INFORMATION SYSTEM DEPARTMENT INFORMATION TECHNOLOGY FACULTY— HCM UNIVERSITY SCIENCE

INTRODUCTION TO DATABASE Chapter 03 Structured Query Language (SQL)

Lecturer- PhD. NGUYEN TRAN MINH THU





- SQL Overview
- Basic SQL
 - SQL Data Definition and Data Types
 - Specifying Constraints in SQL
 - Insert/Update/Delete Statement in SQL
 - Basic Retrieval Queries in SQL
- Advanced SQL
 - Complex SQL Retrieval Queries
 - □ View
 - Ensuring Transaction Integrity
 - ☐ Specifying Constraints as Assertions & Action Trigger
- ☐ Tips for Developing Queries



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- ☐ Structured Query Language
- ☐ The standard for relational database management systems (RDBMS)
- □ RDBMS: A database management system that manages data as a collection of tables in which all relationships are represented by common values in related tables



History of SQL

- ☐ 1970—E. F. Codd develops relational database concept
- ☐ 1974-1979—System R with Sequel (later SQL) created at IBM Research Lab
- 1979—Oracle markets first relational DB with SQL
- 1981 SQL/DS first available RDBMS system on DOS/VSE
- Others followed: INGRES (1981), IDM (1982), DG/SGL (1984),
 Sybase (1986)
- 1986–ANSI SQL standard released
- ☐ 1989, 1992, 1999, 2003, 2006, 2008–Major ANSI standard updates
- ☐ Current—SQL is supported by most major database vendors



Purpose of SQL Standard

- Specify syntax/semantics for data definition and manipulation
- Define data structures and basic operations
- Enable portability of database definition and application modules
- ☐ Specify minimal (level 1) and complete (level 2) standards
- □ Allow for later growth/enhancement to standard (referential integrity, transaction management, user-defined functions, extended join operations, national character sets)



Benefits of a Standardized Relational Language

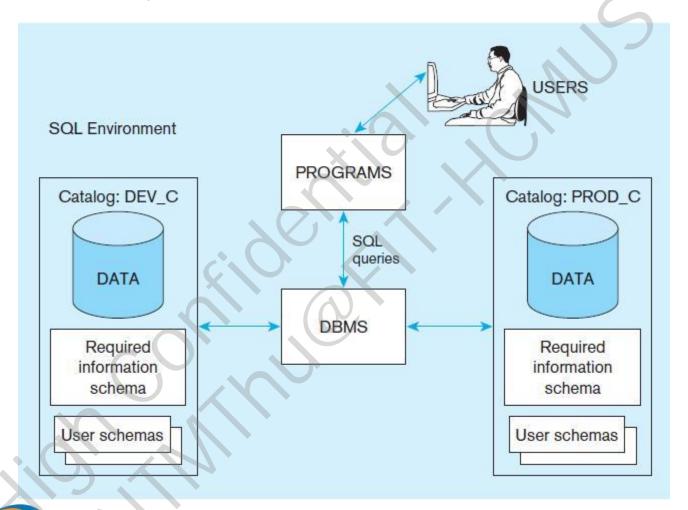
- □ Reduced training costs
- Productivity
- Application portability
- □ Application longevity
- ☐ Reduced dependence on a single vendor
- ☐ Cross-system communication



- ☐ Catalog
 - A set of schemas that constitute the description of a database
- □ Schema
 - The structure that contains descriptions of objects created by a user (base tables, views, constraints)
- □ Data Definition Language (DDL)
 - Commands that define a database, including creating, altering, and dropping tables and establishing constraints
- □ Data Manipulation Language (DML)
 - Commands that maintain and query a database
- □ Data Control Language (DCL)
 - Commands that control a database, including administering privileges and committing data



A simplified schematic of a typical SQL environment, as described by the SQL: 2008 standard





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□ Number

- INTEGER
- SMALLINT
- NUMERIC, NUMERIC(p), NUMERIC(p,s)
- DECIMAL, DECIMAL(p), DECIMAL(p,s)
- REAL
- DOUBLE PRECISION
- FLOAT, FLOAT(p)

□ Character String

- CHARACTER or CHAR
- CHARACTER(n) or CHAR (n)
- CHARACTER VARYING(n) or VARCHAR(n)
- NATIONAL CHARACTER (n) orNCHAR(n)



SQL Data Types

☐ Bit string

- BIT, BIT(x)
- BIT VARYING(x)

□ Datetime

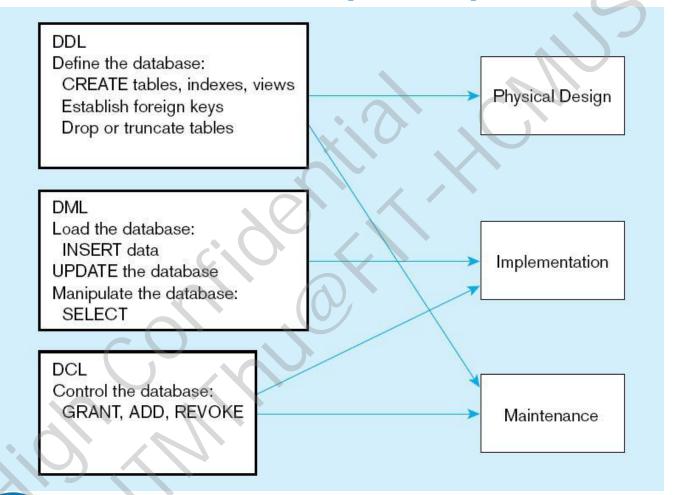
- DATE (day, month, year)
- TIME (hour, minute, second)
- TIMESTAMP (Date & Time)
- DATETIME: DATE and TIME (only in SQLof SQL Server)

□ Boolean

Store truth values: TRUE, FALSE & UNKNOW



DDL, DML, DCL, and the database development process





- ☐ Data Definition Language (DDL)
- Major CREATE statements:
 - ☐ CREATE SCHEMA—defines a portion of the database owned by a particular user
 - ☐ CREATE TABLE—defines a new table and its columns
 - ☐ CREATE VIEW—defines a logical table from one or more tables or views
 - Other CREATE statements: CHARACTER SET, COLLATION, TRANSLATION, ASSERTION, DOMAIN



Steps in Table Creation

- 1. Identify data types for attributes
- 2. Identify columns that can and cannot be null
- 3. Identify columns that must be unique (candidate keys)
- 4. Identify primary key-foreign key mates
- 5. Determine default values
- 6. Identify constraints on columns (domain specifications)
- 7. Create the table and associated indexes



☐ SYNTAX FOR CREATE TABLE

```
CREATE TABLE table_name (
        column1 <datatype> [<column_constraint],
        column2 <datatype> [<column_constraint],
        column3 <datatype> [<column_constraint],
        ....
);</pre>
```



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54.0 Specifying Constraints in SQL

- □ Column Constraints
 - NOT NULL
 - □ NULL
 - UNIQUE
 - DEFAULT
 - PRIMARY KEY
 - ☐ FOREIGN KEY / REFERENCES
 - ☐ CHECK
- □ Naming for column constraint

CONSTRAINT < Name_Constraint > < constraint >



4.0 Create Statements

CREATE TABLE EMPLOYEE

Minit

(Fname VARCHAR(15)

CHAR.

