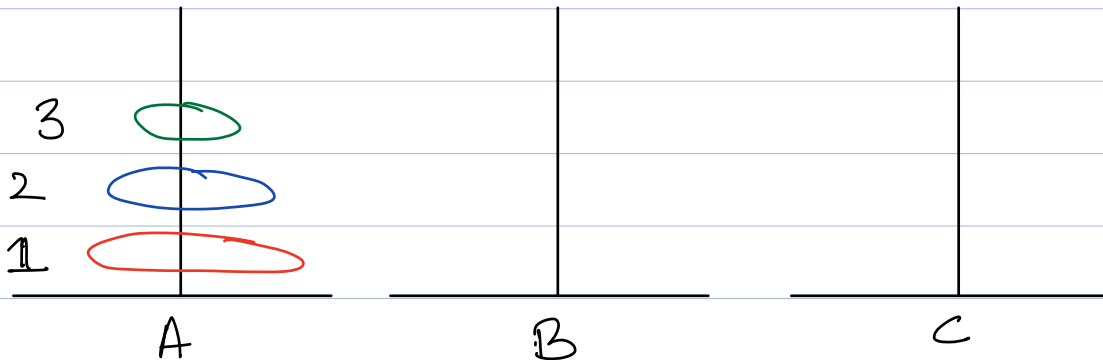
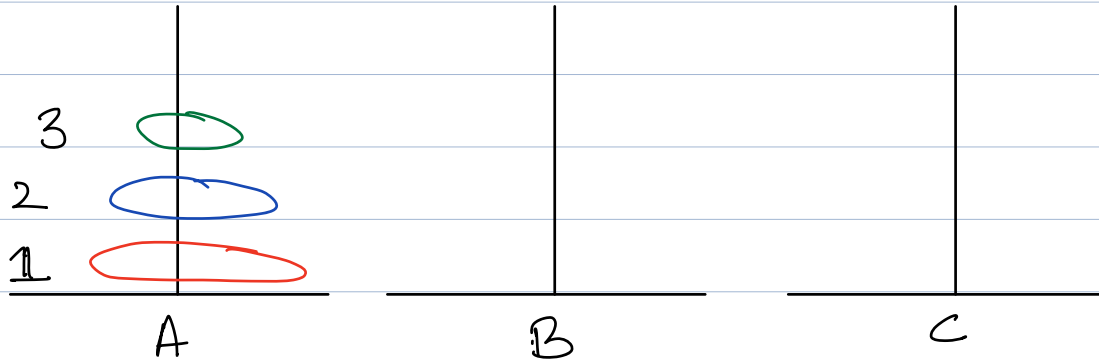


## Tower of Hanoi

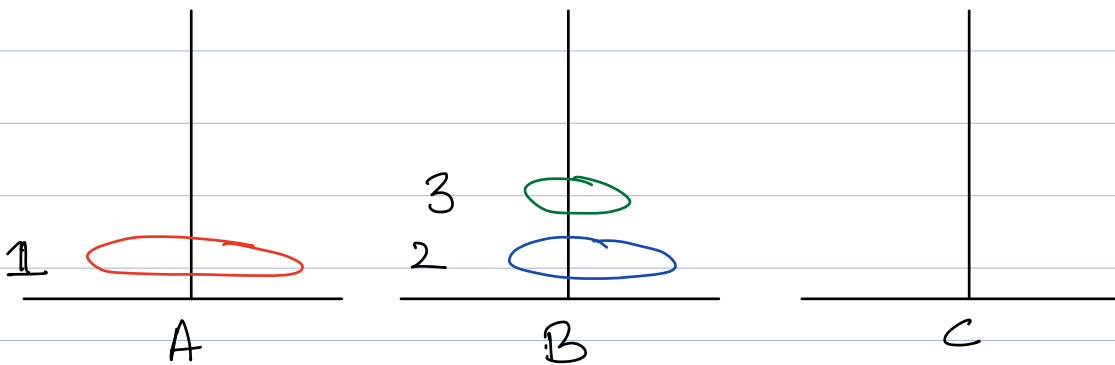
Q Shift the disc from A to C. You can use B.  
Bigger disc cannot be on top of a smaller disc at any point. You can move only 1 disc at a time.



