COMSM1302 Overview of Computer Architecture

Lecture 14

ARM Memory Access Instructions



In the previous lecture

- Multiplication and division instructions.
- Shifting operations.
- Fixed point math.
- Example: Newton and Raphson's method.



In this lecture

LDR/STR

- Load constant and base register
- Addressing modes
- Array access



LDM/STM

- Block data transfer
- Stack

- At the end of this lecture:
 - Learn how to load big constant numbers.
 - How to load and store data to and from memory.
 - How stacks are implemented in ARM.



Load / Store instructions

- The ARM has three sets of instructions which interact with main memory.
 - Single register data transfer (LDR / STR).
 - Block data transfer (LDM/STM).
 - Single Data Swap (SWP).



ARM memory access instructions

LDR/STR

- Load constant and base register
- Addressing modes
- Array access



LDM/STM

- Block data transfer
- Stack



Single register data transfer

- Syntax:
 - -<LDR|STR>{<cond>} Rd, <address>
- Operation
 - Idr loads the content of the memory location at the given address to the register Rd.
 - srt saves the content of register Rd to the memory location at the given address.







- Load constant and base register
- Addressing modes
- Array access

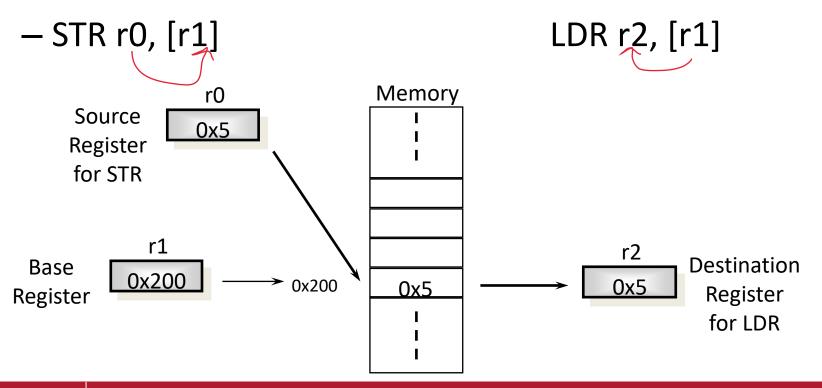


Loading full 32 bit constants

- LDR rd, = numeric constant
 - LDR r0,=0x42
 - Assembler generates MOV r0, #0x42
 - LDR r0,=0x5555555
 - Assembler generate LDR r0, [pc, offset to a literal pool]
- This mechanism will always generate the best instruction for a given case, thus, it is the recommended way of loading constants.

Load and Store : Base register

 The memory location to be accessed is held in a base register







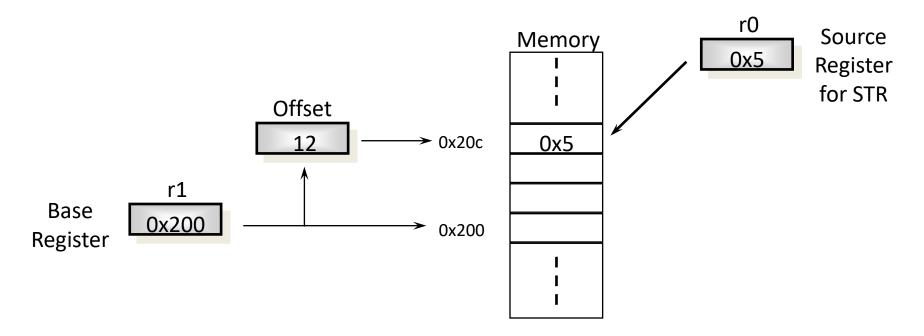


- Load constant and base register
- Addressing modes
- Array access



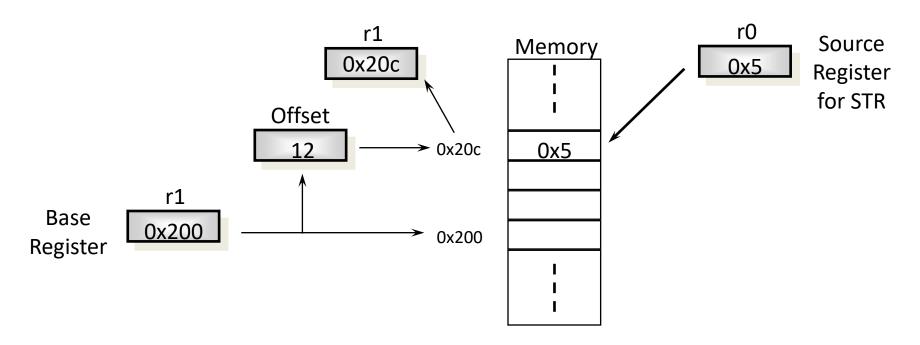
Load and Store – pre-indexed addressing

- Access a location offset from the base register
- Example: STR r0, [r1,#12]



Load and Store - pre-indexed addressing with write back

- Access a location offset from the base register
- Example: STR r0, [r1,#12]!





