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Lindenbaum's lemma

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In this entry, we prove the following assertion known as Lindenbaum's lemma: every consistent set is a subset of a maximally consistent set. From this, we automatically conclude the existence of a maximally consistent set based on the fact that the empty set is consistent (vacuously).

Proposition 1. (Lindenbaum's lemma) Every consistent set can be extended to a maximally consistent set.

We give two proofs of this, one uses Zorn's lemma, and the other does not, and is based on the countability of the set of wff's in the logic. Both proofs require the following:

Lemma 1. The union of a chain of consistent sets, ordered by \subseteq , is consistent.

Proof. Let $C := \{\Gamma_i \mid i \in I\}$ be a chain of consistent sets ordered by \subseteq and indexed by some set I. We want to show that $\Gamma := \bigcup C$ is also consistent. Suppose not. Then $\Gamma \vdash \bot$. Let A_1, \ldots, A_n be a deduction of \bot (which is A_n) from Γ . Then for each $j = 1, \ldots, n-1$, there is some $\Gamma_{i(j)} \in C$ such that $\Gamma_{i(j)} \vdash A_i$. Since C is a chain, take the largest of the $\Gamma_{i(j)}$'s, say Γ_k , so that $\Gamma_k \vdash A_i$ for all $i = 1, \ldots, n-1$. This implies that $\Gamma_k \vdash A_n$, or $\Gamma_k \vdash \bot$, contradicting the assumption that Γ_k is consistent. As a result, Γ is consistent.

First Proof. Suppose Δ is consistent. Let P be the partially ordered set of all consistent supersets of Δ , ordered by inclusion \subseteq . If \mathcal{C} is a chain of elements in P, then $\bigcup \mathcal{C}$ is consistent by the lemma above, so $\bigcup \mathcal{C} \in P$ as each element of \mathcal{C} is a superset of Δ . By Zorn's lemma, P has a maximal element, call it Γ . To see that Γ is maximally consistent, suppose Γ is not maximal. Then there is a wff A such that $\Gamma \not\vdash A$ and $\Gamma \not\vdash \neg A$, the first of which implies that $A \notin \Gamma$, and the second of which implies that $\Gamma \cup \{A\}$ is consistent, and therefore in P. The two imply that $\Gamma \cup \{A\}$ is a consistent proper superset of Γ , contradicting the maximality of Γ in P. Therefore, Γ is maximally consistent.

Second Proof. Let A_1, \ldots, A_n, \ldots be an enumeration of all wff's of the logic in question (this can be achieved if the set of propositional variables can be enumerated). Let Δ be a consistent set of wff's. Define sets $\Gamma_1, \Gamma_2, \ldots, \Gamma_n$

of wff's inductively as follows:

$$\Gamma_{1} := \Delta$$

$$\Gamma_{n+1} := \begin{cases} \Gamma_{n} \cup \{A_{n}\} & \text{if } \Gamma_{n} \vdash A_{n} \\ \Gamma_{n} \cup \{\neg A_{n}\} & \text{otherwise} \end{cases}$$

$$\Gamma := \bigcup_{i=1}^{\infty} \Gamma_{i}.$$

First, notice that each Γ_i is consistent. This is done by induction on i. By assumption, the case is true when i=1. Now, suppose Γ_n is consistent. Then its deductive closure $\operatorname{Ded}(\Gamma_n)$ is also consistent. If $\Gamma_n \vdash A_n$, then clearly $\Gamma_n \cup \{A_n\}$ is consistent since it is a subset of $\operatorname{Ded}(\Gamma)$. Otherwise, $\Gamma_n \not\vdash A_n$, or $\Gamma_n \not\vdash \neg \neg A_n$ by the substitution theorem, and therefore $\Gamma_n \cup \{\neg A_n\}$ by one of the properties of consistency (see http://planetmath.org/PropertiesOfConsistencyhere). In either case, Γ_{n+1} is consistent.

Next, Γ is maximally consistent. Γ is consistent because, by the lemma above, it is the union of a chain of consistent sets. To see that Γ is maximal, pick any wff A. Then A is some A_n in the enumerated list of all wff's. Therefore, either $A \in \Gamma_{n+1}$ or $\neg A \in \Gamma_{n+1}$. Since $\Gamma_{n+1} \subseteq \Gamma$, we have $A \in \Gamma$ or $\neg A \in \Gamma$, which implies that Γ is maximal (see http://planetmath.org/MaximallyConsistenthere). \square .

Given a logic L, let W_L be the set of all maximally L-consistent sets. By Lindebaum's lemma, $W_L \neq \emptyset$. We record two useful corollaries:

- For any consistent set Δ , $\operatorname{Ded}(\Delta) = \bigcap \{ \Gamma \in W_L \mid \Delta \subseteq \Gamma \}$.
- $L = \bigcap W_L$.

Proof. The second statement is a corollary of the first, for $L = \text{Ded}(\emptyset)$. To see the first, let $\mathcal{D} := \{\Gamma \in W_L \mid \Delta \subseteq \Gamma\}$. Then $\mathcal{D} \neq \emptyset$ by Lindebaum's lemma. Also, for any $\Gamma \in \mathcal{D}$, $\text{Ded}(\Delta) \subseteq \Gamma$ since Γ is deductively closed. On the other hand if $A \notin \text{Ded}(\Delta)$, then $\Delta \not\vdash A$, so $\Delta \cup \{\neg A\}$ is consistent, and therefore is contained in a maximally consistent set $\Gamma' \in \mathcal{D}$ by Lindenbaum's lemma. Since $\neg A \in \Gamma'$, $A \notin \Gamma'$, so that $A \notin \bigcap \mathcal{D}$.