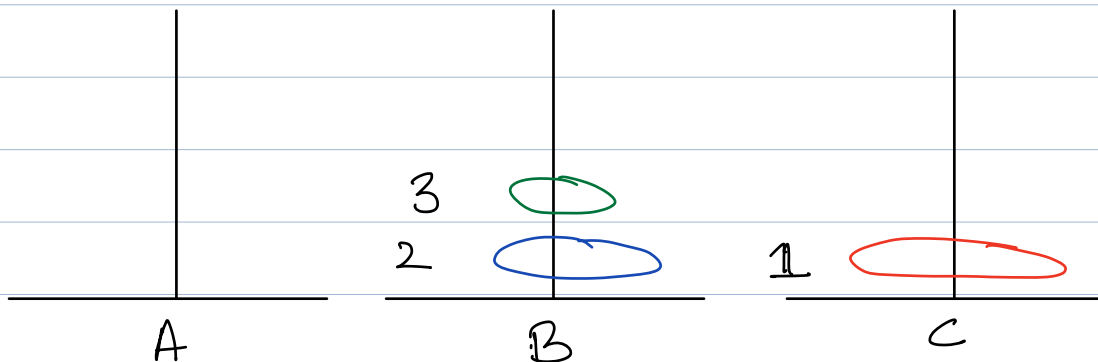
# Approach 1



$3 \rightarrow AC, 2 A \rightarrow B, 3 C \rightarrow B$

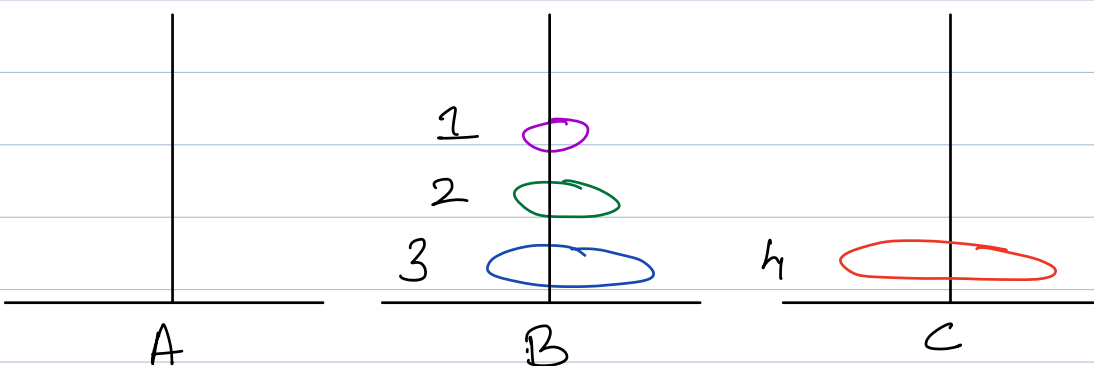
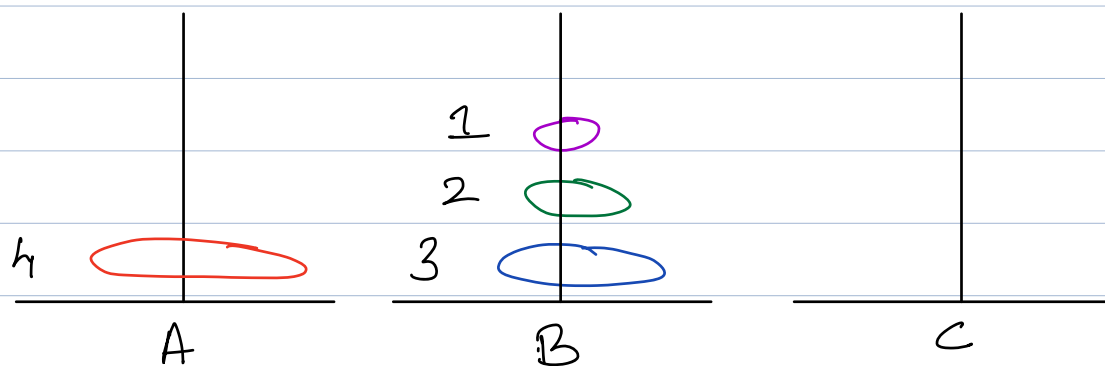


