COMP 348: Principles of Programming Languages Assignment no.1

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1 Example to demonstrate LaTeX

Here is an example for the latex file

```
all_sections(CNAM, CNUM, L) :-
findall(L, takes_course(_, CNAM, CNUM, L), X),
sort(X, L).
```

2 Question 4

2.1 Modified Query

3 Question 6:

Indicate which of the following pairs of terms can be unied together. If they cant be unied, please provide the reason for it. In case of error, indicate the error. If they can be unied successfully, wherever relevant, provide the variable instantiations that lead to successful unication. (Note that: indicates unication)

3.1 food(bread, X) = Food(Y, soup)

In this case, unification will not be successful since **food** is a functor and **Food** is not. For this unification to be successful, both sides would have to contain matching functors such as **food**.

$3.2 \quad \text{Bread} = \text{soup}$

Since we have a variable and an atom, the variable will be instantiated to the value of the atom, thus the terms unify.

$3.3 \quad \text{Bread} = \text{Soup}$

Since we have two variables, the variable will be instantiated to the value of the other variable or vice versa, thus the terms unify.

3.4 food(bread, X, milk) = food(Y, salad, X)

In this case the functors match, but the arguments do not. Y will be instantiated to 'bread', but X will be instantiated to either 'salad' or 'milk', meaning it cannot match with both of the atoms. This means the terms will not unify. If there

were three variables, lets say X,Y,Z instead of X,Y,Y, then the terms would unify.

$3.5 \quad \text{manager}(X) = Y$

In this case unification will be successful and Y will be instantiated to manager(X).

3.6 $\operatorname{meal}(\operatorname{healthyFood}(\operatorname{bread}), \operatorname{drink}(\operatorname{milk})) = \operatorname{meal}(X,Y)$

In this case unification will be successful with **X** instantiated to **healthy-**Food(bread) and **Y** instantiated to **drink(milk)**.

3.7 $\operatorname{meal}(\operatorname{eat}(\mathbf{Z}), \operatorname{drink}(\operatorname{milk})) = [\mathbf{X}]$

In this case, unification is not successful as the arities and functors do not match.

3.8 $[eat(\mathbf{Z}), drink(milk)] = [\mathbf{X}, \mathbf{Y} \mid \mathbf{Z}]$

In this case unification is successful with X being instantiated to $eat([\])$, Y being instantiated to drink(milk), and Z being instantiated to $[\]$.

3.9
$$f(X, t(b, c)) = f(l, t(Z, c))$$

In this case unification is successful with X being instantiated to 1, and and Z being instantiated to b.

3.10 $\operatorname{ancestor}(\operatorname{french}(\operatorname{jean}), B) = \operatorname{ancestor}(A, \operatorname{scottish}(\operatorname{joe}))$

In this case unification is successful with **A** being instantiated to **french(jean)**, and and **B** being instantiated to **scottish(joe)**.

3.11 meal(healthyFood(bread), Y) = meal(X, drink(water))

In this case unification is successful with X being instantiated to **healthy-Food(bread)**, and and Y being instantiated to drink(water).

$$3.12 \quad [H|T] = [a, b, c]$$

In this case unification is successful with \mathbf{H} being instantiated to \mathbf{a} , and \mathbf{T} being instantiated to $[\mathbf{b}, \mathbf{c}]$.

$$3.13 \quad [H, T] = [a, b, c]$$

In this case unification is not successful as the number of elements in the left list does not equal the number of elements in the right list. The list on the left would have to be composed of 3 variables for unification to be successful.

3.14 breakfast(healthyFood(bread), egg, milk) = breakfast (healthyFood(Y), Y, Z)

In this case unification is not successful as **Y** cannot be instantiated to both **bread** and **egg** at the same time.

3.15 $\operatorname{dinner}(X, Y, \operatorname{Time}) = \operatorname{dinner}(\operatorname{jack}, \operatorname{cook}(\operatorname{egg}, \operatorname{oil}), \operatorname{Evening})$

In this case unification is successful with **X** being instantiated to **jack**, **Y** being instantiated to **cook(egg, oil)**, and **Time** being instantiated to **Evening**.

3.16
$$k(s(g), Y) = k(X, t(k))$$

In this case unification is successful with X being instantiated to s(g), and Y being instantiated to t(k).

