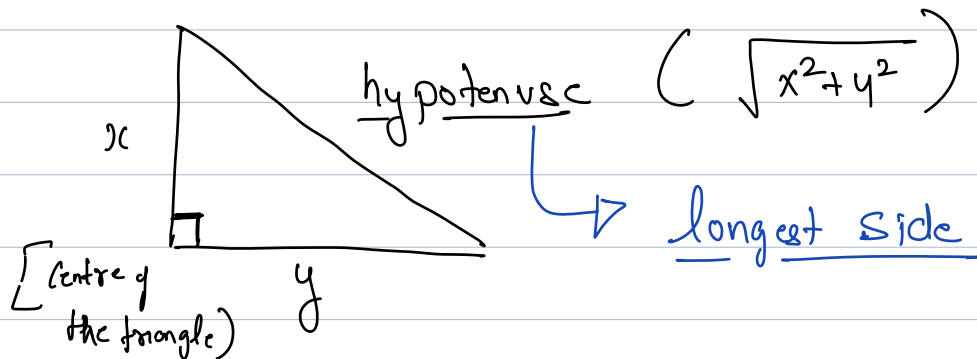
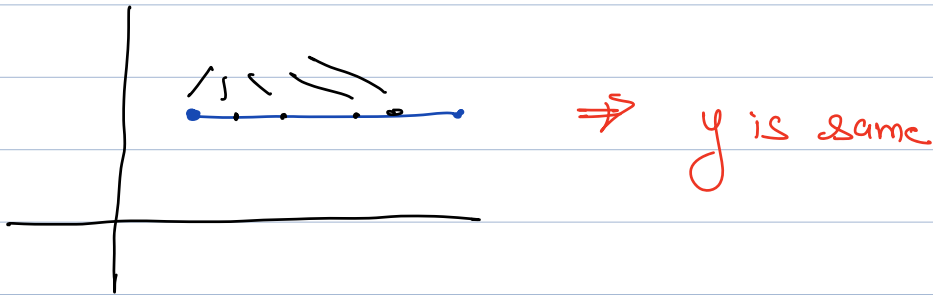