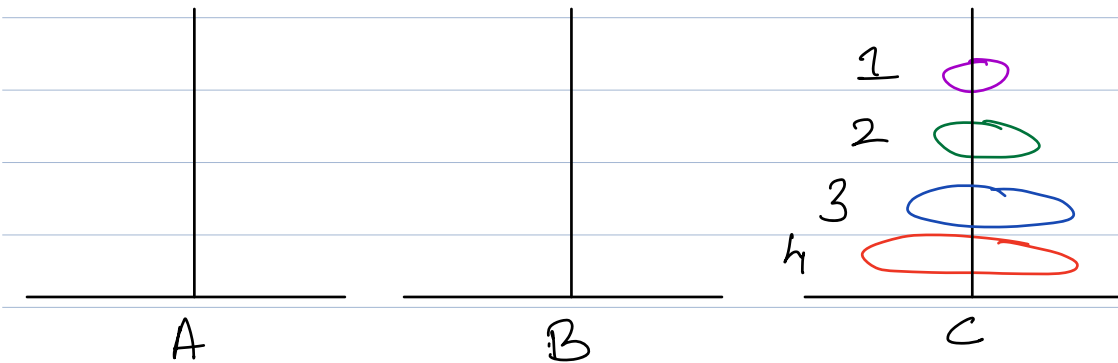
$1 A \rightarrow C$



3  $B \rightarrow A$ , 2  $\rightarrow B \rightarrow C$ , 3  $A \rightarrow C$







Q  $n$  discs. You need to move from Tower A to Tower C using Tower B. Point all the steps.

1 2 3 4 5

i) Manifestation

tower( $n$ )  $\Rightarrow$  solves for  $n$  discs ~~X~~

tower(int  $n$ , String source, String destination, String helper)

## 2) Main logic

1) Move  $(n-1)$  rings/discs to helper.

$\text{tower}(n-1, \text{source}, \text{helper}, \text{destination})$

2) Move  $n^{\text{th}}$  ring to destination.

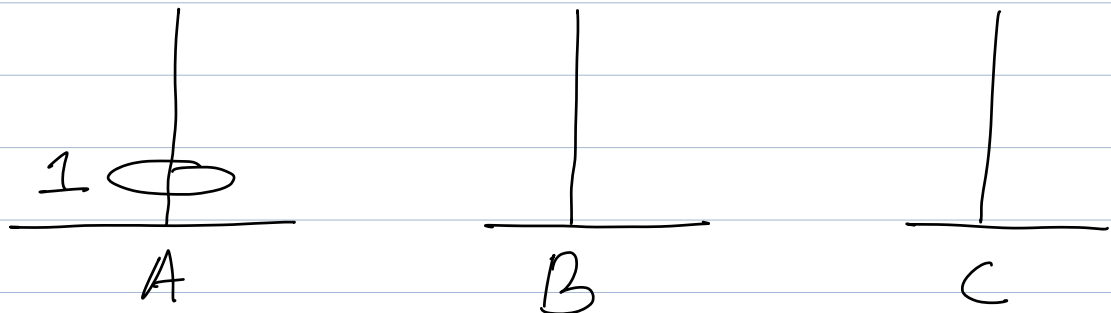
$\text{print}(n: \text{source destination})$

3) Move  $(n-1)$  rings from helper to destination.