3.17 equation(Z,
$$f(x, 17, M), L*M, 17$$
) = equation(C, $f(D, D, y), C, E$)

In this case, unification is not successful as the term f(x, 17, M) does not unify with the term f(D, D, y). These two terms do not unify as D cannot be instantiated to both x and 17 at the same time.

3.18
$$a(X, b(c, d), [H|T]) = a(X, b(c, X), b)$$

In this case unification is not successful since $[\mathbf{H}|\mathbf{T}]$ cannot unify with \mathbf{b} since \mathbf{b} is an atom.

4 Question 7

4.1 field(hit_transfer, engineering)

Query type: ground.

Output: true

Steps:

• Unify the query with the head of rule:

```
field(X, Y) :- course(X, Z), field(Z, Y).
```

- **Instantiate** *X* to hit transfer, *Y* to engineering.
- Resolve to two new queries:

```
course(hit\_transfer, Z), field(Z, engineering).
```

- Both queries are evaluated.
- Both goals prove true \rightarrow succeeds. Thus, the answer to the query is true.

4.2 ? lab_number(fine_arts,X).

Query type: non-ground.

Output:

X = 10; false.

Steps:

• Unify with the head of fact:

```
lab_number(fine_arts,10).
```

• Instantiate X to 10.

4.3 ? field(computer, literature).

Query type: ground.

 $\label{Output:false.} \textbf{Output:} \ \text{false.}$

Steps:

• Unify the query with the head of the rule:

```
field(X, Y) :- course(X, Z), field(Z, Y).
```

- Instantiate computer to X, literature to Y.
- Resolve to two new queries:

```
course(computer, Z), field(Z, literature).
```

- Both queries are evaluated. Once the first goal succeeds, try the next one on the right with the same instantiation.
- After trying all possible cases, no cases prove true. The query is not successful.

4.4? course(X,Y).

Query type: non-ground. Output:

```
X = hit_transfer,
Y = mechanical;
X = web_design,
Y = computer;
X = design_methods,
Y = fine-arts;
X = poetry,
Y = literature;
X = leadership,
Y = management;
X = biology,
Y = medicin.
```

Steps:

- Unify Unify with each statement of course() with 2 arities.
- Instantiate:
 - X is intanciated as hit_transfer, Y is intanciated as mechanical
 - X is intanciated as web_design, Y is intanciated as computer
 - X is intanciated as design_methods, Y is intanciated as fine-arts
 - X is intanciated as poetry, Y is intanciated as literature
 - X is intanciated as leadership, Y is intanciated as management
 - $-\,$ X is intanciated as biology, Y is intanciated as medic in

4.5 ? student(adrian).

```
Query type: ground.
Output: true.
Steps:
```

• Unify the query with the head of the rule:

```
student(X):- student(X,_).
```

• Instantiate adrian to X.

• Resolve into 1 goal:

```
student(adrian, _).
```

 $\bullet\,\,$ The goal will be unified with

```
student(adrian, web_design).
```

• The goal proves true.

4.6 ? student(anna, engineering).

Query type: ground.

Output: true.

Steps:

• Unify the query with the head of the rule:

```
student(X, Y) := field(Z, Y), student(X, Z).
```

- Instantiate anna as X, engineering as Y.
- Resolve to 2 new queries:

```
field(Z, engineering), student(anna, Z).
```

- Both goals are evaluated.
- Both goals prove true with unification of:

field(hit_transfer, engineering), student(anna, hit_transfer).

4.7? student(X, engineering).

Query type: non-ground. Output:

```
X = anna ;
X = daniel ;
X = adrian ;
false.
```

Steps:

• Unify the query with the head of the rule:

```
student(X, Y) :- field(Z, Y), student(X, Z).
```

- Instantiate engineering to Y.
- Resolve to 2 goals:

```
field(Z, engineering), student(X, Z).
```

• The 2 goals are unified with:

```
field(mechanical, engineering), student(X, mechanical).
field(computer, engineering), student(X, computer).
```

• In each of the goal, there are further unifications and instantiations. The final is as belowed:

```
field(mechanical,engineering), student(anna, mechanical).
field(mechanical,engineering), student(daniel, mechanical).
field(computer,engineering), student(adrian, computer).
```

4.8 ? student(X,fine-arts), course(fine_arts, Y).

