

# COMSM1302

## Overview of Computer Architecture

### Lecture 15

### Performance Considerations in Assembly Programming

# In the previous lecture



- Load and store instructions
  - Single register data transfer (LDR / STR).
  - Block data transfer (LDM/STM).
- Pre- and post- addressing modes
- Direct and sequential access of array elements
- Copy data blocks with Block data transfer instructions
- Stack

# In this lecture

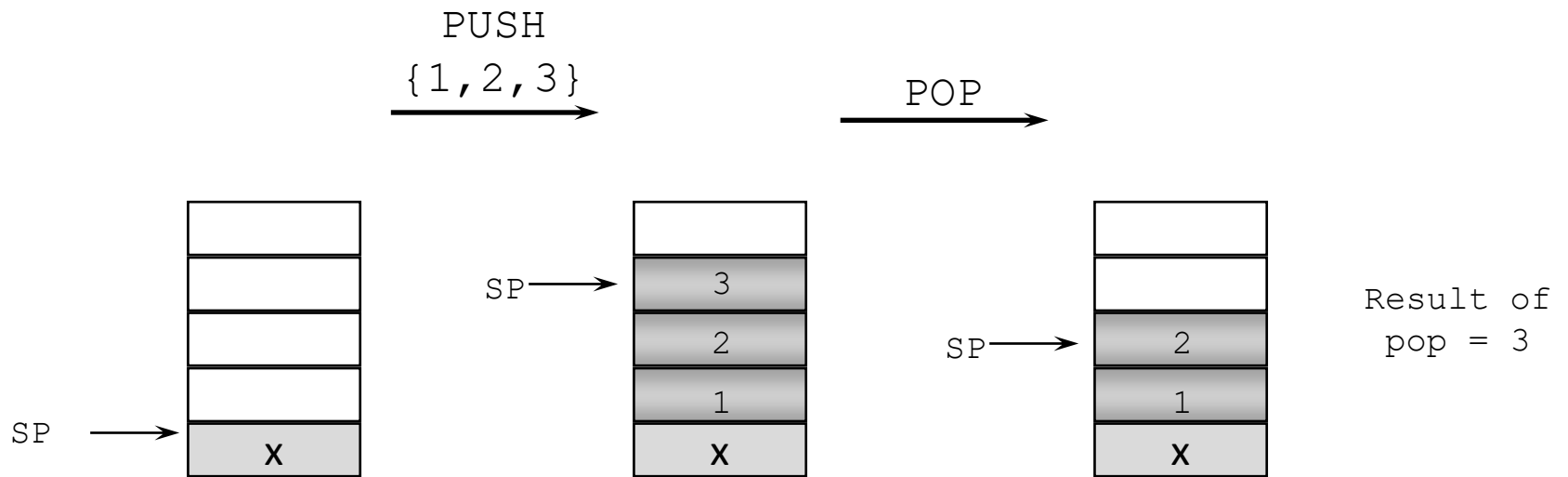


- At the end of this lecture:
  - Understand how the stack help us in tracking function calls.
  - Learn about programs performance measures.
  - Learn how to write efficient assembly programs.

