

CLASS \rightarrow Custom Data type

Roll no
Name \Rightarrow Student

int x \Rightarrow 10

Class Student {

int roll-no;

String name;

Student yash = new Student();

\downarrow
Object

Student (int x
String n)

yash.roll-no \Rightarrow 33;
yash.name = "Yash Raj";

roll-no \Rightarrow x;

name \Rightarrow n;

}

Student yash =
new Student(33, "Yash Raj");

}

yash \rightarrow $\left\{ \begin{array}{l} \text{roll-no} = 33 \\ \text{name} = \text{"Yash Raj"} \end{array} \right.$
rahul \rightarrow

Student rahul = yash.

Static Arrays

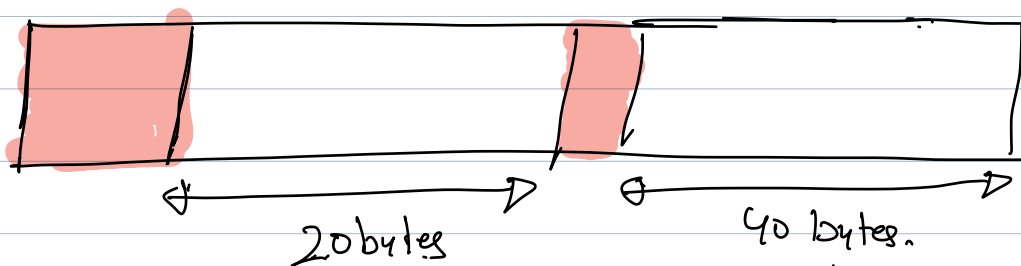
Disadvantage

Advantages

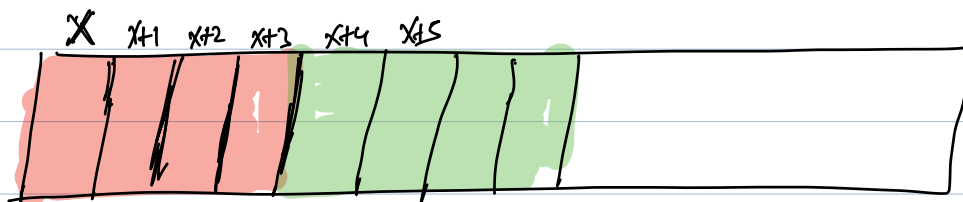
- 1) $O(1)$ access
- 2) $O(1)$ insertion

- 1) No deletion position.
- 2) ~~is~~ fixed size.
- 3) Memory is not utilized optimally.