Lname VARCHAR(15)

Ssn CHAR(9)

Bdate DATE.

Address VARCHAR(30).

Sex CHAR,

Salary DECIMAL(10.2).

Super_ssn CHAR(9),

NOT NULL Dno INT

PRIMARY KEY (Ssn),

CREATE TABLE DEPT LOCATIONS

(Dnumber INT

Dlocation VARCHAR(15)

PRIMARY KEY (Dnumber, Dlocation),

FOREIGN KEY (Dnumber) REFERENCES DEPARTMENT(Dnumber));

CREATE TABLE DEPENDENT

Essn CHAR(9)

Dependent_name VARCHAR(15)

Sex CHAR. Bdate DATE.

VARCHAR(8), Relationship

PRIMARY KEY (Essn, Dependent name),

FOREIGN KEY (Essn) REFERENCES EMPLOYEE(Ssn));

CREATE TABLE DEPARTMENT

NOT NULL.

NOT NULL.

NOT NULL.

NOT NULL

NOT NULL

NOT NULL.

NOT NULL,

VARCHAR(15) (Dname

Dnumber INT

Mgr ssn CHAR(9)

Mgr_start_date DATE.

PRIMARY KEY (Dnumber). UNIQUE (Dname),

FOREIGN KEY (Mgr ssn) REFERENCES EMPLOYEE(Ssn));

NOT NULL.

CREATE TABLE PROJECT

(Pname VARCHAR(15)

Pnumber

Plocation VARCHAR(15),

Dnum

INT

PRIMARY KEY (Pnumber),

UNIQUE (Pname).

FOREIGN KEY (Dnum) REFERENCES DEPARTMENT(Dnumber)):

CREATE TABLE WORKS ON

(Essn CHAR(9)

Pno INT

DECIMAL(3,1) Hours

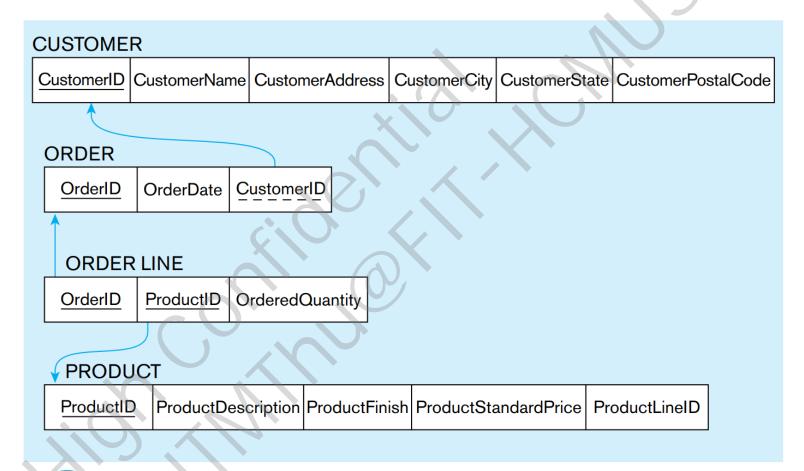
PRIMARY KEY (Essn. Pno).

FOREIGN KEY (Essn) REFERENCES EMPLOYEE(Ssn),

FOREIGN KEY (Pno) REFERENCES PROJECT(Pnumber)):



The following slides create tables for this enterprise data model





SQL Overall table definitions

CREATE TABLE Customer_T

(CustomerID NUMBER(11,0) NOT NULL, CustomerName VARCHAR2(25) NOT NULL,

CustomerAddress VARCHAR2(30),
CustomerCity VARCHAR2(20),
CustomerState CHAR(2),

CustomerPostalCode VARCHAR2(9),

CONSTRAINT Customer_PK PRIMARY KEY (CustomerID));

CREATE TABLE Order_T

(OrderID NUMBER(11,0) NOT NULL,

OrderDate DATE DEFAULT SYSDATE,

CustomerID NUMBER(11,0),

CONSTRAINT Order_PK PRIMARY KEY (OrderID),

CONSTRAINT Order_FK FOREIGN KEY (CustomerID) REFERENCES Customer_T(CustomerID));

CREATE TABLE Product_T

(ProductID NUMBER(11,0) NOT NULL,

ProductDescription VARCHAR2(50),
ProductFinish VARCHAR2(20)

CHECK (ProductFinish IN ('Cherry', 'Natural Ash', 'White Ash',

'Red Oak', 'Natural Oak', 'Walnut')),

ProductStandardPrice DECIMAL(6,2),
ProductLineID INTEGER,

CONSTRAINT Product_PK PRIMARY KEY (ProductID));

CREATE TABLE OrderLine_T

(OrderID NUMBER(11,0) NOT NULL, ProductID INTEGER NOT NULL,

OrderedQuantity NUMBER(11,0),

CONSTRAINT OrderLine_PK PRIMARY KEY (OrderID, ProductID),

CONSTRAINT OrderLine_FK1 FOREIGN KEY (OrderID) REFERENCES Order_T(OrderID),

CONSTRAINT OrderLine_FK2 FOREIGN KEY (ProductID) REFERENCES Product_T(ProductID));





Defining attributes and their data types

CREATE TABLE Product_T

(ProductID NUMBER(11,0)

ProductDescription VARCHAR2(50)

ProductFinish VARCHAR2(20)

CHECK (ProductFinish IN ('Cherry', 'Natural Ash', 'White Ash',

'Red Oak', 'Natural Oak', 'Walnut')),

NOT NULL,

ProductStandardPrice DECIMAL(6,2),

ProductLineID INTEGER,

CONSTRAINT Product_PK PRIMARY KEY (ProductID));