# Assembly programming performance



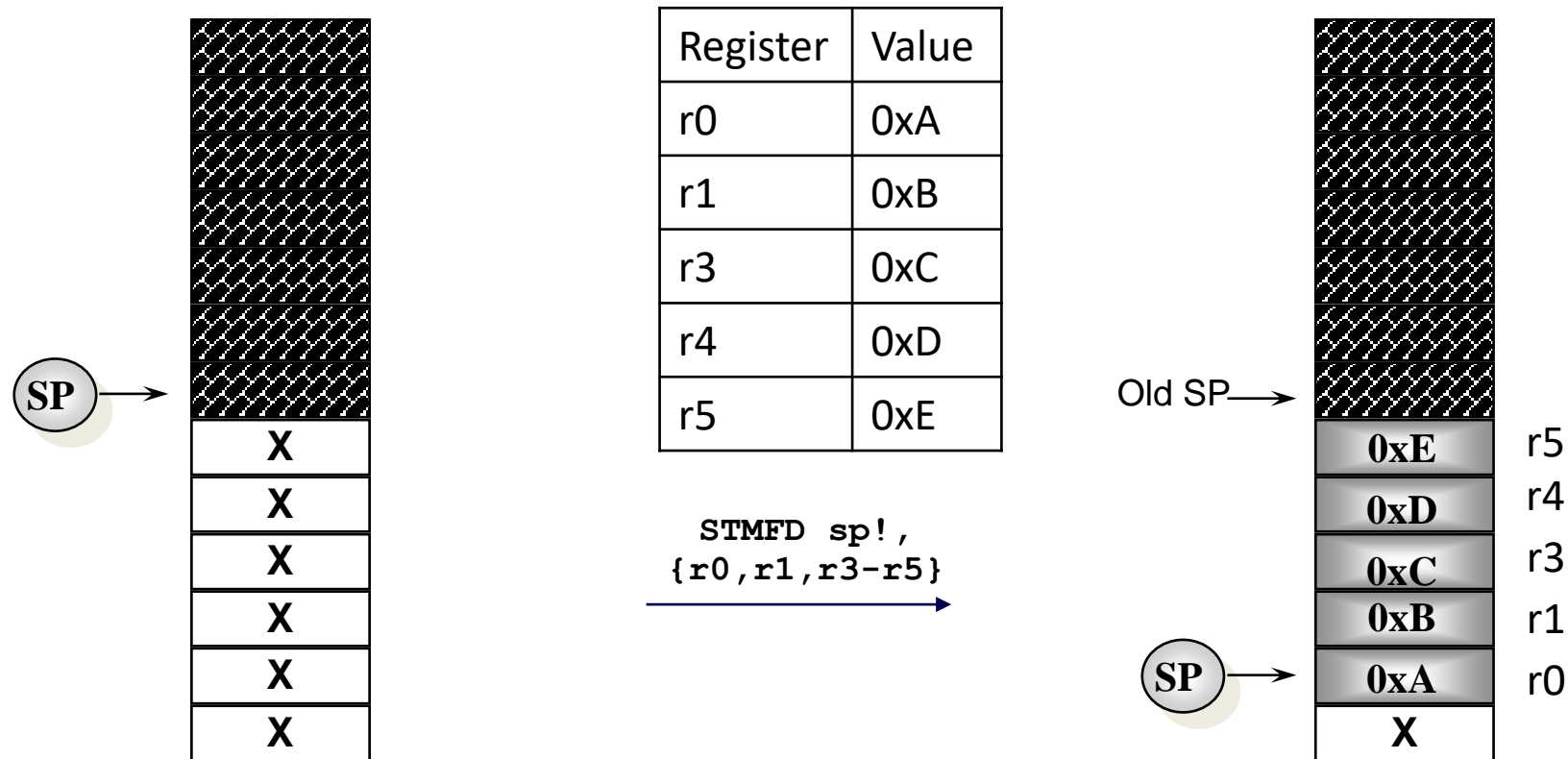
# Stack



# Stacks in ARM

- The stack type to be used is given by the postfix to the instruction:
  - STMFD / LDMFD : Full Descending stack
  - STMFA / LDMFA : Full Ascending stack.
  - STMED / LD MED : Empty Descending stack
  - STMEA / LDMEA : Empty Ascending stack