Query type: non-ground. Output: false. Steps:

• Unify the first query with the head of rule:

```
student(X, Y) :- field(Z, Y), student(X, Z).
```

- Instantiate fine-arts to Y.
- Resolve into 2 goals:

```
field(Z, fine-arts), student(X, Z)
```

• Unify the first goal with the head of rule:

```
field(X, Y) :- course(X, Z), field(Z, Y).
```

- Instantiate Y with fine-arts.
- Resolve into 2 goals:

```
course(X, Z), field(Z, fine-arts).
```

- ...
- This is actually a recursion that will never end. Eventually, Prolog will result in *false* after exausting the knowledge base.

4.9 ? field(_,X).

Query type: non-ground. Output:

```
X = engineering;
X = engineering;
X = art;
X = social;
X = buisiness;
X = engineering;
X = engineering;
X = art;
X = social;
X = buisiness;
false.
```

Steps:

• Unify with facts respectively:

```
field(mechanical, engineering).
field(computer, engineering).
field(fine-arts, art).
field(literature, social).
field(management, buisiness).
```

- Instantiate X to engineering.
- Instantiate X to engineering.
- Instantiate X to art.
- Instantiate X to social.
- Instantiate X to buisiness.
- Resolve
- Unify with the head of rule:

```
field(X, Y) := course(X, Z), field(Z, Y).
```

- Instantiate X to all listed courses (hit_transfer, web_design, design_methods, poetry, leadership, biology respectively).
- Resolve into 2 goals:

```
course(X, Z), field(Z, Y).
```

• Evaluate both goals. The possible results are:

• Output of this part is as followed:

```
X = engineering;
X = engineering;
X = art;
X = social;
X = buisiness;
```

4.10 ? $lab_number(_,X)$, field(X,Y).

```
\label{eq:Query type: non-ground.} \mathbf{Query \ type:} \quad \mathrm{non\text{-}ground.}
```

Output: false.

Steps:

• Unify with the facts respectively:

```
lab_number(mechanical,15).
lab_number(fine_arts,10).
```

- This will **instantiate** X with a number, either 10 or 15.
- However, both will prove false in the second query when field(X, Y) is instantiated and resolved to either field(15, Y) or field(10, Y).
- Thus, Prolog outputs false.

4.11 ? $lab_number(X,15)$, field(X,Y).

Query type: non-ground. Output:

```
X = mechanical,
Y = engineering;
X = hit_transfer,
Y = engineering;
false.
```

Steps:

• Unify the first goal with the fact:

```
lab_number(mechanical,15).
```

- Instantiate mechanical to X.
- Unify the second goal with the fact:

```
field(mechanical, engineering).
```

• Get output: X = mechanical; Y = engineering

Steps when unify with rule:

• Unify with the head of rules:

```
lab_number(X, Z) :-course(X, Y), lab_number(Y, Z).
field(X, Y) :- course(X, Z), field(Z, Y).
```

- Instantiate 15 to Z.
- Resolve into 2 goals:

```
course(X, Y), lab_number(Y, 15).
course(X, Z), field(Z, Y).
```

• Evaluate both goals results in unification as followed:

```
course(hit_transfer, mechanical), lab_number(mechanical, 15),
    field(mechanical, engineering).
```

• Thus, the output for this evaluation is: $X = hit_transfer$, Y = engineering.

4.12 ? $\operatorname{student}(X)$, !, $\operatorname{student}(X, _)$.

Query type: non-ground. Output:

```
X = anna ;
X = anna ;
false.
```

Steps:

• Unify with the head of rule:

```
student(X):- student(X,_).
```

• Resolve the goal:

```
student(X, _).
```

- Evaluate the goal and found the first result of: X = Anna.
- Encounter!, stop evaluating.
- Continue to the thrid goal, Instantiate X with Anna from the first goal.
- Unify with the head of rule:

```
student(X, Y) := field(Z, Y), student(X, Z).
```

- Instantiate X with Anna.
- Resolve into 2 goals:

```
field(Z, Y), student(Anna, Z).
```

• Evaluate the goals, results in X = Anna proves true.

4.13 ? student(X), $student(X, _)$, !.

Query type: non-ground. Output:

```
X = anna.
```

• Unify with the rule:

```
student(X):- student(X,_).
```

• Resolve into 1 goal:

```
student(X, _).
```

- Evaluate and unify with student(anna, hit_transfer)
- Instantiate X with anna. Produce output: X = anna.
- Reached the cut operator, stop.

