely on NAND gates,
  - implement this in Logisim to gain confidence your design is correct, then
  - *(and only then)* transfer to NAND boards and test.

Now what?

- We can add two bits together.
  - By generating a sum and a carry bit.
- We can add three bits together
  - By accommodating a carry-in as well as our regular two inputs.

- How do we add multiple bits together?

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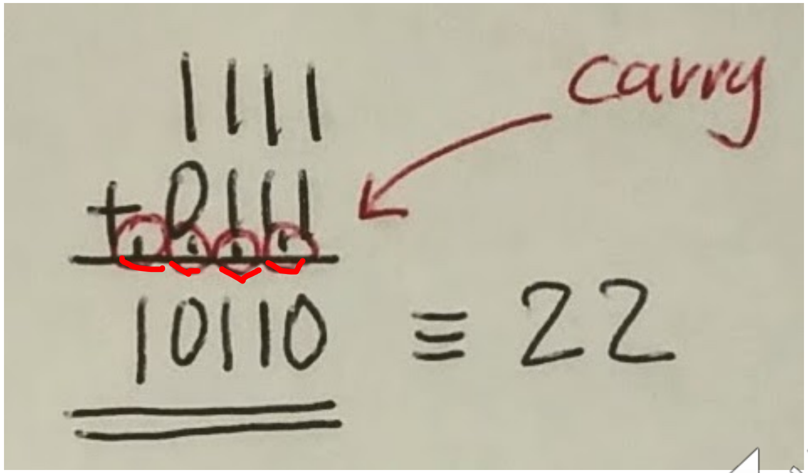
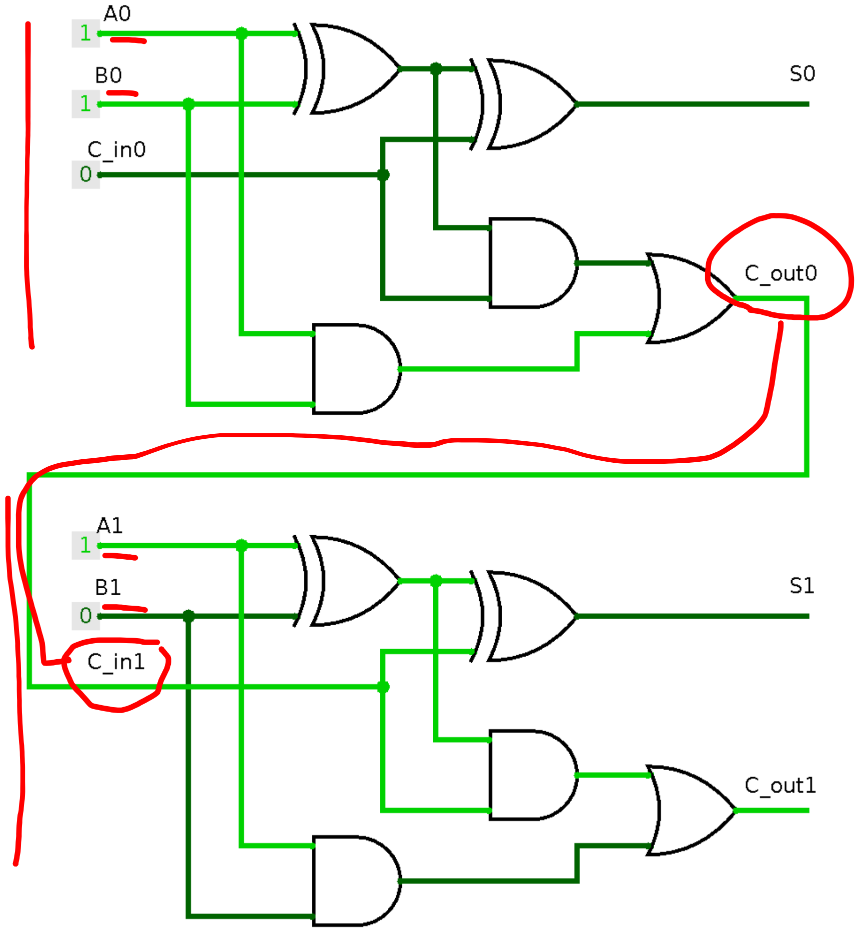
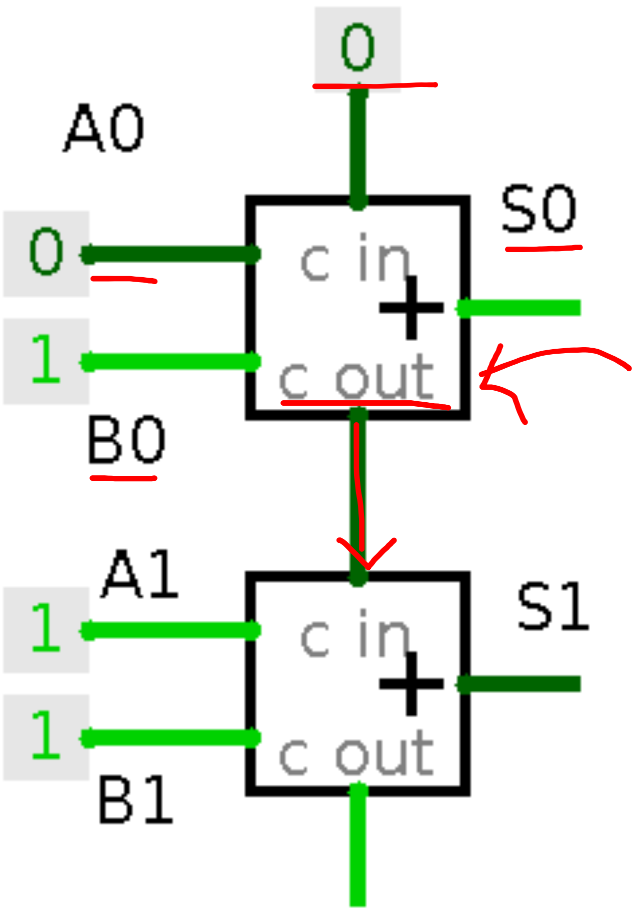


