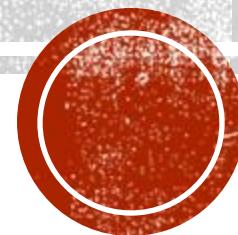


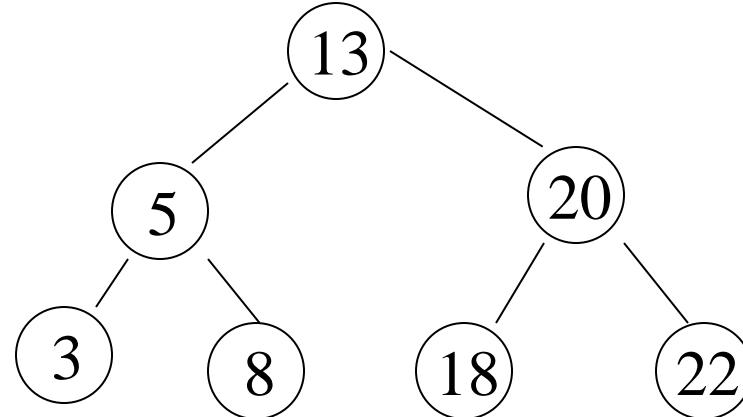
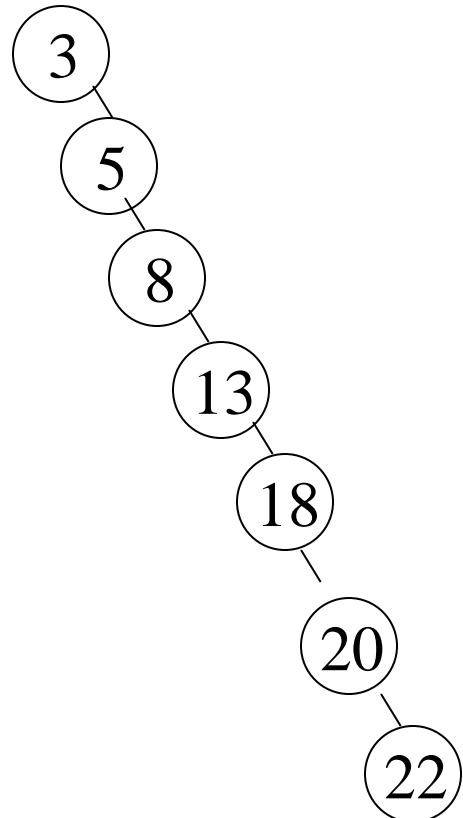
AVL TREES

Self Balancing Binary Search Tree



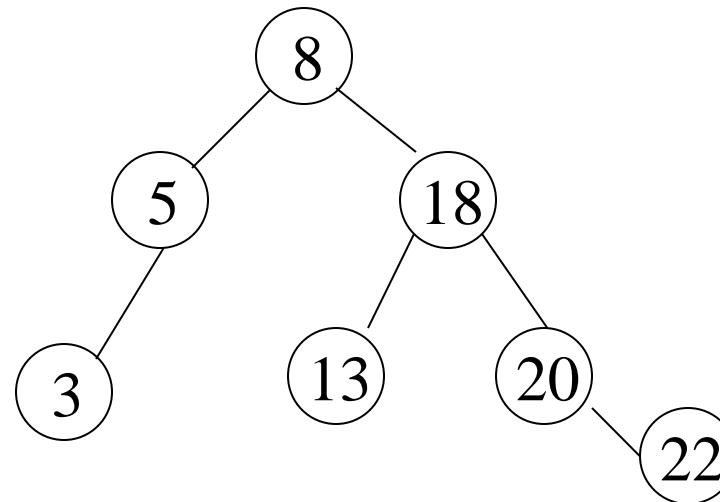
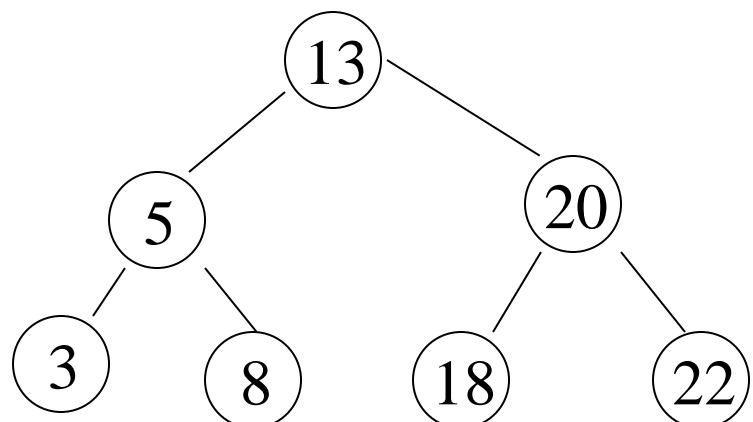
MOTIVATION

When building a binary search tree, what type of trees would we like? Example: 3, 5, 8, 20, 18, 13, 22



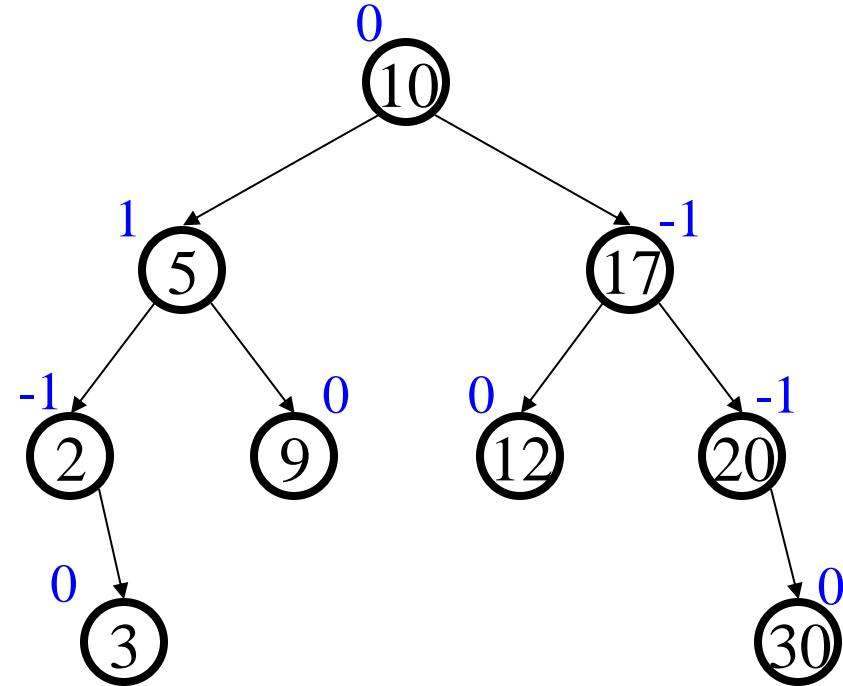
MOTIVATION

- We want a tree that has the following properties
 - Tree height = $O(\log(N))$
 - allows dynamic insert and remove with $O(\log(N))$ time complexity.
- The AVL tree is one of this kind of trees.



AVL (ADELSON-VELSKII AND LANDIS) TREES

- An AVL Tree is a ***binary search tree*** such that for every internal node v of T , the *heights of the children of v can differ by at most 1*.



An example of an AVL tree where the heights are shown next to the nodes:

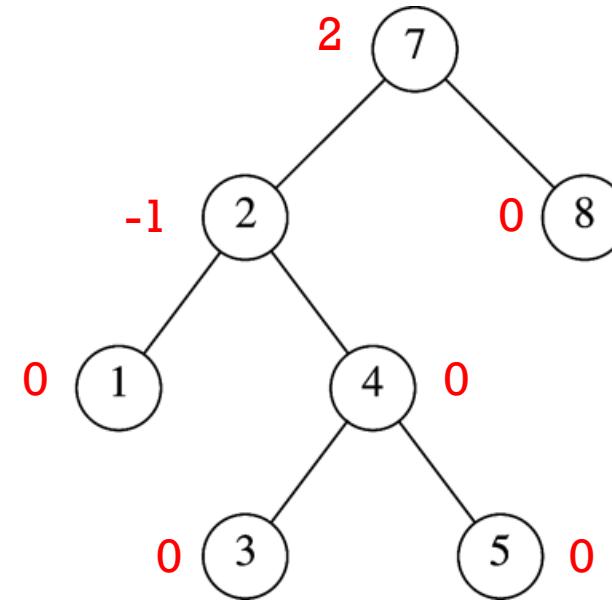
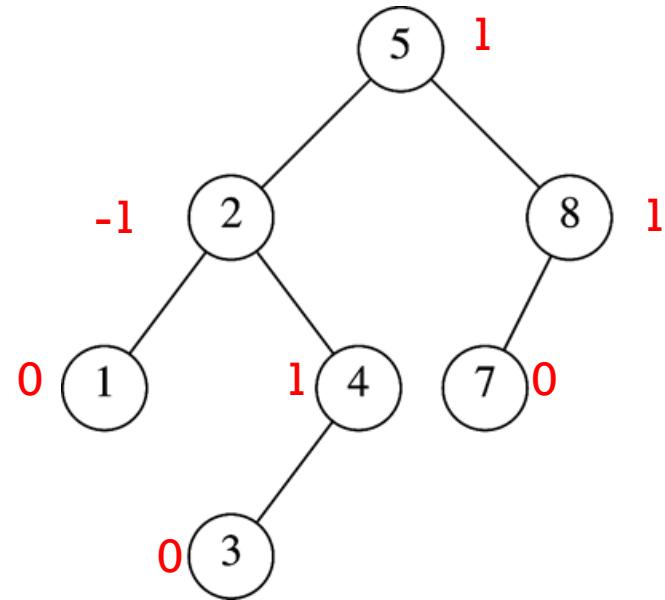
AVL TREE PROPERTIES

- AVL tree is a binary search tree with balance condition
 - To ensure depth of the tree is $O(\log(N))$
 - And consequently, search/insert/remove complexity bound $O(\log(N))$
- Balance condition
 - For **every node** in the tree, height of left and right subtree can differ by at most 1
 - Balance factor of a node = Height of its left subtree – Height of its right subtree
 - Balance factor of a node can be -1,0 or 1

BALANCE FACTOR OF A NODE IN AN AVL TREE

- Height of left subtree > height of right subtree
 - Balance Factor is 1
 - Tree is Left Heavy
- Height of left subtree < height of right subtree
 - Balance Factor is -1
 - Tree is Right Heavy
- Height of left subtree = height if right subtree
 - Balance Factor is 0
 - Tree is balanced

WHICH IS AN AVL TREE?



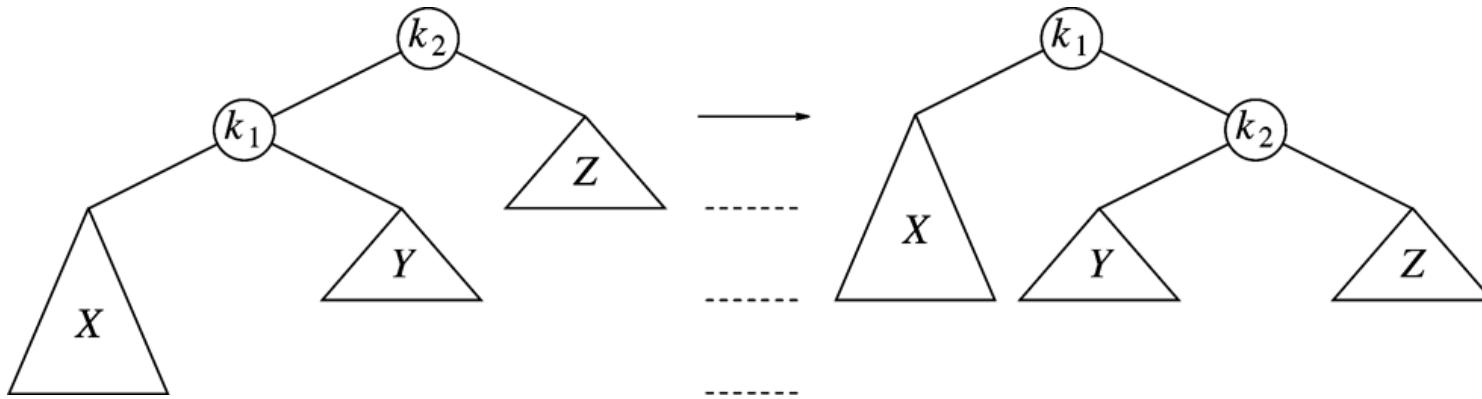
AVL TREE INSERT AND REMOVE

- Do binary search tree insert and remove
- The balance condition can be violated sometimes
 - Do something to fix it : rotations
 - After rotations, the balance of the whole tree is maintained

BALANCE CONDITION VIOLATION

- If condition violated after a node insertion
 - Which nodes do we need to rotate?
 - Only nodes on path from insertion point to root may have their balance altered
- Rebalance the tree through rotation at the **deepest node** with balance violated
 - The entire tree will be rebalanced
- Violation cases at node k (deepest node)
 1. An insertion into left subtree of left child of k
 2. An insertion into right subtree of left child of k
 3. An insertion into left subtree of right child of k
 4. An insertion into right subtree of right child of k
 - Cases 1 and 4 equivalent
 - **Single rotation to rebalance**
 - Cases 2 and 3 equivalent
 - **Double rotation to rebalance**

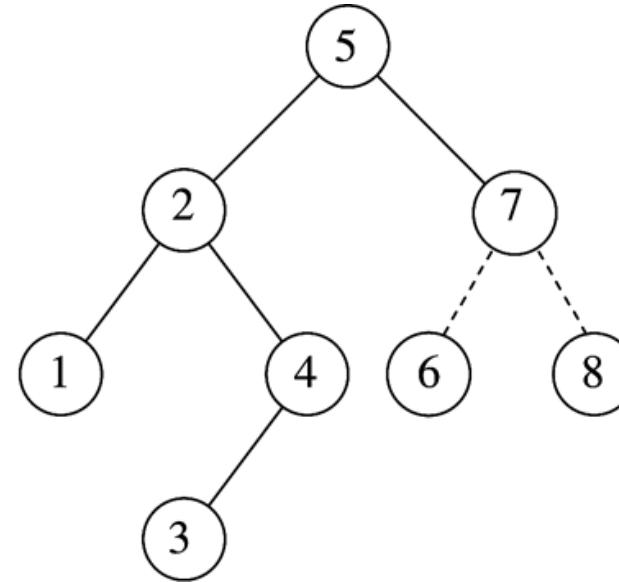
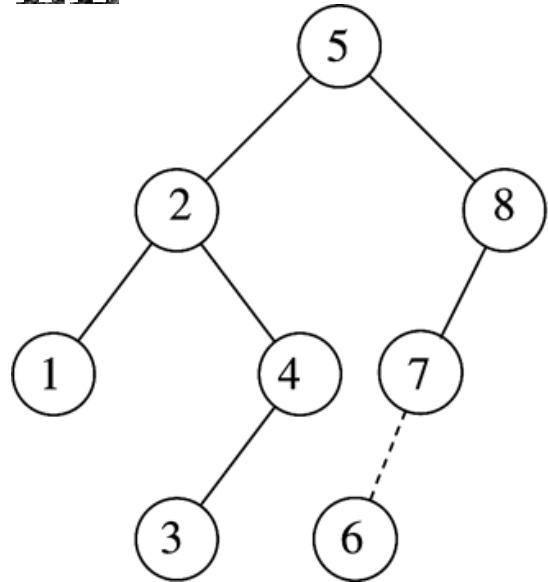
SINGLE RIGHT ROTATION



