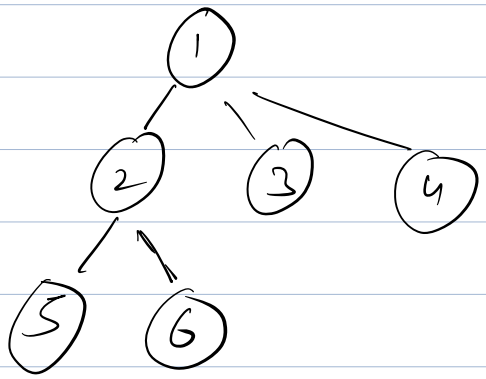
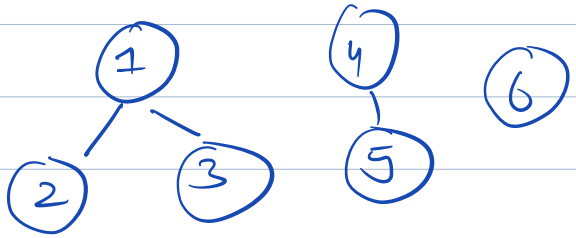


1, 2, 5, 6, 3, 4.

1, 2, 6, 5, 3, 4

1, 2, 6, 5, 4, 3

DFS

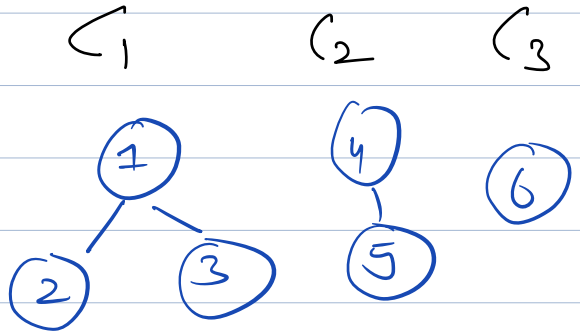


= 1 2 3 4 5 6

```
for (int i=1 ; i <= 6 ; i++) {
    if (!visited[i])
        dfs(i);
}
```

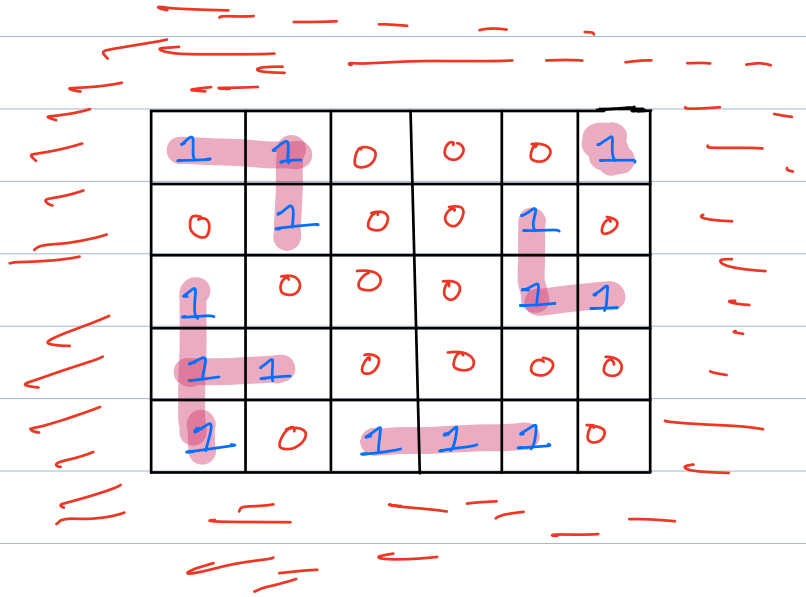
1	2	3	4	5	6
1	1	1	1	1	1

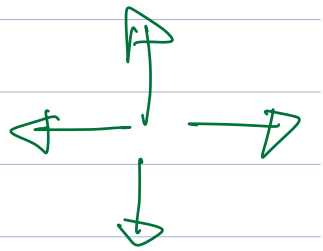
Q Count the number of connected components.




```
cnt = 0;

for (int i = 1 ; i <= 6 ; i++) {
    if (!visited[i]) {
        dfs(i);
        cnt++;
    }
}

return cnt
```