# 🔥 Full Descending stack – 1/3



# 🔥 Full Descending stack – 2/3

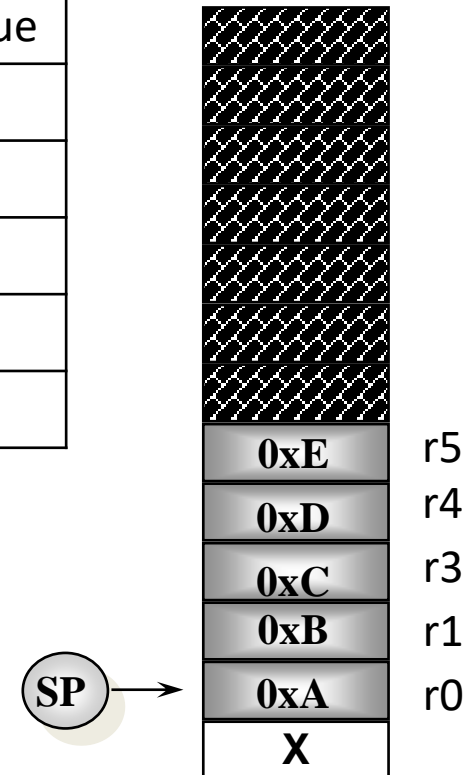


Register	Value
r0	0xA
r1	0xB
r3	0xC
r4	0xD
r5	0xE

Some assembly  
code

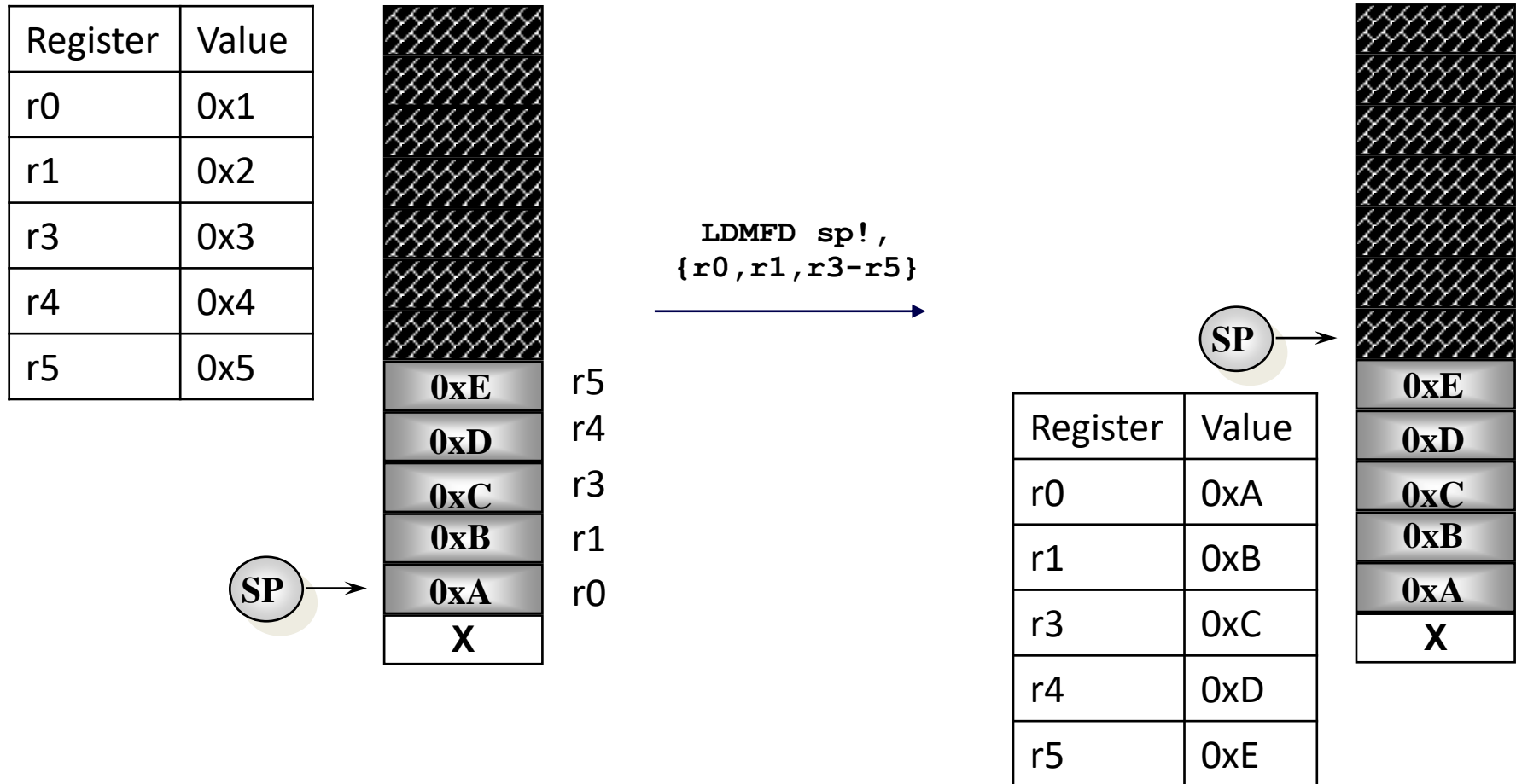


Register	Value
r0	0x1
r1	0x2
r3	0x3
r4	0x4
r5	0x5





# Full Descending stack – 3/3



# ARM Stack Implementations

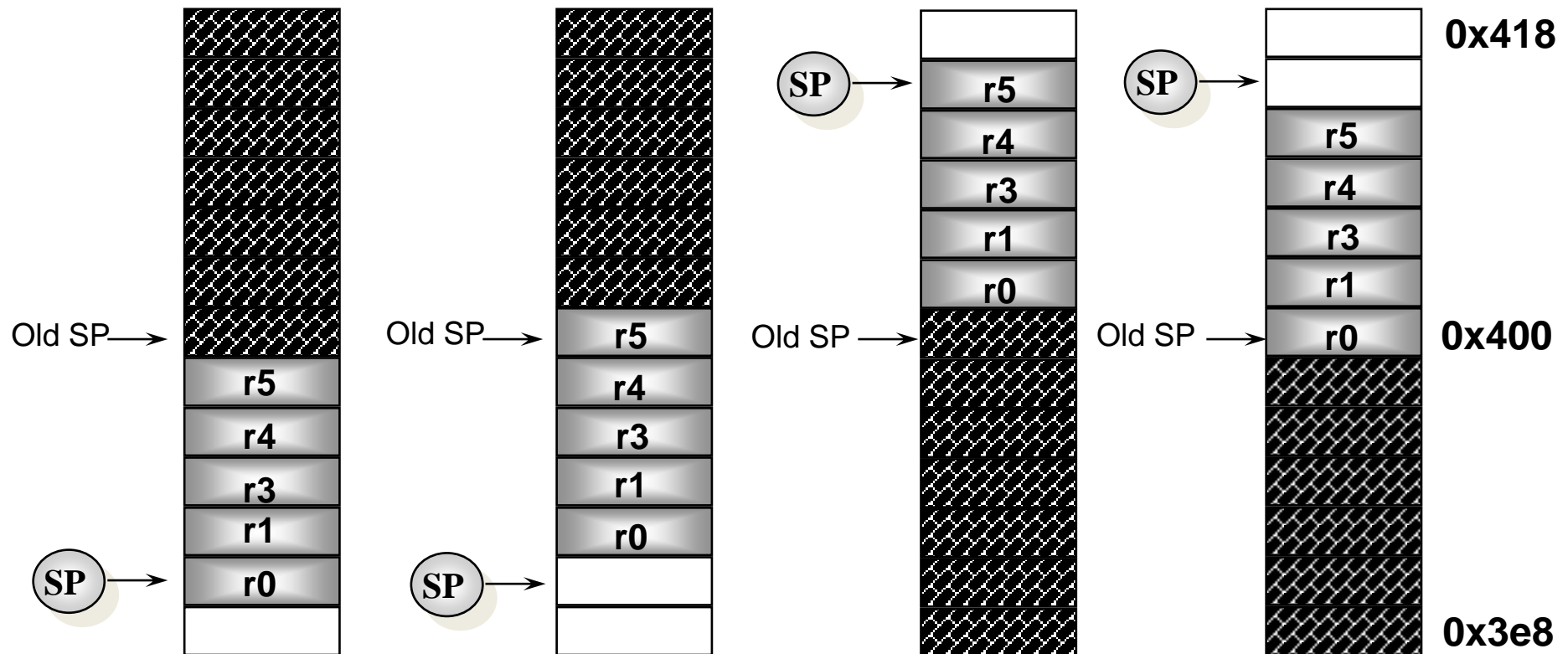


STMFD sp!,  
{r0,r1,r3-r5}

STMED sp!,  
{r0,r1,r3-r5}

STMFA sp!,  
{r0,r1,r3-r5}

STMEA sp!,  
{r0,r1,r3-r5}



# Example : factorial

- Factorial of a non-negative integer, is multiplication of all integers smaller than or equal to  $n$ . For example factorial of 4 is  $4*3*2*1$  which is 24.
- $n! = n*(n-1)*(n-2)*(n-3)*...*1$
- $4! = 4*(4-1)*(4-2)*(4-3) = 24$

# 🔥 Factorial – recursive solution

- $n! = n * (n-1) * (n-2) * (n-3) * \dots * 1$
- $(n-1)! = (n-1) * ((n-1)-1) * ((n-1)-2) * \dots * 1$
- $(n-1)! = (n-1) * (n-2) * (n-3) * \dots * 1$
- $n! = n * (n-1)!$
- $n! = 1$  if  $n = 0$  or  $n = 1$