- **Single rotation:** The basic operation we'll use to rebalance
 - Move child of unbalanced node into parent position
 - Parent becomes the “other” child (always okay in a BST!)
 - Other subtrees move in only way BST allows (we'll see in generalized example)

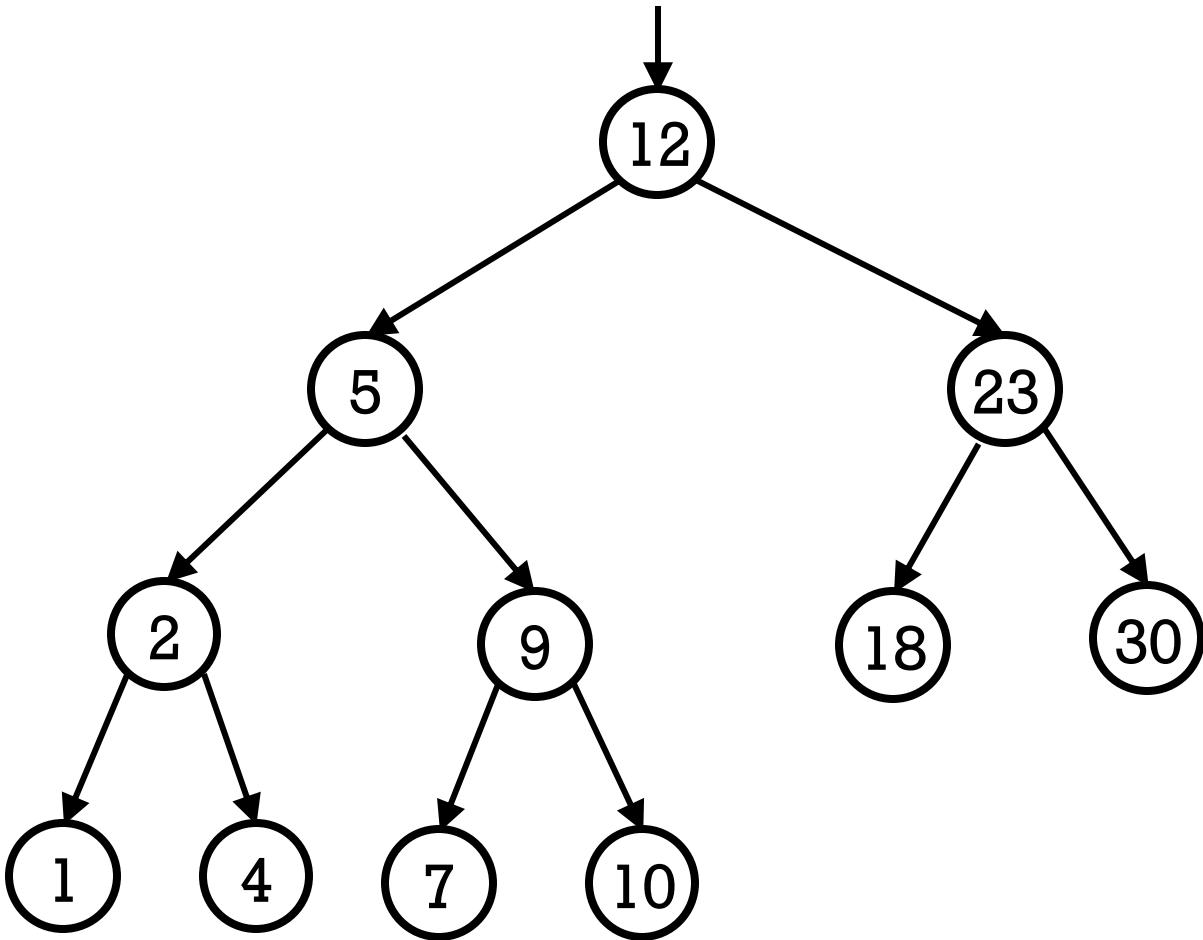
- Replace node k_2 by node k_1
- Set node k_2 to be right child of node k_1
- Set subtree Y to be left child of node k_2
- Case 4 is similar

EXAMPLE

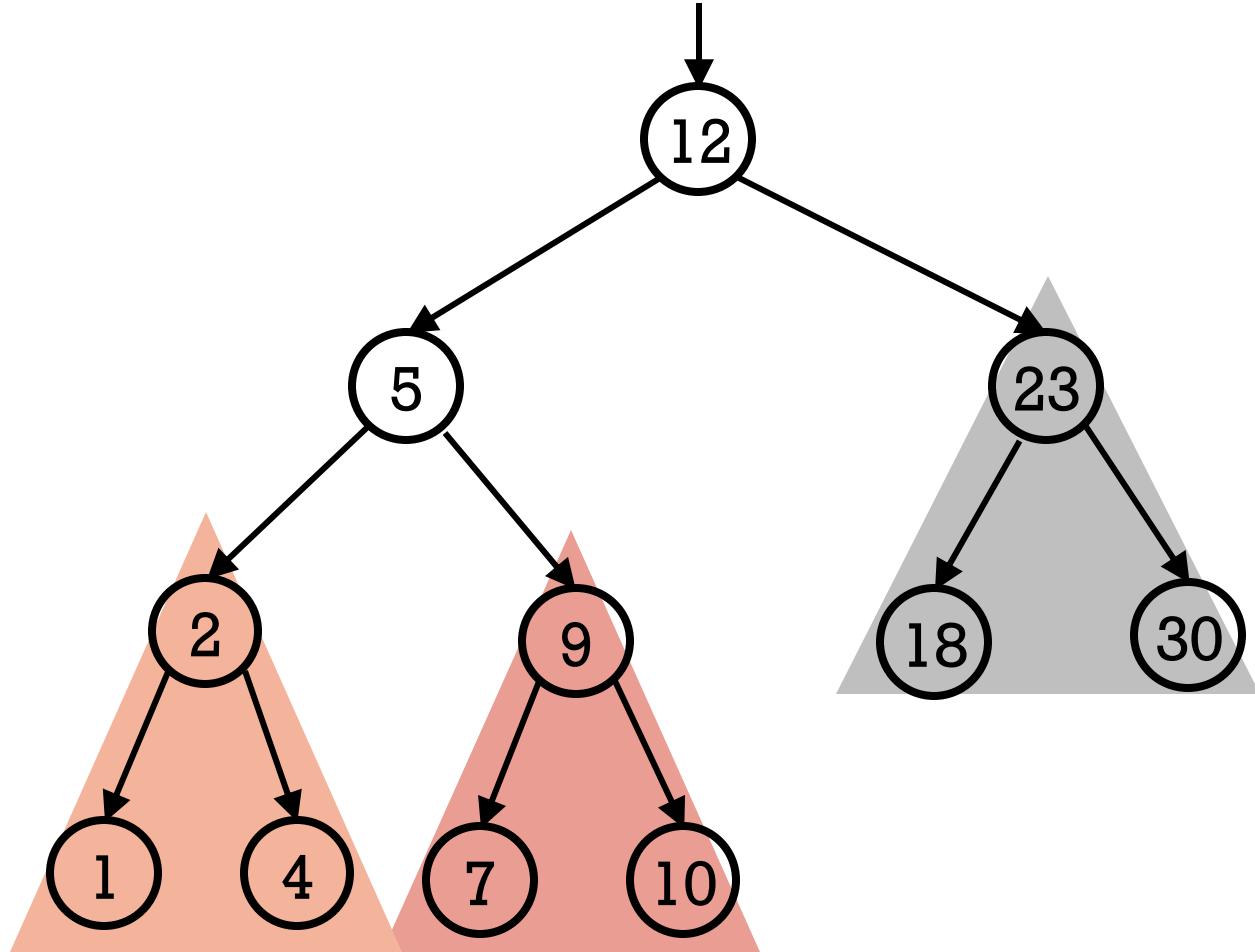


- After inserting 6
 - Balance condition at node 8 is violated
 - Perform Right Rotation on 8

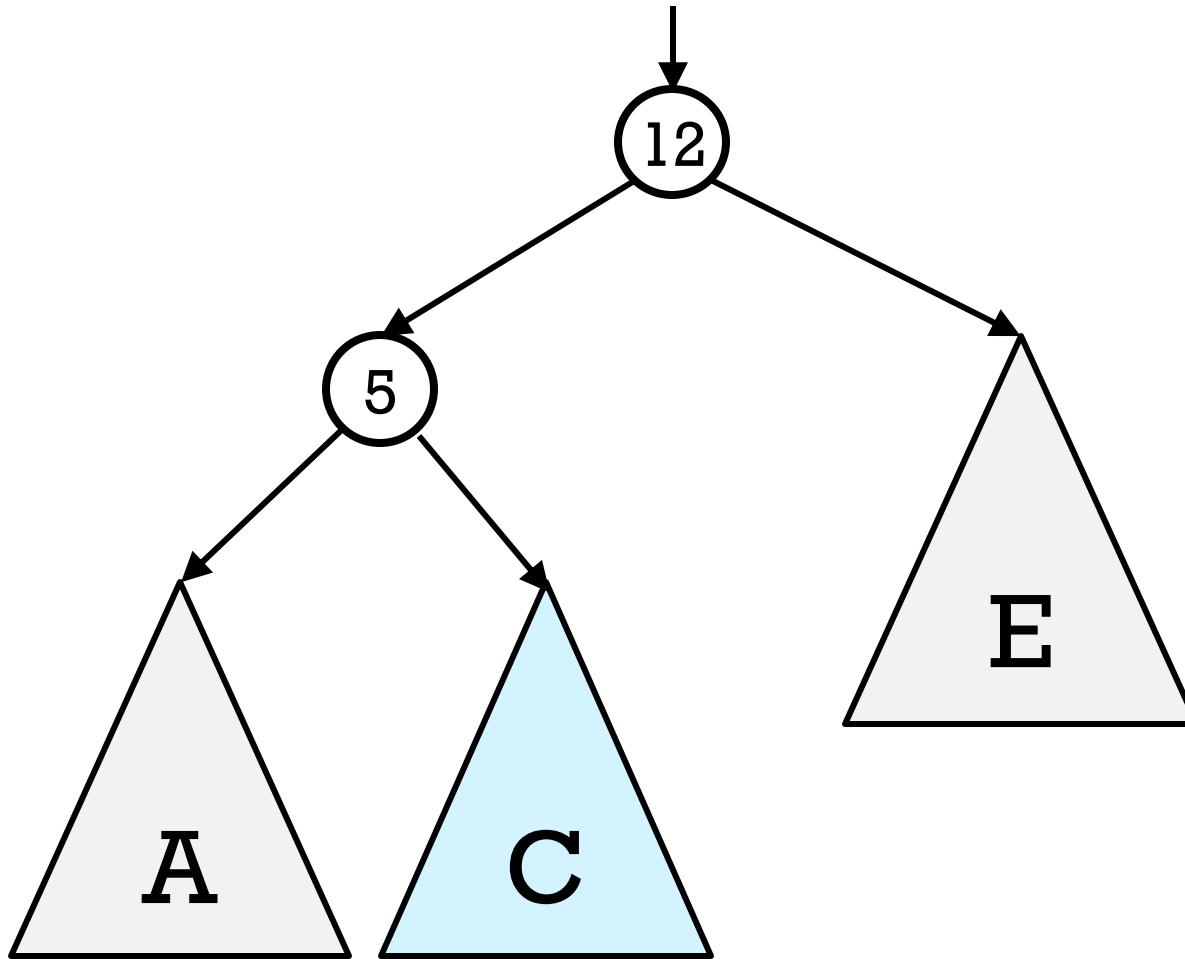
GENERALIZING...



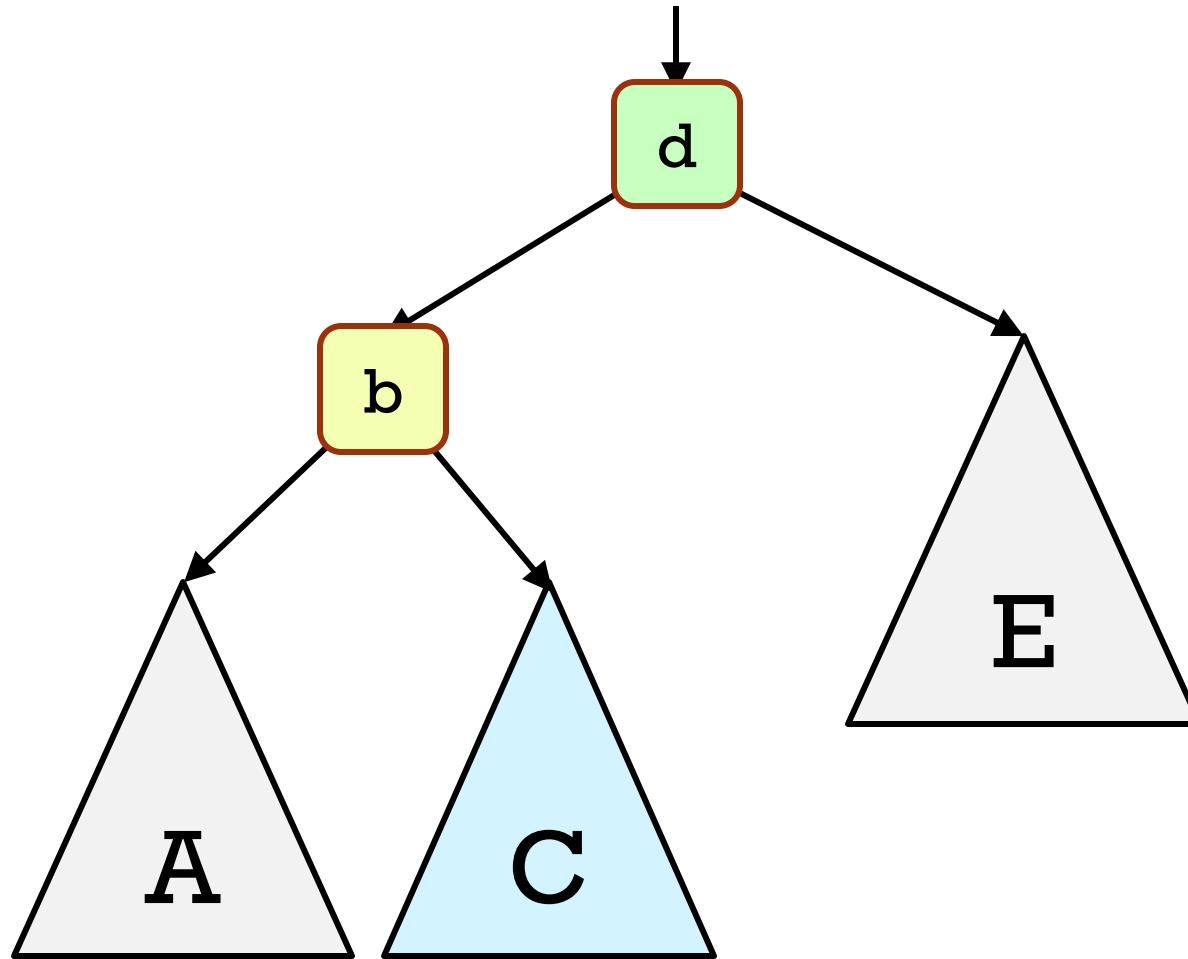
GENERALIZING OUR EXAMPLES...