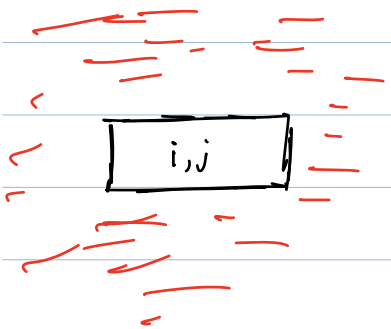
Number of Islands

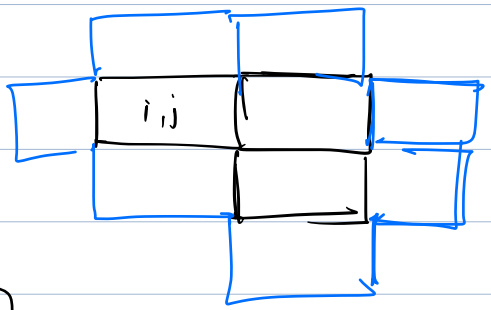


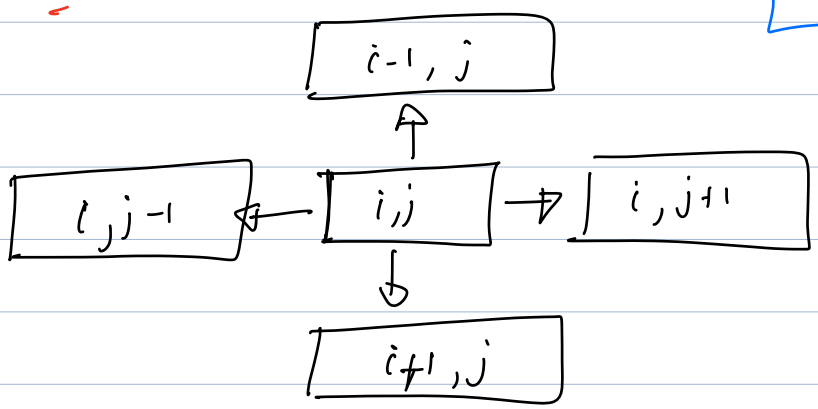


R, D, L, T

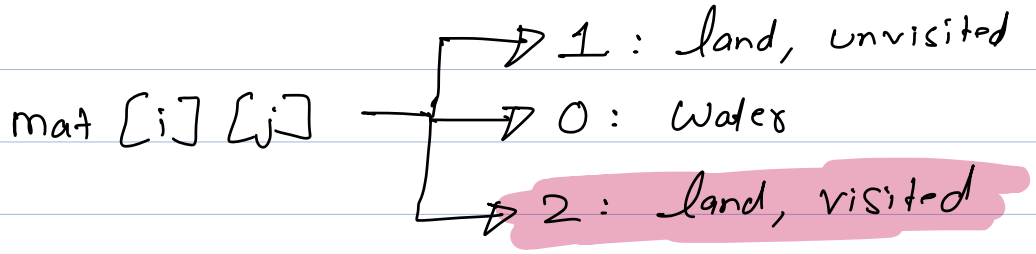
no of islands — 5





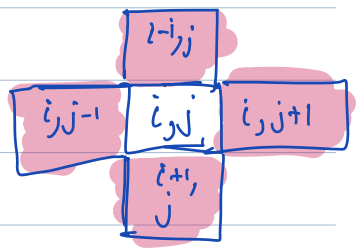


We will use the input 2D matrix to track visited nodes




```
for (int i = 0; i<n ; i++) {
    for (int j=0 ; j<m ; j++ ) {
        if ( mat [i] [j] == 1) {
            dfs (i,j)
            cnt++;
        }
    }
}

return cnt;

bool checkIfValid ( int i , int j, int n, int m) {
    if (i<0 || i≥n || j<0 || j≥m)
        return false
    return true;
}
```

```
void dfs (int i, int j) {

    mat[i][j] = 2;
    int x[4] = {1, -1, 0, 0}
    int y[4] = {0, 0, 1, -1};

    for (int k=0; k<4; k++) {
        int x_n => i + x[k];
        int y_n => j + y[k]

        if (checkIfValid(x_n, y_n) &&
                mat[x_n][y_n] == 1)
            dfs(x_n, y_n);

    }

}
```



$TC: O(V+E)$

$V = mn$

$E = 4mn.$

$TC: O(mn)$

$SC: O(mn)$

Rotten Oranges.