Right Angled Triangle Basics

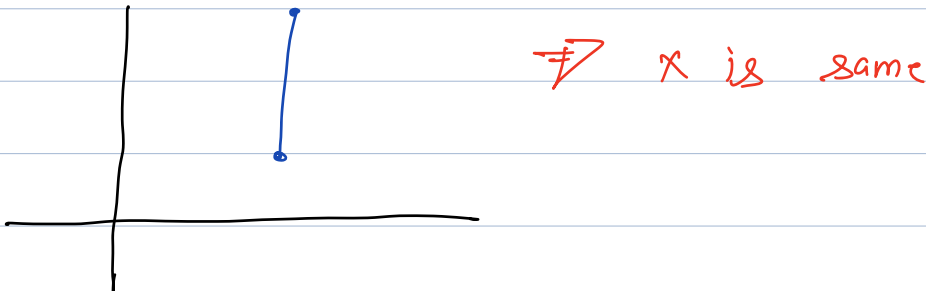
↳ one of the angles = 90°



Line parallel to x axis



Line parallel to y axis

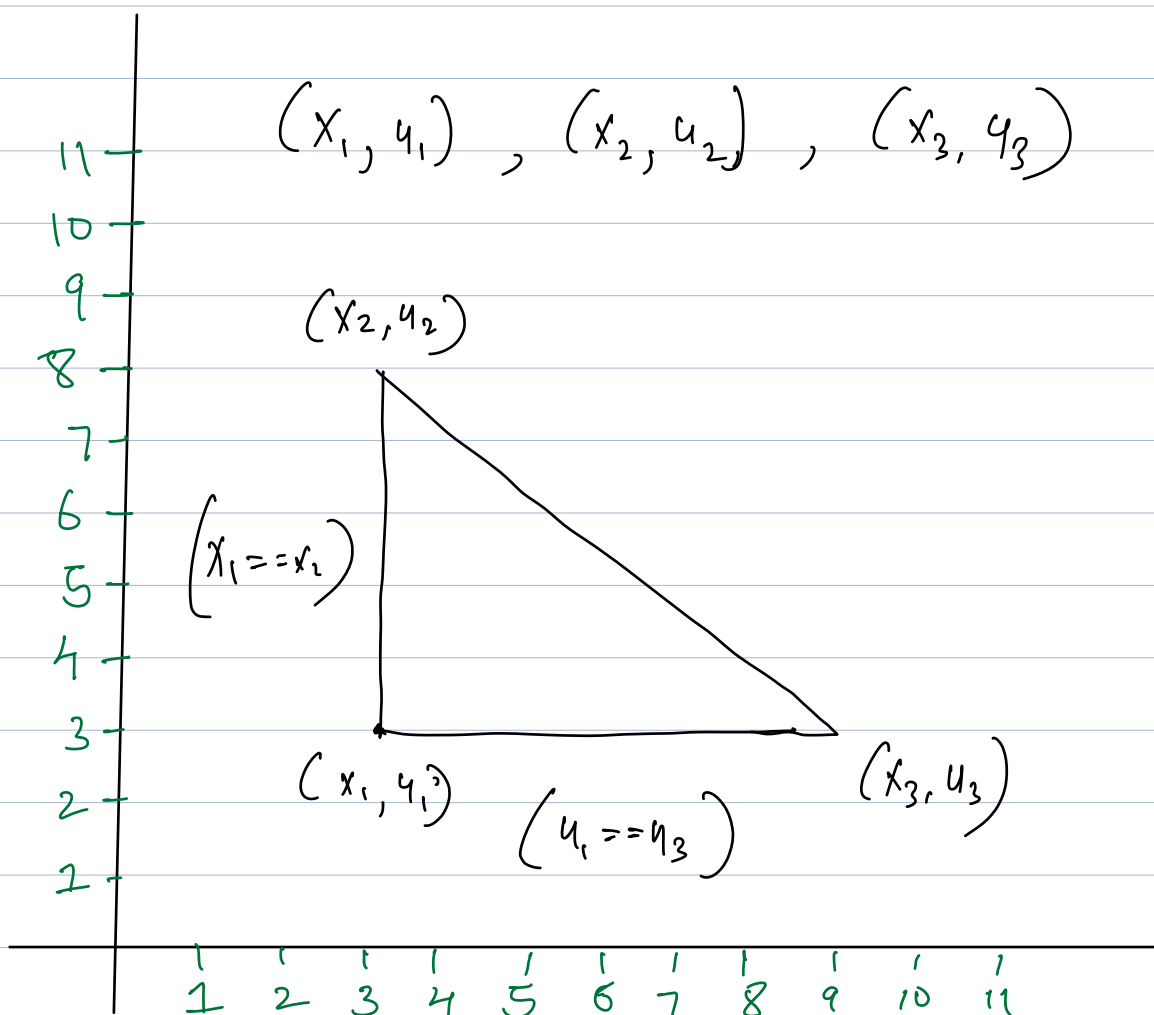


Q1) Given 3 distinct points in a 2D plane check. If they form right angled ^{triangle}, [such that shorter sides are parallel to x-axis & y-axis respectively].

$$\Sigma_{x1} : (1, 8) \quad (1, 4) \quad (5, 4) = \text{YES}$$

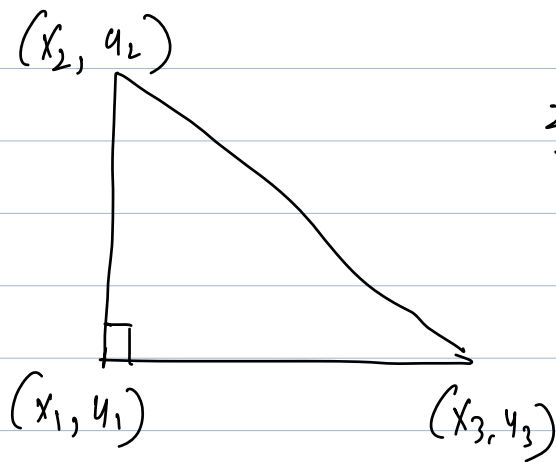
$$\Sigma_{x2} : (5, 10) \quad (1, 3) \quad (5, 3) = \text{YES}$$