Contiguous memory allocation



`int x[10]` \Rightarrow 40 bytes



Zero \Rightarrow x

first \Rightarrow x + 4 \Rightarrow x + (1)4

Second \Rightarrow x + 8 \Rightarrow x + (2)4

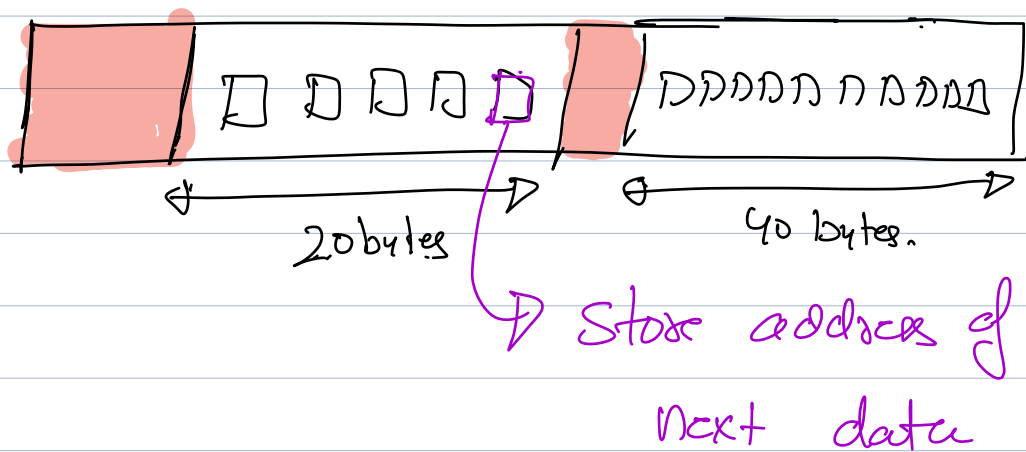
Dynamic Array

Advantages

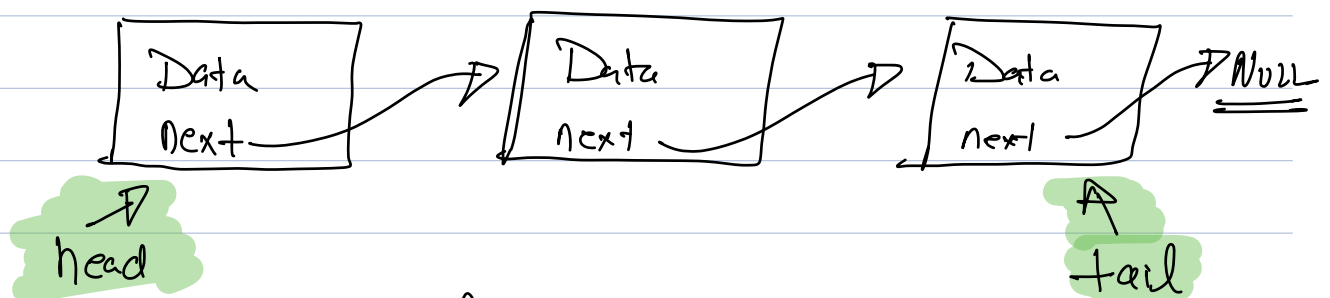
- 1) $O(1)$ access
- 2) $O(1)$ insertion.
- 3) Size can be changed.

Disadvantage.

- 1) Memory is not utilized optimally.

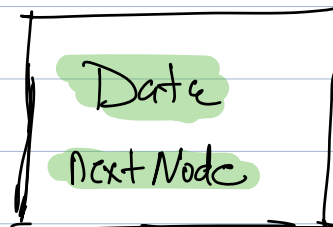


Linked list



class Node {

int data;
Node next;

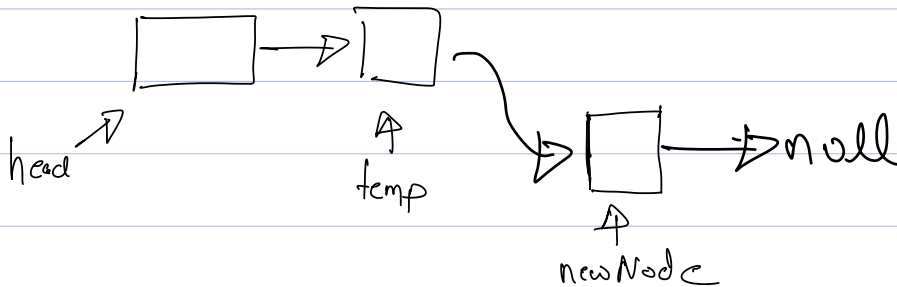


Node (int d) {

data = d
next = null

} }

Q Insert a node at the end of a linked list.



Node insert-at-end (int data , Node head)

Node newNode = new Node (data)

```
if (head == null) {  
    head = newNode;  
    return head;  
}
```

```
Node temp = head;
```

```
while ( temp.next != null ) {  
    temp = temp.next;  
}
```