Chaining full-adders





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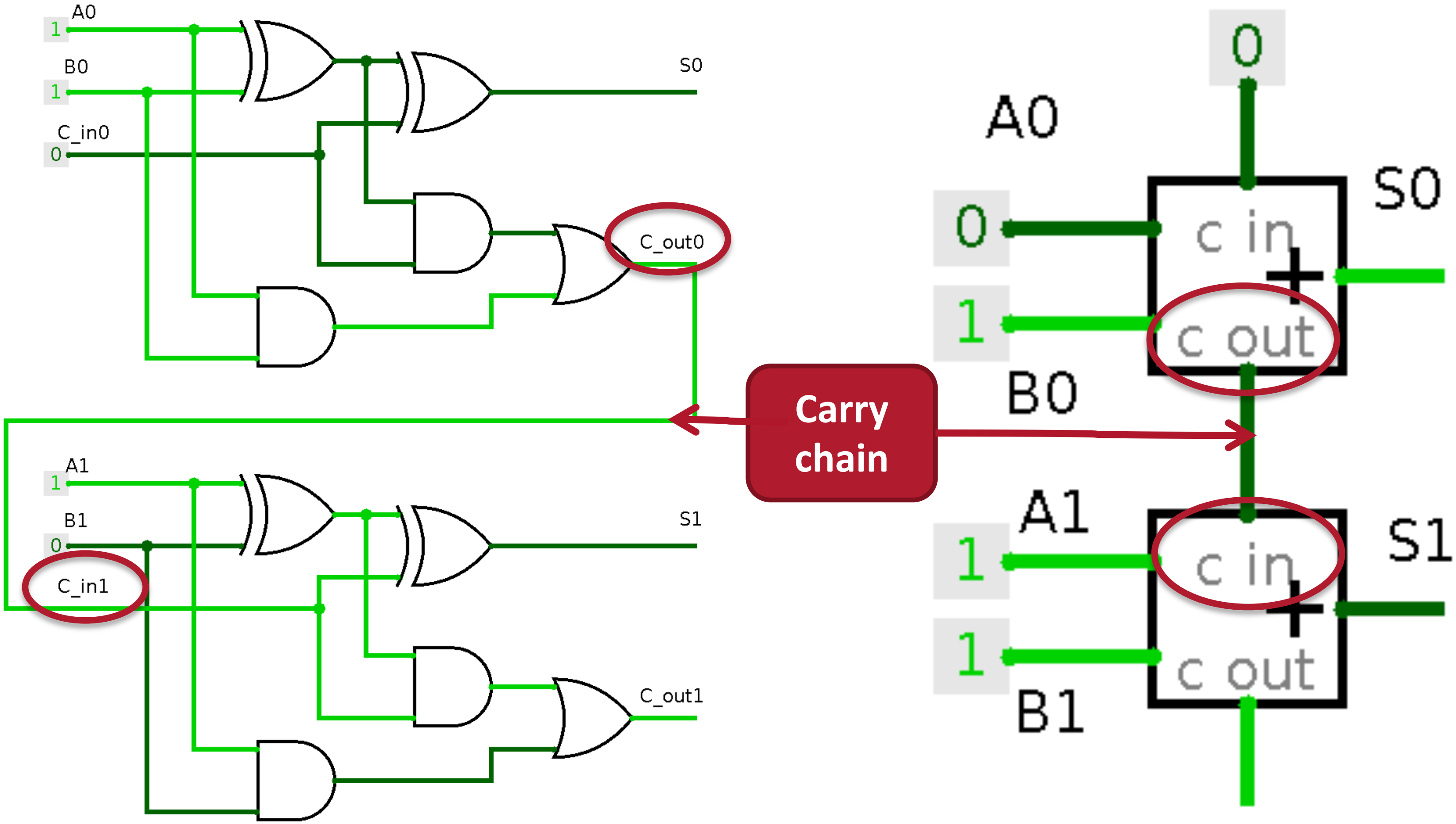
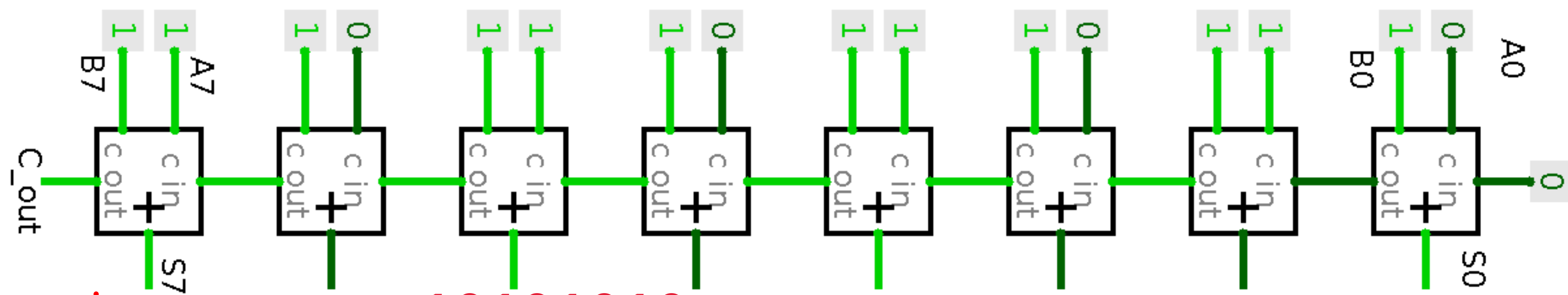


8-bit adder, ripple-carry adder



10101010
+11111111

- Named **ripple-carry** because a carry signal generated at the LSB (Least Significant Bit… bit 0) of the device can affect the result on any/all more significant bits.
- Does this have any implications?

Building blocks

- To recap what we've done:
  – Used Boolean algebra to identify our **building blocks**.
  – Built a unit capable of **adding two bits**.
    - **Half-adder**
  – Extended it to handle **carry-in**.
    - **Full-adder**
  – Chained them together to make an adder of **arbitrary size**.
    - **Ripple-carry adder**
- Now we can add anything!
  – Modern processors typically have adders between 8 and 64 bits.
  – **Why?**

Subtraction

- Subtraction is easy if we think of it as **adding one number to a negative number**.
- So let's represent this subtraction:

  1 – 2 = -1
- As:

  1 + -2 = -1

- How to negate a number?

Subtraction

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- How to negate a number?
  – We use 2s complement!

Reminder: 2s complement

$$Y = -x_{N-1} \cdot 2^{N-1} + \sum_{i=0}^{N-2} x_i \cdot 2^i$$

-128	64	32	16	8	4	2	1	Base 10
0	0	0	0	1	0	1	0	10
0	0	0	0	0	0	0	1	1
1	1	1	1	1	1	1	1	-1
1	1	1	1	1	1	1	0	-2
1	0	0	0	1	0	1	0	-118

Calculating the 2s complement

To calculate the 2's complement of an integer:

1. invert the binary equivalent of the number by changing all of the ones to zeroes and all of the zeroes to ones (also called 1's complement), then
2. add one.

