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source

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| Defines          | source              |
| Defines          | monosource          |
| Defines          | extremal monosource |
| Defines          | sink                |
| Defines          | extremal episink    |

In the whole entry we suppose we are given a category  $\mathbf{A}$ . By an object we always mean an object in  $\mathbf{A}$  and by a morphisms an  $\mathbf{A}$ -morphism.

**Definition 1.** A source in a category  $\mathbf{A}$  is a pair  $(A, (f_i)_{i \in I})$  where  $A$  is an object and  $f_i : A \rightarrow A_i$  are morphisms indexed by a class  $I$ .

The object  $A$  is called the domain of the source and the family  $(A_i)_{i \in I}$  is called the codomain of the source.

A sink is a pair  $((f_i)_{i \in I}, A)$  where  $A$  is an object and  $f_i : A_i \rightarrow A$  are morphisms.

Sources can be composed with morphisms. If  $\mathcal{S} = (A, (f_i)_{i \in I})$  is a source and  $f : B \rightarrow A$  is a morphism, we use the notation  $(B, (f_i \circ f)_{i \in I}) = \mathcal{S} \circ f$ . Similarly, for sinks, we use the notation  $f \circ \mathcal{S} = ((f \circ f_i)_{i \in I}, B)$  if  $\mathcal{S} = ((f_i)_{i \in I}, A)$  is a sink and  $f : A \rightarrow B$  is a morphism.

**Definition 2.** A source  $\mathcal{S} = (A, (f_i)_{i \in I})$  is called a monosource if for any pair  $r, s : B \rightarrow A$  of morphisms from the equality  $\mathcal{S} \circ r = \mathcal{S} \circ s$  follows  $r = s$ .

A sink  $\mathcal{S} = ((f_i)_{i \in I}, A)$  is called an episink if for any pair  $r, s : A \rightarrow B$  of morphisms  $r = s$  whenever  $r \circ \mathcal{S} = s \circ \mathcal{S}$ .

A monosource  $\mathcal{S}$  is called extremal monosource, if the following holds: Whenever  $\mathcal{S} = \overline{\mathcal{S}} \circ e$  for an epimorphism  $e$ , then  $e$  is an isomorphism.

An episink  $\mathcal{S}$  is called extremal episink if the following holds: Whenever  $\mathcal{S} = m \circ \overline{\mathcal{S}}$  pre for a monomorphism  $m$ , tak  $m$  is an isomorphism.

Every limit is an extremal monosource, a colimit is an extremal episink.

## References

- [1] J. Adámek, H. Herrlich, and G. Strecker. *Abstract and Concrete Categories*. Wiley, New York, 1990.