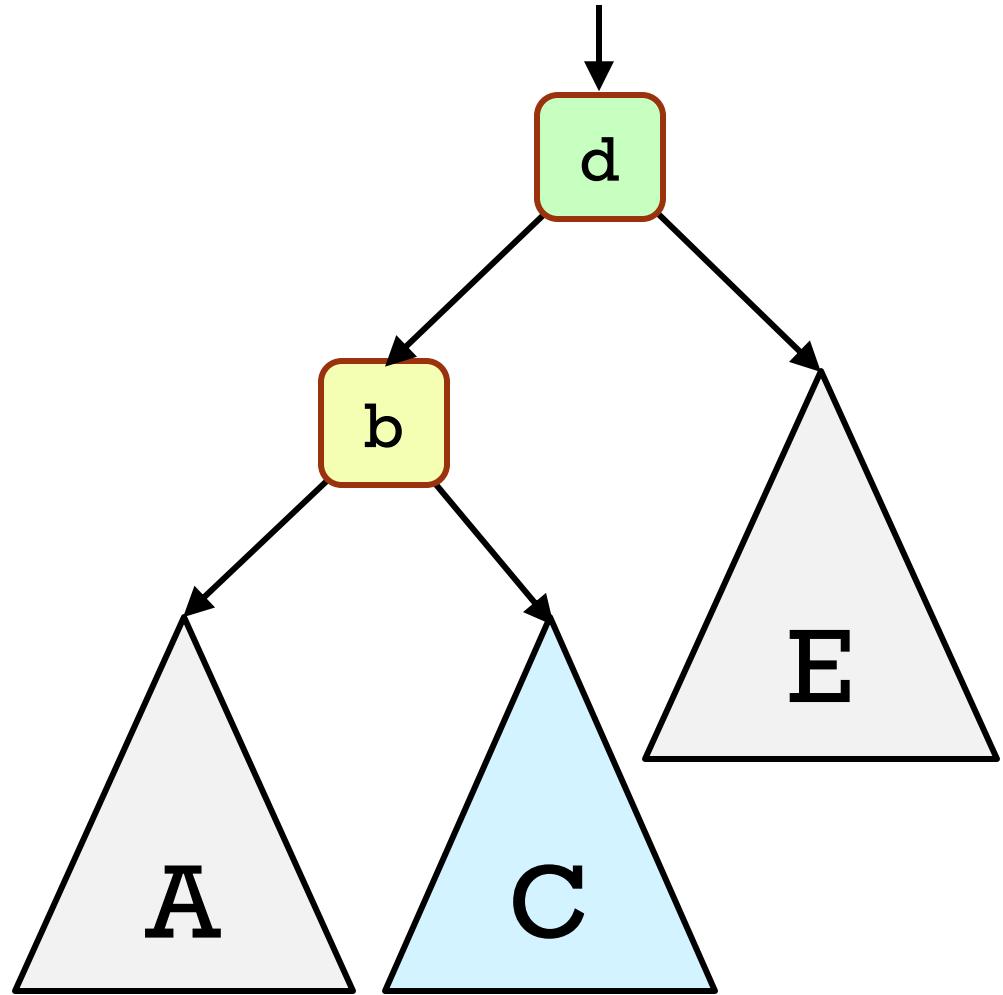
GENERALIZING OUR EXAMPLES...



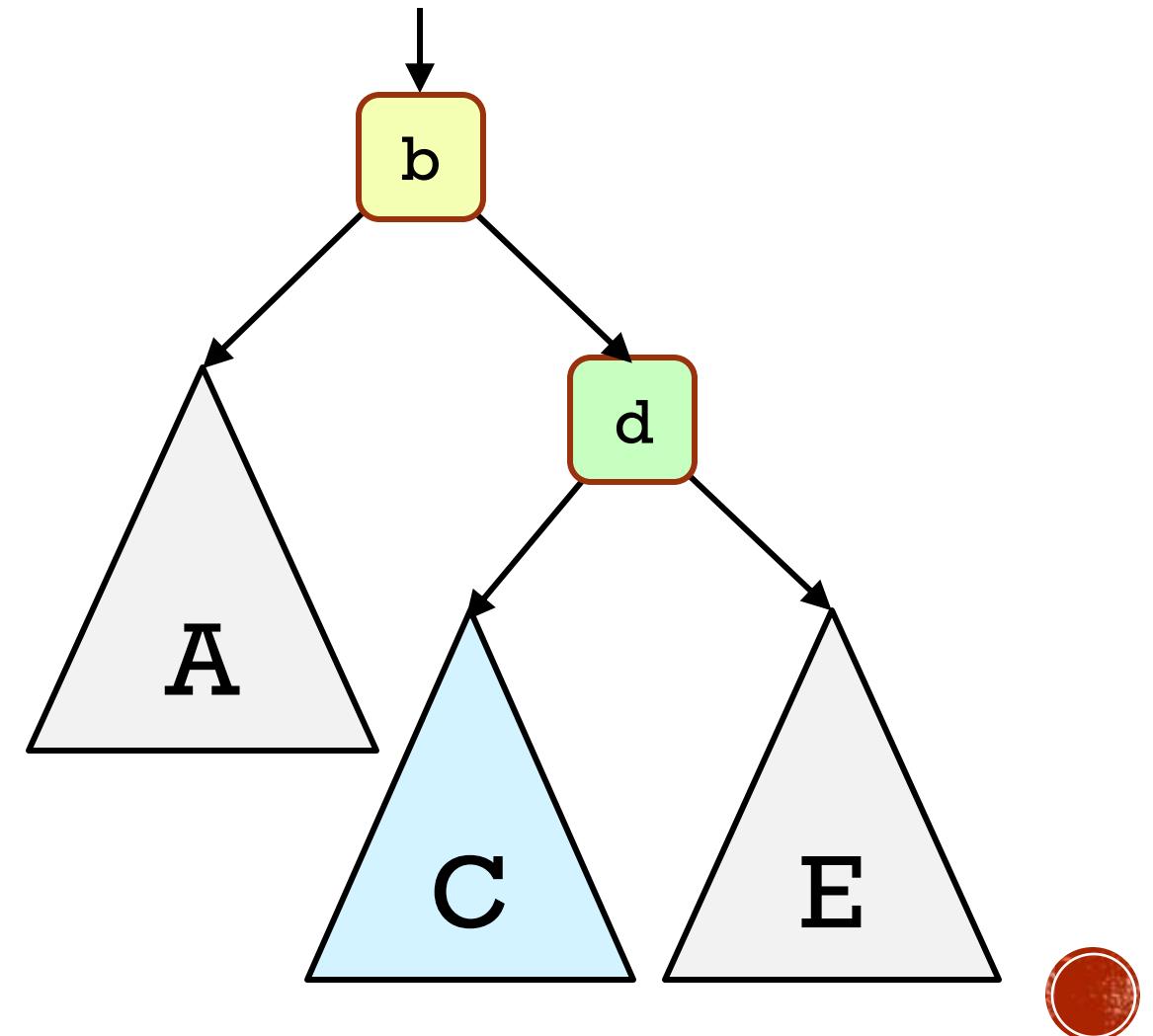
GENERALIZING OUR EXAMPLES...



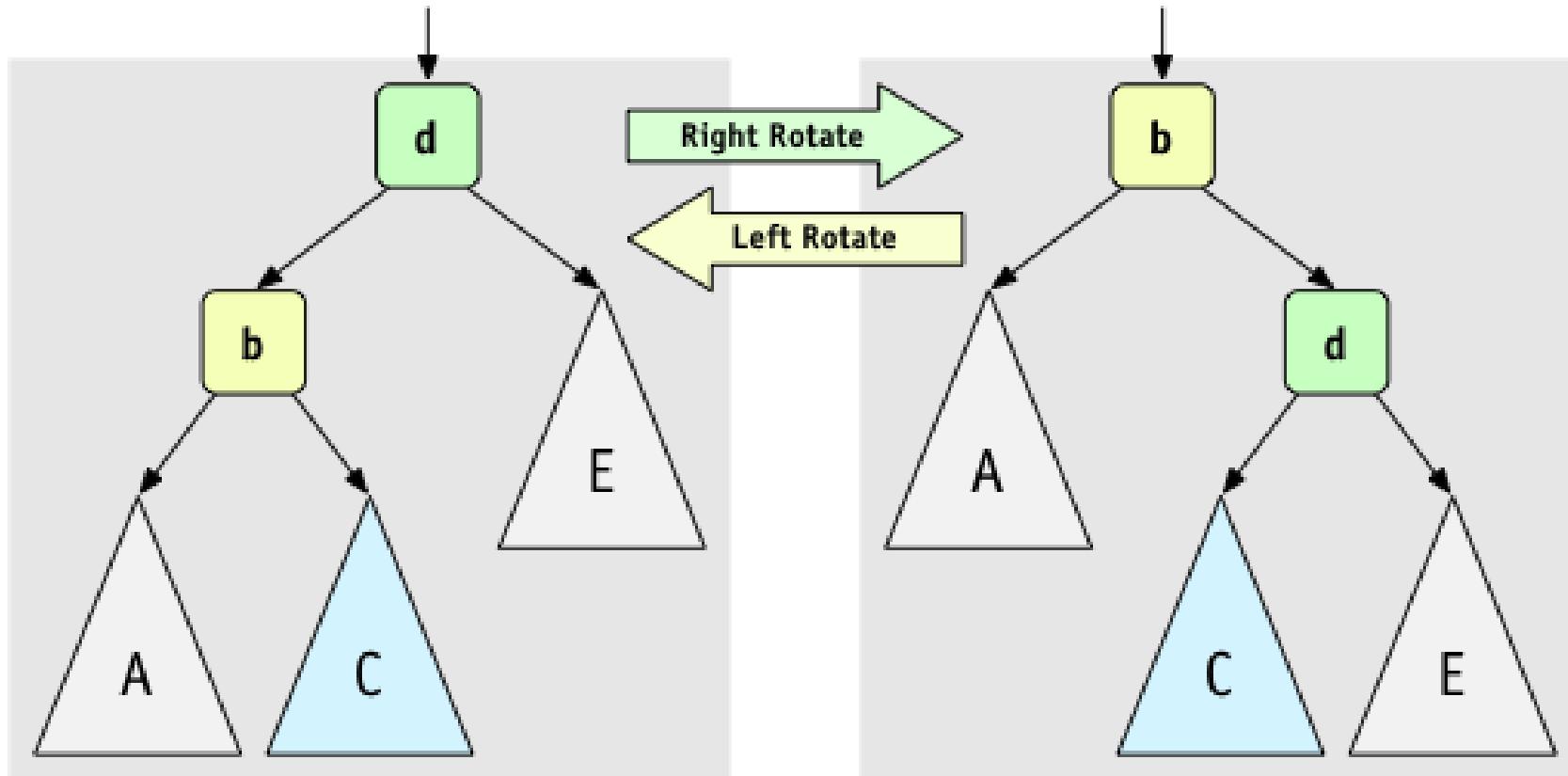
GENERALIZED SINGLE ROTATION



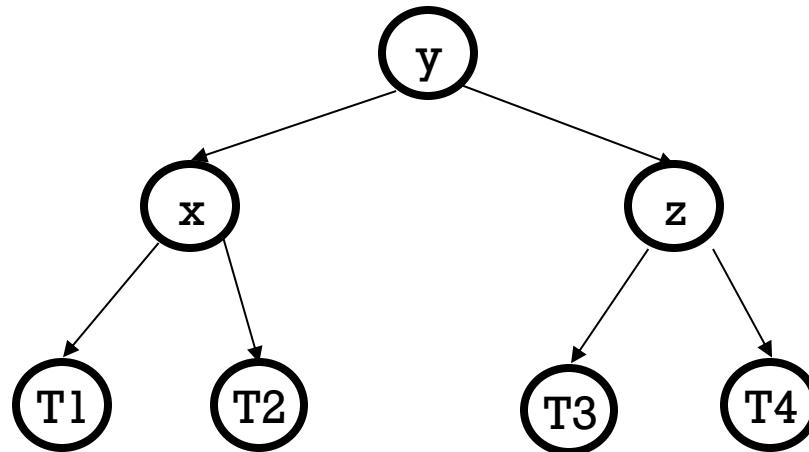
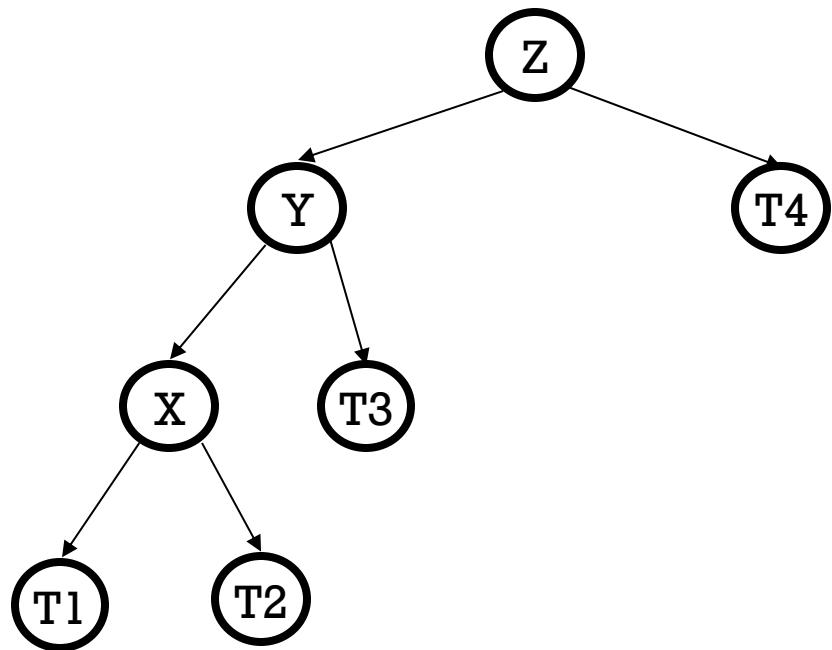
GENERALIZED SINGLE ROTATION