- How do we represent -6?

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$110 \equiv 6$

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0110 ≡ 6

- How do we represent -6?

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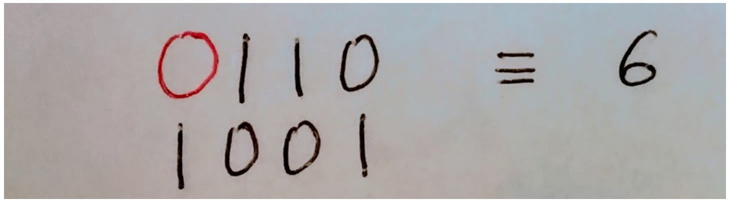

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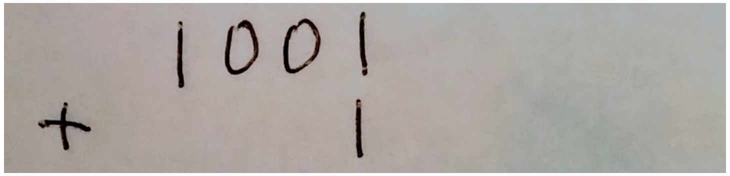

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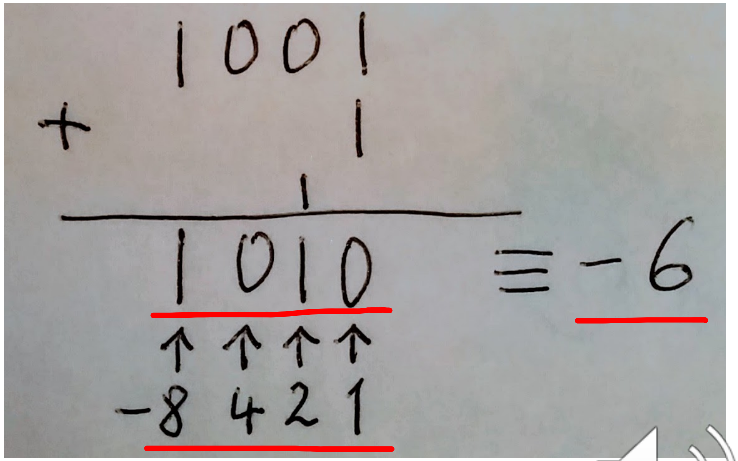

Subtraction with 2s complement

- A – B = A + (-B) = A + (Not(B) + 1)
- We already have all the building blocks we need to implement this!
  - NOT gates to flip bits
  - An unused C_in at the beginning of our adder's carry chain to provide the extra 1.

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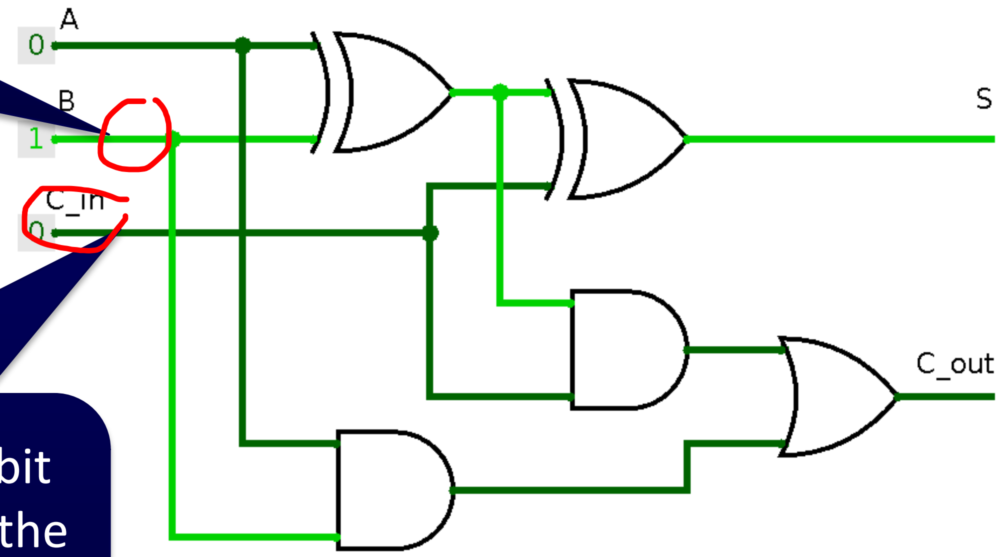
From full-adder to subtractor