Non-nullable specification

CREATE TABLE Product_T	
(ProductID	NUMBER(11,0) NOT NULL,
ProductDescription	VARCHAR2(50),
ProductFinish	VARCHAR2(20)
C	CHECK (ProductFinish IN ('Cherry', 'Natural Ash', 'White Ash',
	'Red Oak', 'Natural Oak', 'Walnut')),
ProductStandardPrice	DECIMAL(6,2),
ProductLineID	INTEGER, Primary keys can never have
CONSTRAINT Product_PK PRIMARY KEY (Pr	oductID)); NULL values

Identifying primary key



Non-nullable specifications

CREATE TABLE OrderLine_T

(OrderID

ProductID

OrderedQuantity

NUMBER(11,0)

INTEGER

NOT NULL,

NUMBER(11,0),

CONSTRAINT OrderLine_PK PRIMARY KEY (OrderID, ProductID),

Primary key

CONSTRAINT OrderLine_FK1 FOREIGN KEY (OrderID) REFERENCES Order_T(OrderID),

CONSTRAINT OrderLine_FK2 FOREIGN KEY (ProductID) REFERENCES Product_T(ProductID));

Some primary keys are composite—composed of multiple attributes



Controlling the values in attributes

```
CREATE TABLE Order T
                                                         NUMBER(11,0)
                                                                            NOT NULL,
                  (OrderID
                   OrderDate
                                                         DATE DEFAULT SYSDATE,
                   CustomerID
                                                         NUMBER(11,0),
                                                                          Default value
CONSTRAINT Order_PK PRIMARY KEY (OrderID),
CONSTRAINT Order_FK FOREIGN KEY (CustomerID) REFERENCES Customer_T(CustomerID));
CREATE TABLE Product T
                  (ProductID
                                                         NUMBER(11,0)
                                                                            NOT NULL,
                   ProductDescription
                                                         VARCHAR2(50),
                   ProductFinish
                                                         VARCHAR2(20)
                                      CHECK (ProductFinish IN ('Cherry', 'Natural Ash', 'White Ash',
     Domain constraint
                                                           'Red Oak', 'Natural Oak', 'Walnut')),
                   ProductStandardPrice
                                                         DECIMAL(6,2),
                   ProductLineID
                                                         INTEGER,
CONSTRAINT Product_PK PRIMARY KEY (ProductID));
```



Identifying foreign keys and establishing relationships

CREATE	TABLE	Customer_	T
--------	-------	-----------	---

(CustomerID		NUMBER(11,0)	NOT NULL,
CustomerName		VARCHAR2(25)	NOT NULL,
CustomerAddress	1,0	VARCHAR2(30),	

Primary key of parent table

CustomerCity

CustomerState

CustomerPostalCode

VARCHAR2(20), CHAR(2),

VARCHAR2(9),

CONSTRAINT Customer_PK PRIMARY KEY (CustomerID)),

CREATE TABLE Order_T

(OrderID NUMBER(11,0)

OrderDate DATE DEFAULT SYSDATE,

CustomerID NUMBER(11,0),

CONSTRAINT Order PK PRIMARY KEY (OrderID),

CONSTRAINT Order_FK FOREIGN KEY (CustomerID) REFERENCES Customer_T(CustomerID));

Foreign key of dependent table

NOT NULL,

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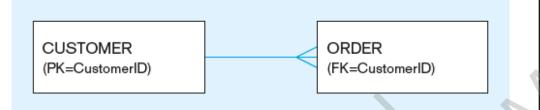


Data Integrity Controls

- □ Referential integrity—constraint that ensures that foreign key values of a table must match primary key values of a related table in 1:M relationships
- ☐ Restricting:
 - □ Deletes of primary records
 - ☐ Updates of primary records
 - ☐ Inserts of dependent records



Ensuring data integrity through updates



Restricted Update: A customer ID can only be deleted if it is not found in ORDER table.

CREATE TABLE CustomerT

(CustomerID CustomerName INTEGER DEFAULT '999' VARCHAR(40) NOT NULL, NOT NULL.

. . .

CONSTRAINT Customer_PK PRIMARY KEY (CustomerID)
ON UPDATE RESTRICT);

Cascaded Update: Changing a customer ID in the CUSTOMER table will result in that value changing in the ORDER table to match.

... ON UPDATE CASCADE);

Set Null Update: When a customer ID is changed, any customer ID in the ORDER table that matches the old customer ID is set to NULL.

... ON UPDATE SET NULL);

Set Default Update: When a customer ID is changed, any customer ID in the ORDER tables that matches the old customer ID is set to a predefined default value.

... ON UPDATE SET DEFAULT);

Relational integrity is enforced via the primary-key to foreign-key match



4.0 Specifying Constraints in SQL

```
CREATE TABLE EMPLOYEE (
Fname VARCHAR(15) NOT NULL,
....,
Dno INT NOT NULL DEFAULT 1,
CONSTRAINT EMPPK PRIMARY KEY (Ssn),
CONSTRAINT EMPSUPERFK
FOREIGN KEY (Super_ssn) REFERENCES EMPLOYEE(Ssn)
ON DELETE SET NULL ON UPDATE CASCADE,
CONSTRAINT EMPDEPTFK
FOREIGN KEY(Dno) REFERENCES DEPARTMENT(Dnumber)
ON DELETE SET DEFAULT ON UPDATE CASCADE)
```

Specifying Key and Referential Integrity Constraints



3.0 Schema Change Statements in SQL

- DROP Command
 - DROP SCHEMA: drop a whole schema
 - □ DROP TABLE: remove tables from schema
 - ☐ There are two drop behavior options: CASCADE and RESTRICT

DROP SCHEMA COMPANY CASCADE;

DROP TABLE DEPENDENT CASCADE;;



3.0 Schema Change Statements in SQL

□ The ALTER Command

■ Alter table actions include adding or dropping a column (attribute), changing a column definition, and adding or dropping table constraints

ALTER TABLE COMPANY. EMPLOYEE ADD COLUMN Job VARCHAR (12);

ALTER TABLE COMPANY. EMPLOYEE DROP COLUMN Address CASCADE;

ALTER TABLE COMPANY DEPARTMENT ALTER COLUMN Mgr_ssn DROP DEFAULT;

ALTER TABLE COMPANY.DEPARTMENT ALTER COLUMN Mgr_ssn SET DEFAULT '333445555';

ALTER TABLE COMPANY.EMPLOYEE

DROP CONSTRAINT EMPSUPERFK CASCADE;





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- ☐ The INSERT Command: Adds one or more rows to a table
 - ☐ Inserting into a table

Inserting a record that has some null attributes requires identifying the fields that actually get data

```
INSERT INTO EMPLOYEE (Fname, Lname, Dno, Ssn)
VALUES ('Richard', 'Marini', 4, '653298653');
```