4.14 ? $course(X, _)$, $+ student(_, X)$.

Query type: non-ground. Output:

X = biology.

Steps:

• Unify the first query firstly with the fact:

course(hit_transfer, mechanical).

- Instantiate X as hit_transfer,
- Resolve the querry to be true.
- The second query will be unified with the fact:

student(anna, hit_transfer).

- the goals prove false. Backtracking..
- Unify the first query firstly with the fact:

course(web_design, computer).

- Instantiate X as web_design,
- Resolve the querry to be true.
- The second query will be unified with the fact:

student(adrian, web_design).

- the goals prove false. Backtracking..
- Unify the first query firstly with the fact:

course(design_methods, fine-arts).

- Instantiate X as design_methods.
- Resolve the querry to be true.
- The second query will be unified with the fact:

student(ava, design_methods).

• the goals prove false. Backtracking..

• Unify the first query firstly with the fact:

course(poetry, literature).

- Instantiate X as poetry.
- Resolve the querry to be true.
- The second query will be unified with the fact:

student(jack, poetry).

- the goals prove false. Backtracking..
- Unify the first query firstly with the fact:

course(leadership, management).

- Instantiate X as leadership.
- Resolve the querry to be true.
- The second query will be unified with the fact:

student(lee, leadership).

- the goals prove false. Backtracking..
- Unify the first query firstly with the fact:

course(biology, medicin).

- Instantiate X as biology.
- Resolve the querry to be true.
- The second query could not be unified with the any fact, produce false.
- Not false = true.
- the goals prove true. Hence, output: X = biology.

5 Question 8

5.1 ? magic(Hermione)

In this non ground query, we will find all atoms such that they satisfy magic (Hermione). First, magic (Hermione) is unified with the head of the rule $\operatorname{magic}(X)$:- house_elf(X). So X is instantiated to Hermione. Through resolution, we go to the body of the rule and house_elf(Hermione) is now our goal to satisfy and it is treated as a new non ground query. We have house_elf(dobby) in our knowledge base so house_elf(Hermione) unifies with house_elf(dobby) and Hermione is instantiated to dobby so dobby satisfies our goal and thus it is one answer to the initial query.

There are no more answers for house_elf(Hermione), so using backtracking, we reach the next rule for magic(X). magic(Hermione) is unified with the head of the rule magic(X):- wizard(X). So X is instantiated to Hermione. Then in the resolution step, wizard(Hermione) is now our goal to satisfy. We have wizard(dobby) in our knowledge base so wizard_elf(Hermione) unifies with wizard(dobby) and Hermione is instantiated to dobby so dobby is another answer of the query.

There are no more answers for wizard(Hermione), so using backtracking, we reach the next rule for magic(X). magic(Hermione) is unified with the head of the rule magic(X):- witch(X). X is instantiated to Hermione. Then in the resolution step, we make a transition to the body of the rule and witch(Hermione) is now our goal to satisfy. We have witch(hermione), witch(mcGonagall) and witch(rita_skeeter) in our knowledge base so Hermione is first instantiated hermione and when we continue the search, there are other answers mcGonagall and rita_skeeter to the non groud query witch(Hermione). Through backtracking, we can see that there is no other answer for our initial query. The final answer to the query is:

```
Hermione = dobby;
Hermione = dobby;
Hermione = hermione;
Hermione = mcGonagall;
Hermione = rita_skeeter;
```

5.2 ? magic(hermione)

This is a ground query, so we expect either true or false as an answer. First, Prolog searches the database from top to bottom. It matches the query with head of the rule magic(X):- house_elf(X). So X is instantiated to hermione. Now our goal is house_elf(hermione) (through resolution) and we treat it as a new ground query. We can see that house_elf(hermione) cannot match any clause in our database so we go to the next rule through backtracking: magic(X):- wizard(X). Now X is instantiated to hermione. In the resolution step, we make a transition to the body of the rule and our goal is to satisfy (find) wizard(hermione). We can clearly see in the database that hermione is a wizard so the result of our

goal query (?- wizard(hermione).) is true. Therefore we can tell that the initial query also returns true.

6 Question 11

The prolog file is in the file $\it question\mbox{-}11.pl$ Run query:

lucas(N, X)

to call the rule and execute with:

- N: the desired index of the lucas sequence.
- X: the return list.