Load and Store - pre-indexed addressing - questions

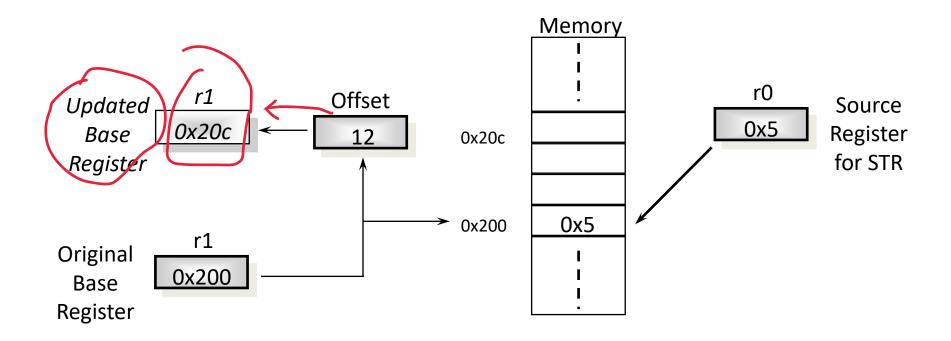
 If r1 = 0x200, what is the address where r0 will be stored, and what is the value of r1 after executing the following instructions?

- 1. STR r0, [r1,#-12] O_x | f + O_x 200
- 3. If r2 contains 12, STR r0, [r1, r2] 0_{\times} 20_{\circ}
- 4. If r2 contains 3, STR r0, [r1, r2, LSL #2]



Load and Store - post-indexed addressing

• Example: STR r0, [r1], #12



Load and Store - post-indexed addressing - questions

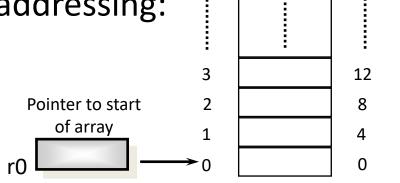
- If r1 = 0x200, what is the address where r0 will be stored, and what is the value of r1 after executing the following instructions?

 - STR r0, [r1], #-12 $O_x 200$, $Y_1 = O_x 1f +$ If r2 contains 12, STR r0, [r1], r2 $O_x 200$, $Y_1 = O_x 200$
 - If r2 contains 3, STR r0, [r1], r2, LSL #2 0_{κ}^{200} ,
- Why we do not have "STR r0, [r1], #12!"?

Because we already saved the result in r1

Usage of addressing modes

- If r0 points to the first element of an array in the memory.
- How can we access a particular element?
- Let r1 contains the index of the required element, and r2 be the destination register.
- Then we can use pre-indexed addressing:
 - LDR r2, [r0, r1, LSL #2]



Memory

Offset



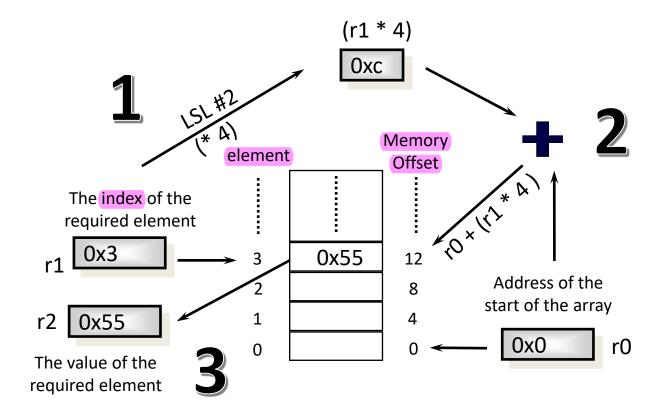


- Load constant and base register
- Addressing modes
- Array access



Usage of addressing modes – direct access

• LDR r2, [r0, r1, LSL #2] : r2 = [r0 + (r1 * 4)]





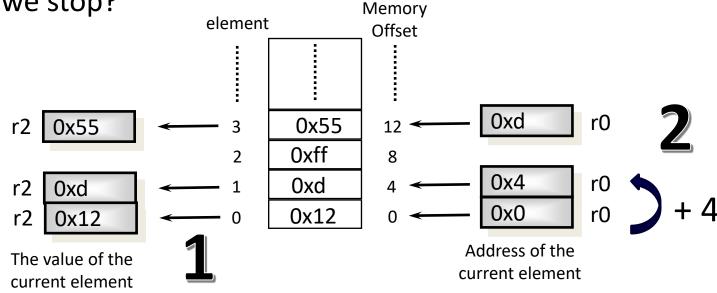
Usage of addressing modes – sequential access

- LDR r2, [r0], #4:
 - r2 = [r0]
 - r0 = r0 + 4

No

Do we still have the address of the first item at the end?