Add a NOT gate here to obtain the 1s complement of B, i.e. NOT(B)



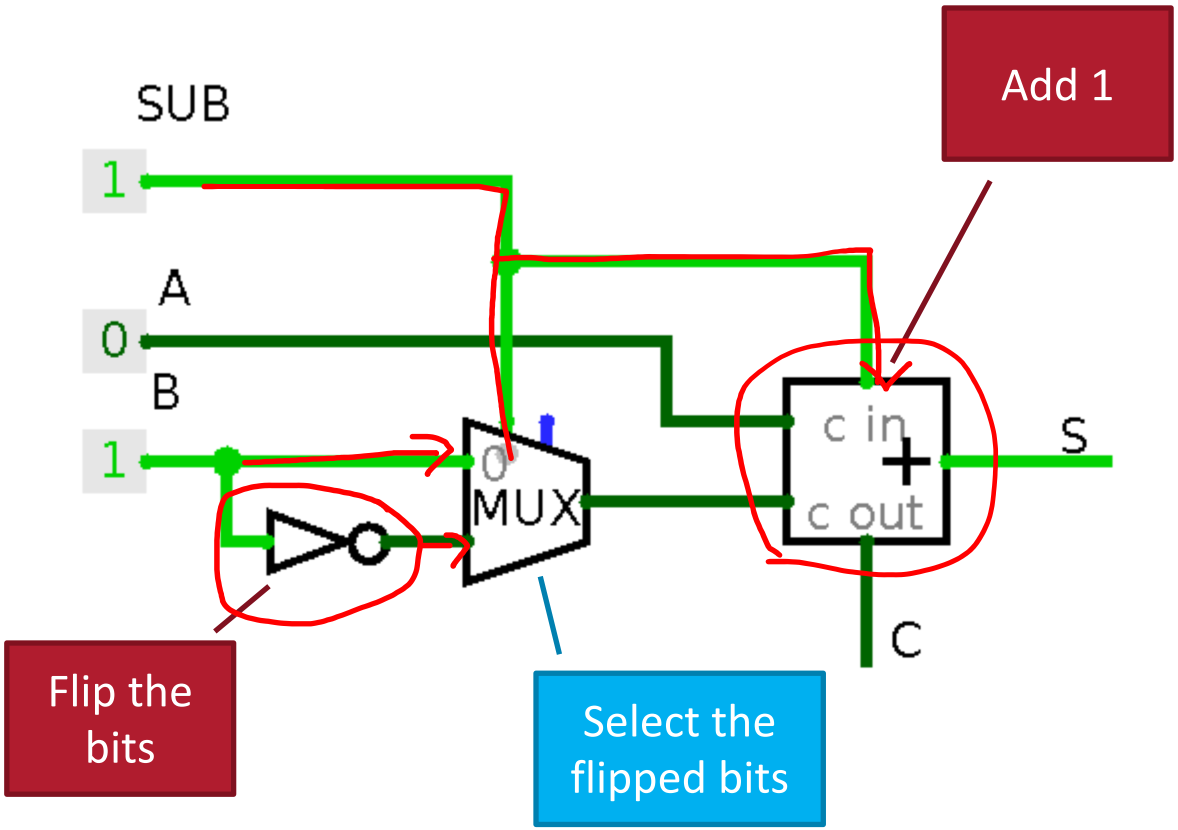
Set the C_in (at the lowest bit position) to 1; this provides the “extra” 1 we need to obtain the 2s complement of B.

The adder-subtractor?