SINGLE ROTATIONS

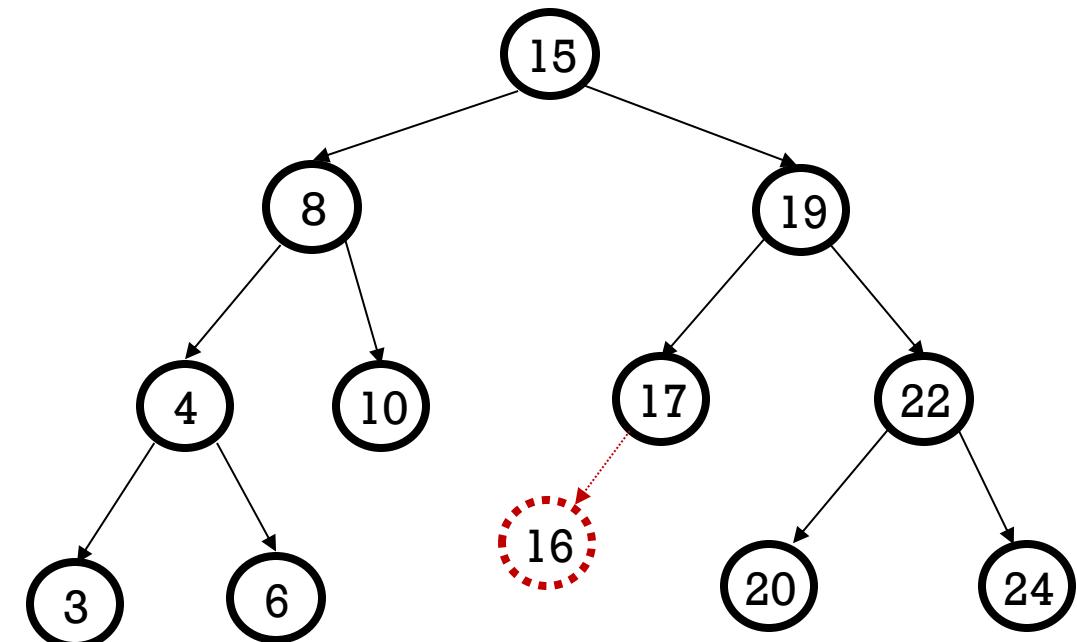
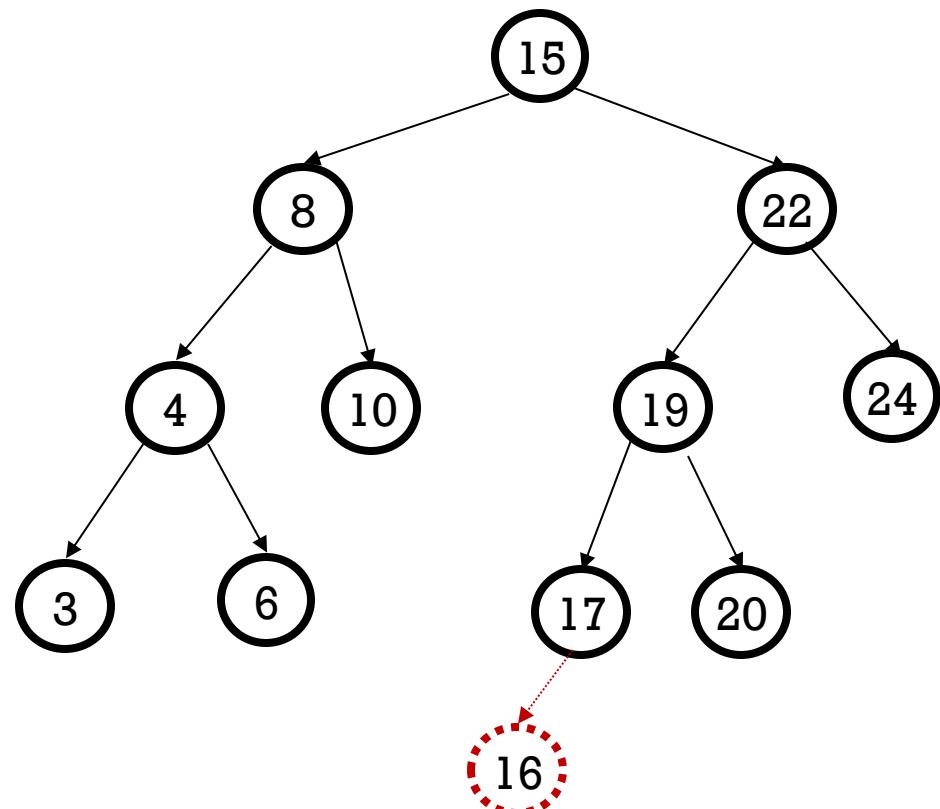


INSERTION CASE #1 LEFT LEFT

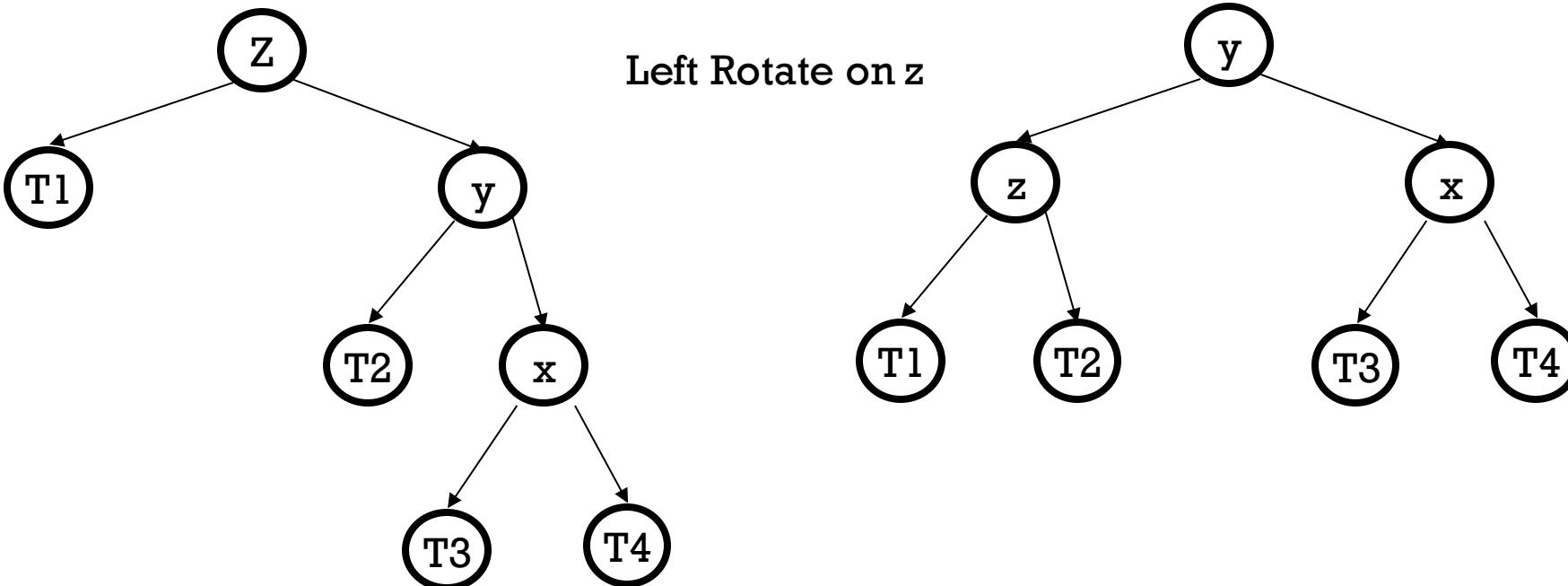


Right Rotate on z

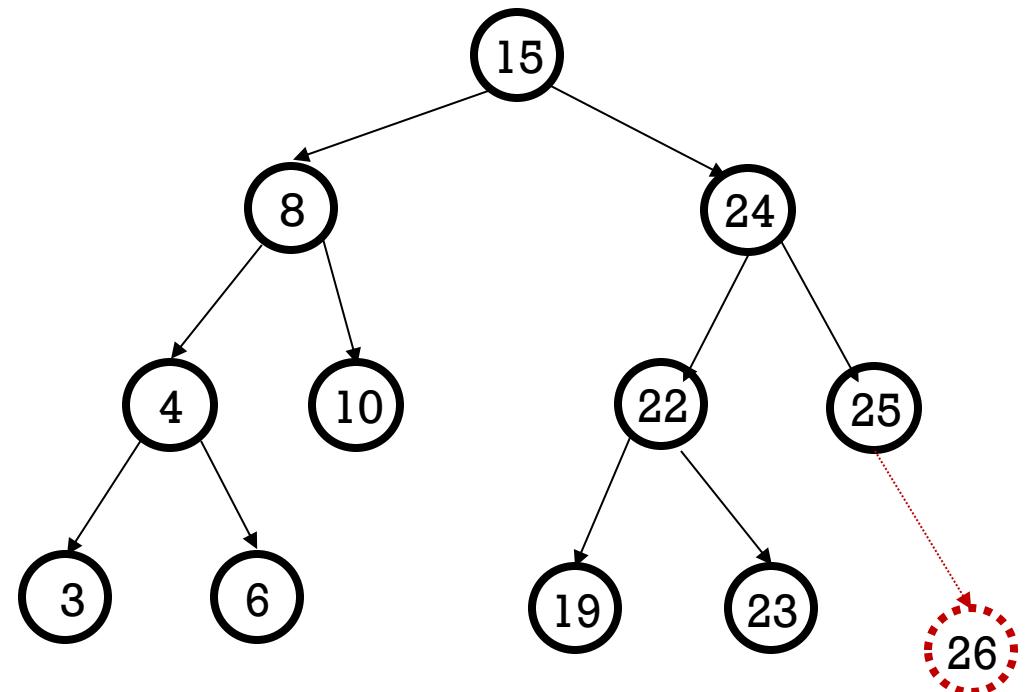
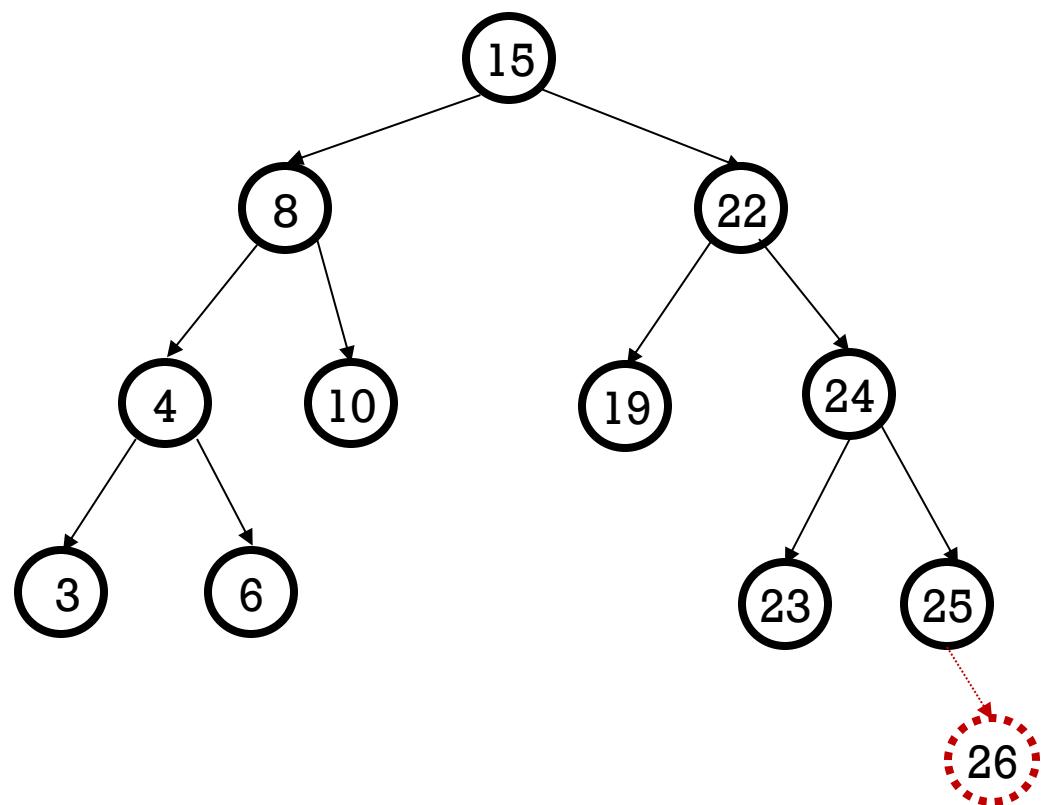
INSERT CASE # 1 EXAMPLE



