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categorical pullback

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Pullbacks

Let $f : X \rightarrow B$ and $g : Y \rightarrow B$ be morphisms in a category \mathcal{C} . Then a *pullback diagram*, or *pullback square* of f and g is a commutative diagram

$$\begin{array}{ccc} A & \xrightarrow{q} & Y \\ p \downarrow & & \downarrow g \\ X & \xrightarrow{f} & B. \end{array}$$

such that if we have another commutative diagram

$$\begin{array}{ccc} Z & \xrightarrow{s} & Y \\ r \downarrow & & \downarrow g \\ X & \xrightarrow{f} & B. \end{array}$$

then there is a unique morphism $h : Z \rightarrow A$ with the commutative diagram

$$\begin{array}{ccccc} Z & & & & \\ & \searrow s & & & \\ & & A & \xrightarrow{q} & Y \\ & \searrow h & \downarrow p & & \downarrow g \\ & & X & \xrightarrow{f} & B. \end{array}$$

A *pullback* of (f, g) is the ordered triple (A, q, p) . We also say that p is a *pullback of g along f* , and q a *pullback of f along g* . When a pullback of (f, g) exists, it is unique up to isomorphism. The object A , and the morphisms p and q in the diagram above are often denoted by

$$X \times_B Y, \quad 1_X \times_B g \quad \text{and} \quad f \times_B 1_Y$$

respectively, and the uniquely determined morphism h by

$$\begin{pmatrix} r \\ s \end{pmatrix}.$$

It is easy to see that

$$X \times_B Y \cong Y \times_B X,$$

whenever one (and hence the other) exists.

Remarks.

- The pullback of f and g can be equivalently defined as a limiting cone over the diagram $X \rightarrow B \leftarrow Y$. In other words, a pullback diagram is a terminal object in the category of commutative squares of the form

$$\begin{array}{ccc} Z & \longrightarrow & Y \\ \downarrow & & \downarrow g \\ X & \xrightarrow{f} & B. \end{array}$$

- A category \mathcal{C} is said to *have pullbacks* if every diagram $X \rightarrow B \leftarrow Y$ can be completed into a pullback diagram.
- The notion of pullbacks can be generalized: let $\{x_i : C_i \rightarrow C \mid i \in I\}$ be a collection of morphisms indexed by set I , considered as a small diagram. The *generalized pullback* of the x_i 's is just the limiting cone of the diagram. Using this definition, the generalized pullback of one morphism f is the identity morphism of $\text{dom}(f)$, the domain of f , and the generalized pullback of the empty set is a terminal object.
- A pullback is sometimes known as an *amalgamated sum*.

Pushouts

Dually, given morphisms $f : B \rightarrow X$ and $g : B \rightarrow Y$, a *pushout square* of f and g is a commutative diagram

$$\begin{array}{ccc} B & \xrightarrow{g} & Y \\ f \downarrow & & \downarrow s \\ X & \xrightarrow{t} & A \end{array}$$

such that if we have another commutative diagram

$$\begin{array}{ccc} B & \xrightarrow{g} & Y \\ f \downarrow & & \downarrow u \\ X & \xrightarrow{v} & Z \end{array}$$

Then there is a unique morphism $h : A \rightarrow Z$ with the following commutative diagram:

$$\begin{array}{ccc}
 B & \xrightarrow{g} & Y \\
 f \downarrow & & \downarrow s \\
 X & \xrightarrow{t} & A \\
 & \searrow v & \nearrow u \\
 & & Z
 \end{array}$$

(Note: A dotted arrow labeled h points from A to Z in the original diagram.)

The pair (s, t) of morphisms is a *pushout* of (f, g) . s is the *pushout of f along g* , and t the *pushout of g along f* . Like pullbacks, pushouts are unique up to unique isomorphism when they exist. The object A and the morphisms s and t are typically written

$$X \amalg_B Y, \quad f \amalg_B 1_Y \quad \text{and} \quad 1_X \amalg_B g,$$

and the unique morphism h is denoted by

$$(f \amalg g).$$

Remark. The pushout of f and g can be thought of as the limiting cocone under the diagram $X \leftarrow B \rightarrow Y$. Equivalently, they are initial objects in the category of commutative squares whose top edge is $B \rightarrow Y$ and left edge is $B \rightarrow X$. A category is said to *have pushouts* if every diagram $X \leftarrow B \rightarrow Y$ can be completed to a pushout diagram. The *generalized pushout* is defined as the limiting cocone under the diagram consisting of morphisms $y_i : B \rightarrow B_i$, where i belongs to some set I .