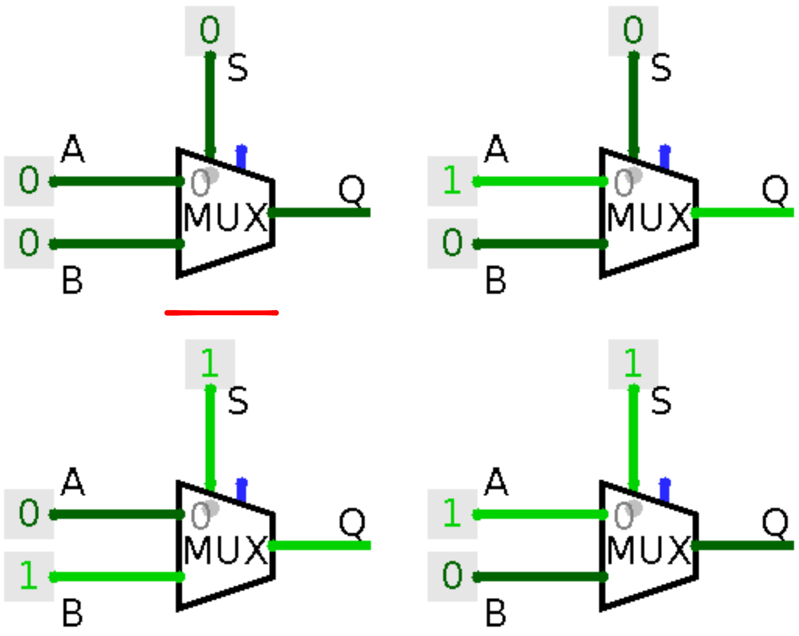
- We can build an adder **or** a subtractor.
- They are **very similar**.
- Can we build one unit that does **both**?

- Almost…

Adder-subtractor



Selecting a signal



S *(Select)*	A	B	Q
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

Selecting a signal

- Q = (A ∧ ¬S) ∨ (B ∧ S)
- Consider S = 0
- Consider S = 1

S	A	B	Q
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

Selecting a signal

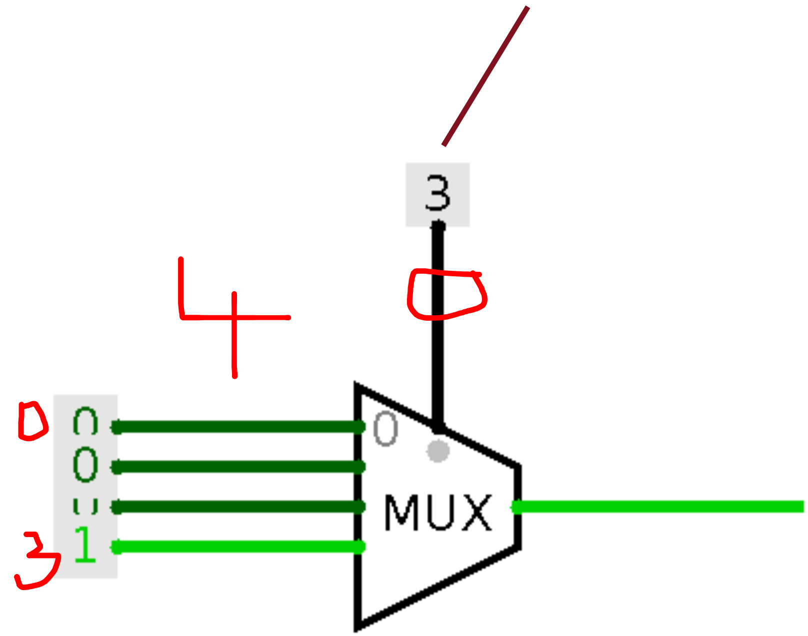
- $Q = (A \land \neg \underline{S}) \lor (B \land \underline{S})$
- Consider $\underline{S = 0}$
  - $Q = (A \land \underline{1}) \lor (B \land \underline{0})$
  - $Q = A \lor 0$
  - $\underline{Q = A}$
- Consider $\underline{S = 1}$
  - $\underline{Q} = (A \land 0) \lor (B \land 1)$
  - $Q = 0 \lor B$
  - $\underline{Q = B}$