INSERT CASE #2 RIGHT RIGHT

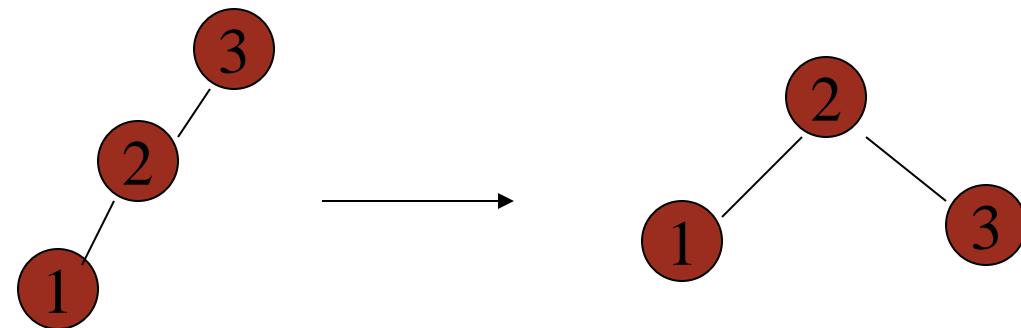


INSERT CASE # 2 EXAMPLE



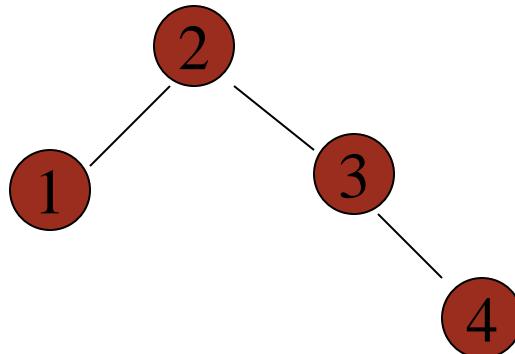
EXAMPLE

- Inserting 3, 2, 1, and then 4 to 7 sequentially into empty AVL tree

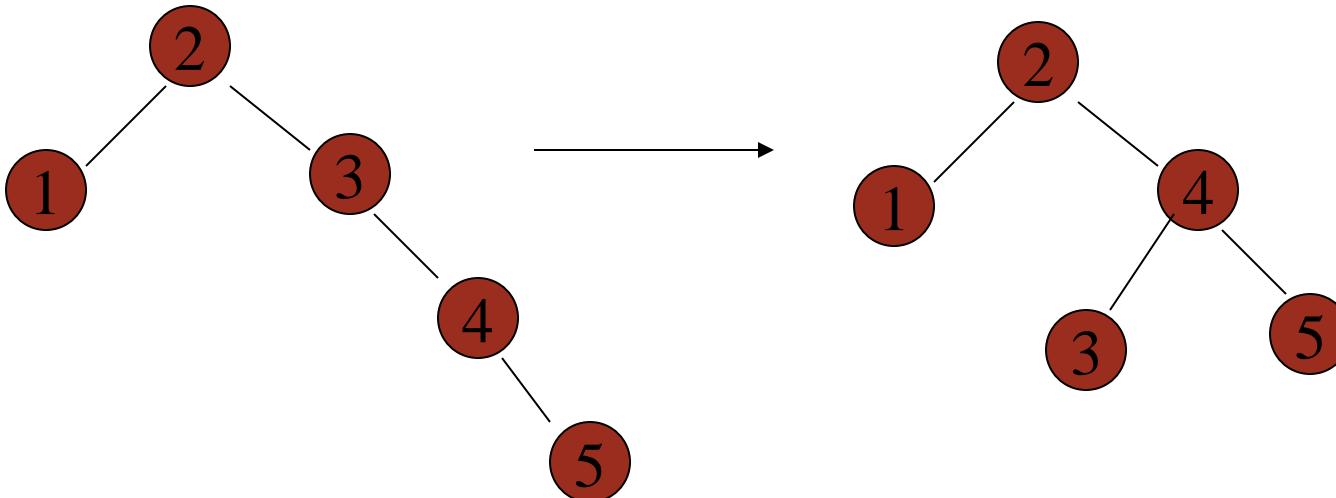


EXAMPLE

- Inserting 4

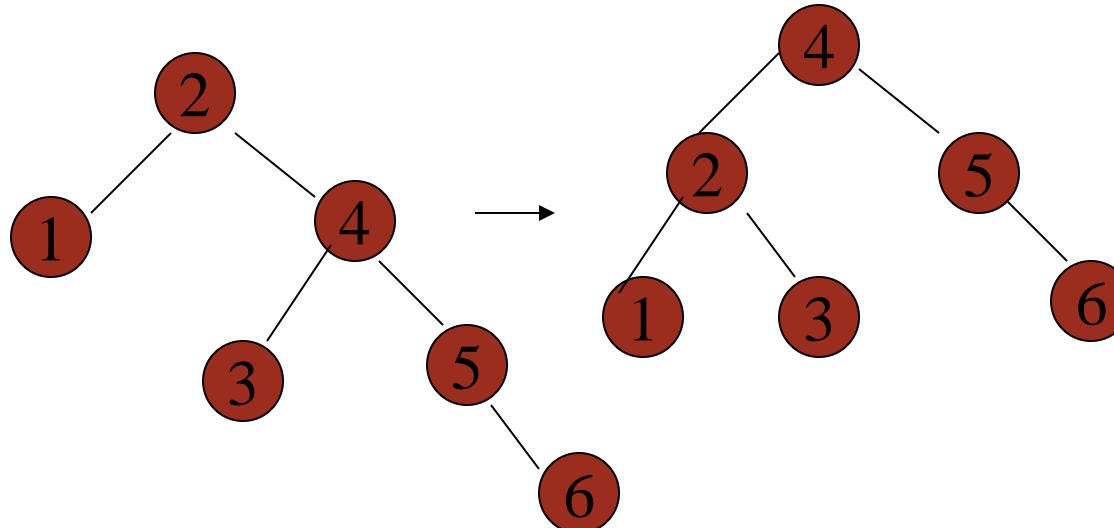


- Inserting 5

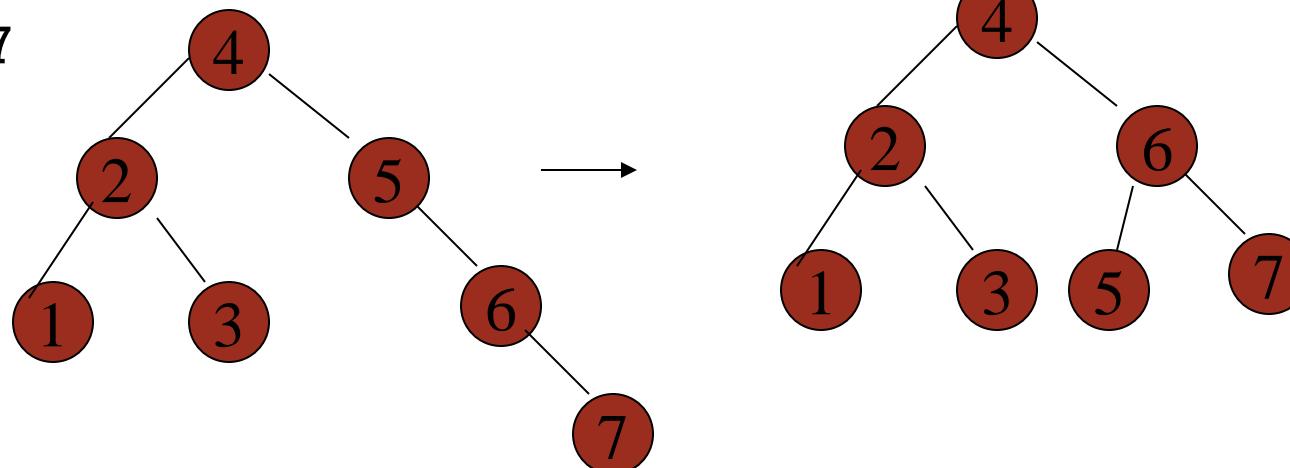


EXAMPLE

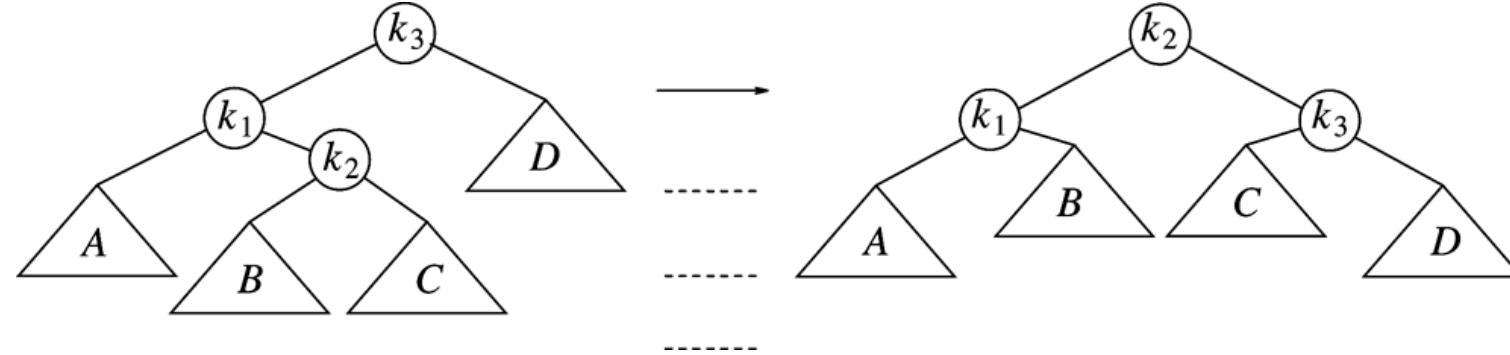
- Inserting 6



- Inserting 7

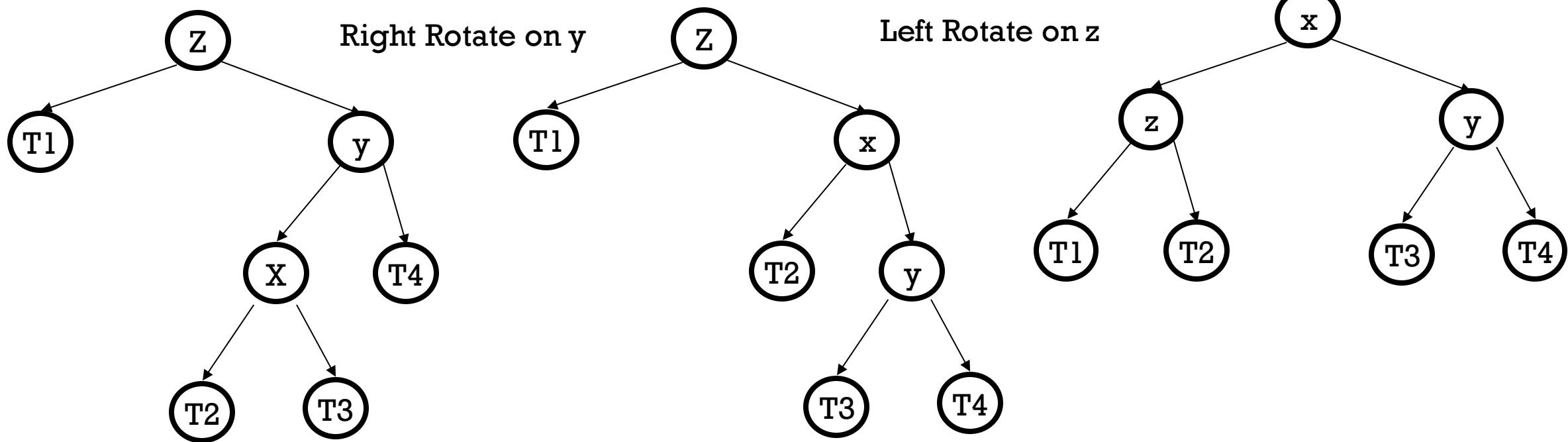


DOUBLE ROTATION

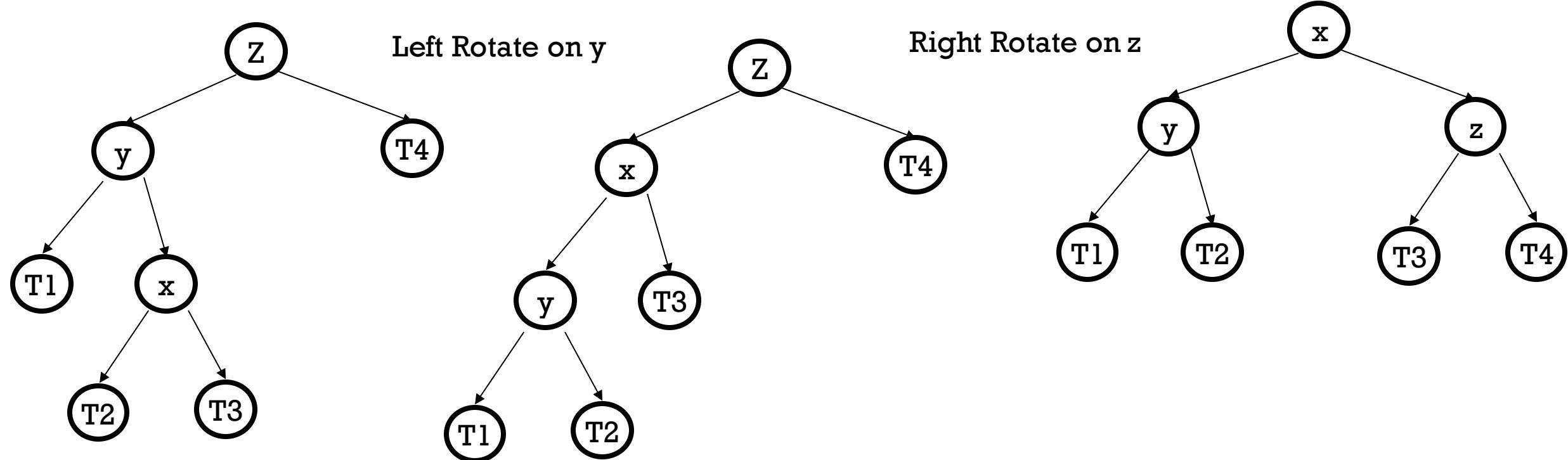


- Left-right double rotation to fix
- First rotate between \mathbf{k}_1 and \mathbf{k}_2
- Then rotate between \mathbf{k}_2 and \mathbf{k}_3