$$\Sigma_{x3} : (4, 3) \quad (8, 3) \quad (4, 8) = \text{YES}$$

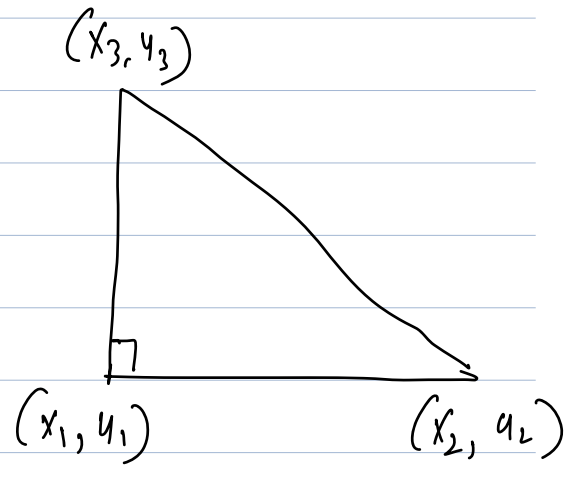


6 conditions

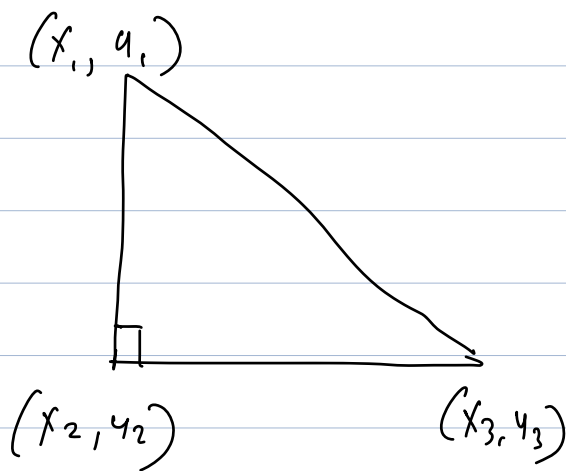
1)



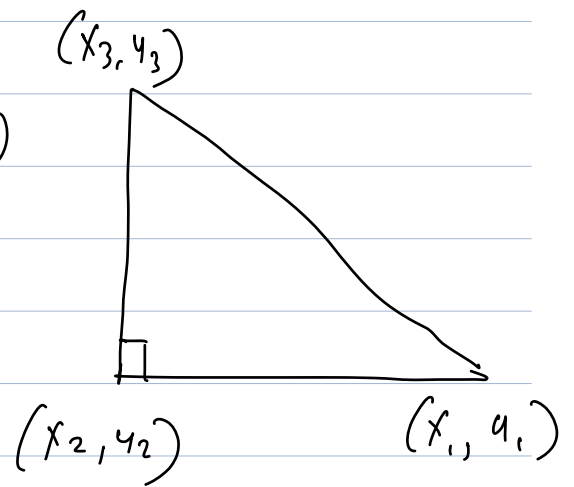
2)



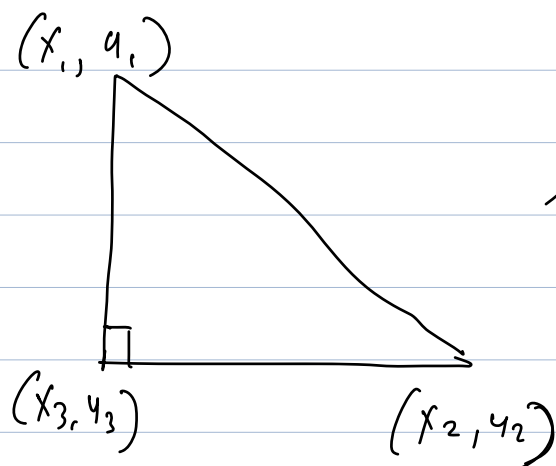
3)



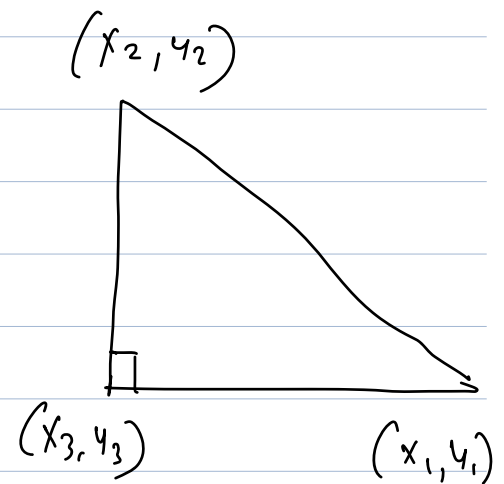
4)



5)



6)



Pseudo Code

Input : $(x_1, y_1), (x_2, y_2), (x_3, y_3)$

bool isRightAngleTriangle($x_1, y_1, x_2, y_2, x_3, y_3$) {

if ($x_1 == x_3$ && $y_1 == y_2$) {

return true;

} else if ($x_1 == x_2$ && $y_1 == y_3$) {

return true;