When do we stop?

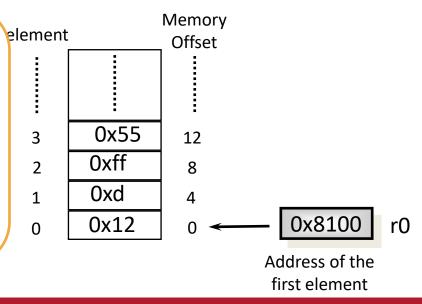




Load and store example

- Assume an array of 4 items is saved in the memory.
 The first element is in the address 0x8100.
- Write an assembly code to sum the elements of this array and save the sum in the address 0x8110.

The sum should be stored just after the last item. The address of this location is 0x8110 i.e. 4 words away from the first address. That is 16 bytes away; this why we added 0x1 to the base address 0x8100.





Load and store example - loop

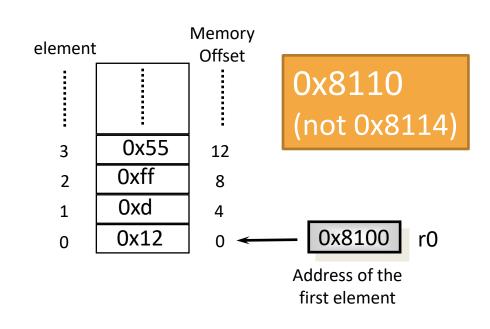
Design a loop to iterate four times.

```
MOV r1,#4
_loop:
@ Code in here will be repeated four times
SUBS r1,r1,#1
BNE _loop
_end: B _end
```



Load and store example – data access

```
MOV r1,#4
MOV r3, #0
LDR r0, = data
loop:
LDR r2, [r0], #4
ADD r3, r3, r2
SUBS r1, r1, #1
BNE loop
STR r3, [r0]
end: B end
```



Load and store example - code

```
.section .text
.align 2
.global _start
start:
MOV r1,#4
MOV r3, #0
LDR r0, = data
loop:
LDR r2, [r0], #4
ADD r3, r3, r2
SUBS r1, r1, #1
```

```
BNE _loop
STR r3, [r0]
_end: B _end

.section .data

data: .word 0x12,0xd,0xff,0x55
```

ARM memory access instructions

LDR/STR

- Load constant and base register
- Addressing modes
- Array access



LDM/STM

- Block data transfer
- Stack



Multiple registers data transfer

- Syntax:
 - -<LDM|STM>{<cond>} Rd, <address>
- Operation

 LDM / STM allow between 1 and 16 register to be transferred to or from memory.







- Block data transfer
- Stack



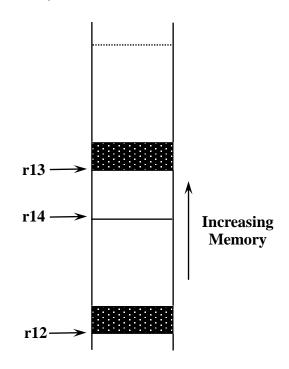
Direct functionality of block data transfer

- STMIA / LDMIA : Increment After
- STMIB / LDMIB : Increment Before
- STMDA / LDMDA : Decrement After
- STMDB / LDMDB : Decrement Before



Example: block copy – 1/2

Copy a block of memory, which is an exact multiple of 12 words long from the location pointed to by r12 to the location pointed to by r13. r14 points to the end of block to be copied.





Example: block copy – 2/2

- r12 points to the start of the source data
- r14 points to the end of the source data
- r13 points to the start of the destination data

```
_loop: write back new location to r12

LDMIA r12!, {r0-r11} @ load 48 bytes

STMIA r13!, {r0-r11} @ and store them

CMP r12, r14 @ check for the end

BNE _loop: @ and loop until done
```





