Probabilistic Languages

Some Models of Translation

- IBM Models 1-5
- Hidden Markov Model
- Phrase-Based Models

Q: What do all of these things have in common?

Definition 1: Natural language

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Gettin Stairtit

Walcome page — Writin Scots — Editin Lessons



Scotland

Aw the airts — Cultur — Economy — Fowk — Law — Leids — Politics



Applee'd sciences

Agricultur — Airchitectur — Communication — Electronics — Ingineerin — Heal — Industry — Medicine — Transport — Wather



Naitural Sciences an Maths

Astronomy - Biology - Chemistry - Computers - Yird science - Mathematics - Pheesics



Fowk an Social Studies

Anthropology — Airchaeology — Geography — Eddication — History — Leids — Philosophy — Psychology — Releegion — Sociology



Govrenment an Law

Economics — Govrenment — Law — Military — Politics



Airt an Cultur

Airt - Fuid & Farin — Cultur — Dance — Habbies — Media — Muisic — Gemmes & Sports — Theatre

Tends to not be well-defined.

Definition 1: Natural language





Tends to not be well-defined.

Definition 1: Formal language

Well-defined, so that a computer can process it: a (possibly infinite) set of strings.

- All of the English words in a dictionary.
- All sequences of any length over those words.
- All English sentences with non-zero $p(e \mid f)$ for some French sentence f, according to your model.

Desiderata

- We need efficient algorithms and data structures to:
 - Encode all of the strings in the language.
 - Assign probabilities to all of those strings.
 - Via products such as p(e)p(f|e).
 - Find the string with the highest probability.
 - Compute expectations over substrings.
 - Compute mappings between strings.

$$\mathcal{L}_1 = \left\{ \begin{array}{l} a \ a \ a \ a \\ a \ b \ a \\ a \ b \ b \end{array} \right\}$$

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$$\mathcal{L}_2 = a^* = \{a, aa, aaa, ...\}$$