} else if ($x_2 == x_3$ && $y_2 == y_1$) {

return true;

} else if ($x_2 == x_1$ && $y_2 == y_3$) {

return true;

} else if ($x_3 == x_1$ && $y_3 == y_2$) {

return true;

} else if ($x_3 == x_2$ && $y_3 == y_1$) {

return true;

} else

return false;

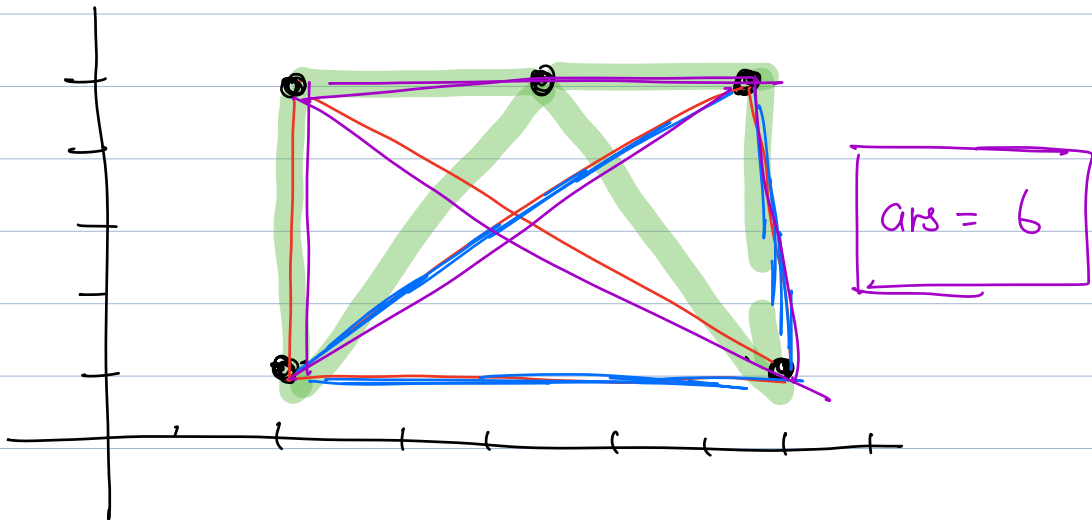
Tc: $O(1)$

Sc: $O(1)$

}

Q2) Given N distinct points in a 2D plane, calculate no of triangles that can be formed such that shorter sides are parallel to x -axis & y -axis.

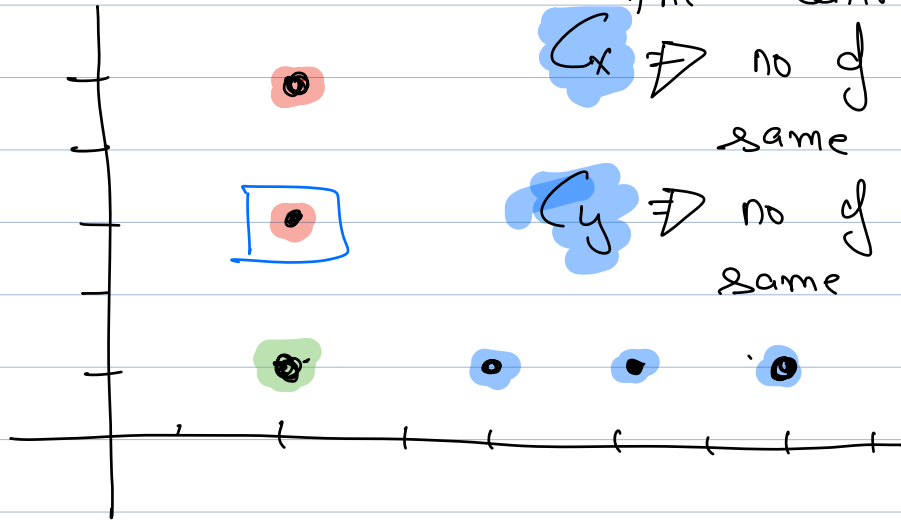
Note: Input will be in form of 2 arrays $x[n]$ & $y[n]$, where i^{th} coordinate = $(x[i], y[i])$



Brute Force: for all triplets, pass it through the code for the last question

$$Tc: O(n^3)$$

Assuming Green point to be the centre.



$C_x \Rightarrow$ no of points having same x axis

$C_y \Rightarrow$ no of points having same y axis.