temp.next → newNode

```
return head;
```

```
}
```

```
|
```

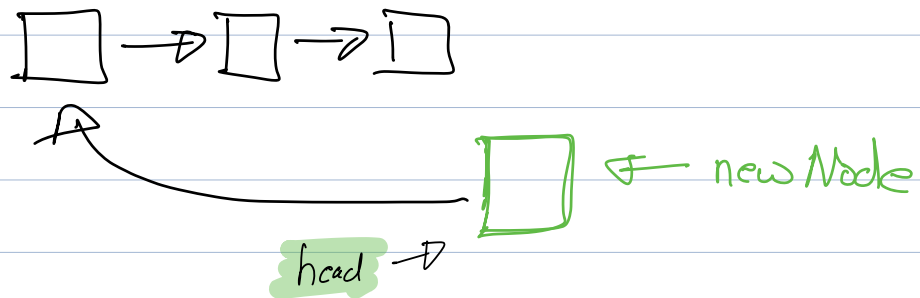
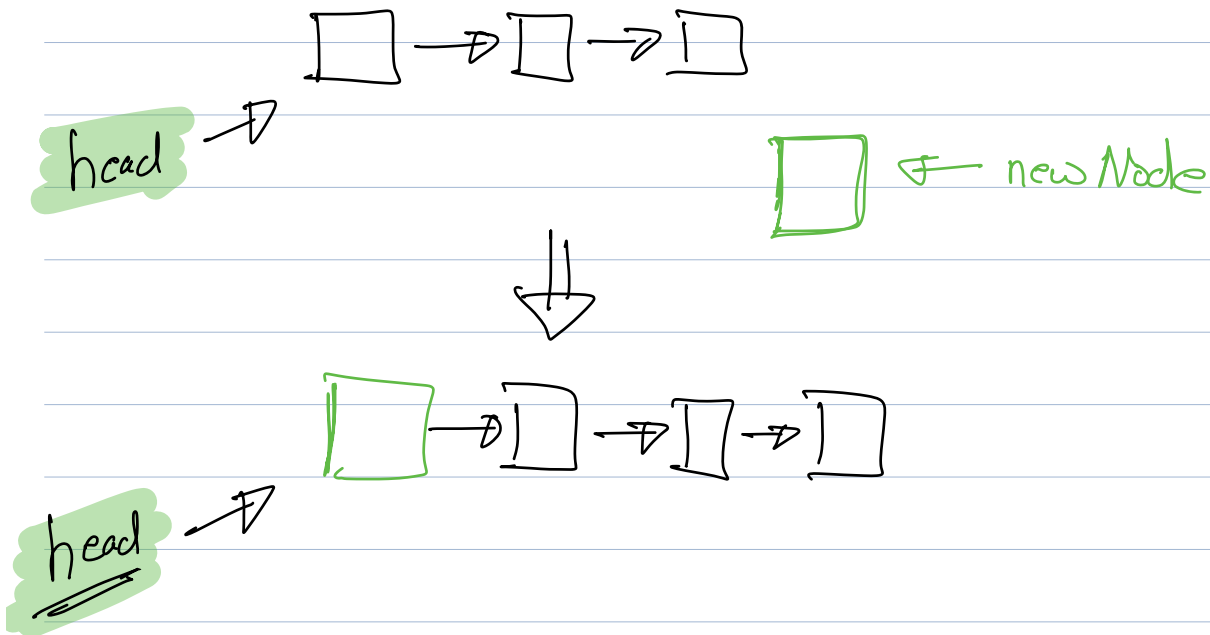
Edge Cases

1) If head was pointing to null.

head \rightarrow null

2) When linked list has only 1 node.

Q Insert at beginning



Node insert-at-beginning (int d, Node head)

