

ALGEBRAIC FORMULATION AND GRAPHICAL REPRESENTATION

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Linear Programming (LP)

- Linearity in objective function and constraints
- Continuous decision variables

$$\begin{array}{ll}\max & x_1 + x_2 \\ \text{st.} & \\ & x_1 \leq 5 \\ & x_2 \leq 4\end{array}$$

$$\begin{array}{ll}\min & x_1 + x_2 \\ \text{st.} & \\ & x_1 \geq 1 \\ & x_2 \geq 2\end{array}$$

$$\begin{array}{ll}\max & x_1 + x_2 \\ \text{st.} & \\ & x_1 \geq 1 \\ & x_2 \geq 2\end{array}$$

Unbounded

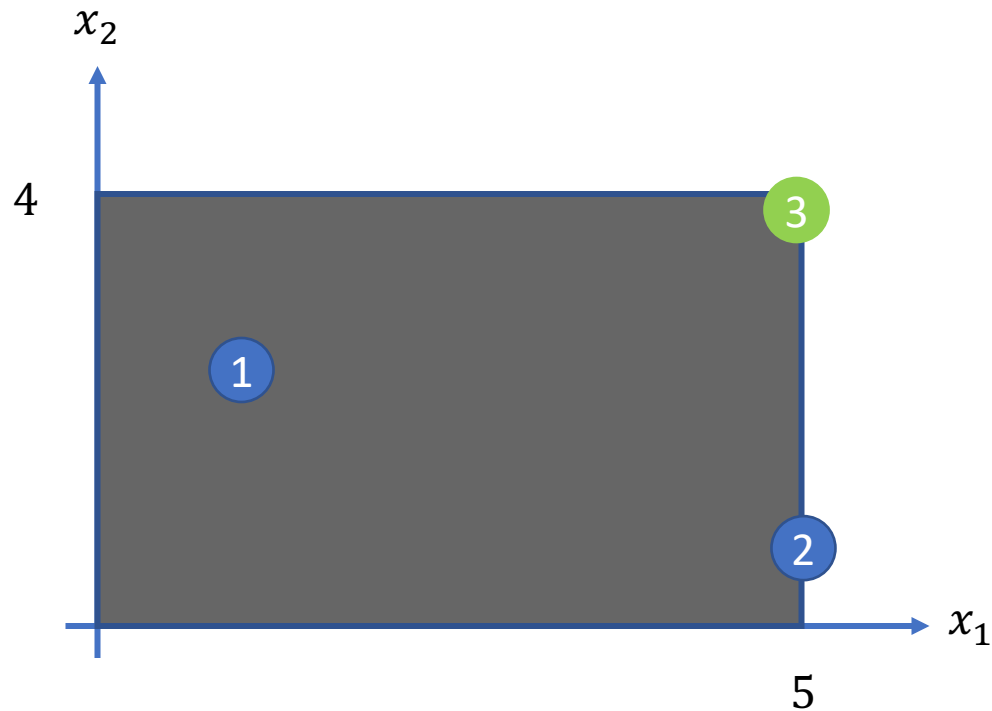
$$\begin{array}{ll}\min & x_1 \\ \text{st.} & \\ & x_1 \geq 1 \\ & x_2 \geq 2\end{array}$$

Degenerate

$$\begin{array}{ll}\min & x_1 \\ \text{st.} & \\ & x_1 \leq 1 \\ & x_1 \geq 2\end{array}$$

Infeasible

Linear Programming



$$\begin{aligned} \max z &= x_1 + x_2 \\ \text{st.} \end{aligned}$$

$$x_1 \leq 5$$

$$x_2 \leq 4$$

$$x_i \geq 0 \quad \forall i$$

Solution 1:

$$x_1 = 1.2$$

$$x_2 = 2.9$$

$$z = 4.1$$

Feasible
Sub-Optimum

Solution 2:

$$x_1 = 5$$

$$x_2 = .9$$

$$z = 5.9$$

Feasible
Sub-Optimum

Solution 3:

$$x_1 = 5$$

$$x_2 = 4$$

$$z = 9$$

Optimum

Mixed Integer Linear Programming (MILP)