Total Δ triangles having Green point as centre.

$$\Rightarrow \underline{\underline{(C_x - 1) \times (C_y - 1)}}$$

Note: There will be no duplicate right angled triangle as the centres would be different.

Pseudo Code

HashMap <int, int> x_count;

HashMap <int, int> y_count;

for (int i=0; i<n; i++) {

 x_count[x[i]]++;

 y_count[y[i]]++;

}

int ans = 0;

for (int i=0; i<n; i++) {

 int c_x = x_count[x[i]];

 int c_y = y_count[y[i]];

 ans = ans + (c_x - 1) * (c_y - 1);

}

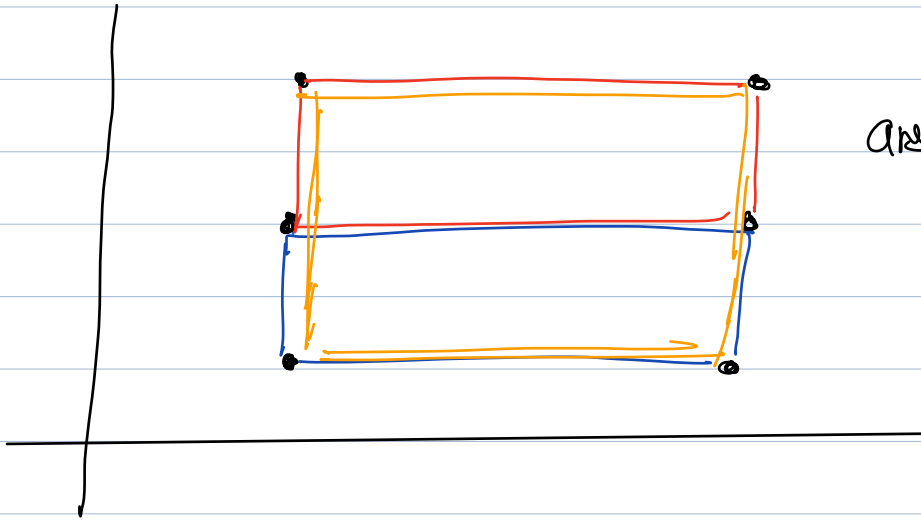
return ans;

Tc: $O(n)$

Sc: $O(n)$

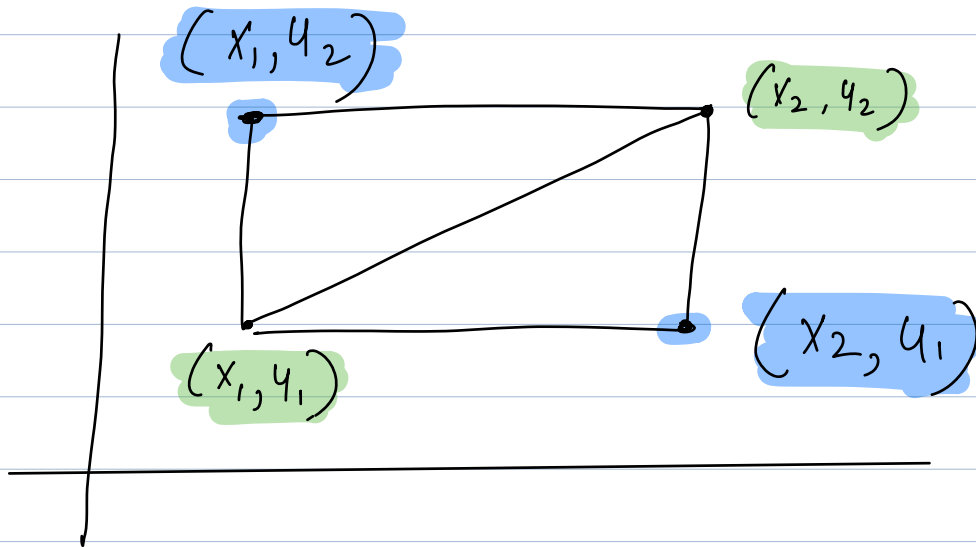
Q2 Given N distinct points on a 2D plane.
Find no of rectangles such that their sides are parallel to x-axis & y-axis.

Ex 1



Brute Force : for all 4 points

$$\underline{\underline{T.C: O(n^4)}}$$



1) We will store a pair in Set

```
class Pair {
    int x;
    int y;
}
```

→ $O(1)$
~~#~~ unordered_set
 ordered_set → $O(\log n)$

→ $O(1)$
~~#~~ unordered_map
 ordered_map → $O(\log n)$

→ $O(1)$
unordered_set

Int, String, float, double, ~~Object~~, ~~Pair~~

ordered_set ⇒ Can take complex keys.