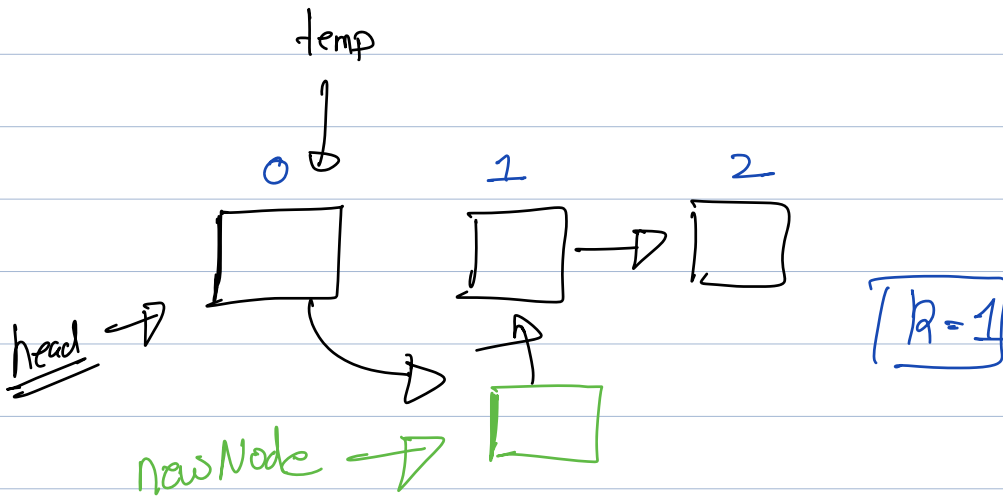
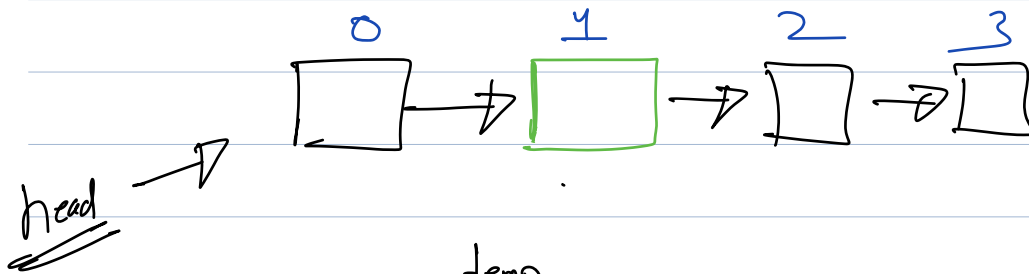
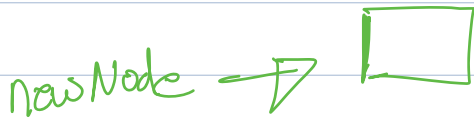
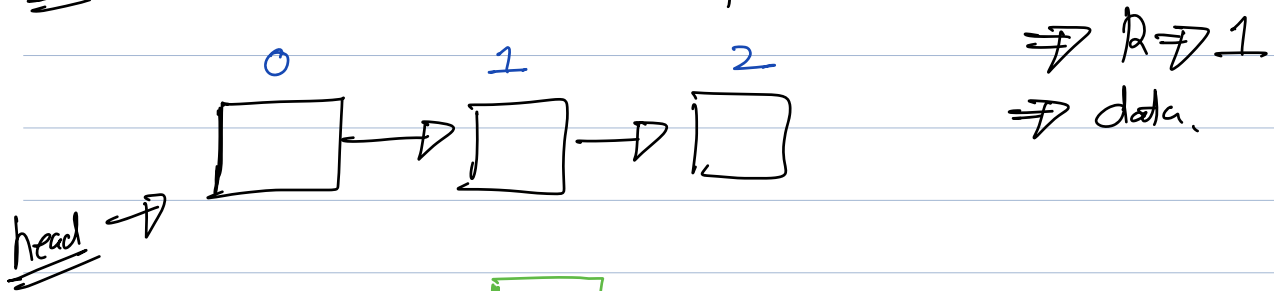
Node newNode = new Node(d)

newNode.next = head

head \Rightarrow newNode

return head;

Q Insert at k^{th} position.



$\text{newNode.next} \Rightarrow \text{temp.next}$
 $\text{temp.next} \Rightarrow \text{newNode.}$

Node insertAtK (int d, int K, Node head) {

if (K == 0)
return insert-at-beginning (d, head);

if (head == null)
return head;

Node newNode = new Node (d);

Node temp = head;

for (int i = 1; i < K; i++) {

if (temp.next == null)
return head;

temp = temp.next;

}

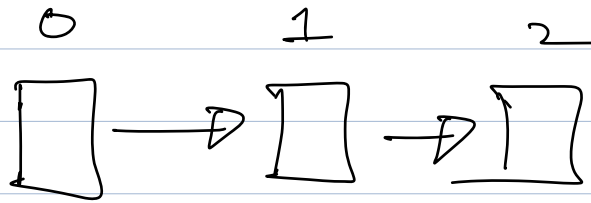
newNode.next \Rightarrow temp.next

temp.next \Rightarrow newNode.

return head;

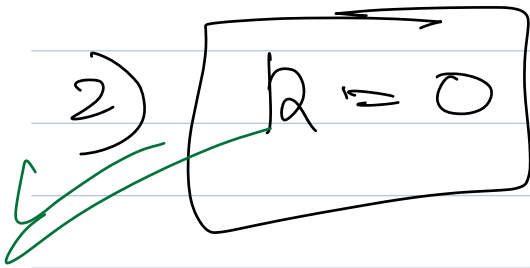
}

1) If k is larger than size of linked list.