- Linearity in objective function and constraints
- Continuous and integer/binary decision variables

$$\max z = x_1 + x_2$$

st.

$$x_1 \leq 3.3$$

$$x_2 \leq 3$$

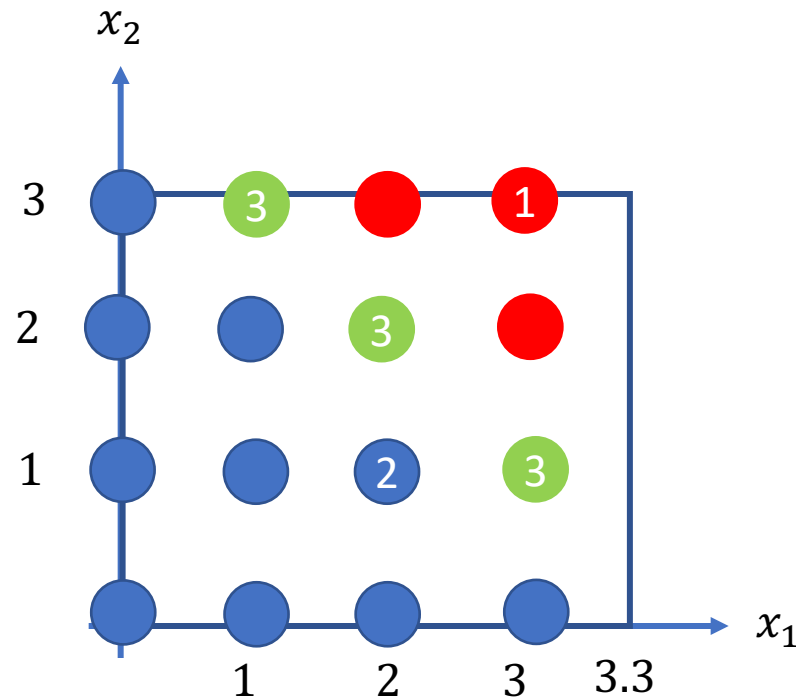
$$x_1 + x_2 \leq 4$$

$$x_i \geq 0 \quad \forall i$$

$$x_i \in \mathbb{I} \quad \forall i$$

Mixed Integer Linear Programming

$$\begin{aligned} \max z &= x_1 + x_2 \\ \text{st.} \\ x_1 &\leq 3.3 \\ x_2 &\leq 3 \\ x_1 + x_2 &\leq 4 \\ x_i &\geq 0 \quad \forall i \\ x_i &\in \mathbb{I} \quad \forall i \end{aligned}$$



Solution 1:

$$\begin{aligned} x_1 &= 3 \\ x_2 &= 3 \\ x_1 + x_2 &= 6 \\ z &= 6 \end{aligned}$$

Infeasible

Solution 2:

$$\begin{aligned} x_1 &= 2 \\ x_2 &= 1 \\ x_1 + x_2 &= 3 \\ z &= 3 \end{aligned}$$

Sub-Optimum

Solutions 3:

$$\begin{aligned} x_1 + x_2 &= 4 \\ z &= 4 \end{aligned}$$

Optimum &
Degenerate

Goal Programming

- Objective is to achieve a certain target value
- Minimize penalty associated of not achieving the value

