

# Homework 6 - 5/21/2019

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- 1) 5.32 (you can use the Java wait-notify constructs to solve this problem)
- 2) 8.11
- 3) 8.21
- 4) 8.23

## Exercise 5.32

A file is to be shared among different processes, each of which has a unique number. The file can be accessed simultaneously by several processes, subject to the following constraint: the sum of all unique numbers associated with all the processes currently accessing the file must be less than  $n$ . Write a monitor to coordinate access to the file.

```
import java.io.*;
import java.util.*;
public class FileShare {
    int total, nSum;

    FileShare(int n)
    {
        this.nSum = n;
    }

    public synchronized void access(int i)
    {
        if(i + total >= nSum)
        {
            try{
                wait();
            }
            catch (InterruptedException e) {}
        }
        total += i;
    }

    public synchronized void release(int i){
        total -= i;
        notify();
    }

    public static void main(String[] args){
    }
}
```

## Exercise 8.11

Given six memory partitions of 300KB, 600KB, 350KB, 200KB, 750KB, and 125KB(in order), how would the first-fit, best-fit, and worst-fit algorithms place processes of size 115KB, 500KB, 358KB, 200KB ,and 375KB (in order)? Rank the algorithms in terms of how efficiently they use memory.

Ranking of Memory Efficiency:

Both First & Best Fit have same memory allocation, First Fit should have slightly better performance due to not crawling through each memory partition.

**1. First Fit:**

P1: 115 KB -> 300 KB  
P2: 500 KB -> 600 KB  
P3: 358 KB -> 750 KB  
P4: 200 KB -> 350 KB  
P5: 375 KB -> not able to allocate

**2. Best Fit:**

P1: 115 KB -> 125 KB  
P2: 500 KB -> 600 KB  
P3: 358 KB -> 750 KB  
P4: 200 KB -> 200 KB  
P5: 375 KB -> not able to allocate

**3. Worst Fit:**

P1: 115 KB -> 750 KB  
P2: 500 KB -> 600 KB  
P3: 358 KB -> not able to allocate  
P4: 200 KB -> 350 KB  
P5: 375 KB -> not able to allocate

**Exercise 8.21**

The BTV operating system has a 21-bit virtual address, yet on certain embedded devices, it has only a 16-bit physical address. It also has a 2-KB page size. How many entries are there in each of the following?

- A. A conventional, single-level page table

- 2KB page size =  $2^{11}$ , or 11 bits for page size. Using a 21 bit virtual address or  $2^{21}$  for 2MB.  
Subtract  $\frac{2^{21}}{2^{11}} = 2^{10}$  to get the remainder of  $2^{10}$  or 1024 entries in the page table .

- B. An inverted page table

- Using the physical address,  $\frac{2^{16}}{2^{11}} = 2^5$  or 32 entries.

**Exercise 8.23**

Consider a logical address space of 256 pages with a 4-KB page size, mapped onto a physical memory of 64 frames.

- A. How many bits are required in the logical address?

- 256 pages =  $2^8$  or 8 page bits, with a 4KB page size =  $2^{12}$ , thus we require  $2^{12} * 2^8 = 2^{20}$  addresses or 20 bits for the logical address.

- B. How many bits are required in the physical address?

- 64 frames =  $2^6$  or 6 frame bits, with a 4KB page size =  $2^{12}$ , thus we require  $2^{12} * 2^6 = 2^{18}$  addresses or 18 bits for the physical address.