- ☐ The INSERT Command: Adds one or more rows to a table
 - ☐ Inserting from another table

```
CREATE TABLE WORKS_ON_INFO

( Emp_name VARCHAR(15),
    Proj_name VARCHAR(15),
    Hours_per_week DECIMAL(3,1));

INSERT INTO WORKS_ON_INFO
SELECT E.Lname, P.Pname, W.Hours
FROM PROJECT P, WORKS_ON W, EMPLOYEE E
WHERE P.Pnumber = W.Pno AND W.Essn = E.Ssn;
```



- ☐ The DELETE Command: Removes rows from a table
 - Delete certain rows

```
DELETE FROM EMPLOYEE

WHERE Lname = 'Brown';

DELETE FROM EMPLOYEE

WHERE Ssn = '123456789';

DELETE FROM EMPLOYEE

WHERE Dno = 5;
```

Delete all rows

DELETE FROM EMPLOYEE



☐ Update Statement: Modifies data in existing rows

```
UPDATE PROJECT
SET Plocation = 'Bellaire', Dnum = 5
WHERE Pnumber = 10;

UPDATE EMPLOYEE
SET Salary = Salary * 1.1
WHERE Dno = 5;
```



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☐ The SELECT-FROM-WHERE Structure of Basic SQL Queries

```
SELECT <attribute list>
FROM 
WHERE <condition>;
```

where

- <attribute list> is a list of attribute names whose values are to be retrieved by the query.
- is a list of the relation names required to process the query.
- <condition> is a conditional (Boolean) expression that identifies the tuples to be retrieved by the query

Logical comparison operators: =, <, <=, >, >=, and <>



Query 0. Retrieve the birth date and address of the employee(s) whose name is 'John B. Smith'

```
SELECT Bdate, Address
FROM EMPLOYEE
WHERE Fname = 'John' AND Minit = 'B' AND
Lname = 'Smith';
```

Bdate	Address	
1965-01-09	731Fondren, Houston, TX	



Query 1. Retrieve the name and address of all employees who work for the 'Research' department

```
SELECT Fname, Lname, Address
FROM EMPLOYEE, DEPARTMENT
WHERE Dname = 'Research'
AND Dnumber = Dno;
```

Fname	Lname	<u>Address</u>	
John	Smith	731 Fondren, Houston, TX	
Franklin	Wong	638 Voss, Houston, TX	
Ramesh	Narayan	975 Fire Oak, Humble, TX	
Joyce English		5631 Rice, Houston, TX	



Query 2. For every project located in 'Stafford', list the project number, the controlling department number, and the department manager's last name, address, and birth date

```
SELECT Pnumber, Dnum, Lname, Address, Bdate
FROM PROJECT, DEPARTMENT, EMPLOYEE
WHERE Dnum = Dnumber AND Mgr_ssn = Ssn AND
Plocation = 'Stafford
```

Pnumber	Dnum	Lname	<u>Address</u>	<u>Bdate</u>
10	4	Wallace	291Berry, Bellaire, TX	1941-06-20
30	4	Wallace	291Berry, Bellaire, TX	1941-06-20



□ Ambiguous Attribute Names, Aliasing, Renaming, and Tuple Variables.

```
SELECT Fname, Name, Address
FROM EMPLOYEE, DEPARTMENT
WHERE Name = 'Research' AND Dnumber = Dnumber;

SELECT E.Fname, E.Name, E.Address
FROM EMPLOYEE E, DEPARTMENT D
WHERE D.Name = 'Research' AND D.Dnumber = E.Dnumber;
```

```
SELECT E.Fname, E.Lname, S.Fname, S.Lname
FROM EMPLOYEE AS E, EMPLOYEE AS S
WHERE E.Super ssn = S.Ssn;
```



☐ Unspecified WHERE Clause and Use of the Asterisk

■ A missing WHERE clause indicates no condition on tuple selection; hence, all tuples of the relation specified in the FROM clause qualify and are selected for the query result.

SELECT Ssn
FROM EMPLOYEE;

■ If more than one relation is specified in the FROM clause and there is no WHERE clause, then the CROSS PRODUCT—all possible tuple combinations—of these relations is selected.

SELECT Ssn, Dname **FROM** EMPLOYEE, DEPARTMENT;



□ Tables as Sets in SQL

- Duplicate elimination is an expensive operation. One way to implement it is to sort the tuples first and then eliminate duplicates.
- The user may want to see duplicate tuples in the result of a query.
- When an aggregate function is applied to tuples, in most cases we do not want to eliminate duplicates.

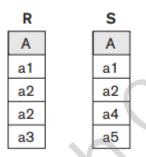
```
SELECT Salary
FROM EMPLOYEE;
```

SELECT DISTINCT Salary **FROM** EMPLOYEE;

Salary
30000
40000
25000
43000
38000
25000
25000
55000



- □ Set union (UNION), set difference (EXCEPT), and set intersection (INTERSECT) operations
 - □ SQL also has corresponding multiset operations, which are followed by the keyword ALL (UNION ALL, EXCEPT ALL, INTERSECT ALL). Their results are multisets (duplicates are not eliminated)



Two tables: R(A) and S(A)

T	
Α	
a1	
a1	
a2	
a2	
a2	
a3	
a4	
a5	

R(A)UNION ALL S(A)



R(A) EXCEPT ALL S(A)



R(A) INTERSECT ALL S(A)



Query 4. Make a list of all project numbers for projects that involve an employee whose last name is 'Smith', either as a worker or as a manager of the department that controls the project

```
(SELECT DISTINCT Pnumber
FROM PROJECT, DEPARTMENT, EMPLOYEE
WHERE Dnum = Dnumber AND Mgr_ssn = Ssn
AND Lname = 'Smith')
UNION
(SELECT DISTINCT Pnumber
FROM PROJECT, WORKS_ON, EMPLOYEE
WHERE Pnumber = Pno AND Essn = Ssn
AND Lname = 'Smith');
```



■ Substring Pattern Matching and Arithmetic Operators **LIKE** comparison operator (string pattern matching): % replaces an arbitrary number of zero or more characters underscore (_) replaces a single character ESCAPE (\): escape character For example: 'AB_CD\%EF' ESCAPE '\' represents the literal string 'AB_CD%EF' because \ is specified as the escape character Query: Retrieve all employees whose address is in Houston, Texas SELECT Fname, Lname FROM EMPLOYEE WHERE Address LIKE '%Houston, TX%'; Query: retrieve all employees who were born during the 1970s SELECT Fname, Lname FROM EMPLOYEE WHERE Bdate LIKE '



