Lexical Analysis

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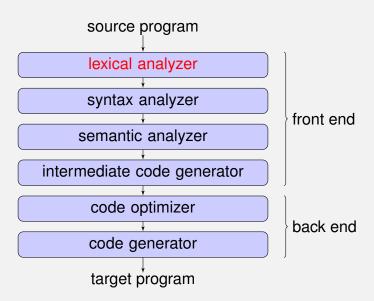
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Outline

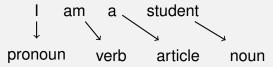
- Introduction
- 2 Roles
- Implementation
- Use ANTLR to generate Lexer

Compilation Phases



Lexical Analysis

- Like a word extractor in ⇒ in ⇒ in
- Like a spell checker
 l ogog to socholsochol
- Like a classification



Lexical Analysis Roles

- Identify lexemes: substrings of the source program that belong to a grammar unit
- Return tokens: a lexical category of lexemes
- Ignore spaces such as blank, newline, tab
- Record the **position** of tokens that are used in next phases

Example on Lexeme and Token

r	е	S	u		t	, ,	=	, ,	0	ldsum - value / 100;
---	---	---	---	--	---	-----	---	-----	---	----------------------

Kind of Tokens
IDENT
ASSIGN_OP
IDENT
SUBSTRACT_OP
IDENT
DIV_OP
INT_LIT
SEMICOLON

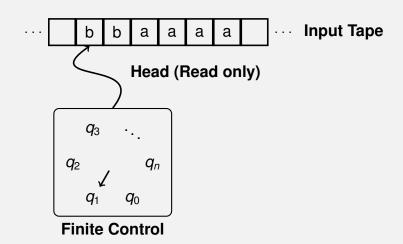
How to build a lexical analyzer?

- How to build a lexical analysis for English?
 - 65000 words
 - Simply build a dictionary: {(I,pronoun);(We,pronoun);(am,verb);...}
 - Extract, search, compare
- But for a programming language?
 - How many words?
 - Identifiers: abc, cab, Abc, aBc, cAb, ...
 - Integers: 1, 10, 120, 20, 210, ...
 - ...
 - Too many words to build a dictionary, so how?
 - Apply rules for each kind of word (token)

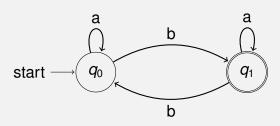
Rule Representations

- Finite Automata
 - Deterministic Finite Automata
 - Nondeterministic Finite Automata
- Regular Expressions

Finite Automata



State Diagram



Input: abaabb

Current state	Read	New State
q_0	а	q_0
q_0	b	q_1
$oldsymbol{q}_1$	а	q_1
$oldsymbol{q}_1$	а	q_1
q_1	b	q_0
q_0	b	q_1

Deterministic Finite Automata

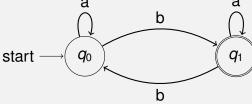
Definition

Deterministic Finite Automaton(DFA) is a 5-tuple $M = (K, \Sigma, \delta, s, F)$ where

- K = a finite set of state
- Σ = alphabet
- $s \in K = the initial state$
- $F \subseteq K$ = the set of final states
- δ = a transition function from K $\times \Sigma$ to K

$$\begin{array}{ll} \mathsf{M}=(\mathsf{K},\Sigma,\delta,\mathsf{s},\mathsf{F})\\ \mathsf{where}\;\mathsf{K}=\{q_0,\,q_1\} & \Sigma=\{\mathsf{a},\mathsf{b}\} & \mathsf{s}\!=\!q_0 & \mathsf{F}\!=\!\{q_1\}\\ \mathsf{and}\;\delta & \end{array}$$