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$$a \ a \ b$$

$$a \ b \ b$$

$$a \ b \ b$$

$$a \ b \ b$$

$$b \ a \ b$$

$$c \ a \ b$$

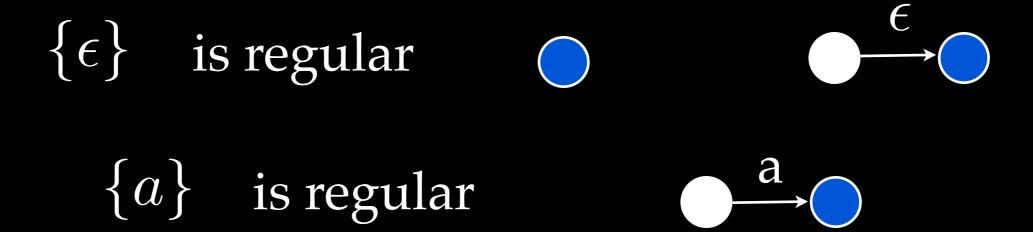
$$c \ a \ b$$

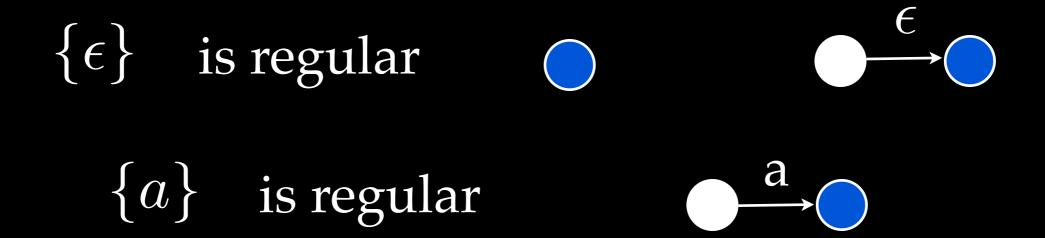
$$d \ b \ b$$

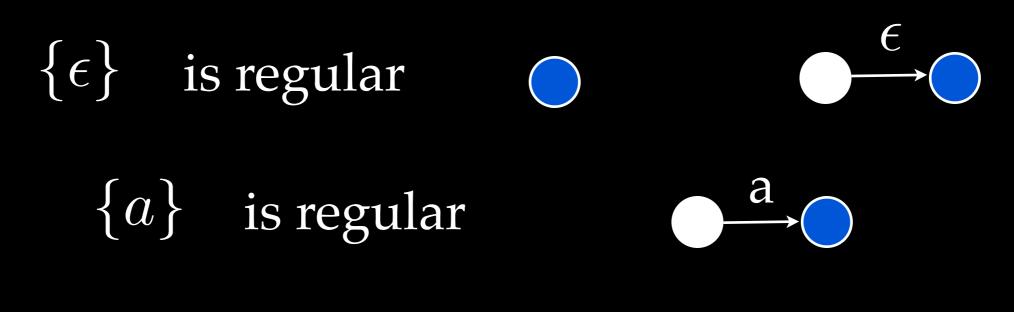
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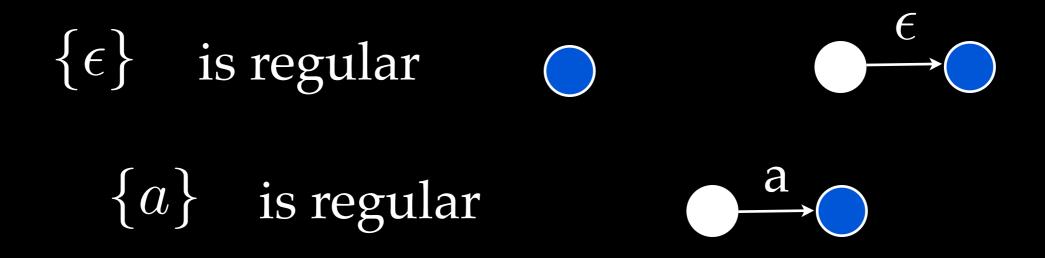
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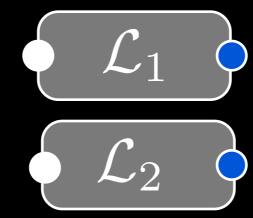


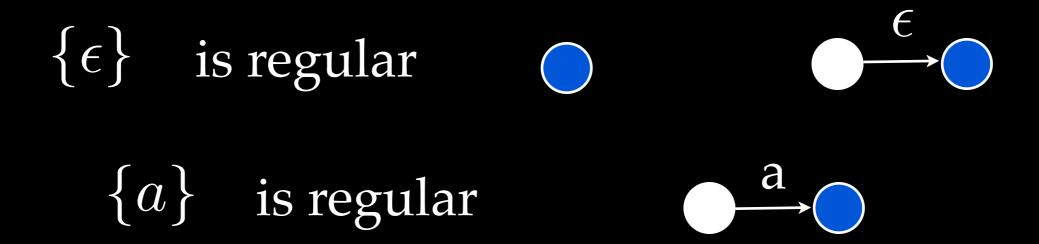


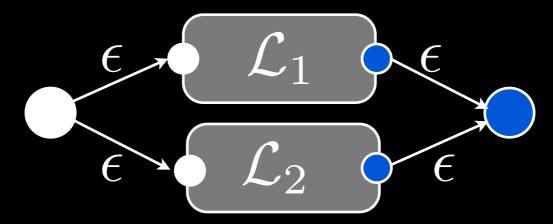




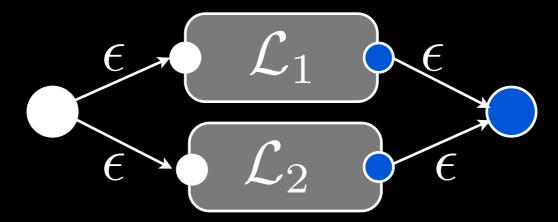




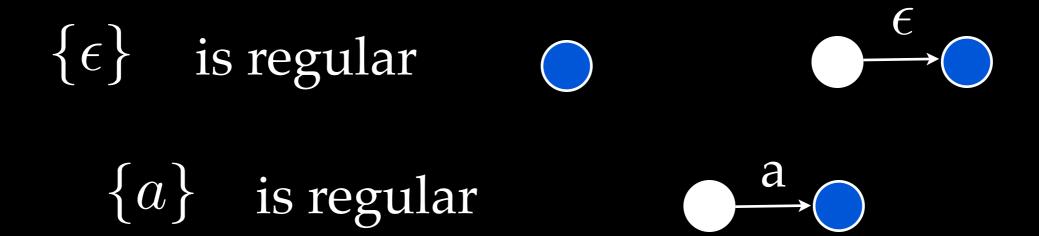


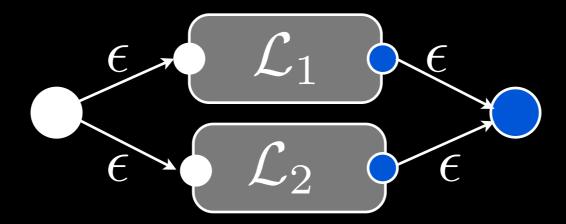


 $\mathcal{L}_1 \cup \mathcal{L}_2$ is regular if \mathcal{L}_1 and \mathcal{L}_2 are regular