- Another feature allows the use of arithmetic in queries.
 - ☐ The standard arithmetic operators for addition (+), subtraction (−), multiplication (*), and division (/) can be applied to numeric values or attributes with umeric domains
 - we can rename an attribute in the query result using AS in the SELECT clause

Query: Show the resulting salaries if every employee working on the 'ProductX' project is given a 10% raise

```
SELECT E.Fname, E.Lname, 1.1 * E.Salary AS Increased_sal
FROM EMPLOYEE AS E, WORKS_ON AS W, PROJECT AS P
WHERE E.Ssn = W.Essn AND W.Pno = P.Pnumber AND
P.Pname = 'ProductX';
```



- Boolean Operators: AND, OR, and NOT Operators for customizing conditions in WHERE clause
 - Boolean query a without use of parentheses

SELECT ProductDescription, ProductFinish, ProductStandardPrice

FROM Product_T

WHERE ProductDescription LIKE '%Desk'

OR ProductDescription LIKE '%Table' AND ProductStandardPrice > 300:

By default, processing order of Boolean operators is NOT, then

AND, then OR

■ With parentheses:

SELECT ProductDescription, ProductFinish, ProductStandardPrice

FROM Product T;

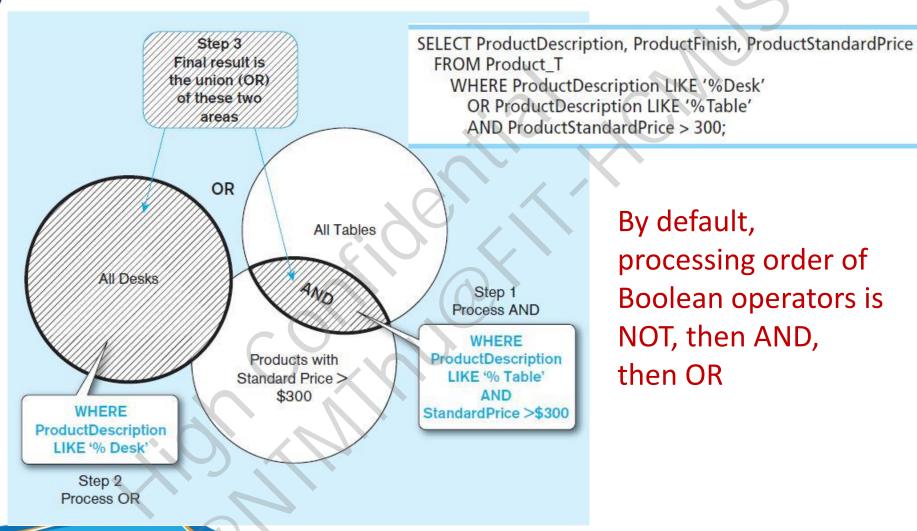
WHERE (ProductDescription LIKE '%Desk' OR ProductDescription LIKE '%Table')

AND ProductStandardPrice > 300;

These override the normal precedence of Boolean operators. this case parentheses make the OR take place before the AND



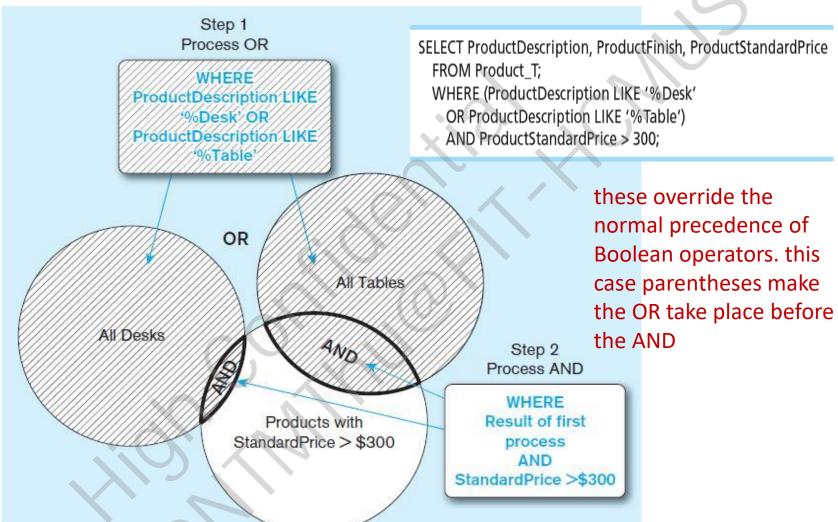
4.0 Boolean query A without use of parentheses



By default, processing order of Boolean operators is NOT, then AND, then OR



4.0 Boolean query B with use of parentheses





☐ Comparison operator: BETWEEN, NOT BETWEEN

Query: Retrieve all employees in department 5 whose salary is between \$30,000 and \$40,000.

```
SELECT *
FROM EMPLOYEE
WHERE (Salary BETWEEN 30000 AND 40000) AND Dno = 5;
```

☐ Ordering of Query Results: ORDER BY

Query: Retrieve a list of employees and the projects they are working on, ordered by department and, within each department, ordered alphabetically by last name, then first name

```
FROM DEPARTMENT AS D, EMPLOYEE AS E,

WORKS_ON AS W, PROJECT AS P

WHERE D.Dnumber = E.Dno AND E.Ssn = W.Essn

AND W.Pno =P.Pnumber
```

ORDER BY D. Dname, E. Lname, E. Fname; DESC:descending order

ASC: ascending order



Summary of **Basic SQL Retrieval Queries**

```
SELECT <attribute list
FROM 
[WHERE <condition>
[ORDER BY <attribute list> ];
```



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□ Comparisons Involving NULL

```
SELECT Fname, Lname
FROM EMPLOYEE
WHERE Super_ssn IS NULL;
```

```
SELECT Fname, Lname
FROM EMPLOYEE
WHERE Super ssn IS NOT NULL;
```



■ Multiset Comparisons: ■ IN/NOT IN: to match a list of values, consider using IN ■ EXISTS/NOT EXISTS □ ALL ANY/SOME **SELECT** column name(s) FROM table name WHERE column name IN (value1, value2, ...); SELECT column name(s) FROM table name WHERE column name IN (SELECT STATEMENT); **SELECT** column name(s) FROM table name WHERE EXISTS (SELECT STATEMENT);



■ Set/Multiset Comparisons:

```
SELECT DISTINCT Pnumber
FROM PROJECT
WHERE Pnumber IN
      (SELECT Pnumber
       FROM PROJECT, DEPARTMENT, EMPLOYEE
       WHERE Dnum = Dnumber AND Mgr_ssn = Ssn AND
             Lname = 'Smith' )
      OR
      Pnumber IN
      (SELECT Pno
       FROM WORKS ON, EMPLOYEE
       WHERE Essn = Ssn AND Lname = 'Smith');
```