$\text{tower}(n-1, \text{helper}, \text{destination}, \text{source})$

## 3) Base Case

$n=1$



Point 1: A C

```
void tower (int n , string source , string destination ,  
            string helper )
```

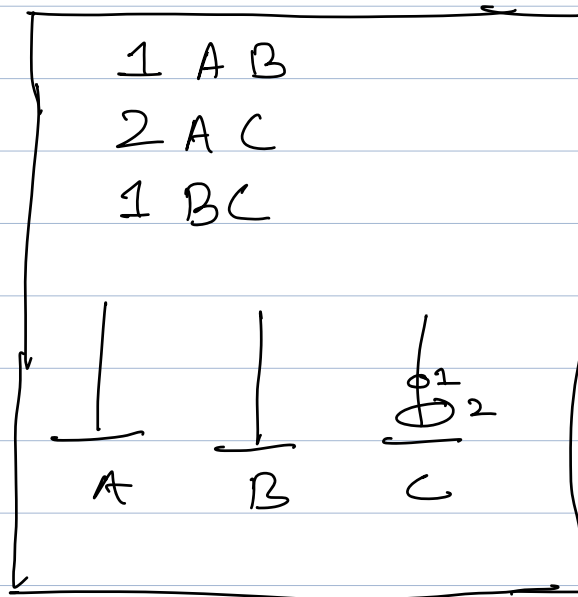
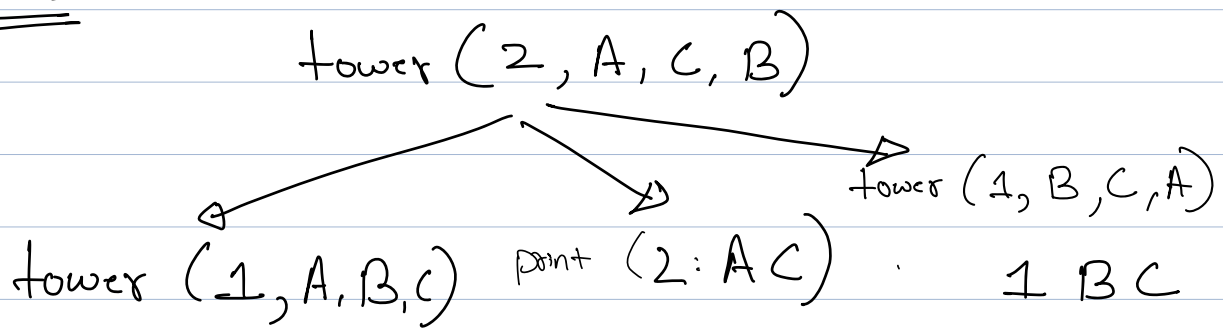
```
    if (n == 1) {  
        print (1: source destination) // Base  
        return ;                      Case  
    }
```

```
    tower (n-1 , source , helper , destination)
```

```
    print ( n: source destination)
```