# Factorial – recursive solution – C code

```
// function to find factorial of given number
unsigned int factorial(unsigned int n)
{
    if (n == 0)
        return 1;
    return n * factorial(n - 1);
}
```

# Factorial – recursive solution – assembly code – 1/6

```
Ldr r0, =#3  
bl _factorial  
b _end
```

```
_factorial:  
    cmp r0, #0  
    moveq r0, #1          if (n == 0)  
    moveq pc, lr          return 1;
```

# Factorial – recursive solution – assembly code – 2/6

```
Ldr r0, =#3
bl _factorial
b _end
```

```
_factorial:
```

```
    cmp r0, #0
    moveq r0, #1          if (n == 0)
                           return 1;
```

```
    moveq pc, lr
```

```
    mov r1, r0           Save n and calculate n-1 then call factorial for n-1
```

```
    sub r0, r0, #1
```

```
    stmfd sp!, {r1, lr}
```

```
    bl _factorial
```

# Factorial – recursive solution – assembly code – 3/6



```
Ldr r0, =#3
bl _factorial
b _end                <- 0x2

_factorial:
    cmp r0, #0
    moveq r0, #1
    moveq pc, lr
    mov r1, r0
    sub r0, r0, #1
    stmfd sp!, {r1, lr}
    bl _factorial

                                <- 0xA
```

0x2
3



# 🔥 Factorial – recursive solution – assembly code – 4/6



```
Ldr r0, =#3
bl _factorial
b _end                <- 0x2

_factorial:
    cmp r0, #0
    moveq r0, #1
    moveq pc, lr
    mov r1, r0
    sub r0, r0, #1
    stmfd sp!, {r1, lr}
    bl _factorial
                                <- 0xA
```

0x2
3
0xA
2

# Factorial – recursive solution – assembly code – 5/6



```
Ldr r0, =#3
bl _factorial
b _end                <- 0x2

_factorial:
    cmp r0, #0
    moveq r0, #1
    moveq pc, lr
    mov r1, r0
    sub r0, r0, #1
    stmfd sp!, {r1, lr}
    bl _factorial
                                <- 0xA
```

0x2
3
0xA
2
0xA
1

# Factorial – recursive solution – assembly code – 6/6



```
Ldr r0, =#3
bl _factorial
b _end                                <- 0x2

_factorial:
    cmp r0, #0
    moveq r0, #1
    moveq pc, lr
    mov r1, r0
    sub r0, r0, #1
    stmfd sp!, {r1, lr}
    bl _factorial
    lmdfd sp!, {r1, lr}    <- 0xA
    mul r0, r1, r0
    mov pc, lr
```

0x2
3
0xA
2
0xA
1

# Assembly programming performance



# Program performance measures

- Program Execution Time
  - Worst-case execution time (WCET):
  - Best-case execution time (BCET)
  - Average-case execution time (ACET)
- Program Size: number of instructions in ARM
- Program Energy Consumption

# Assembly programming performance



# ARM assembly programming performance issues



- LDM /STM
- Conditional execution
- Using the barrel shifter
- Optimising register usage
- Loop unrolling
- Initiate a register with zero
- Addressing modes
- Multiplication / Division

# LDM / STM

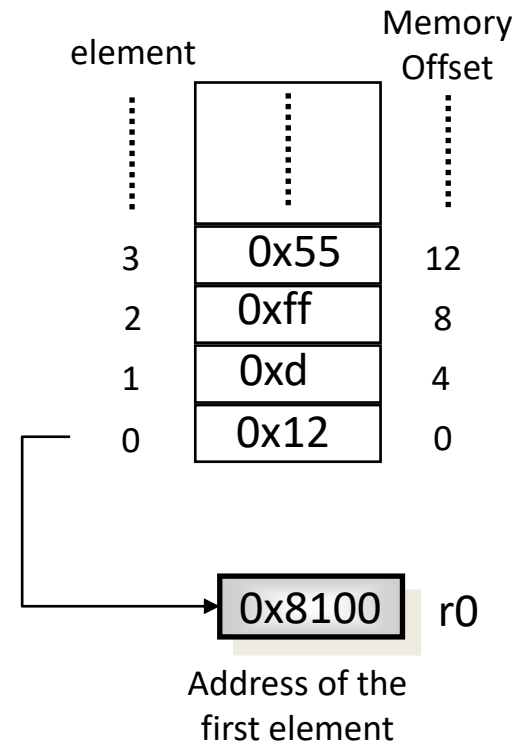
- Use LDM and STM instead of a sequence of LDR or STR instructions wherever possible.
  - The code is smaller.
  - An instruction fetch cycle can be saved for each eliminated LDR or STR.
  - Can turn non-sequential memory cycles into faster memory sequential cycles.