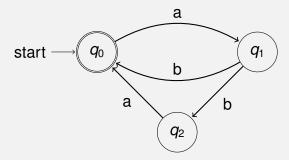
K	Σ	$\delta(K,\Sigma)$			
q_0	a	q_0			
$ \begin{array}{c} q_0 \\ q_0 \\ q_1 \\ q_1 \end{array} $	b	q_1			
q_1	а	<i>q</i> ₁ <i>q</i> ₀			
q_1	b	q_0			
a a b					



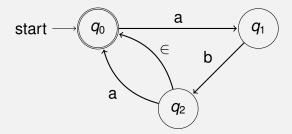
Nondeterministic Finite Automata

- Permit several possible "next states" for a given combination of current state and input symbol
- Accept the empty string \in in state diagram
- Help simplifying the description of automata
- Every NFA is equivalent to a DFA

Language L = $({ab} \cup {aba})^*$



Language $L = (\{ab\} \cup \{aba\})^*$



Regular Expression (regex)

- Describe regular sets of strings
- Symbols other than () | * stand for themselves
- Use \in for an empty string
- Concatenation α β = First part matches α , second part β
- Union $\alpha \mid \beta$ = Match α or β
- Kleene star α^* = 0 or more matches of α
- Use () for grouping

RE		Language
0	=>	{ 0 }
01	=>	{ 01 }
0 1	=>	{0,1}
0(0 1)	=>	{00,01}
(0 1)(0 1)	=>	{00,01,10,11}
0*	=>	$\{\in,0,00,000,0000,\}$
(0 1)*	=>	$\{\in,0,1,00,01,10,11,000,001,\}$

(i|I)(f|F)

Keyword if of language Pascal

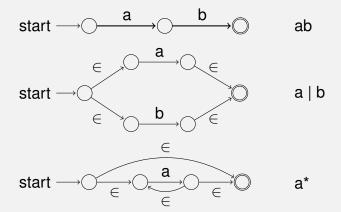
- if
- IF
- If
- iF

E(0|1|2|3|4|5|6|7|8|9)*

An E followed by a (possibly empty) sequence of digits

- E123
- E9
- E

Regular Expression and Finite Automata



Convenience Notation

- α + = one or more (i.e. $\alpha\alpha*$)
- α ? = 0 or 1 (i.e. $(\alpha | \in)$)
- [xyz]= x|y|z
- [x-y]= all characters from x to y, e.g. [0-9] = all ASCII digits
- [^x-y]= all characters other than [x-y]
- matches any character

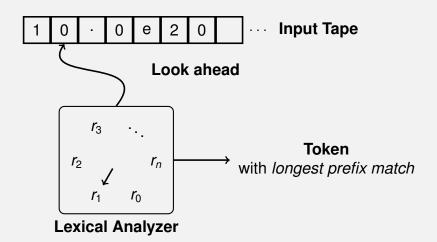
```
\begin{array}{lll} (0|1|2|3|4) & => & [0-4] \\ (a|g|h|m) & => & [aghm] \\ (0|1|2|3|4|5|6|7|8|9)(0|1|2|3|4|5|6|7|8|9)^* & => & [0-9]+ \\ (E|e)(+|-|\in)(0|1|2|3|4|5|6|7|8|9)+ & => & [Ee][+-]?[0-9]+ \end{array}
```

ANTLR [1]

- ANother Tool for Language Recognition
- Terence Parr, Professor of CS at the Uni. San Francisco
- powerful parser/lexer generator

```
/**
 * Filename: Hello.g4
 */
lexer grammar Hello:
// match any digits
INT: [0-9]+:
// Hexadecimal number
HEX: 0[Xx][0-9A-Fa-f]+:
// match lower-case identifiers
ID : [a-z]+ ;
// skip spaces, tabs, newlines
WS: [ \t \r \n] + -> skip ;
```

Lexical Analyzer



Summary

- A lexical analyzer is a pattern matcher that isolates small-scale parts of a program
- Lexical rules are represented by Regular expressions or Finite Automata.
- How to write a lexical analyzer (lexer) in ANTLR

References I

[1] ANTLR, http:antlr.org, 19 08 2016.