$\min pen$
st.
 $pen \geq 10 - x$
 $pen \geq x - 10$
 $pen \geq 0$

Goal = 10

Solution 1:
if $x = 9$
then $pen = 1$

Sub-Optimum

Solution 2:
if $x = 12$
then $pen = 2$

Sub-Optimum

Solution 3:
if $x = 10$
then $pen = 0$

Optimum

Multicriteria Programming

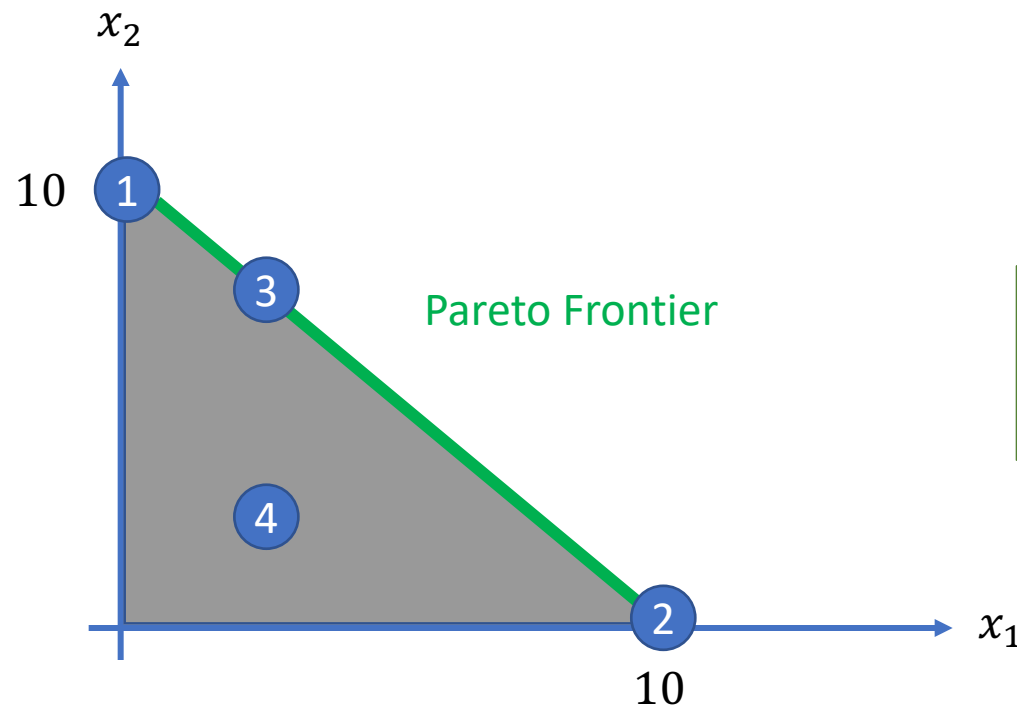
- Optimize several objectives at the same time

$$\begin{array}{ll} \max & x_1 \\ \max & x_2 \\ \text{st.} & \\ & x_1 + x_2 \leq 10 \\ & x_i \geq 0, \forall i \end{array}$$

Multicriteria Programming

$\max x_1$
 $\max x_2$
st.