■ Set/Multiset Comparisons:

```
SELECT DISTINCT Essn
FROM WORKS_ON
WHERE Pno IN ( 1,2,3)
```



- Nested Queries:
 - Noncorrelated—executed once for the entire outer query
 - □ Correlated—executed once for each row returned by the outer query

(Outer query)

SELECT <danh sách các cột>

FROM <danh sách các bảng>

WHERE <so sánh tập hợp> (

SELECT <danh sách các cột>

FROM <danh sách các bảng>

WHERE <điều kiện>)

(Subquery)



4.0 Correlated vs. Noncorrelated Subqueries

- **☐** Noncorrelated subqueries:
 - Do not depend on data from the outer query
 - Execute once for the entire outer query
- **□** Correlated subqueries:
 - Make use of data from the outer query
 - Execute once for each row of the outer query
 - □ Can use the EXISTS operator



- Nested Queries:
 - □ Noncorrelated—executed once for the entire outer query

```
FROM EMPLOYEE E The IN operator will test to see if the E.Dno value of a row where E.Dno IN is included in the list returned from the subquery

(SELECT D.Dnumber
FORM DEPARTMENT D
WHERE D.Name = 'Research');
```

Subquery is embedded in parentheses. In this case it returns a list that will be used in the WHERE clause of the outer query



- Nested Queries:
 - Noncorrelated—executed once for the entire outer query

```
FROM EMPLOYEE E The IN operator will test to see if the E.SSN value of a row WHERE E.SSN IN is included in the list returned from the subquery

(SELECT W.ESSN
FORM WORK ON W);
```

Subquery is embedded in parentheses. In this case it returns a list that will be used in the WHERE clause of the outer query



■ Nested Queries:

```
Correlated—executed once for each row returned by the The EXISTS operator will return a TRUE value if the subquery resulted in a non-empty set, otherwise it returns a FALSE

SELECT E.Fname, E.Name, E.Address

FROM EMPLOYER E

WHERE EXISTS

(SELECT *
FORM WORK ON W
WHERE W.ESSN = E.SSN );
```

→ A correlated subquery always refers to an attribute from a table referenced in the outer query

The subquery is testing for a value that comes from the outer query



■ Nested Queries:

☐ Correlated—executed once for each row returned by the outer query



■ Nested Queries:

□ Correlated—executed once for each row returned by the outer query



■ Nested Queries:

□ Correlated—executed once for each row returned by the outer query

```
SELECT E.Fname, E.Lname
FROM EMPLOYEE AS E
WHERE NOT EXISTS

(SELECT *
FROM DEPENDENT AS D
WHERE E.Ssn = D.Essn
AND E.Sex = D.Sex
AND E.Fname = D.Dependent name)
```

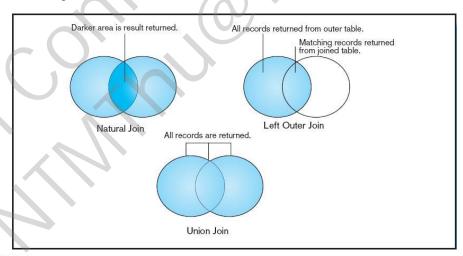


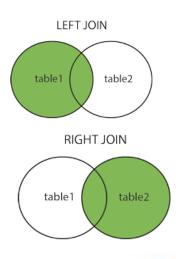
- ☐ Joined Tables in SQL and Outer Joins
 - Join—a relational operation that causes two or more tables with a common domain to be combined into a single table or view
 - **Equi-join**—a join in which the joining condition is based on equality between values in the common columns; common columns appear redundantly in the result table
 - Natural join—an equi-join in which one of the duplicate columns is eliminated in the result table

The common columns in joined tables are usually the primary key of the dominant table and the foreign key of the dependent table in 1:M relationships.



- Joined Tables in SQL and Outer Joins
 - Outer join (Left/Right Join)—a join in which rows that do not have matching values in common columns are nonetheless included in the result table (as opposed to inner join, in which rows must have matching values in order to appear in the result table)
 - ☐ Union join/Cross join—includes all columns from each table in the join, and an instance for each row of each table







☐ Joined Tables in SQL and Outer Joins

```
SELECT Fname, Lname, Address
FROM (EMPLOYEE [inner] JOIN DEPARTMENT
          ON Dno = Dnumber)
WHERE Dname = 'Research';

SELECT Fname, Lname, Address
FROM (EMPLOYEE NATURAL JOIN
  (DEPARTMENT AS DEPT (Dname, Dno, Mssn, Msdate)))
WHERE Dname = 'Research';
```



☐ Joined Tables in SQL and Outer Joins

```
SELECT Pnumber, Dnum, Lname, Address, Bdate
FROM ((PROJECT JOIN DEPARTMENT ON Dnum = Dnumber)
        JOIN EMPLOYEE ON Mgr_ssn = Ssn)
WHERE Plocation = 'Stafford';

SELECT E.Lname AS Employee_name,
        S.Lname AS Supervisor_name
FROM (EMPLOYEE AS E LEFT OUTER JOIN EMPLOYEE AS S
        ON E.Super ssn = S.Ssn);
```



Aggregate Functions in SQL

Query 19. Find the sum of the salaries of all employees, the maximum salary, the minimum salary, and the average salary

```
SELECT SUM (Salary), MAX (Salary),
MIN (Salary), AVG (Salary)
FROM EMPLOYEE;
```



□ Aggregate Functions in SQL

Query 20: Find the sum of the salaries of all employees of the 'Research' department, as well as the maximum salary, the minimum salary, and the average salary in this department

```
SELECT SUM (Salary), MAX (Salary), MIN (Salary),
AVG (Salary)
FROM (EMPLOYEE JOIN DEPARTMENT ON Dno = Dnumber)
WHERE Dname = 'Research';
Query 21: Retrieve the total number of employees in the company
SELECT COUNT (*)
FROM EMPLOYEE;
```

Query 22: The number of employees in the 'Research' department

```
SELECT COUNT (*)
FROM EMPLOYEE, DEPARTMENT
WHERE DNO = DNUMBER
AND DNAME = 'Research';
```



More Complex SQL Retrieval Queries

□ Aggregate Functions in SQL

Query 23. Count the number of distinct salary values in the database

```
SELECT COUNT (DISTINCT Salary)
FROM EMPLOYEE;
```