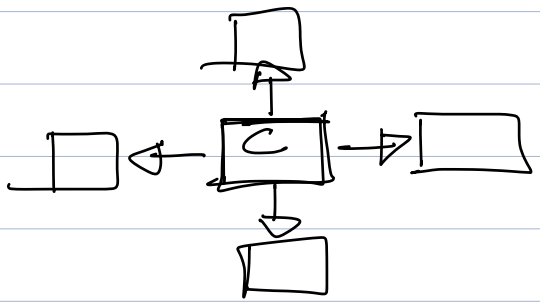
In how many days every living area is covid affected

C	1	1	0	1	0
1	1	0	1	1	0
0	0	1	C	0	1
C	1	1	0	1	1
1	1	0	0	1	C

C ⇒ affected by covid.

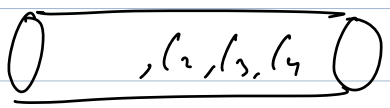
1 ⇒ not affected area

0 ⇒ no life present.



DAY 1

C	1	1	0	1	0
1	1	0	1	1	0
0	0	1	C	0	1
C	1	1	0	1	1
1	1	0	0	1	C



DAY 2

C	1	1	0	1	0
1	1	0	1	1	0
0	0	1	C	0	1
C	1	1	0	1	1
1	1	0	0	1	C

DAY 3

C	1	1	0	1	0
1	1	0	1	1	0
0	0	1	C	0	1
C	1	1	0	1	1
1	1	0	0	1	C

Ans ⇒ 3 days

C	1	1	0	1	0
1	1	0	1	1	0
0	0	1	C	0	1
C	1	1	0	1	1
1	1	0	0	1	C

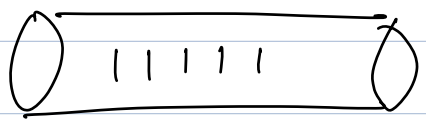
1) Traverse the matrix and insert covid cells in the Queue.

2) Perform BFS and maintain level.

Whenever from (i,j) you visit a neight (x,y)

level[x][y] = level[i][j] + 1

C	1	1	0	1	0
1	1	0	1	1	0
0	0	1	C	0	1
C	1	1	0	1	1
1	1	0	0	1	C

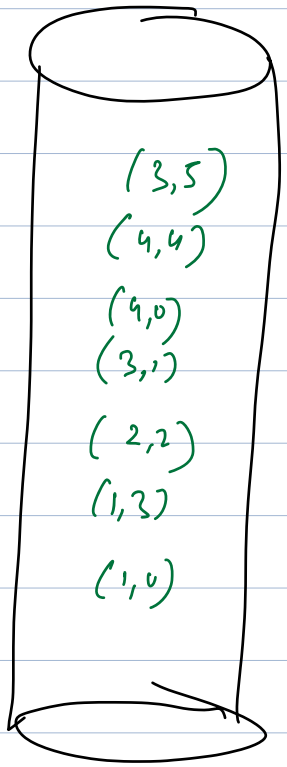


Size = 8

max-level = 2

10:32 pm

	0	1	2	3	4	5
0	C	1	1	0	1	0
1	1	1	0	1	1	0
2	0	0	1	C	0	1
3	C	1	1	0	1	1
4	1	1	0	0	1	C

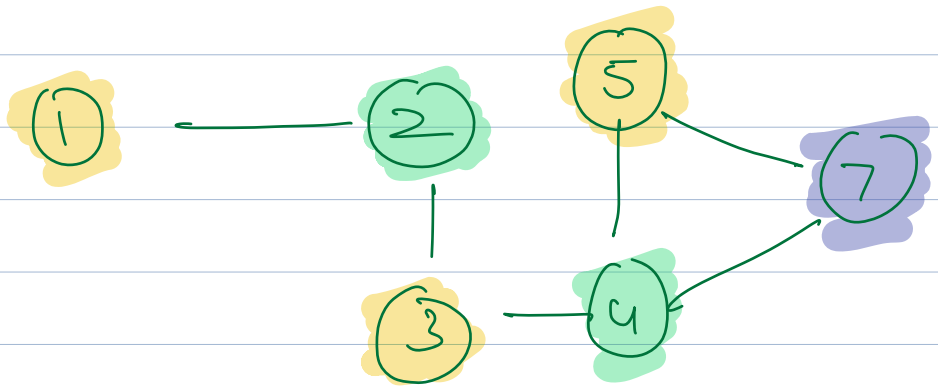


	0	1	2	3	4	5
0	0	1	2		3	
1	1	2		1	2	
2			1	0		2
3	0	1	2		2	1
4	1	2			1	0

(0,1)