INSERT CASE # 3 RIGHT LEFT



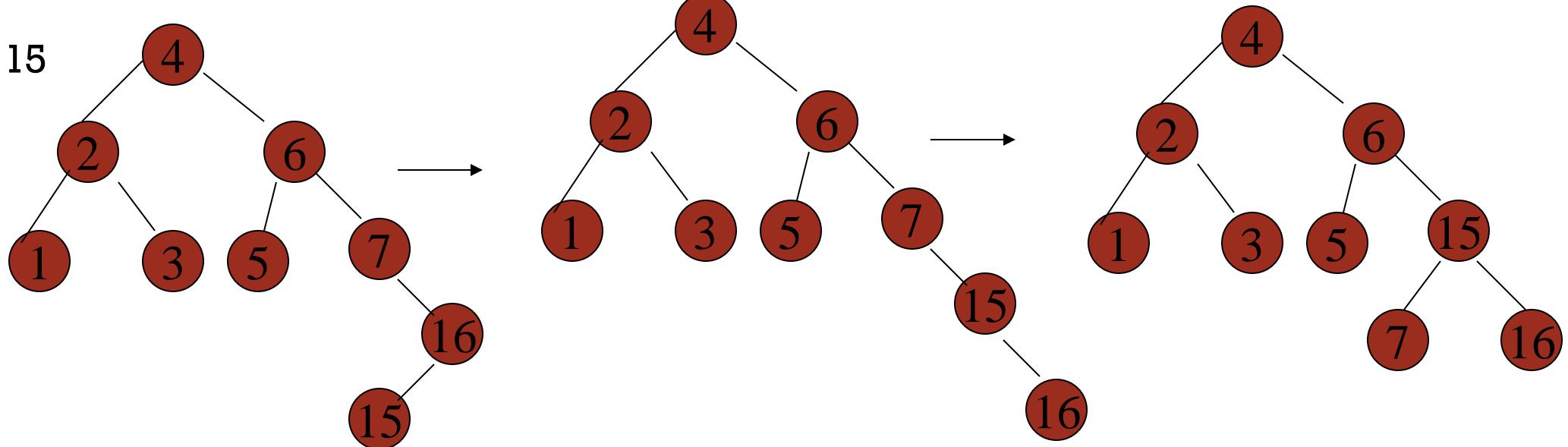
INSERT CASE #4 LEFT RIGHT



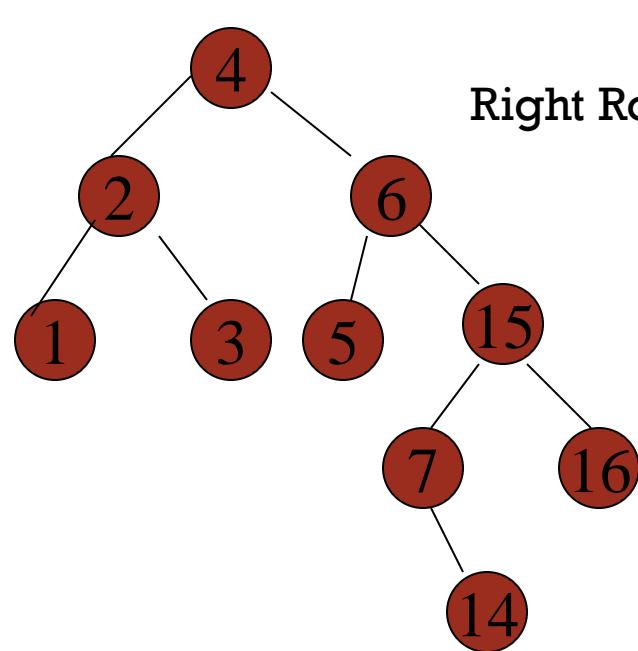
EXAMPLE

- Continuing the previous example by inserting
 - 16 down to 10, and then 8 and 9

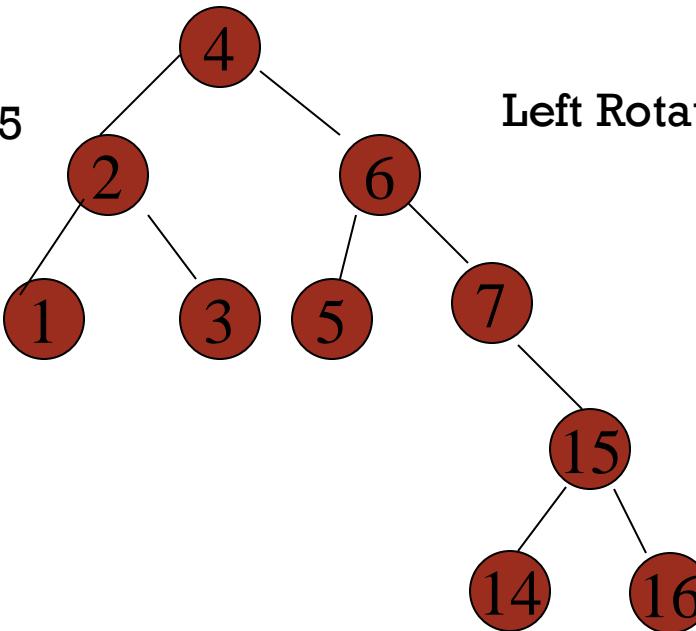
- Inserting 16 and 15



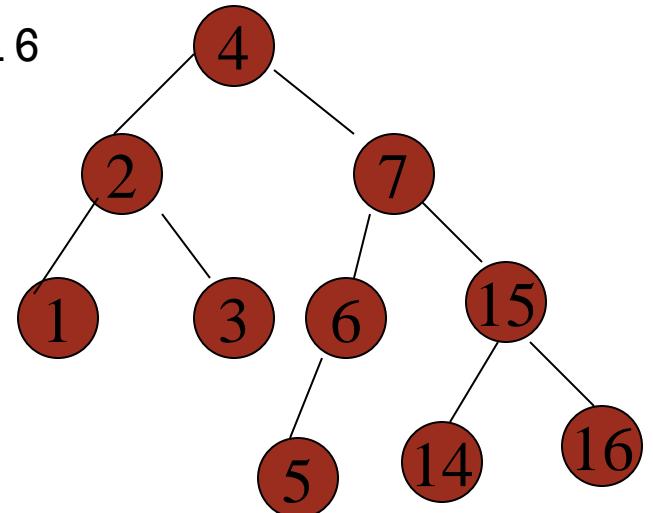
- Inserting 14



Right Rotate on 15



Left Rotate on 6



AVL DELETION ALGORITHM

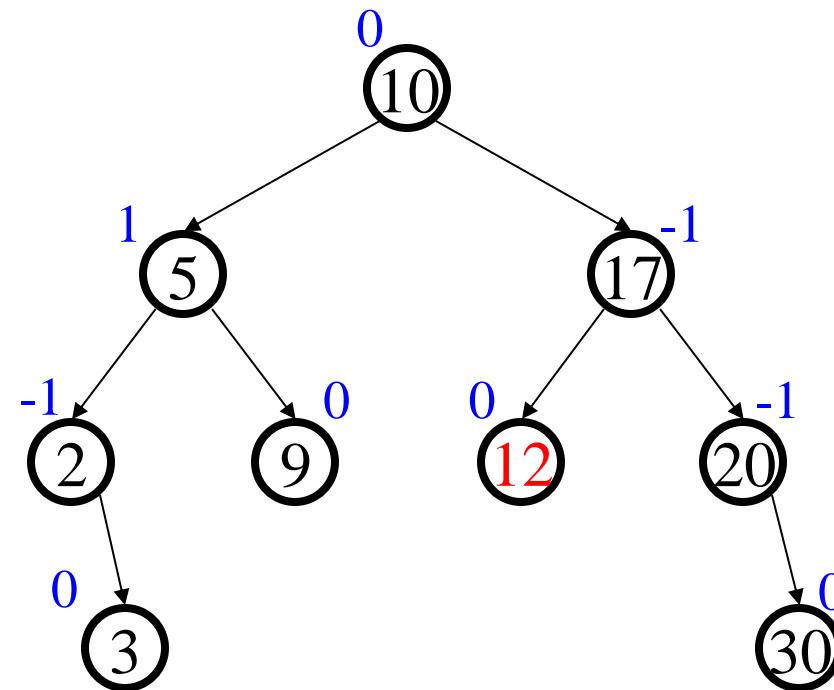
Recursive

1. Search downward for node
2. Delete node
3. Unwind, correcting heights as we go
 - a. If imbalance is left-left or right-right,
single rotate
 - b. If imbalance right-left or left-right,
double rotate

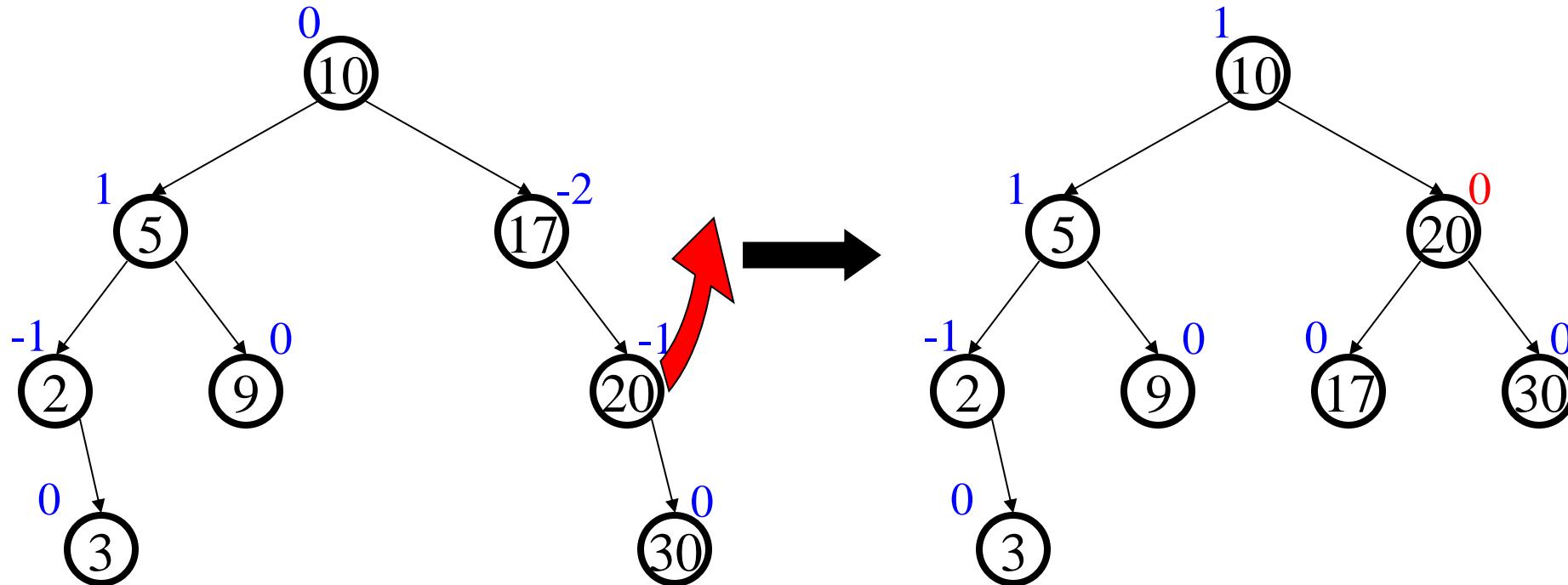


DELETION CASE #1

Delete(12)



SINGLE ROTATION ON DELETION



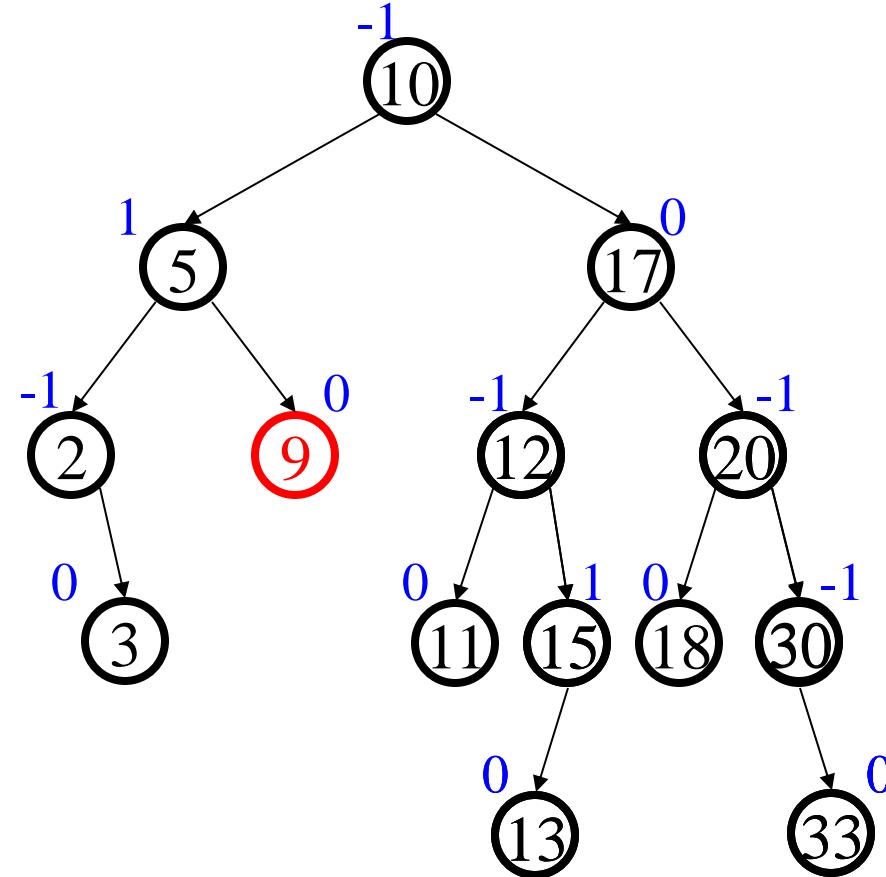
Deletion can differ from insertion – ***How?***