Query 24: Retrieve the names of all employees who have two or more dependents



More Complex SQL Retrieval Queries

☐ Grouping: The GROUP BY and HAVING Clauses

Query 25. For each department, retrieve the department number, the number of employees in the department, and their average salary.

```
SELECT Dno, COUNT (*), AVG (Salary)
FROM EMPLOYEE
GROUP BY Dno;
```

Query 26. For each project, retrieve the project number, the project name, and the number of employees who work on that project.

```
SELECT Pnumber, Pname, COUNT (*)
FROM PROJECT, WORKS_ON
WHERE Pnumber = Pno
GROUP BY Pnumber, Pname;
```



More Complex SQL Retrieval Queries

☐ Grouping: The GROUP BY and HAVING Clauses

Query 27. For each project on which more than two employees work, retrieve the project number, the project name, and the number of employees who work on.

```
SELECT Pnumber, Pname, COUNT (*)
FROM PROJECT, WORKS_ON
WHERE Pnumber = Pno
GROUP BY Pnumber, Pname
HAVING COUNT (*) > 2;
```



• R÷S

R	A	В	С	D	Е
	α	a	α	a	1
	α	a	γ	a	1
	α	a	γ	b	1
	β	a	γ	a	1
	β	a	γ	b	3
	γ	a	γ	a	1
	γ	a	γ	b	1
	γ	a	β	b	1

S	D	E
b _i	a	1
	b	1

R÷S	A	В	С
a _i	α	a	γ
•	γ	a	γ



Using EXCEPT

```
SELECT R1.A, R1.B, R1.C
FROM R R1
WHERE NOT EXISTS (
                    (SELECT S.D, S.E FROM S)
                    EXCEPT
                    (SELECT R2.D, R2.E
                    FROM R R2
                    WHERE R1.A=R2.A AND R1.B=R2.B
                    AND R1.C=R2.C)
```



Using NOT EXISTS

```
SELECT R1.A, R1.B, R1.C
FROM R R1
WHERE NOT EXISTS (
       SELECT*
       FROM S
       WHERE NOT EXISTS
              SELECT*
              FROM R R2
              WHERE R2.D=S.D AND R2.E=S.E
              AND R1.A=R2.A AND R1.B=R2.B AND R1.C=R2.C ))
```



Using Count Function

```
SELECT R.A

FROM R

[WHERE R.B IN (SELECT S.B FROM S [WHERE <ĐK>]]

GROUP BY R.A

HAVING COUNT(DISTINCT R.B) = ( SELECT COUNT(S.B)

FROM S

[WHERE <ĐK>])
```



Using EXCEPT

Query: Find the names of employees who work on all the projects controlled by department number 5

```
SELECT E.Fname, E.Lname
FROM EMPLOYEE E, WORK_ON W1
WHERE E.Ssn = W1.Essn
        AND NOT EXISTS (
                            (SELECT P.Pnumber
                             FROM PROJECT P
                             WHERE P.Dnum = 5)
                             EXCEPT
                            ( SELECT W2.Pno
                             FROM WORK_ON W2
                              WHERE W1.SSN=W2.Essn)
```



Using EXISTS

Query: Find the names of employees who work on all the projects controlled by department number 5

```
SELECT E.Fname, E.Lname

FROM EMPLOYEE E ,WORK_ON W1,

WHERE E.Ssn = W1.Essn

AND NOT EXISTS (SELECT *

FROM PROJECT P

WHERE P.Dnum = 5 AND

NOT EXISTS (SELECT *

FROM WORK_ON W2

WHERE W1.ESsn=W2.Essn

AND P.Dnum = W2.Pno ))
```



Using Count Function

Query: Find the names of employees who work on all the projects controlled by department number 5

```
SELECT E.Fname, E.Lname

FROM EMPLOYEE E, WORK_ON W

WHERE E.Ssn = W.Essn

AND W.Pno IN (SELECT P.Pnumber

FROM PROJECT P

WHERE P.Dnum = 5)

GROUP BY W.Essn

HAVING COUNT(DISTINCT W.Pno) = (SELECT COUNT(P.Pnumber)

FROM PROJECT P

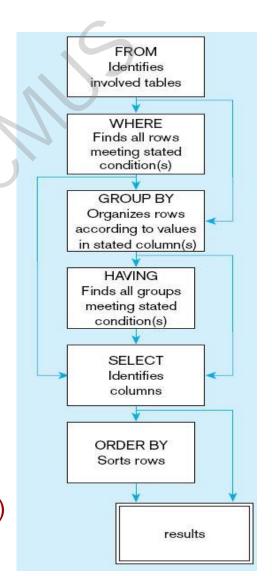
WHERE P.Dnum = 5)
```



Summary of SQL Queries

```
SELECT <attribute and function list>
FROM 
[ WHERE <condition> ]
[ GROUP BY <grouping attribute(s)> ]
[ HAVING <group condition> ]
[ ORDER BY <attribute list> ];
```

SQL statement processing order (based on van der Lans, 2006 p.100)





- SQL Overview
- Basic SQL
 - SQL Data Definition and Data Types
 - Specifying Constraints in SQL
 - ☐ Insert/Update/Delete Statement in SQL
 - Basic Retrieval Queries in SQL
- □ Advanced SQL
 - Complex SQL Retrieval Queries
 - □ View
 - Ensuring Transaction Integrity
 - ☐ Specifying Constraints as Assertions & Action Trigger
- ☐ Tips for Developing Queries



Views (Virtual Tables) in SQL

- □ Concept of a View in SQL
 - □ Views provide users controlled access to tables
 - Base Table—table containing the raw data
 - Dynamic View
 - A "virtual table" created dynamically upon request by a user
 - No data actually stored; instead data from base table made available to user
 - Based on SQL SELECT statement on base tables or other views
 - Materialized View
 - Copy or replication of data
 - Data actually stored
 - ☐ Must be refreshed periodically to match corresponding base tables



4.0 Views (Virtual Tables) in SQL

- □ Specification of Views in SQL
 - □ CREATE VIEW
 - DROP VIEW
 - We can now specify SQL queries on a view—or virtual table—in the same way we specify queries involving base tables

```
CREATE VIEW WORKS_ON1
```

```
AS SELECT Fname, Lname, Pname, Hours
FROM EMPLOYEE, PROJECT, WORKS_ON
WHERE Ssn = Essn AND Pno = Pnumber;
```

WORKS_ON1

```
Fname Lname Pname Hours
```

```
CREATE VIEW DEPT_INFO(Dept_name, No_of_emps, Total_sal)

AS SELECT Dname, COUNT (*), SUM (Salary)

FROM DEPARTMENT, EMPLOYEE

WHERE Dnumber = Dno

GROUP BY Dname;

DEPT_INFO

Dept_name No_of_emps
```