S	A	B	Q
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

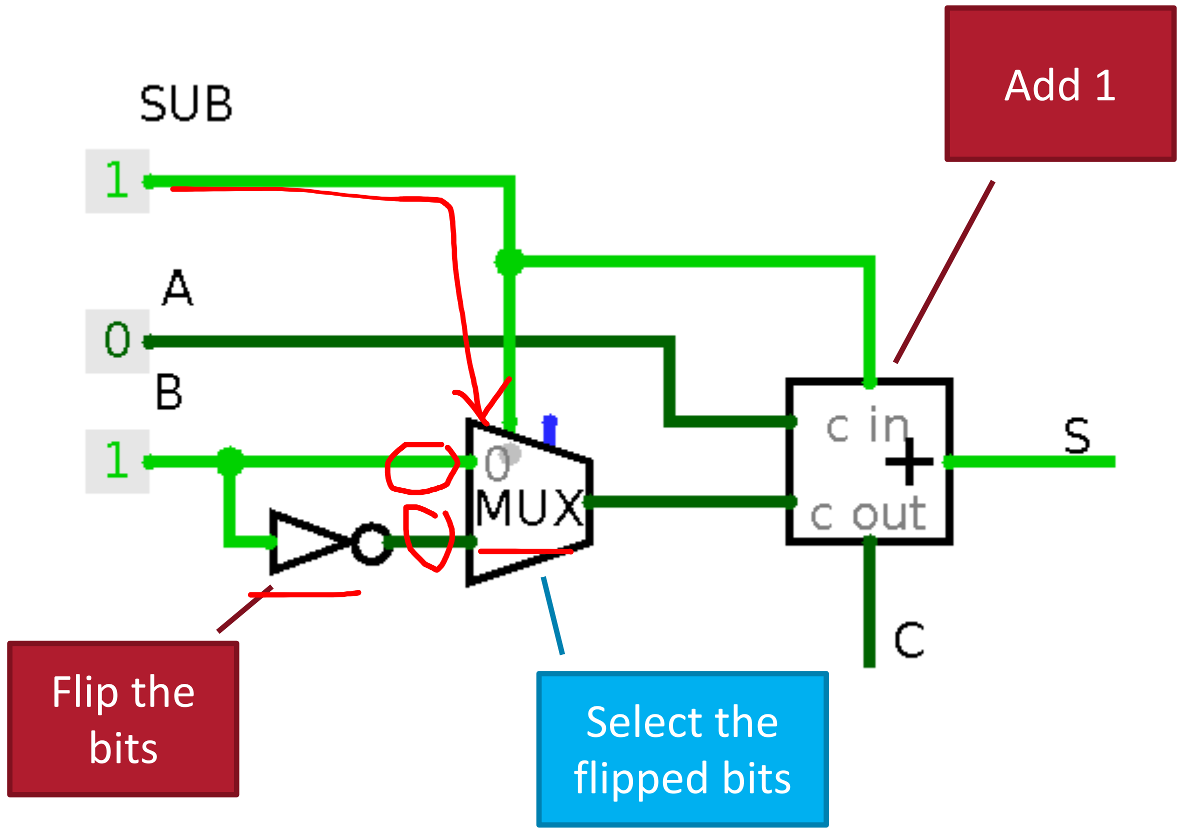
The multiplexer

- S selects which input to propagate to the output.
- 2-to-1 multiplexer
  - 1 select bit
- **N**-to-1 multiplexers are also possible
  - $\log_2(N)$ select bits needed
    - Why?

2-bit signal (imagine a bundle of two wires)