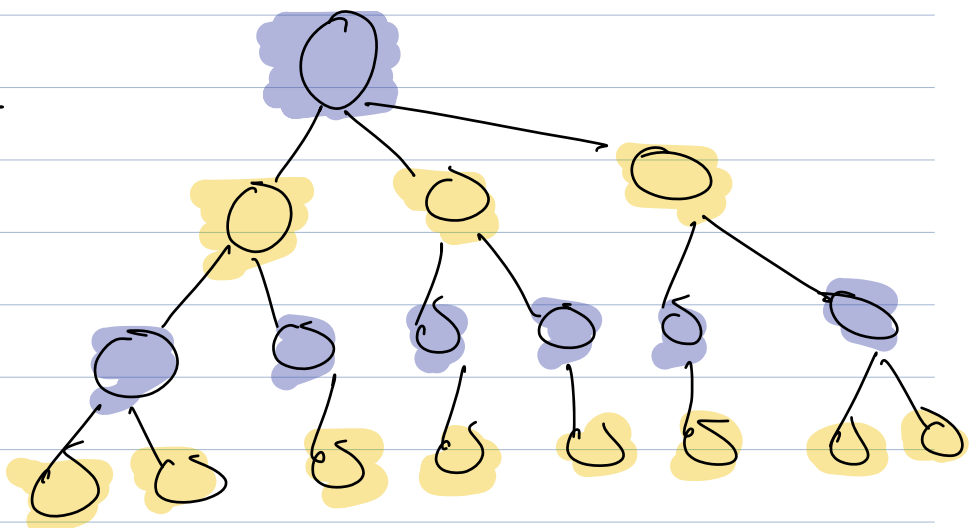
Chromatic numbers ⇒ np hard = [no polynomial time solution]

Min number of colors required to color the graph such that no 2 adjacent nodes have the same color.