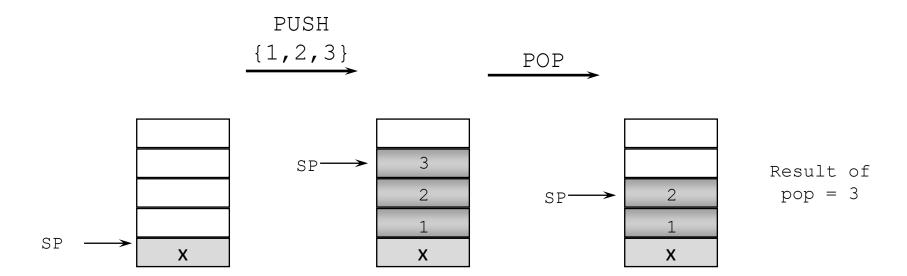


- Block data transfer
- Stack









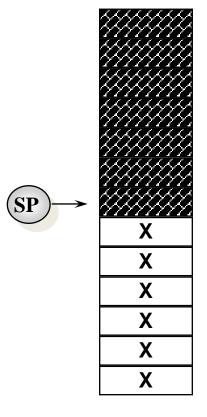
Stacks in ARM

- The stack type to be used is given by the postfix to the instruction:
 - STMFD / LDMFD : Full Descending stack
 - STMFA / LDMFA : Full Ascending stack. sp points to a value
 - STMED / LDMED : Empty Descending stack
 - STMEA / LDMEA : Empty Ascending stack

sp points to a empty address

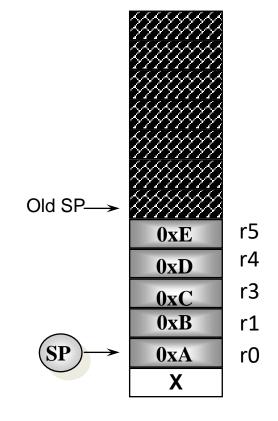


Full Descending stack – 1/3



Register	Value				
r0	0xA				
r1	ОхВ				
r3	0xC				
r4	0xD				
r5	0xE				
STMFD sp!					

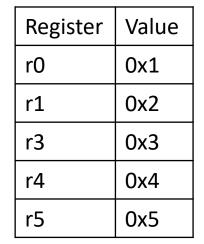
STMFD sp!, {r0,r1,r3-r5}

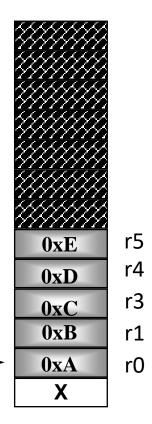


Full Descending stack − 2/3

Register	Value		
r0	0xA		
r1	0xB		
r3	0xC		
r4	0xD		
r5	0xE		

Some assembly code



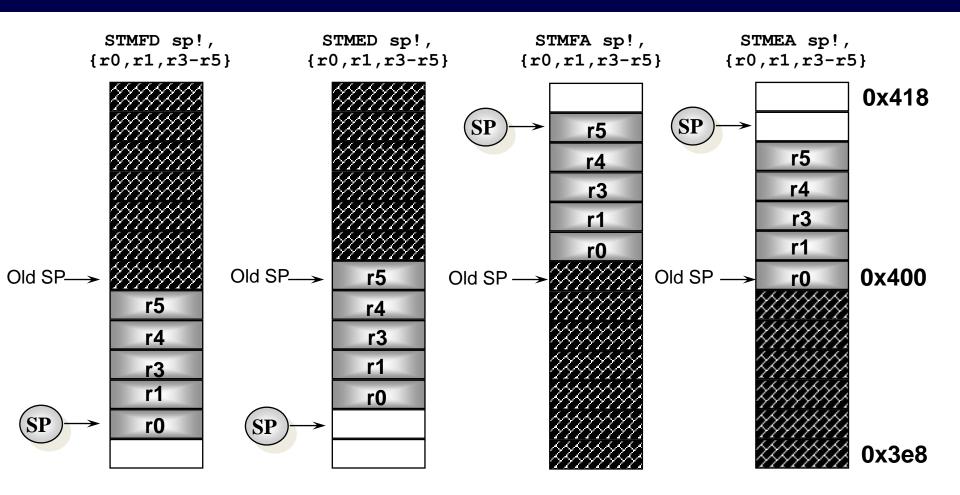


₩ Full Descending stack — 3/3

Register r0 r1 r3 r4	Value 0x1 0x2 0x3 0x4		LDMFD sp!, {r0,r1,r3-r5}		(SP)	
r5	0x5	0xE	r5			0xE
		0xD	r4	Register	Value	0xD
		0xC	r3	r0	0xA	0xC
	(SD)	OxB	r1	r1	0xB	0xB 0xA
	$(SP) \longrightarrow$	0xA X	r0 V	r3	0xC	X
				r4	0xD	
				r5	OxE	



ARM Stack Implementations



Stacks and subroutines

 One use of stacks is to create temporary register workspace for subroutines. Any registers that are needed can be pushed onto the stack at the start of the subroutine and popped off again at the end so as to restore them before return to the caller

```
STMFD sp!, {r0-r12, lr}

@ stack registers and the return address
......

LDMFD sp!, {r0-r12, pc}

@ load all the registers and return
automatically
```





• LDR/STR

- Load constant and base register
- Addressing modes
- Array access
- LDM/STM
 - Block data transfer
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