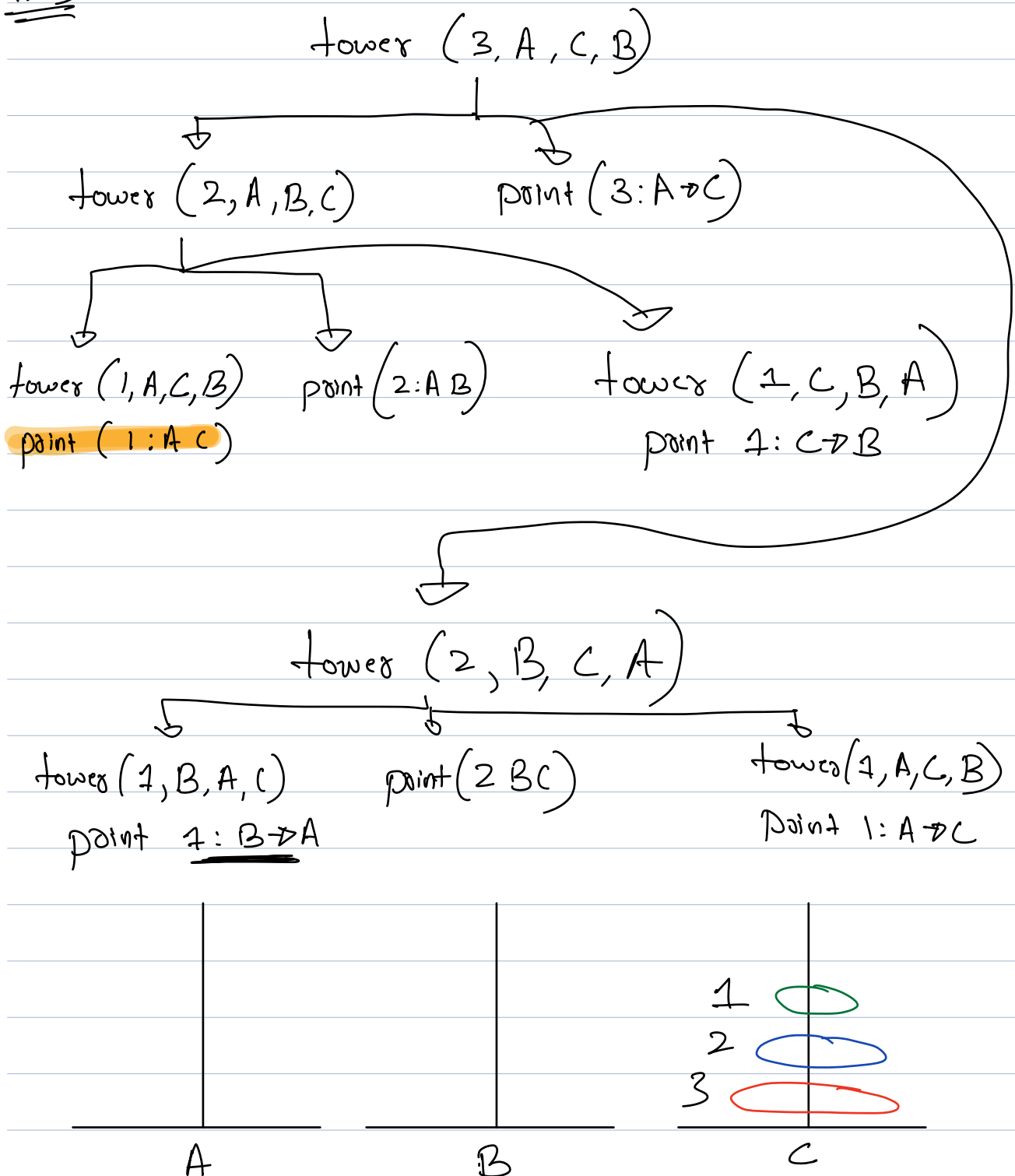
```
    tower (n-1 , helper , destination , source)
```

```
}
```

$$\underline{\underline{n = 2}}$$




N=3



1: A C , 2: A B , 1: C B , 3: A → C , 1: B → A

2BC

1AC

Time Complexity.

$$T(n) \Rightarrow 1 + 2T(n-1)$$

$$\begin{aligned} T(n) &\Rightarrow 1 + 2(1 + 2T(n-2)) \\ &\Rightarrow 3 + 4T(n-2) \end{aligned}$$