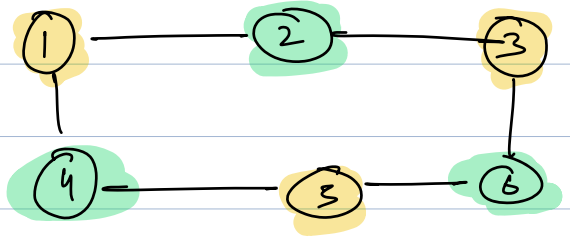


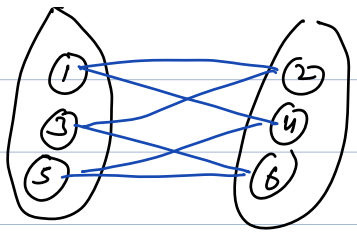
chromatic number = 3

Tree



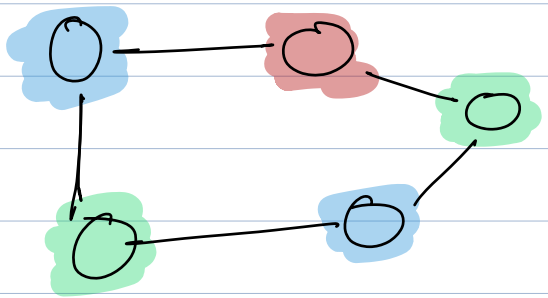
Chromatic number for a tree = 2





Even Cycles

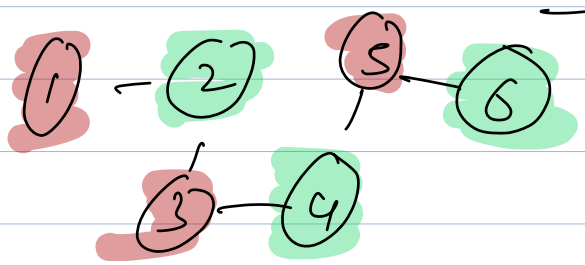
chromatic number = 2



Odd Cycles

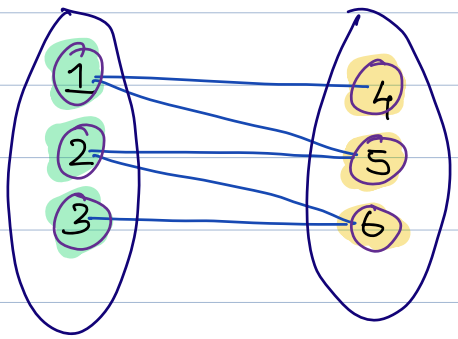
chromatic number = 3

For any acyclic graphs the chromatic number ≤ 2

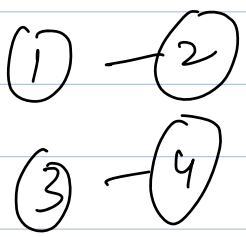


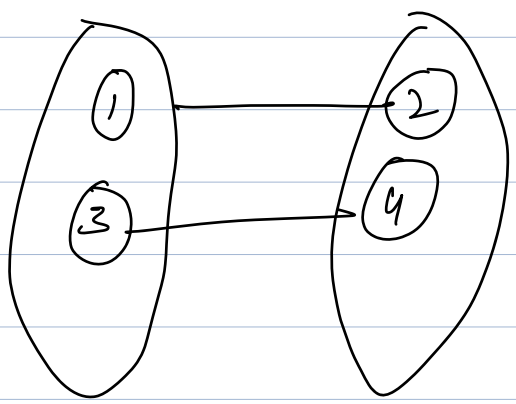
Bipartite graphs

A graph which can be broken down into two distinct sets U & V such that no two nodes in U & V have an edge b/w them.

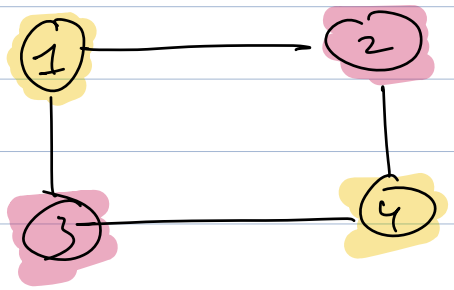


Chromatic number of a bipartite graph is 2.





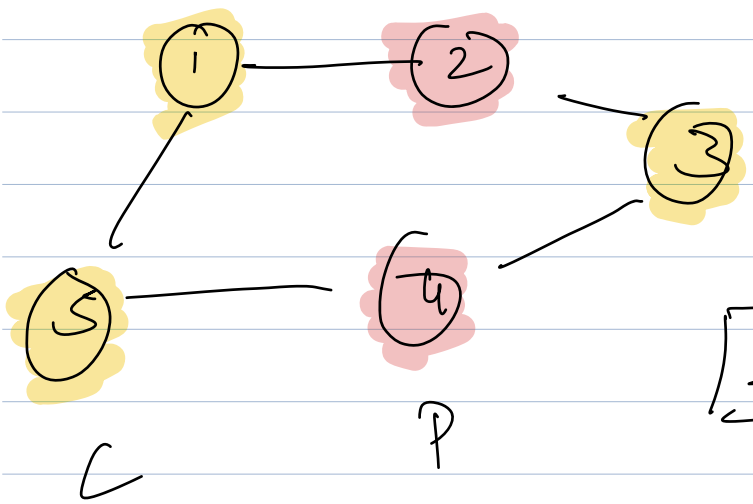
Even length cycles are bipartite.
Trees are bipartite.
Acyclic graphs are bipartite.



= 1

= 0

1	0	0	1
1	2	3	4



1	0	1	0	1
1	2	3	4	5

Pseudo Code : (1, 0) Colors → 0 → 1

```
int visited[n]; int color[n], isBipartite = true;

void dfs(int node, int c) {

    visited[node] = 1;
    color[node] = c;

    for(int i=0 ; i< adj[node].size() ; i++) {

        int child = adj[node][i];

        if(visited[child]) {
            if(color[child] == color[node])
                isBipartite = false;
        }else {
            dfs(child, 1-c)
        }
    }
}
```

0 → 1

1 → 0

1-c

1^c

TC : O(V+E)

SC : O(V+E)