# 🔥 LDM / STDM - example

```
MOV r1, #4
MOV r3, #0
LDR r0, =_data
_loop:
LDR r2, [r0], #4
ADD r3, r3, r2
SUBS r1, r1, #1
BNE _loop
STR r3, [r0]
_end: B _end
```

```
MOV r3, #0
LDR r0, =_data
LDMIA r0!, {r3-r6}
ADD r3, r3, r4
ADD r3, r3, r5
ADD r3, r3, r6
STR r3, [r0]
_end: B _end
```



# Conditional execution

- Using conditionally executed instructions can avoid branches around short pieces of code.
  - This reduces the size of code

```
TST r0, #1
BLEQ _even:
_end: B _end
```

```
_even:
LSL r0, r0, #1
MOV r1, r0
MOV pc, lr
```

```
TST r0, #1
LSLEQ r0, r0, #1
MOVEQ r1, r0
_end: B _end
```

# Using the barrel shifter

- Combining shift operations with other operations can significantly increase the code density and thus performance.

```
TST  r0,#1
LSLEQ r0,r0,#1
MOVEQ r1,r0
_end: B _end
```

```
TST  r0,#1
MOVEQ r1,r0,LSL #1
_end: B _end
```

# Conditional execution and barrel shifter example

```
TST r0,#1
BLEQ _even:
_end: B _end
```

```
_even:
LSL r0,r0,#1
MOV r1,r0
MOV pc, lr
```

```
TST r0,#1
LSLEQ r0,r0,#1
MOVEQ r1,r0
_end: B _end
```

```
TST r0,#1
MOVEQ r1,r0,LSL #1
_end: B _end
```

# Optimising register usage

- Register spillage happens when we have more variables than the number of available registers.
  - a value has to be reloaded.
  - an intermediate value saved and then reloaded.

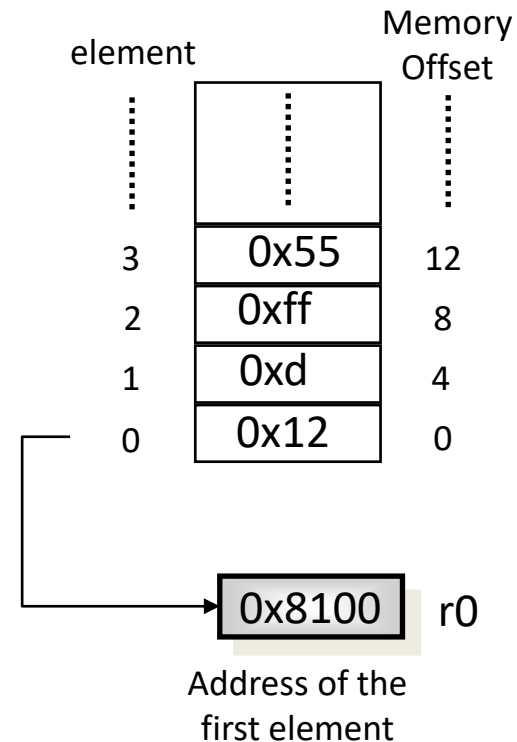
# Loop unrolling

- Loop unrolling involves using more than one copy of the inner loop of an algorithm
  - The branch back to the beginning of the loop is executed less frequently.
  - it may be possible to combine some of one iteration with some of the next iteration, and thereby significantly reduce the cost of each iteration.