Pseudo Code

ordered-set $\langle \text{Pair} \langle \text{int}, \text{int} \rangle \rangle$ Ts;

for (int i = 0 ; i < n ; i++) {

 Ts.insert (make_pair (x[i], y[i]));

}

int ans \Rightarrow 0;

for (int i = 0 ; i < n ; i++) {

 for (int j = i + 1 ; j < n ; j++) {

$x_1, y_1 \Rightarrow x[i], y[i]$

$x_2, y_2 \Rightarrow x[j], y[j]$

 if ($x_1 == x_2$ || $y_1 == y_2$)
 continue

 Pair $\langle \text{int}, \text{int} \rangle$ first_point = make_pair (x_1, y_1);

 Pair $\langle \text{int}, \text{int} \rangle$ second_point = make_pair (x_2, y_2);

 if (Ts.contains (first_point) && Ts.contains (second_point))
 ans++;

 }

}

return ans/2;

Tc: $n^2 \log n$

Sc: $O(n)$

Example

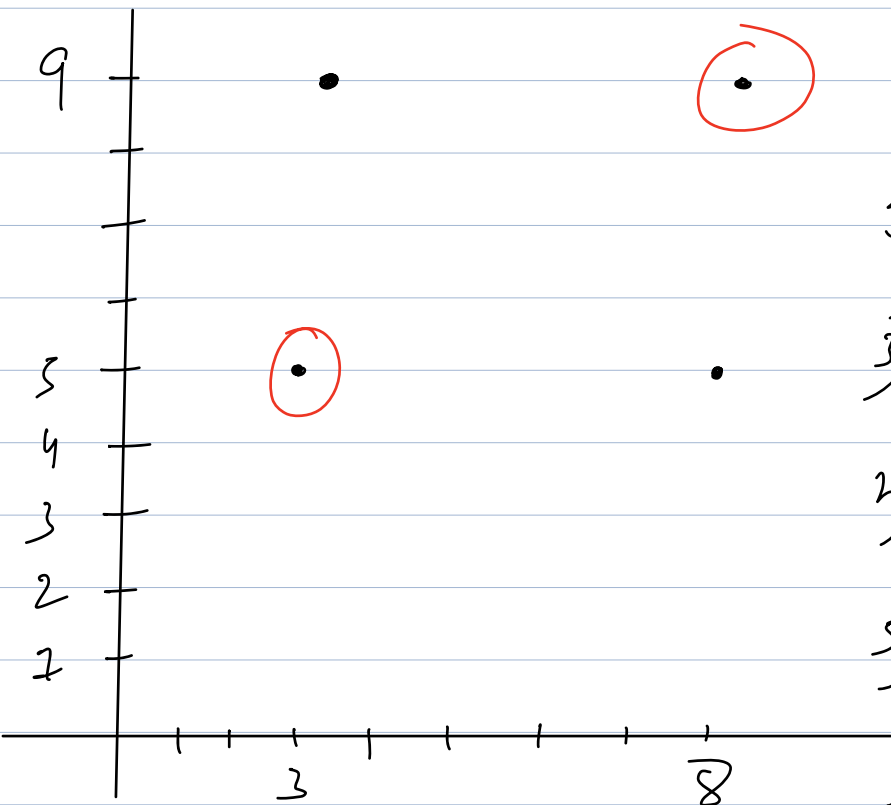
$$x[] = \{3, 3, 8, 8\}$$

$$y[] = \{9, 5, 9, 5\}$$

\Rightarrow

$$(3, 9), (3, 5)$$

$$(8, 9), (8, 5)$$



1) $(3, 9) (3, 5)$ ~~X~~

2) $(3, 9) (8, 9)$ ~~X~~

3) $(3, 9) (8, 5)$ ✓

4) $(3, 5) (8, 9)$ ✓

5) $(3, 5) (8, 5)$ ~~X~~

6) $(8, 9) (8, 5)$ ~~X~~

Total	⇒	168	
Sleep	⇒	49	(7 × 7)
Office	⇒	45	(9 × 5)
normal chores	⇒	28	(4 × 7)
School class	⇒	7.5	(2.5 × 3)
School revise	⇒	15	(5 × 3)

20 hours

⇒ 145

backlog.

1 week off from office

Sat - Sun ⇒