Lớp: Phân tích và thiết kế thuật toán – CS112.N22

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Bài học: Design and Analysis of Algorithms (Độ phức tạp thuật toán không đệ quy)

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Ex1:

Câu 1:

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//3  \begin{aligned} & \text{sum = 0;} \\ & \text{for (i = 0; i < n; i++)} \\ & \text{for (j = i + 1; j <= n; j++)} \\ & \text{for (k = 1; k < 10; k++)} \\ & \text{sum = sum + i * j * k;} \end{aligned}   T_1 = O(1)   T_5 = O(1)   T_4 = O(1)   T_{45} = T_4 * T_5 = O(1) * O(1) = O(1)   T_{3} = O(n)
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$$\begin{split} T_{345} &= T_3 * T_{45} = O(n) * O(1) = O(n) \\ T_2 &= O(n) \\ T_{2345} &= T_2 * T_{345} = O(n) * O(n) = O(n^2) \\ T_{12345} &= T_1 + T_{2345} = O(1) + O(n^2) = O(n^2) \\ \hline Câu 2: \\ //5 \\ sum &= 0; \\ thisSum &= 0; \\ for & (i = 0; i < n; i++) \ \{ \\ thisSum > sum \} \\ sum &= thisSum; \\ else \\ thisSum &= sum; \\ \} \\ \hline T_1 &= O(1) \\ T_2 &= O(1) \\ T_{12} &= T_1 + T_2 = O(1) + O(1) = O(1) \\ \hline T_6 &= O(1) \\ T_8 &= O(1) \\ \hline T_5 &= O(1) \\ \hline T_4 &= O(1) \\ \hline T_4 &= O(1) \\ \hline T_4 &= O(1) \\ \hline T_{3678} &= T_4 + T_{5678} = O(1) + O(1) = O(1) \\ \hline T_3 &= O(n) \\ \hline T_{3456789} &= T_3 * T_{45678} = O(n) * O(1) = O(n) \\ \hline \end{split}$$

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T_{123456789} = T_{12} + T_{3456789} = O(1) + O(n) = O(n)
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Ex2:

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Câu 1:
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```
//3
sum = 0;
i = 1;
    while(i<=n) {
        j = n-i;
        while(j<=i) {
            sum = sum+j;
            j=j+1;
        }
        i=i+1;
}</pre>
```

1: 1 gán

2: 1 gán

3: n + 1 so sánh

4: n gán

5:
$$\left[\sum_{i=1}^{n} 0 \text{ if } n-i > i \text{ else } i-(n-i)+1\right] + n \text{ so sánh}$$

6:
$$\sum_{i=1}^{n} 0 \text{ if } n-i > i \text{ else } i-(n-i)+1 \text{ gán}$$

7:
$$\sum_{i=1}^{n} 0 \text{ if } n-i > i \text{ else } i-(n-i)+1 \text{ gán}$$

8: không tính('}')

9: n gán

10: không tính('}')

$$\begin{split} &\sum_{i=1}^{n} 0 \ if \ n-i > i \ else \ i-(n-i)+1 \\ &= \sum_{i=1}^{n} 2i-n+1 \ if \ i \geq \frac{n}{2} \\ &= \sum_{i=\frac{n}{2}}^{n} 2i-n+1 \ if \ n \% \ 2 == 0 \ else \sum_{i=\frac{n+1}{2}}^{n} 2i-n+1 \\ &= \frac{\binom{n}{2}+1)*(2*n+\frac{n}{2}*2)}{2} - (\frac{n}{2}+1)*n+(\frac{n}{2}+1) \ if \ n \% \ 2 == 0 \\ &\text{Else} \frac{\binom{(n+1)}{2}*[2*(n+1)+(\frac{n-1}{2})*2]}{2} - (\frac{n+1}{2})*n+(\frac{n+1}{2}) \\ &= \frac{n^2}{4}+n+1 \ if \ n\%2 == 0 \ else \frac{n^2}{4}+n+\frac{3}{4} \\ &\text{Gán}(n) = 2n+2+2*(\frac{n^2}{4}+n+1 \ if \ n\%2 == 0 \ else \frac{n^2}{4}+n+\frac{3}{4}) \\ &= \frac{n^2}{2}+4n+4 \ if \ n\%2 == 0 \ else \frac{n^2}{2}+4n+\frac{7}{2} \\ &\text{So sánh}(n) = n+1+(\frac{n^2}{4}+n+1 \ if \ n\%2 == 0 \ else \frac{n^2}{4}+n+\frac{3}{4})+n \\ &= \frac{n^2}{4}+3n+2 \ if \ n\%2 == 0 \ else \frac{n^2}{4}+3n+\frac{7}{4} \\ &\text{T}(n) = \text{Gán}(n) + \text{So Sánh}(n) \\ &= (\frac{n^2}{2}+4n+4 \ if \ n\%2 == 0 \ else \frac{n^2}{2}+4n+\frac{7}{2}) \\ &+ (\frac{n^2}{4}+3n+2 \ if \ n\%2 == 0 \ else \frac{n^2}{4}+3n+\frac{7}{4}) \\ &= \frac{3n^2}{4}+7n+6 \ if \ n\%2 == 0 \ else \frac{3n^2}{4}+7n+\frac{21}{4} \\ &\Rightarrow \text{T}(n) = \text{O}(n^2) \end{split}$$

Câu 2:

```
//4
          5 = 0;
          i =1;
          while(i<=n){
          j = 1;
                    while(j<=i*i) {
                              j = j + 1;
               i = i + 1;
          }
1: 1 gán
2: 1 gán
3: n + 1 so sánh
4: n gán
5: (\sum_{i=1}^{n} i^2) + n so sánh
6: \sum_{i=1}^{n} i^2 gán
7: \sum_{i=1}^{n} i^2 gán
8: không tính('}')
9: n gán
10: không tính('}')
\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6} = \frac{n^3}{3} + \frac{n^2}{3} + \frac{n}{6}
Gán(n) = 2n + 2 + 2*(\frac{n^3}{3} + \frac{n^2}{3} + \frac{n}{6}) = \frac{2n^3}{3} + n^2 + \frac{7n}{3} + 2
So sanh(n) = n+1 + (\frac{n^3}{3} + \frac{n^2}{2} + \frac{n}{6}) + n = \frac{n^3}{3} + \frac{n^2}{2} + \frac{13n}{6} + 1
```

T(n) = Gán(n) + So sánh(n)

$$= \left(\frac{2n^3}{3} + n^2 + \frac{7n}{3} + 2\right) + \left(\frac{n^3}{3} + \frac{n^2}{2} + \frac{13n}{6} + 1\right)$$

$$= n^3 + \frac{3n^2}{2} + \frac{9n}{2} + 3$$

$$\Rightarrow T(n) = O(n^3)$$

Câu 3:

```
//5
sum = 0;
i = 1;
while (i<=n) {
    j = n - i*i;
    while (j<=i*i) {
        sum = sum + i*j;
        j=j+1;
    }
    i=i+1;
}</pre>
```

- 1: 1 gán
- 2: 1 gán
- 3: n + 1 so sánh
- 4: n gán

5:
$$(\sum_{i=1}^{n} i^2 - (n - i^2) + 1 if i^2 \ge n - i^2) + n$$
 so sánh

6:
$$\sum_{i=1}^{n} i^2 - (n - i^2) + 1 if i^2 \ge n - i^2$$
 gán

7:
$$\sum_{i=1}^{n} i^2 - (n - i^2) + 1 if i^2 \ge n - i^2 gán$$

8: không tính('}')

9: n gán

10: không tính('}')

$$\sum_{i=1}^{n} i^{2} - (n - i^{2}) + 1 if i^{2} \ge n - i^{2}$$

$$\begin{split} &= \sum_{i=1}^{n} \ 2i^{2} - n \ + 1 \ if \ i \geq \ \sqrt[2]{\frac{n}{2}} \\ &= \sum_{i=1}^{n} \ 2i^{2} - n \ + \ 1 - \sum_{i=1}^{\sqrt[3]{n/2}} \ 2i^{2} - n \ + \ 1 \\ &= \ 2 * \left(\frac{n^{3}}{3} + \frac{n^{2}}{2} + \frac{n}{6}\right) - \ 2 * \left(\frac{\sqrt[2]{n/2}^{3}}{3} + \frac{\sqrt[2]{n/2}^{2}}{2} + \frac{\sqrt[2]{n/2}}{6}\right) \\ &- (n - \sqrt[2]{n/2})(n - 1) \\ &= \frac{2n^{3}}{3} + \frac{2n^{1.5}}{3\sqrt{2}} + \frac{5n}{6} - \frac{4n^{0.5}}{3\sqrt{2}} \\ &G{\acute{a}}{n}(n) = 2n + 2 + 2 * \left(\frac{2n^{3}}{3} + \frac{2n^{1.5}}{3\sqrt{2}} + \frac{5n}{6} - \frac{4n^{0.5}}{3\sqrt{2}}\right) \\ &= \frac{4n^{3}}{3} + \frac{4n^{1.5}}{3\sqrt{2}} + \frac{11n}{3} - \frac{8n^{0.5}}{3\sqrt{2}} + 2 \\ &So \ s{\acute{a}}{n}{n}(n) = n + 1 + \left(\frac{2n^{3}}{3} + \frac{2n^{1.5}}{3\sqrt{2}} + \frac{5n}{6} - \frac{4n^{0.5}}{3\sqrt{2}}\right) + n \\ &= \frac{2n^{3}}{3} + \frac{2n^{1.5}}{3\sqrt{2}} + \frac{17n}{6} - \frac{4n^{0.5}}{3\sqrt{2}} + 1 \\ &T(n) = G{\acute{a}}{n}(n) - So \ s{\acute{a}}{n}{n}(n) \\ &= \left(\frac{4n^{3}}{3} + \frac{4n^{1.5}}{3\sqrt{2}} + \frac{11n}{3} - \frac{8n^{0.5}}{3\sqrt{2}} + 2\right) + \left(\frac{2n^{3}}{3} + \frac{2n^{1.5}}{3\sqrt{2}} + \frac{17n}{6} - \frac{4n^{0.5}}{3\sqrt{2}} + 1\right) \\ &= 2n^{3} + \frac{2n^{1.5}}{\sqrt{2}} + \frac{13n}{2} - \frac{4n^{0.5}}{\sqrt{2}} + 3 \end{split}$$

 \Rightarrow T(n) = O(n^3)