# 🔥 Loop unrolling - example

```
MOV r1, #4
MOV r3, #0
LDR r0, =_data
_loop:
LDR r2, [r0], #4
ADD r3, r3, r2
SUBS r1, r1, #1
BNE _loop
STR r3, [r0]
_end: B _end
```

```
MOV r3, #0
LDR r0, =_data
LDMIA r0!, {r3-r6}
ADD r3, r3, r4
ADD r3, r3, r5
ADD r3, r3, r6
STR r3, [r0]
_end: B _end
```

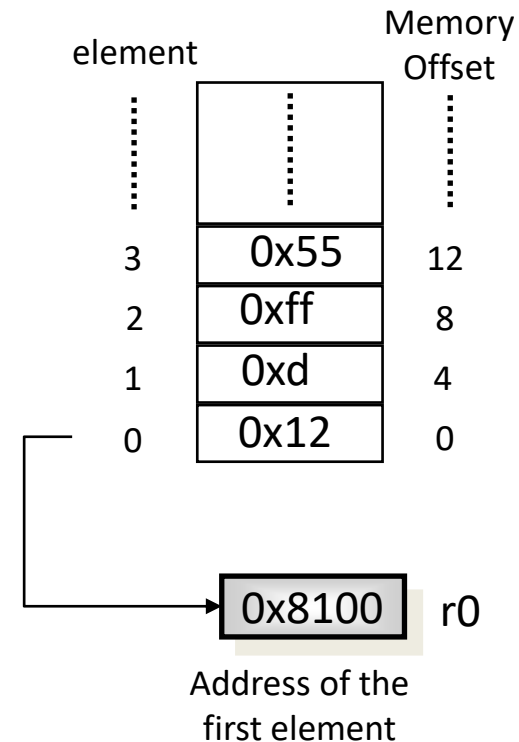


# 🔥 Initiate a register with zero

- What registers to reset and how?

```
MOV r3, #0
LDR r0,=_data
LDMIA r0!, {r3-r6}
ADD r3, r3, r4
ADD r3, r3, r5
ADD r3, r3, r6
STR r3, [r0]
_end: B _end
```

```
EOR r3, r3, r3
LDR r0,=_data
LDMIA r0!, {r3-r6}
ADD r3, r3, r4
ADD r3, r3, r5
ADD r3, r3, r6
STR r3, [r0]
_end: B _end
```



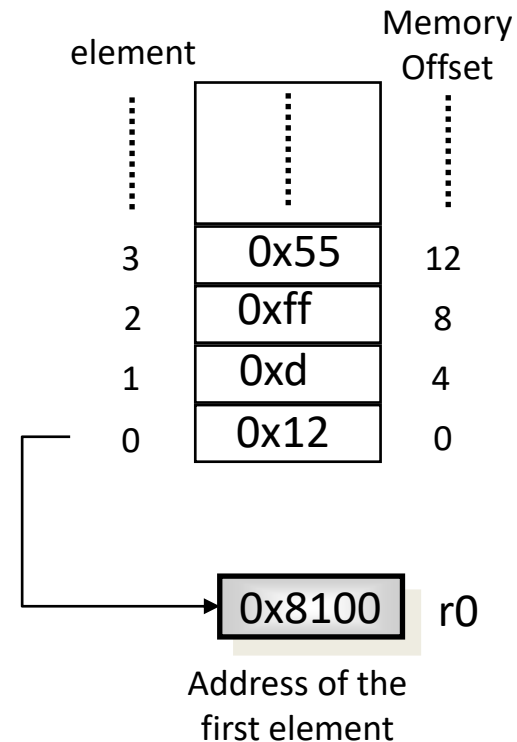


# 🔥 Initiate a register with zero

- What registers to reset and how?

```
MOV r3, #0
LDR r0,=_data
LDMIA r0!, {r3-r6}
ADD r3, r3, r4
ADD r3, r3, r5
ADD r3, r3, r6
STR r3, [r0]
_end: B _end
```

```
EOR r3, r3, r3
LDR r0,=_data
LDMIA r0!, {r3-r6}
ADD r3, r3, r4
ADD r3, r3, r5
ADD r3, r3, r6
STR r3, [r0]
_end: B _end
```



# Other performance issues

- **Addressing modes**
  - Using LDR or STR pre- or post-indexed with a non-zero offset increments the base register and performs the data transfer
- **Multiplication / Division**
  - Be aware of the time taken by the ARM multiply and division instructions

# Summary

- Stack operation
- Recursive solution to compute the factorial of a given number.
- Programs performance measures
- How to write efficient assembly programs.