Total sal



4.0 Views (Virtual Tables) in SQL

□ Specification of Views in SQL

- We can now specify SQL queries on a view—or virtual table
- DROP VIEW

```
SELECT Fname, Lname
FROM WORKS_ON1
WHERE Pname = 'ProductX';
```

```
DROP VIEW WORKS ON1;
```



Views (Virtual Tables) in SQL

- Advantages of Views
 - □ Simplify query commands
 - Assist with data security (but don't rely on views for security, there are more important security measures)
 - Enhance programming productivity
 - □ Contain most current base table data
 - Use little storage space
 - Provide customized view for user
 - ☐ Establish physical data independence
- Disadvantages of Views
 - ☐ Use processing time each time view is referenced
 - May or may not be directly updateable



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Ensuring Transaction Integrity

- ☐ Transaction = A discrete unit of work that must be completely processed or not processed at all
 - May involve multiple updates
 - ☐ If any update fails, then all other updates must be cancelled
- ☐ SQL commands for transactions
 - □ BEGIN TRANSACTION/END TRANSACTION
 - Marks boundaries of a transaction
 - □ COMMIT
 - Makes all updates permanent
 - □ ROLLBACK
 - Cancels updates since the last COMMIT



An SQL Transaction sequence (in pseudocode)

BEGIN transaction INSERT OrderID, Orderdate, CustomerID into Order T; INSERT OrderID, ProductID, OrderedQuantity into OrderLine_T; INSERT OrderID, ProductID, OrderedQuantity into OrderLine_T; INSERT OrderID, ProductID, OrderedQuantity into OrderLine_T; **END** transaction Invalid ProductID entered. Valid information inserted. COMMIT work. Transaction will be ABORTED. ROLLBACK all changes made to Order_T. All changes to data All changes made to Order_T are made permanent. and OrderLine T are removed. Database state is just as it was before the transaction began.

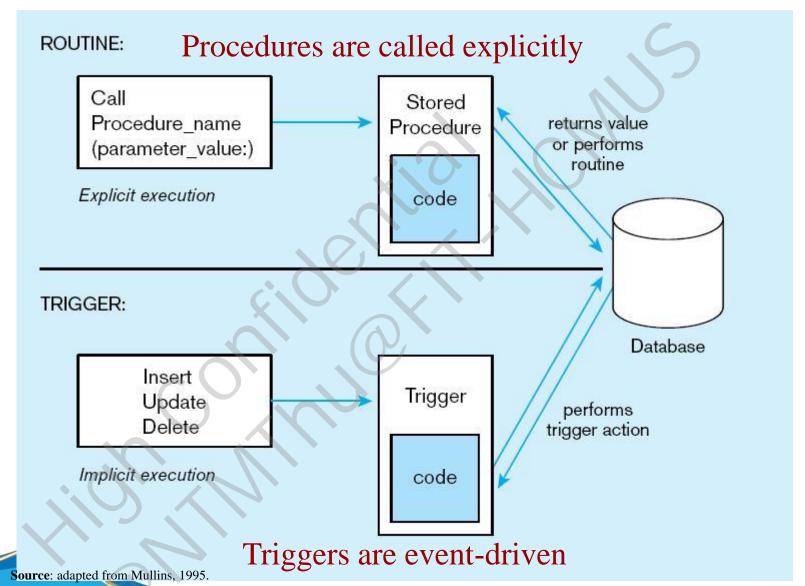


□ Routines and Triggers

- □ Routines
 - Program modules that execute on demand
- ☐ Functions—routines that return values and take input parameters
- □ Procedures—routines that do not return values and can take input or output parameters
- ☐ Triggers—routines that execute in response to a database event (INSERT, UPDATE, or DELETE)



Triggers contrasted with stored procedures (based on Mullins 1995)



Introduction to Database



Specifying Constraints as Assertions and Actions as Triggers

- Specifying General Constraints as Assertions in SQL
 - □ CREATE ASSERTION statement
 - Each assertion is given a constraint name and is specified via a condition similar to the WHERE clause of an SQL query

For example, to specify the constraint that the salary of an employee must not be greater than the salary of the manager of the department that the employee works for in SQL, we can write the following assertion:



Specifying Constraints as Assertions and Actions as Triggers

- Introduction to Triggers in SQL
 - □ CREATE TRIGGER statement
 - Triggers are used when you need to perform, under specified conditions, a certain action as the result of some database event (e.g., the execution of a DML statement such as INSERT, UPDATE, or DELETE or the DDL statement ALTER TABLE)
 - Thus, a trigger has three parts—the event, the condition, and the action—and these parts are reflected in the coding structure for triggers

Events are: inserting a new employee record, changing an employee's salary, or changing an employee's supervisor

CREATE TRIGGER SALARY VIOLATION

BEFORE INSERT OR UPDATE OF SALARY, SUPERVISOR SSN ON EMPLOYEE

FOR EACH ROW

WHEN (NEW.SALARY >

Condition that determines whether the rule action should be executed

SELECT SALARY FROM EMPLOYEE

WHERE SSN = NEW.SUPERVISOR SSN))

INFORM_SUPERVISOR(NEW.Supervisor_ssn, NEW.Ssn);

The action to be taken:

- The action is usually a sequence of SQL statements
 - The action is to execute the stored procedure

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Tips for Developing Queries

- Be familiar with the data model (entities and relationships)
- Understand the desired results
- Know the attributes desired in results
- Identify the entities that contain desired attributes
- □ Review ERD
- ☐ Construct a WHERE equality for each link
- ☐ Fine tune with GROUP BY and HAVING clauses if needed
- Consider the effect on unusual data



Query Efficiency Considerations

- □ Instead of SELECT *, identify the specific attributes in the SELECT clause; this helps reduce network traffic of result set
- ☐ Limit the number of subqueries; try to make everything done in a single query if possible
- ☐ If data is to be used many times, make a separate query and store it as a view



Guidelines for Better Query Design

- Understand how indexes are used in query processing
- □ Keep optimizer statistics up-to-date
- ☐ Use compatible data types for fields and literals
- □ Write simple queries
- □ Break complex queries into multiple simple parts
- ☐ Don't nest one query inside another query
- Don't combine a query with itself (if possible avoid self-joins)



- ☐ Create temporary tables for groups of queries
- ☐ Combine update operations
- □ Retrieve only the data you need
- ☐ Don't have the DBMS sort without an index
- ☐ Learn!
- ☐ Consider the total query processing time for ad hoc queries