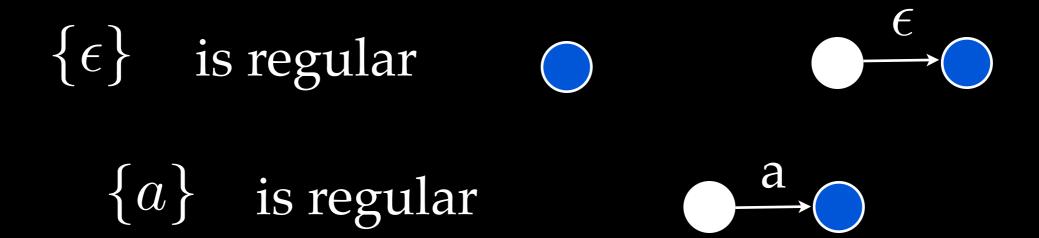


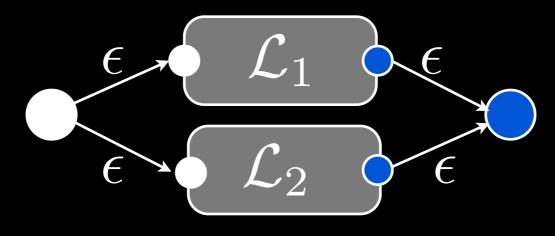
 $\mathcal{L}_1 \cdot \mathcal{L}_2$



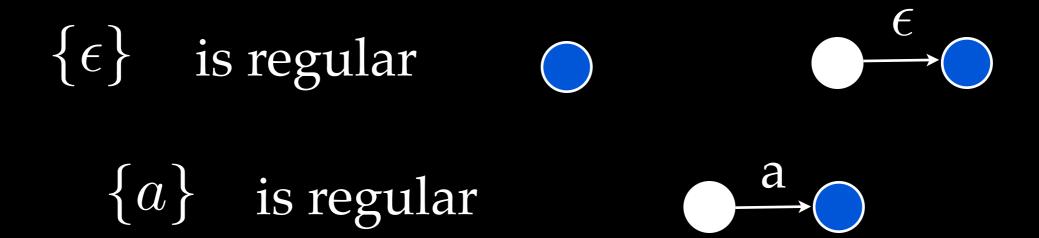


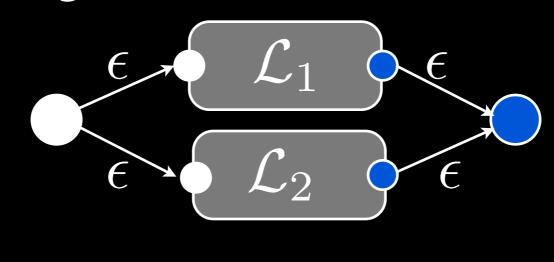
$$\mathcal{L}_1 \cdot \mathcal{L}_2 \bullet \mathcal{L}_1$$



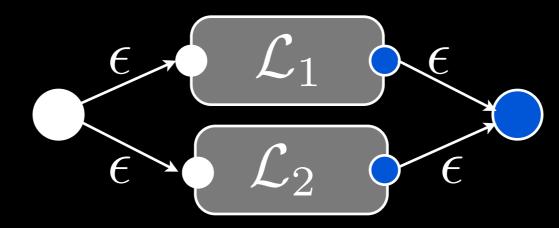


$$\mathcal{L}_1 \cdot \mathcal{L}_2 \bullet \mathcal{L}_1 \diamond \bullet \mathcal{L}_2 \diamond$$