$$T(n) \Rightarrow 7 + 8T(n-3)$$

$$T(n) \Rightarrow 2^k - 1 + 2^k T(n-k)$$

$$\Rightarrow 2^{n-1} - 1 + 2^{n-1}$$

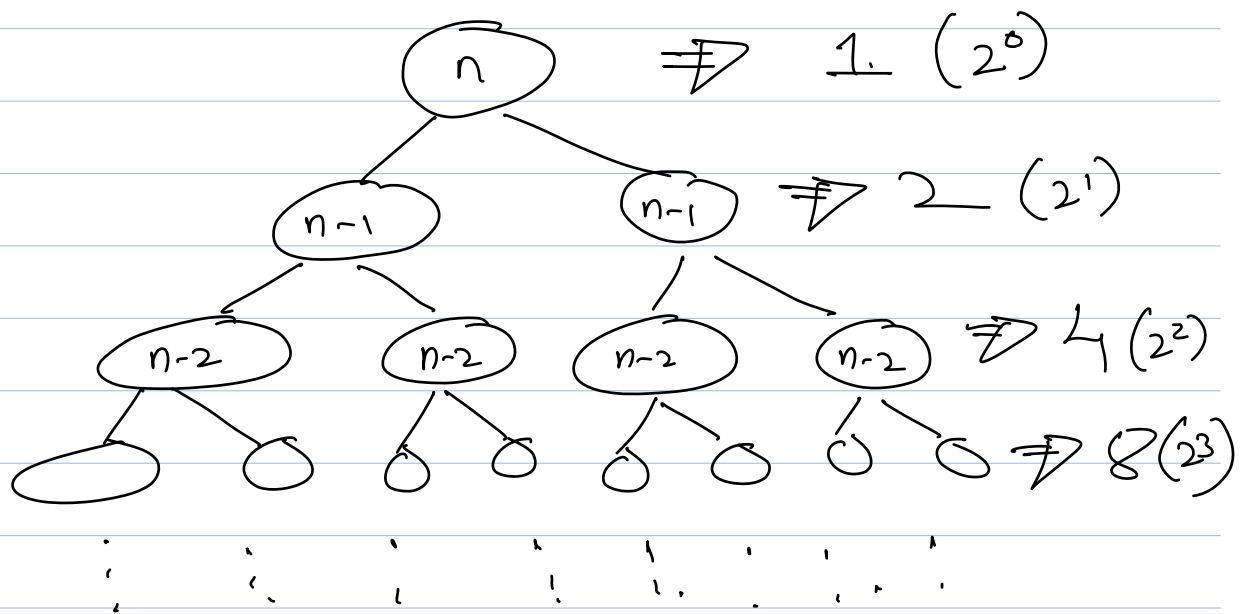
$$\Rightarrow 2^n$$

$$T(1) = 1$$

$$n-k = 1$$

$$\boxed{n-1} = k$$

$$TC : \underline{\underline{O(2^n)}}$$



$$\textcircled{1} \textcircled{2} \textcircled{1} \textcircled{1} \textcircled{1} \textcircled{1} \textcircled{2} \textcircled{1} \Rightarrow 2^{n-1}$$

$$2^0 + 2^1 + 2^2 \dots 2^{n-1} \Rightarrow 2^n - 1$$

$$TC: 2^n \times O(1) \\ \underline{\underline{= 2^n}}$$

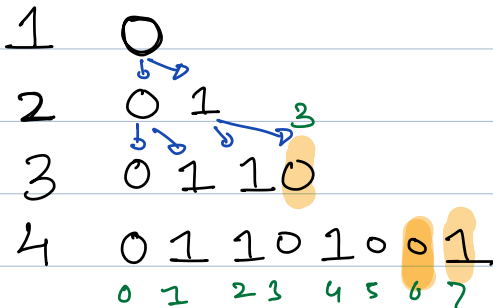
Height of the tree  $\Rightarrow n$

$$SC: O(n)$$

10:33pm

$k^{\text{th}}$  Symbol

0 1 2 3



$0 \rightarrow 01$

$1 \rightarrow 10$

Q Return the value at  $n^{\text{th}}$  row and  $k^{\text{th}}$  column.

Ex:1  $n \rightarrow 3, k \rightarrow 2$

$\Rightarrow 1$

Ex2:  $n \rightarrow 4, k \rightarrow 7$

$\Rightarrow 1$

Constraints

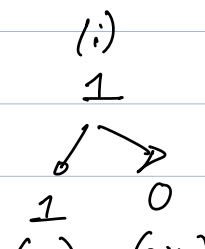
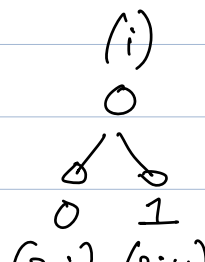
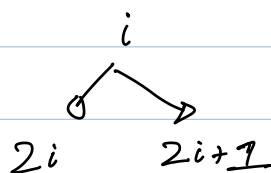
$1 \leq n \leq 10^5$