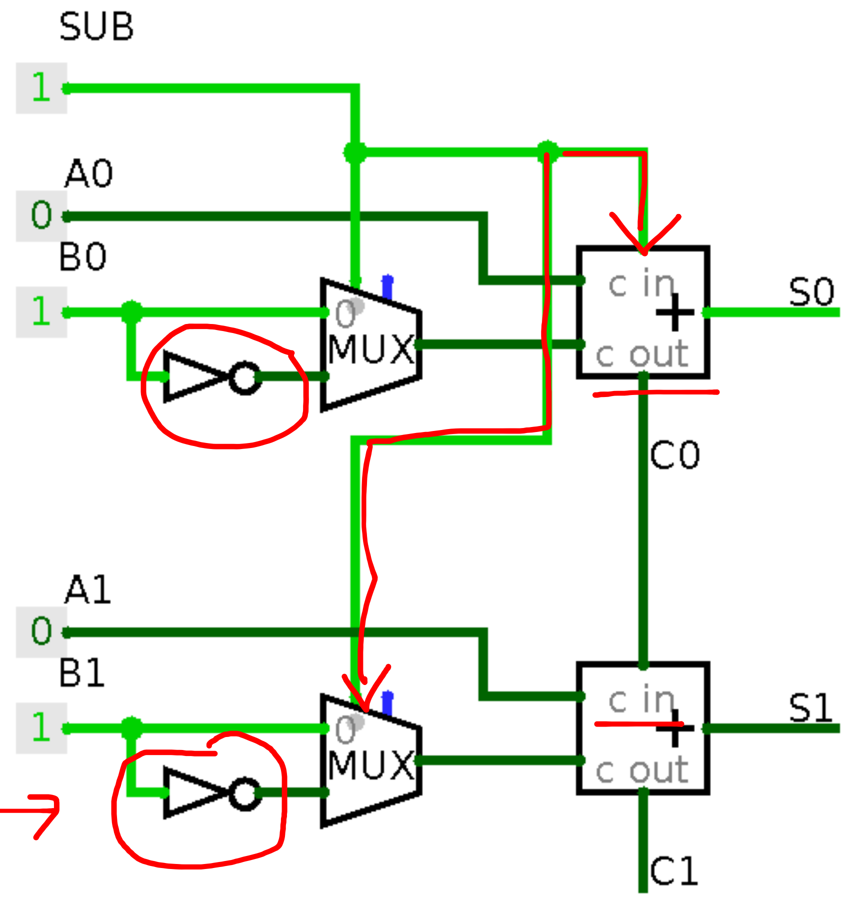


Adder-subtractor



Ripple-carry adder-subtractor

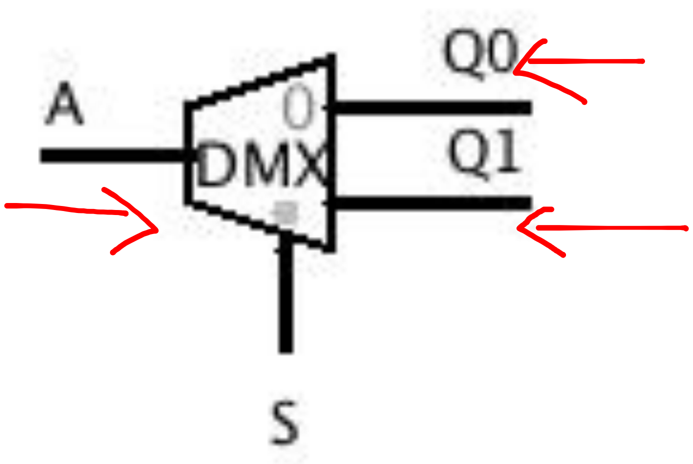
Only a one is needed at the beginning



Multi bits: like 8-bit 64-bit

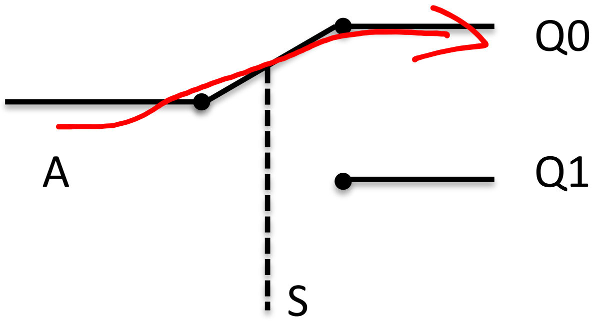
Attention: If we want to add a + b + c, calculate the first item, then add the result and c

Demultiplexer



- 1-to-2 demultiplexer
  - Choose **which of 2 wires** to **propagate** the input signal onto.
  - Q0 = A ∧ ¬S
  - Q1 = A ∧ S
- Instead of choosing which signal to select, we choose **where to send** a signal to, i.e. which output carries the input.

Demultiplexer



The demultiplexer acts like a switch.

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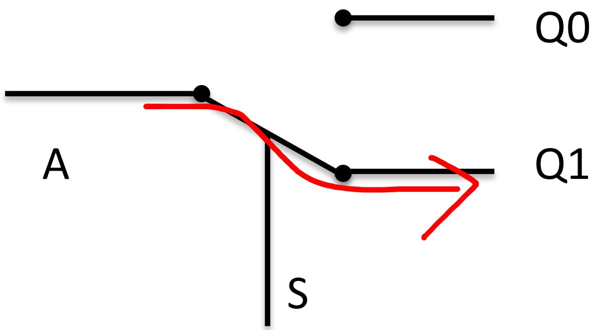

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Demultiplexer

S	A	Q0	Q1
0	**0**	**0**	0
0	**1**	**1**	0
1	0	0	0
1	1	0	1

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- Instead of choosing which signal to select, we choose **where to send** a signal to, i.e. which output carries the input.

Summary

- Used Boolean algebra to build **four simple devices**:
  - Demux, Mux, Adder, Subtractor.
- Combined Mux, NOT gate and adder to build **adder-subtractor**.
- We can do **basic arithmetic** with a **bunch of NAND gates**!

Imagine if we could **store** the results of that arithmetic, somehow…

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In the next lecture

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