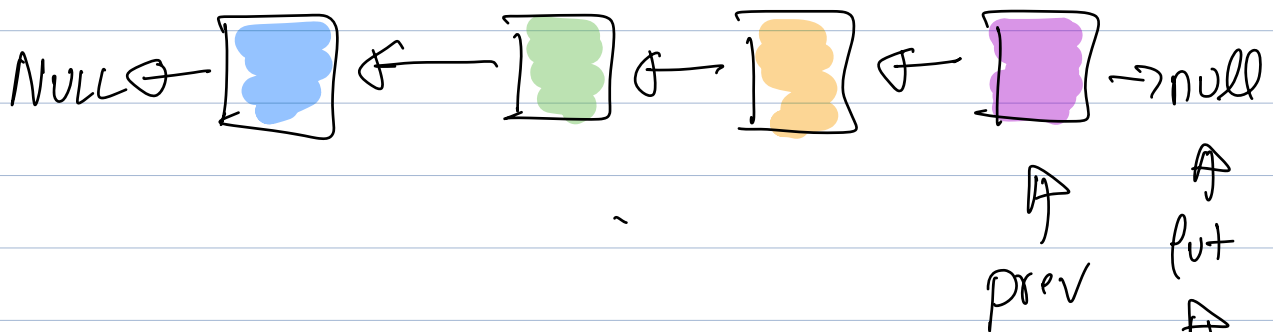
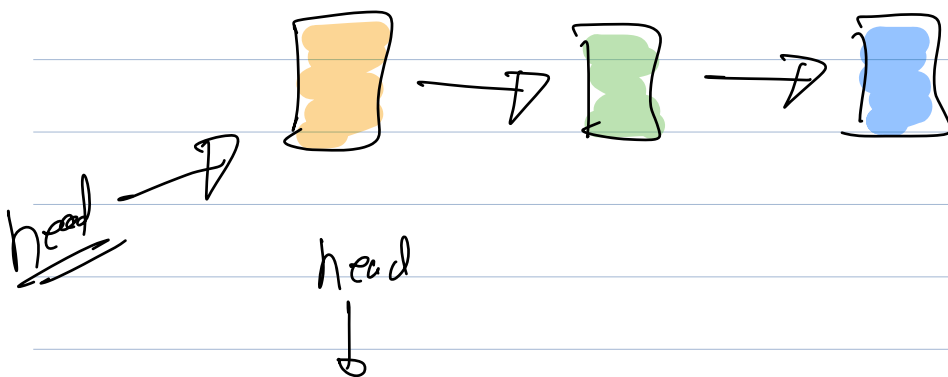
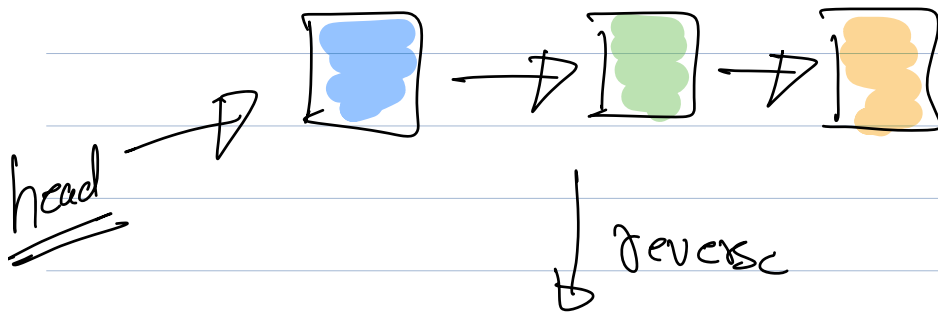
$k=3$

if ($\text{size} < k$)
return head,



3) head == null

Q Reverse a linked list. Sc: $O(1)$



if ($head == null$)
return head;

curr \rightarrow head

prev \rightarrow null

ptr \rightarrow head.next

```
while (curr != null) {
```

```
    curr.next = prev.
```

```
    prev = curr.
```

```
    curr = fut
```

```
    if (fut != null) {
```

```
        fut = fut.next
```

```
}
```

```
head = prev
```

```
return head;
```

```
}
```

2nd Approach

curr = head

prev = null

while (curr != null)

next = curr.next

curr.next = prev

prev = curr

curr = next

}

return prev