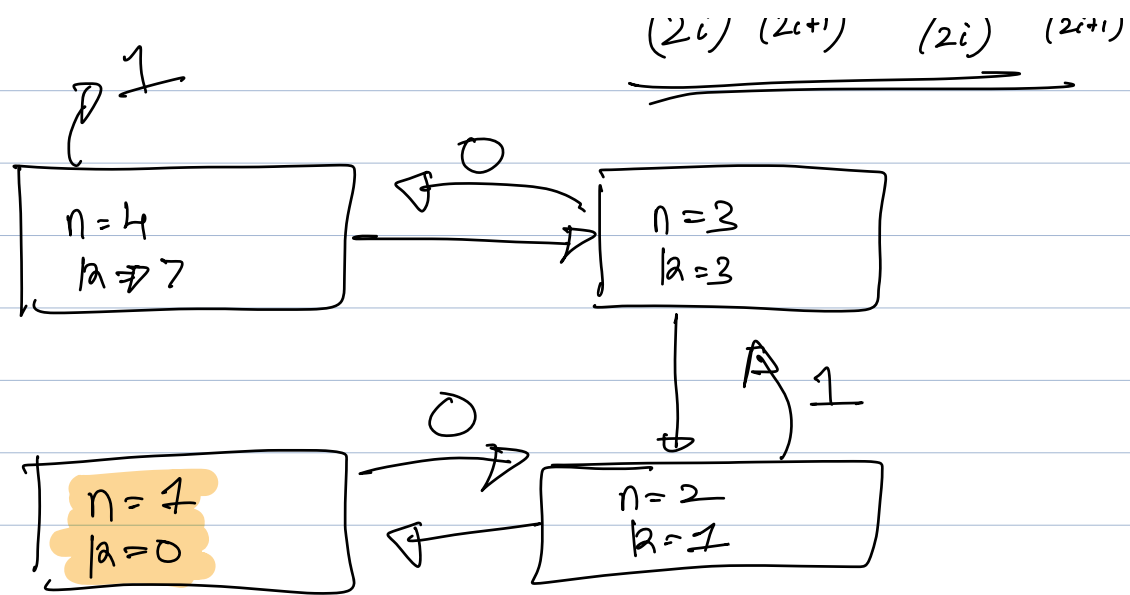
$0 \leq k \leq 10^{18}$

$\text{string}(n) \Rightarrow \text{string}(n-1) + \text{toggle}(\text{string}(n-1))$

Observation

,  $\text{val}(n, k)$   
 $\downarrow$  parent  
 $\text{val}(n-1, k/2)$





1) Manipulation.

```
int solve (int n, int k)
```

↳ returns the value at  $n^{\text{th}}$  row &  $k^{\text{th}}$  column.

2) Main logic

1) Find parent

```
int parent => solve (n-1, k/2)
```

2) Find the value.

```
if (k%2 == 0) {
```

```
    return parent;
```

```
} else {
```

```
    return (1-parent);
```

```
}
```

↳ ! parent

↳ 1 ^ parent

↳ ~ parent

3) Base Case !

(i) if ( $n == 1$ )  
return 0;

(ii) if ( $k == 0$ )  
return 0;

```
int solve (int n, int k) {  
    if ( $n == 1$  ||  $k == 0$ )  
        return 0;
```

```
    int parent = solve ( $n-1$ ,  $k/2$ );  
    if ( $k \% 2 == 0$ ) {  
        return parent;  
    } else {  
        return (1 + parent);  
    }  
}
```

```
}
```

~~$$T(n) \Rightarrow T(n-1) + 1$$~~

$$T(k) \Rightarrow T(k/2) + 1$$

$$T_C: O(\log k)$$

$$S_C \Rightarrow O(\log k)$$

Ex 1

$$\text{Time comp} = O(n)$$

$$\boxed{\begin{array}{l} n = 10^5 / 10^6 / 10^7 \\ k \neq 10 \end{array}}$$

$$\Rightarrow O(n)$$

$$O(\log k)$$