Deletion can propagate

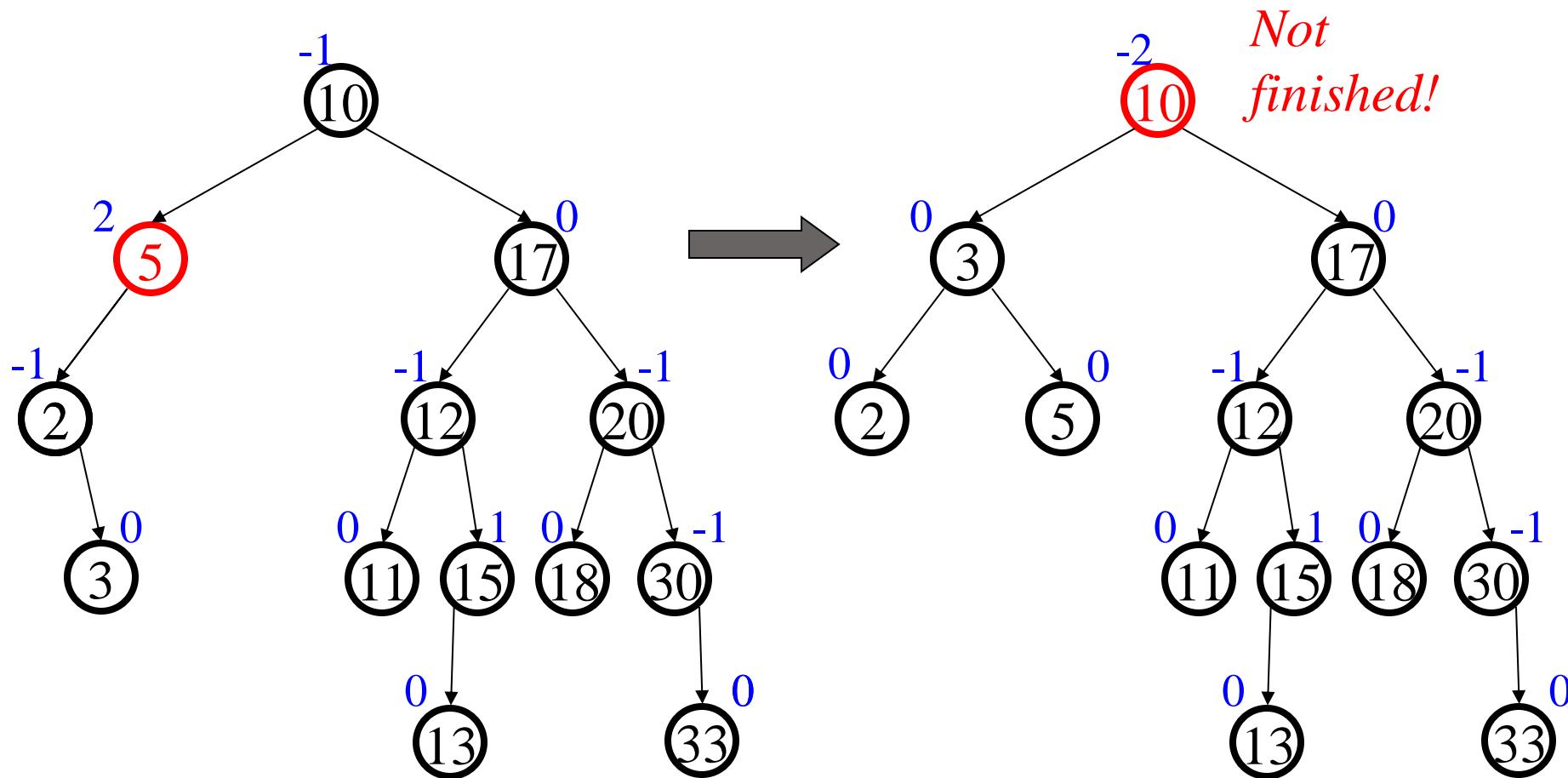


DELETION CASE #2 - PROPAGATION

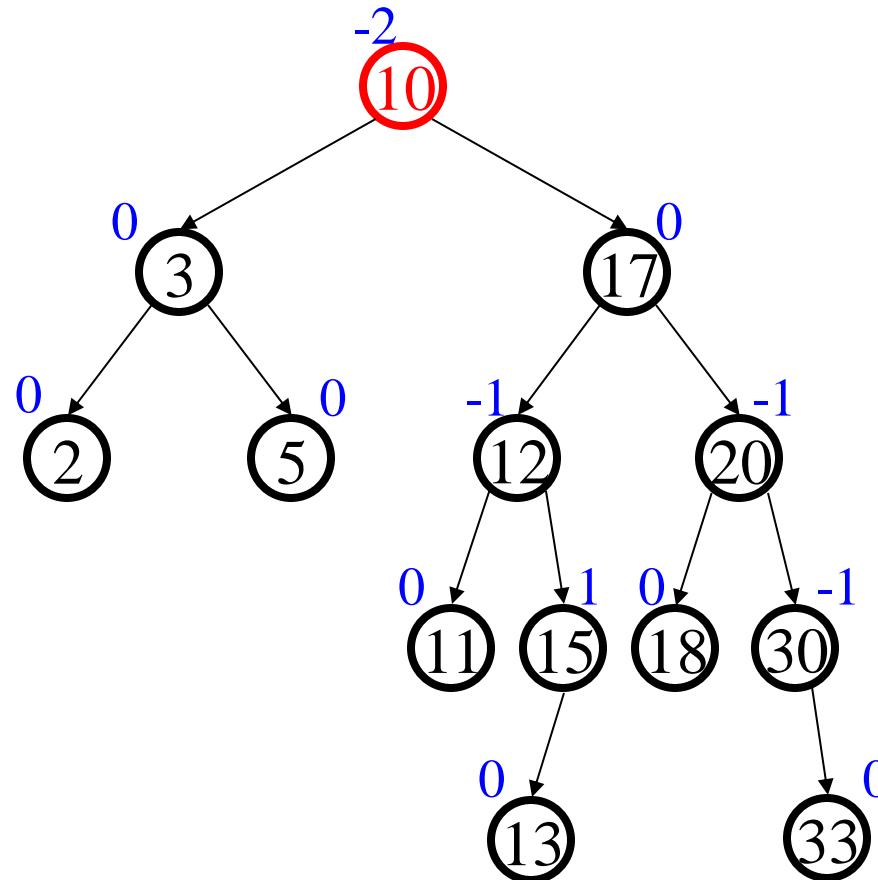
Delete(9)



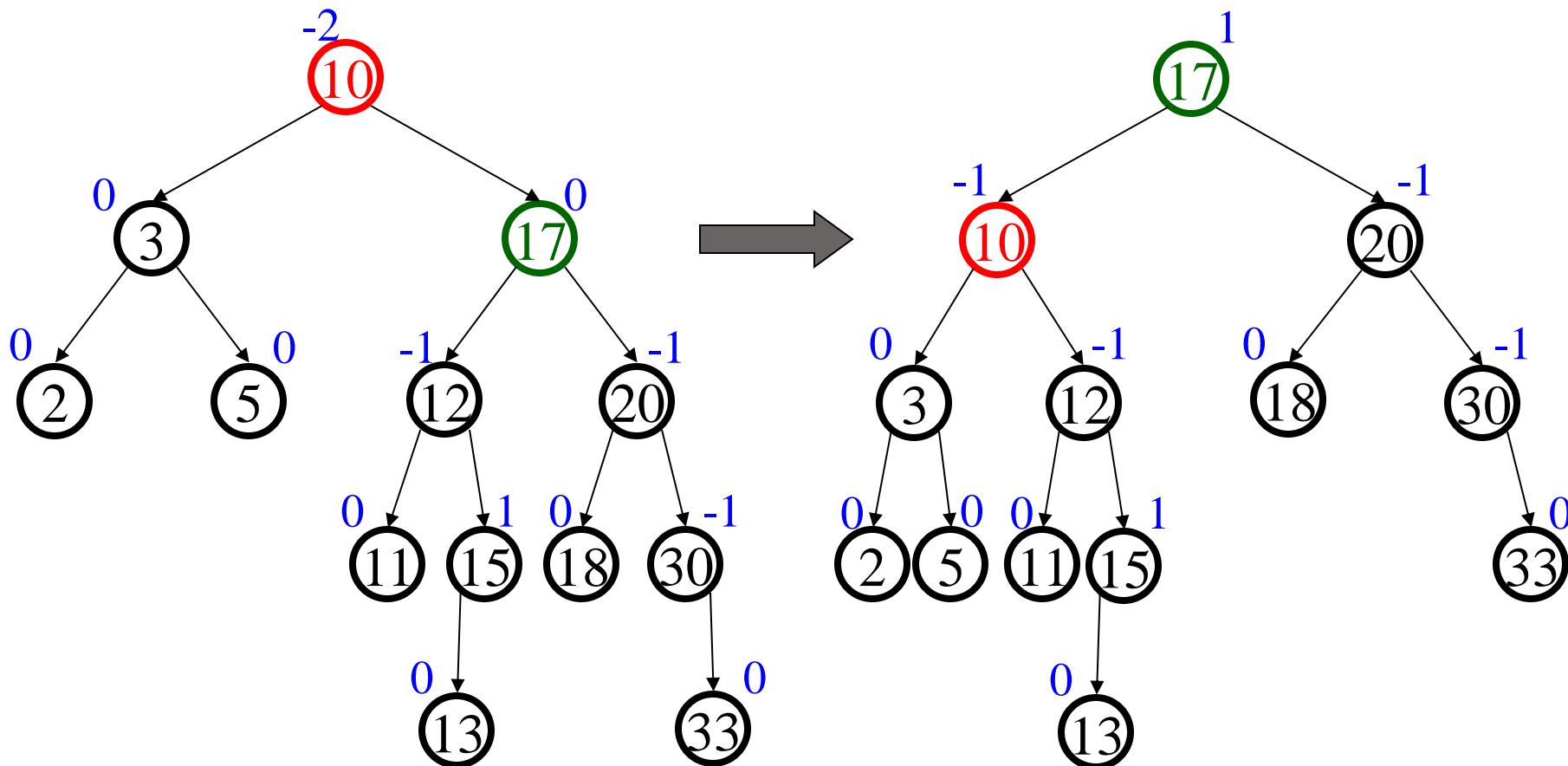
DOUBLE ROTATION ON DELETION



DELETION WITH PROPAGATION



PROPAGATED SINGLE ROTATION



PROS & CONS

Arguments for AVL trees:

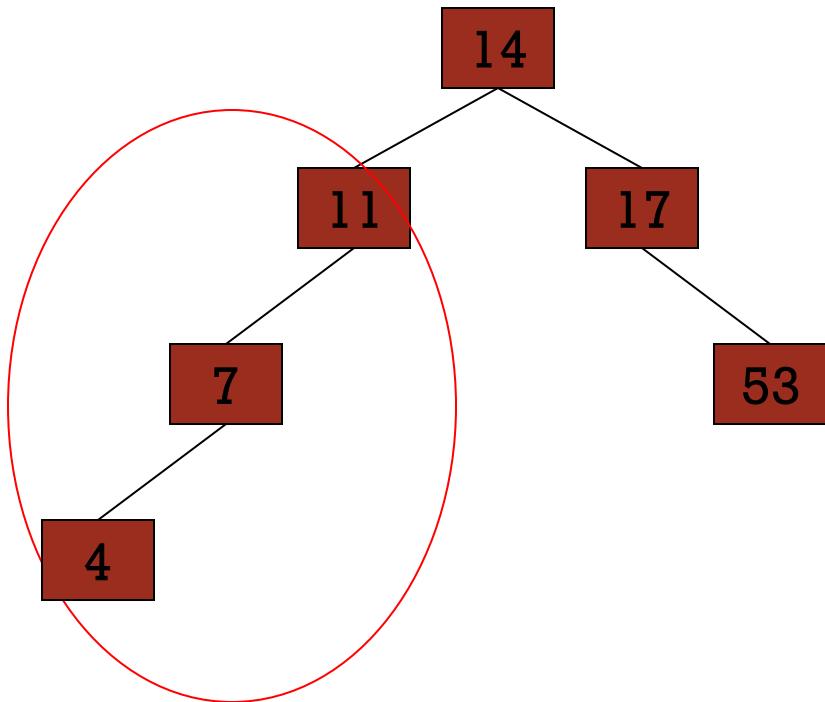
- All operations logarithmic worst-case because trees are *always* balanced
- Height balancing adds no more than a constant factor to the speed of insert and delete

Arguments against AVL trees:

- Difficult to program & debug
- More space for height field
- Asymptotically faster but rebalancing takes time

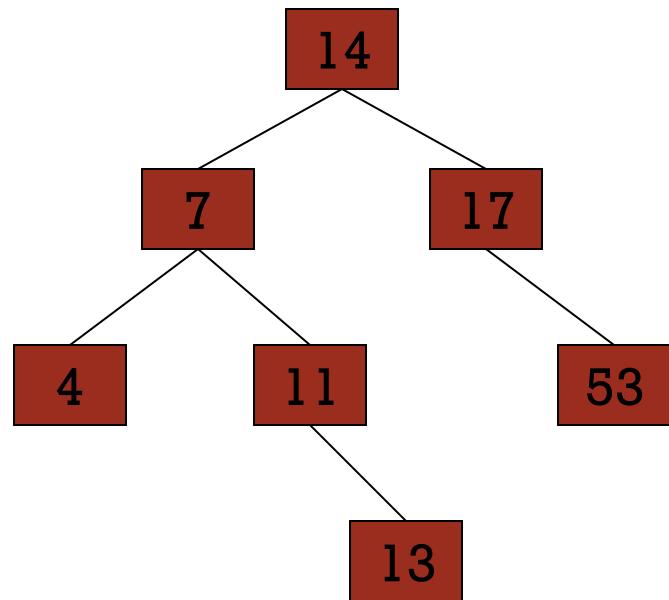
AVL Tree Example:

- Insert 14, 17, 11, 7, 53, 4, 13 into an empty AVL tree



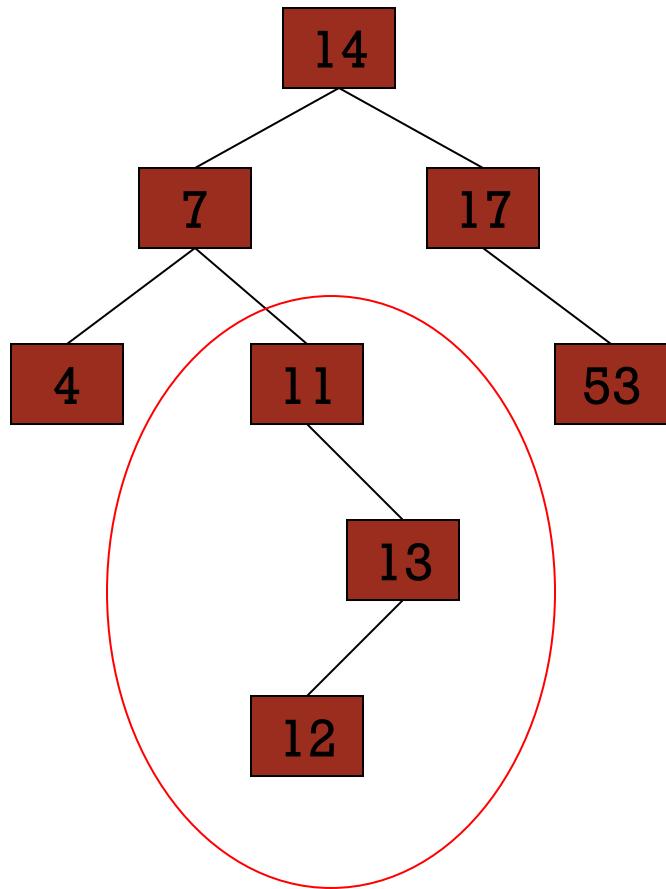
AVL Tree Example:

- Insert 14, 17, 11, 7, 53, 4, 13 into an empty AVL tree



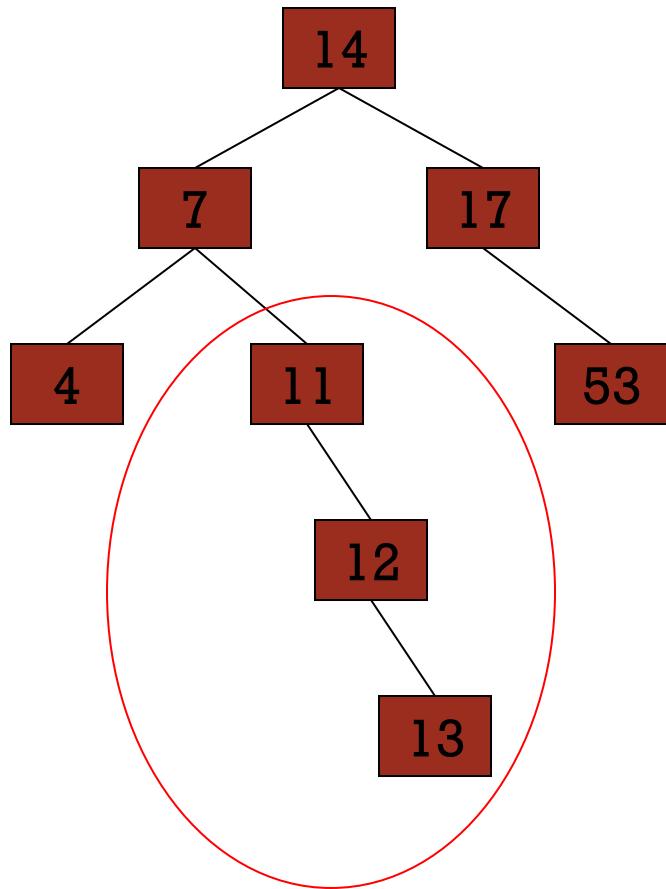
AVL Tree Example:

- Now insert 12



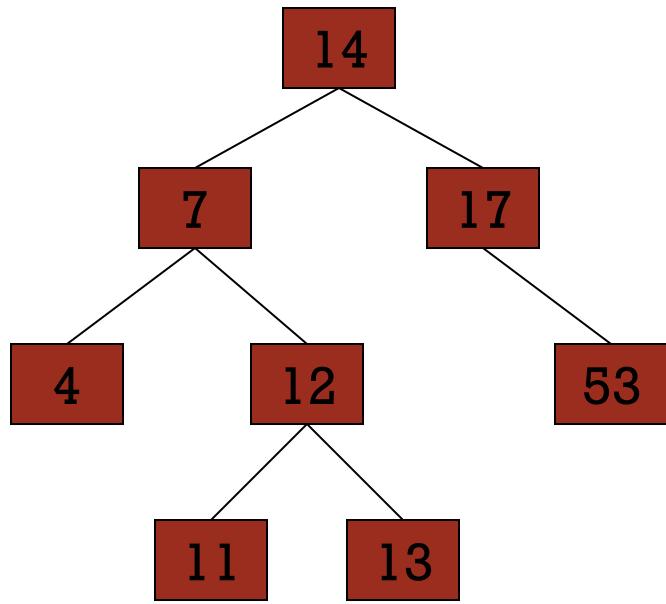
AVL Tree Example:

- Now insert 12



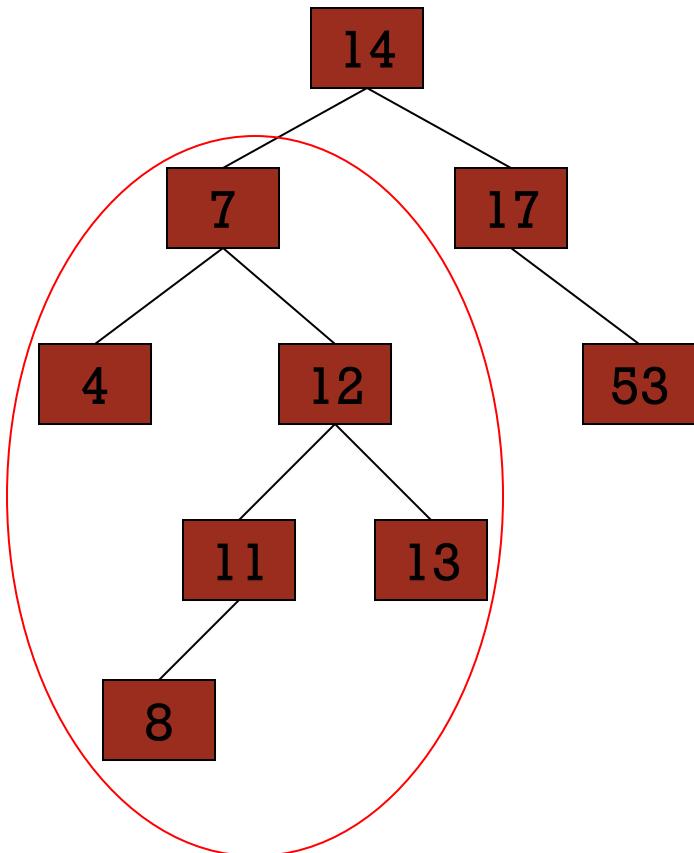
AVL Tree Example:

- Now the **AVL tree is balanced.**



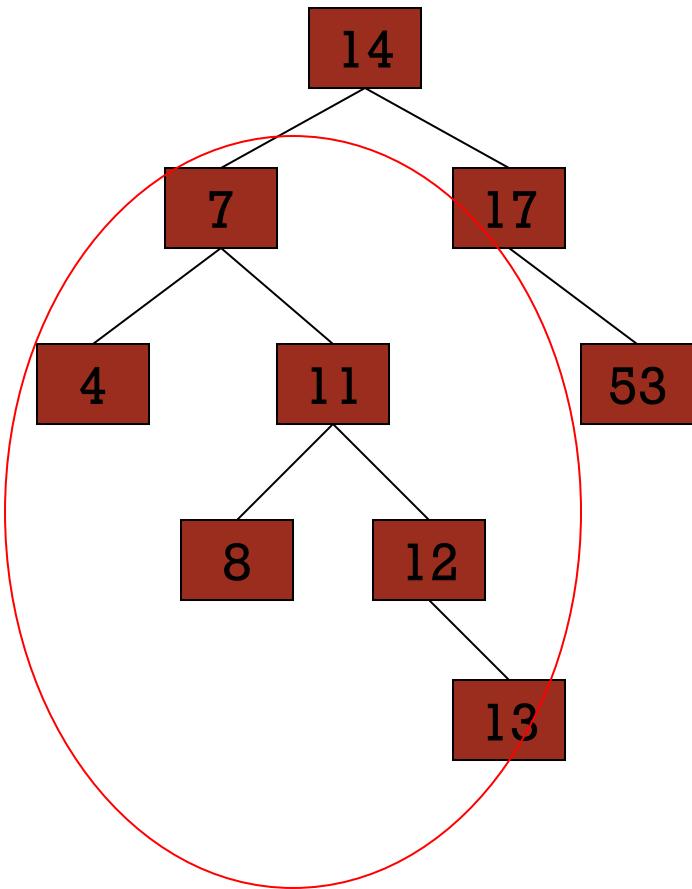
AVL Tree Example:

- Now insert 8



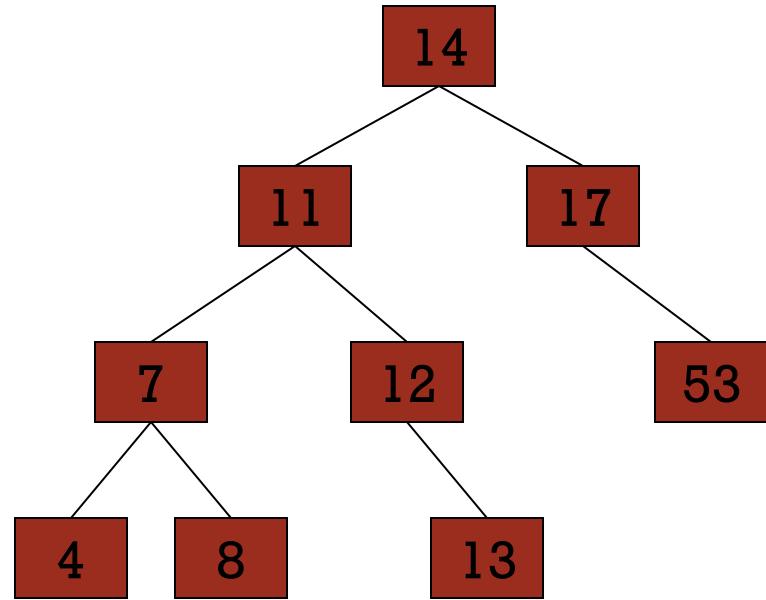
AVL Tree Example:

- Now insert 8



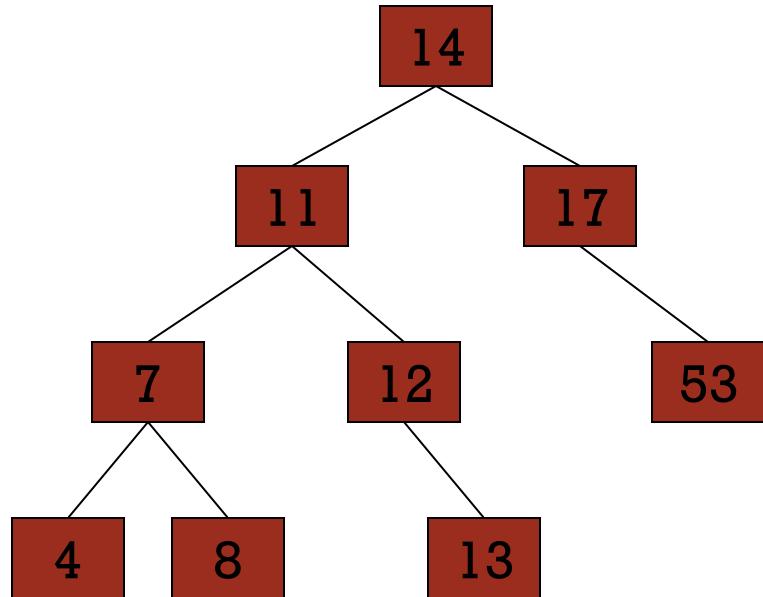
AVL Tree Example:

- Now the **AVL tree is balanced.**



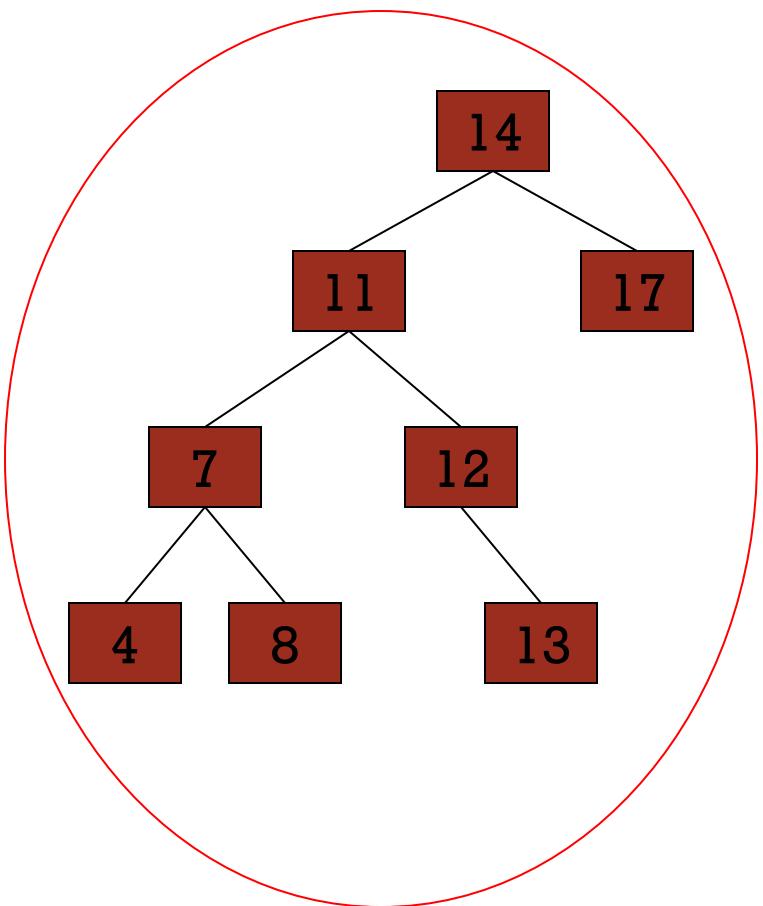
AVL Tree Example:

- Now remove 53



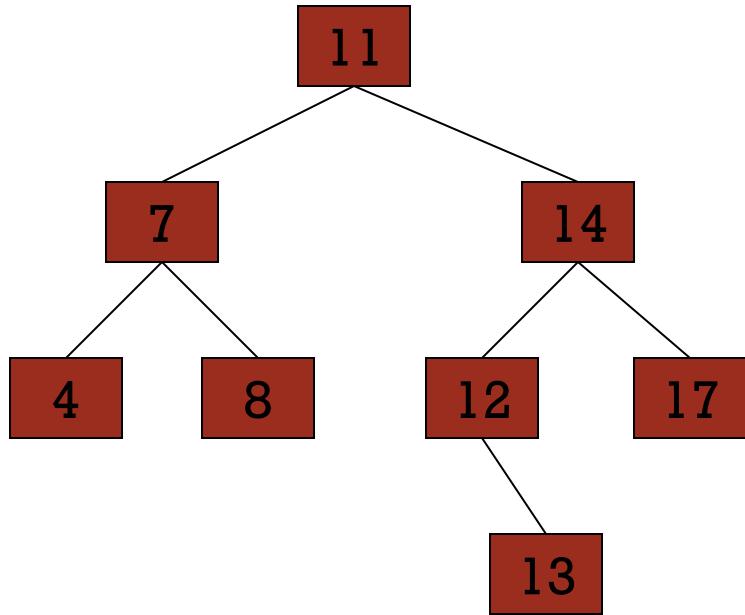
AVL Tree Example:

- Now remove 53, unbalanced



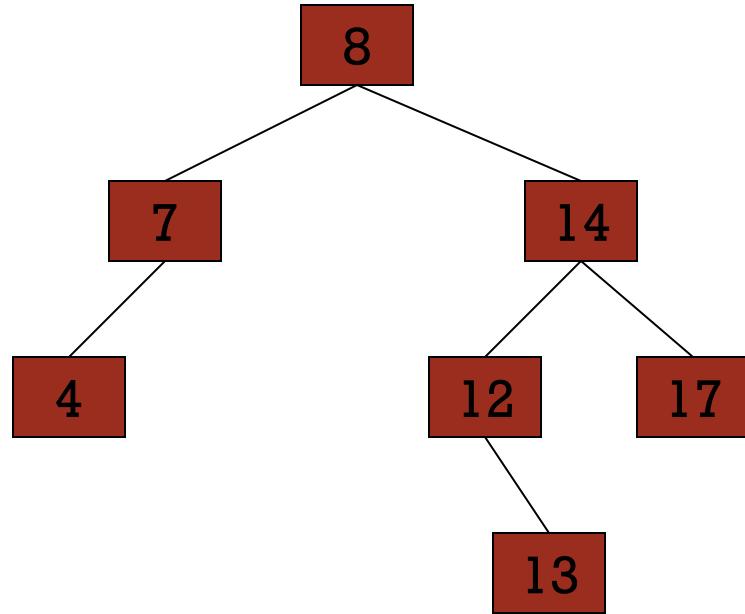
AVL Tree Example:

- Balanced! Remove 11



AVL Tree Example:

- Remove 11, replace it with the largest in its left branch



AVL Tree Example:

- Remove 8, replace with 12. Balanced.