$$x_1 + x_2 \leq 10$$
$$x_i \geq 0, \forall i$$



Solution 1:
 $x_1 = 0$
 $x_2 = 10$

Pareto
Optimum

Solution 2:
 $x_1 = 10$
 $x_2 = 0$

Pareto
Optimum

Solution 3:
 $x_1 = 8$
 $x_2 = 2$

Pareto
Optimum

Solution 4:
 $x_1 = 2$
 $x_2 = 2$

Dominated

Non-Linear Programming

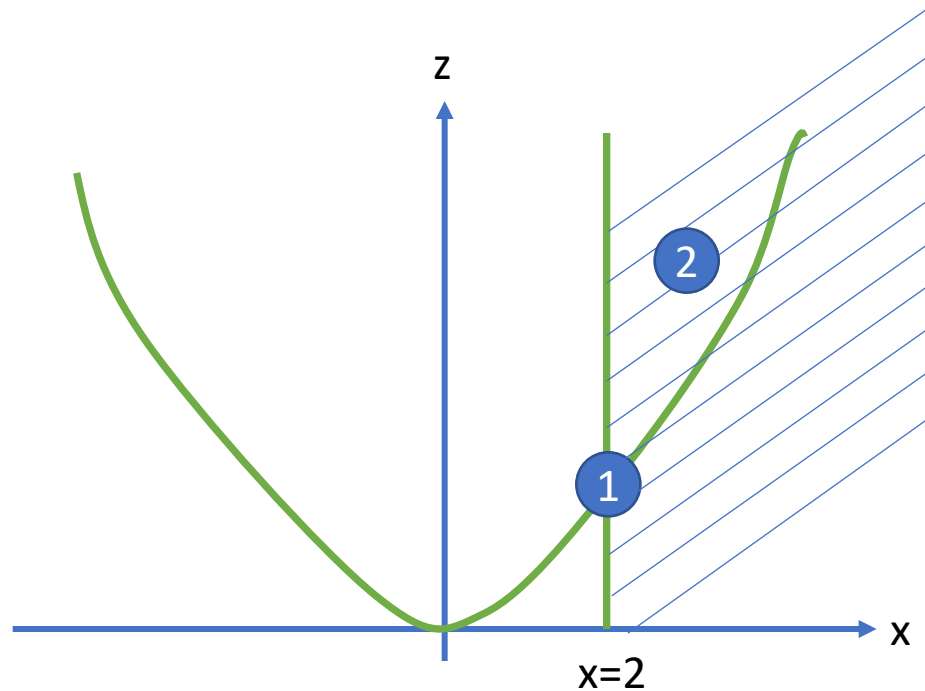
- Nonlinearity either in objective function or in constraints

$$\begin{array}{ll}\min z = x^2 \\ \text{st.} \\ x \geq 2\end{array}$$

Non-Linear Programming

$$\begin{array}{ll} \min z = x^2 \\ \text{st.} \\ x \geq 2 \end{array}$$

- Nonlinearity either in objective function or in constraints



Solution 1:
 $x = 2$
 $z = 4$

Optimum

Solution 2:
 $x = 2.5$
 $z = 6.25$

Feasible
Sub-Optimum

Mathematical Formulation General Form - LP

$$\max(\min) \quad c_1x_1 + c_2x_2 + \cdots + c_nx_n$$

s. t.

$$a_{11}x_1 + a_{21}x_2 + \cdots + a_{n1}x_n \leq (\geq) b_1$$

$$a_{12}x_1 + a_{22}x_2 + \cdots + a_{n2}x_n \leq (\geq) b_2$$

$$\vdots$$

$$a_{1m}x_1 + a_{2m}x_2 + \cdots + a_{nm}x_n \leq (\geq) b_m$$

$$x_1 \geq 0$$

$$x_2 \geq 0$$

$$\vdots$$

$$x_n \geq 0$$

$$\max(\min) \quad \sum_{i=1}^n c_i x_i$$

s. t.

$$\sum_{i=1}^n a_{ij}x_i \leq (\geq) b_j \quad \forall j = 1 \dots m$$

$$x_i \geq 0 \quad \forall i = 1 \dots n$$

[Refresher on Linear Algebra](#)

Building Blocks for Problem Solving:

- Sets
- Summations

Mathematical Formulation General Form - MILP

$$\max(\min) \quad c_1x_1 + c_2x_2 + \dots + c_nx_n$$

s. t.

$$a_{11}x_1 + a_{21}x_2 + \dots + a_{n1}x_n \leq (\geq) b_1$$

$$a_{12}x_1 + a_{22}x_2 + \dots + a_{n2}x_n \leq (\geq) b_2$$

\vdots

$$a_{1m}x_1 + a_{2m}x_2 + \dots + a_{nm}x_n \leq (\geq) b_m$$

$$x_i \in \mathbb{I} \quad \forall i \in I$$

$$x_i \in \{0,1\} \quad \forall i \in B$$

$$x_i \geq 0 \quad \forall i \in C$$

$$\max(\min) \quad \sum_{i=1}^n c_i x_i$$

s. t.

$$\sum_{i=1}^n a_{ij}x_i \leq (\geq) b_j \quad \forall j = 1 \dots m$$

$$x_i \in \mathbb{I} \quad \forall i \in I$$

$$x_i \in \{0,1\} \quad \forall i \in B$$

$$x_i \geq 0 \quad \forall i \in C$$

Steps to Formulate Optimization Problems