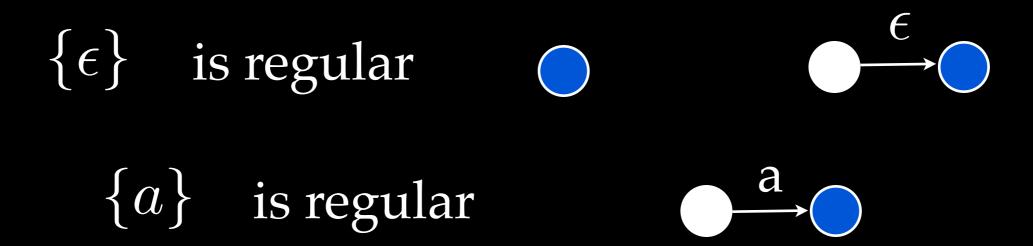


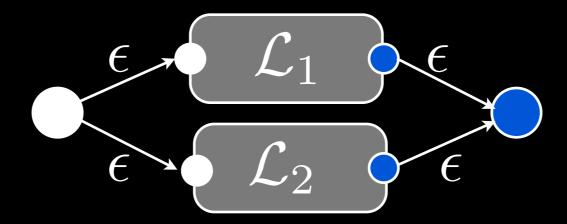






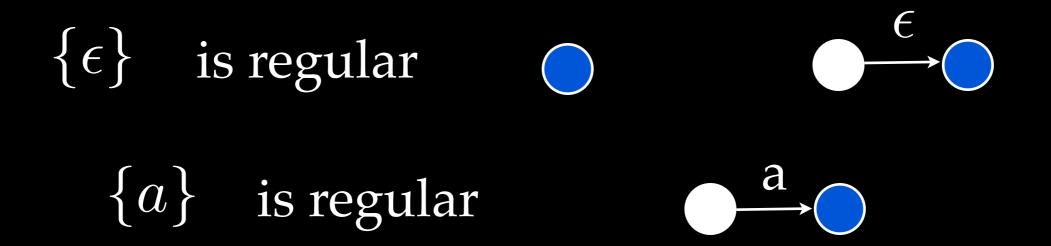
$$\mathcal{L}_1 \cdot \mathcal{L}_2 \bullet \mathcal{L}_1 \stackrel{\epsilon}{\triangleright} \bullet \mathcal{L}_2 \stackrel{\epsilon}{\triangleright}$$

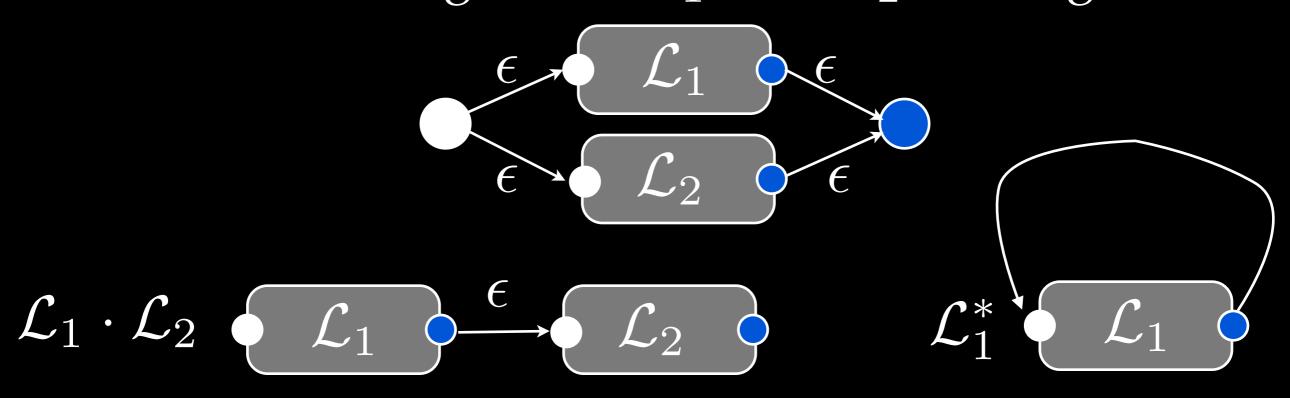




$$\mathcal{L}_1 \cdot \mathcal{L}_2 \bullet \mathcal{L}_1 \bullet \bullet \mathcal{L}_2 \bullet$$

$$\mathcal{L}_1 \bullet \mathcal{L}_1 \bullet \mathcal{L}_$$





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But not all languages are context-free, either!

We want a function:

$$f:\mathcal{L} \to \mathbb{R}^+$$

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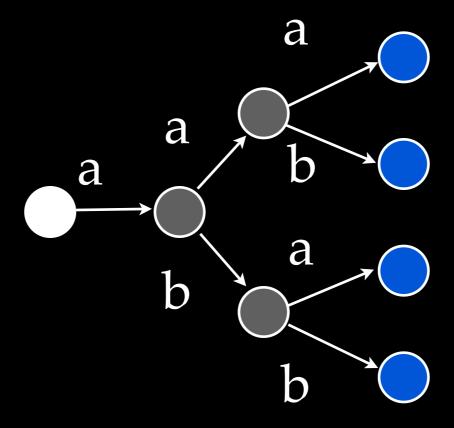
such that:

$$f(w) \in [0, 1]$$

$$\sum_{w} f(w) \in [0, 1]$$

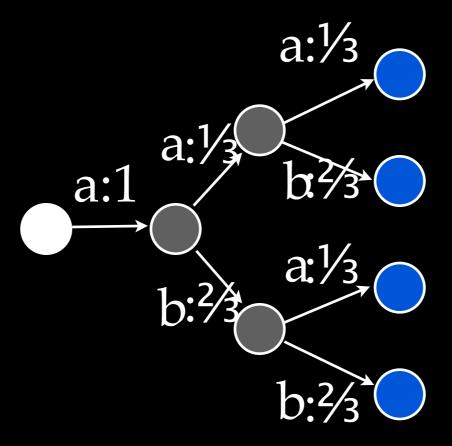
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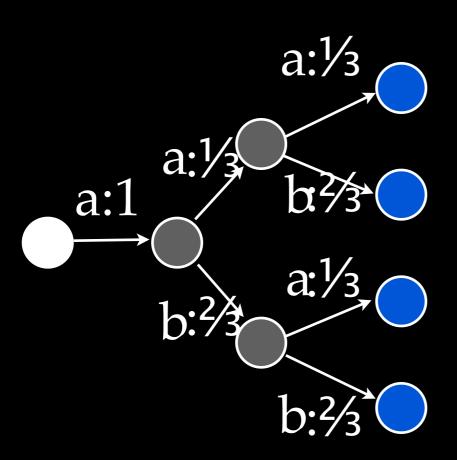


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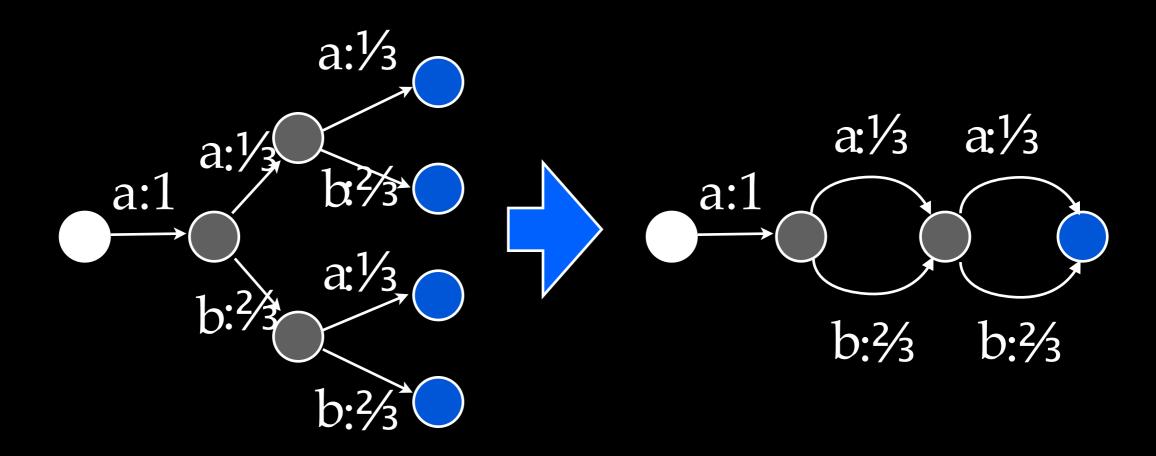
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Minimization



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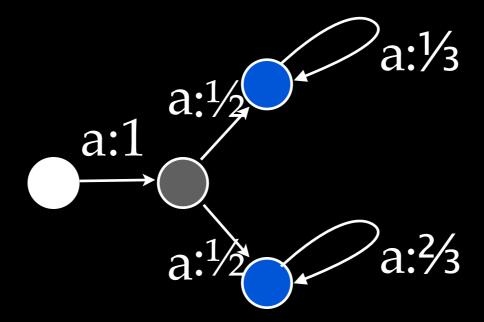


Other Algorithms

- Shortest path (e.g. Dijkstra, A*): most probable
- Determinization (not all can be determinized)
- Epsilon-removal
- Lazy composition (e.g. intersection): p(e)p(f|e)

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Practical Issues

- OpenFST (openfst.org)
 - Efficient C++ implementation.
 - Used in speech recognition (Google, Kaldi @ JHU)
 - Machine translation (JHU → Cambridge, Google)

Some Models of Translation

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Q: What do all of these things have in common?

A: They all define weighted regular languages over a set of output sentences. Details Thursday.

Questions

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Are natural languages regular?

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Does it matter for MT if they aren't?