

Standard Code Library

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Chapter 1

Algorithms and Datastructures

1.1 *High Precision in C Plus Plus

Comming Soon

1.2 *Fraction Class

Comming Soon

1.3 Splay Tree

注意初始化内存池和 null 节点，以及根据需要修改 update 和 relax，区间必须是 1-based

```
1  const int MAX_NODE = 50000 + 10;
2  const int INF = 2000000000;
3
4  struct Node *null;
5
6  struct Node
7  {
8      int rev, add;
9      int val, maxv, size;
10     Node *ch[2], *p;
11
12     void set(Node *t, int _d) {
13         ch[_d] = t;
14         t->p = this;
15     }
16     int dir() {
17         return this == p->ch[1];
18     }
19     void update() {
20         maxv = max(max(ch[0]->maxv, ch[1]->maxv), val);
21         size = ch[0]->size + ch[1]->size + 1;
22     }
23     void relax() {
24         if (add) {
25             ch[0]->appAdd(add);
26             ch[1]->appAdd(add);
27             add = 0;
28         }
29         if (rev) {
30             ch[0]->appRev();
31             ch[1]->appRev();
32             rev = false;
33         }
34     }
35     void appAdd(int x) {
```

```

36         if (this == null) return;
37         add += x;
38         val += x;
39         maxv += x;
40     }
41     void appRev() {
42         if (this == null) return;
43         rev ^= true;
44         swap(ch[0], ch[1]);
45     }
46 };
47
48 Node nodePool[MAX_NODE], *curNode;
49
50 Node *newNode(int val = 0)
51 {
52     Node *t = curNode++;
53     t->maxv = t->val = val;
54     t->rev = t->add = 0;
55     t->size = 1;
56     t->ch[0] = t->ch[1] = t->p = null;
57     return t;
58 }
59
60 struct Splay
61 {
62     Node *root;
63
64     Splay() {
65         root = newNode();
66         root->set(newNode(), 0);
67         root->update();
68     }
69
70     Splay(int *a, int N) { //sequence is 1-based
71         root = build(a, 0, N + 1);
72     }
73
74     Node* build(int *a, int l, int r) {
75         if (l > r) return null;
76         int mid = l + r >> 1;
77         Node *t = newNode(a[mid]);
78         t->set(build(a, l, mid - 1), 0);
79         t->set(build(a, mid + 1, r), 1);
80         t->update();
81         return t;
82     }
83
84     void rot(Node *t)
85     {
86         Node *p = t->p; int d = t->dir();
87         p->relax(); t->relax();
88         if (p == root) root = t;
89         p->set(t->ch[! d], d);
90         p->p->set(t, p->dir());
91         t->set(p, ! d);
92         p->update();
93     }
94
95     void splay(Node *t, Node *f = null)
96     {
97         for(t->relax(); t->p != f; ) {
98             if (t->p->p == f) rot(t);

```

```

99         else t->dir() == t->p->dir() ? (rot(t->p), rot(t)) : (rot(t), rot(t));
100     }
101     t->update();
102 }
103
104 Node* getKth(int k) {
105     Node *t = root;
106     int tmp;
107     for( ; ; ) {
108         t->relax();
109         tmp = t->ch[0]->size + 1;
110         if (tmp == k) return t;
111         if (tmp < k) {
112             k -= tmp;
113             t = t->ch[1];
114         } else
115             t = t->ch[0];
116     }
117 }
118
119 //make range[l,r] root->ch[1]->ch[0]
120 //make range[x+1,x] to add something after position x
121 void getRng(int l, int r) {
122     r += 2;
123     Node *p = getKth(l);
124     Node *q = getKth(r);
125     splay(p); splay(q, p);
126 }
127
128 void addRng(int l, int r, int x) {
129     getRng(l, r);
130     root->ch[1]->ch[0]->appAdd(x);
131 }
132
133 void revRng(int l, int r) {
134     getRng(l, r);
135     root->ch[1]->ch[0]->appRev();
136 }
137
138 int maxvRng(int l, int r) {
139     getRng(l, r);
140     return root->ch[1]->ch[0]->maxv;
141 }
142 };
143
144 void initNull()
145 {
146     curNode = nodePool;
147     null = curNode++;
148     null->size = 0;
149     null->maxv = - INF;
150 }

```

1.4 Link/Cut Tree

根据需求修改 Node 中的 relax 和 update 函数, 修改 access, 以及 Node 的构造函数, 注意初始化内存池和 null 节点

```

1 struct Node
2 {
3     Node *ch[2], *p;
4     int isroot;
5     bool dir();
6     void set(Node*, bool);

```

```

7     void update();
8     void relax();
9 } *null;
10
11 void rot(Node *t)
12 {
13     Node *p = t->p; bool d = t->dir();
14     p->relax(); t->relax();
15     p->set(t->ch[! d], d);
16     if (p->isroot) t->p = p->p, swap(p->isroot, t->isroot);
17     else p->p->set(t, p->dir());
18     t->set(p, ! d);
19     p->update();
20 }
21
22 void Splay(Node *t)
23 {
24     for(t->relax(); ! t->isroot; ) {
25         if (t->p->isroot) rot(t);
26         else t->dir() == t->p->dir() ? (rot(t->p), rot(t)) : (rot(t), rot(t));
27     }
28     t->update();
29 }
30
31 void Access(Node *t)
32 {
33     for(Node *s = null; t != null; s = t, t = t->p) {
34         Splay(t);
35         t->ch[1]->isroot = true;
36         s->isroot = false;
37         t->ch[1] = s;
38         t->update();
39     }
40 }
41 bool Node::dir()
42 {
43     return this == p->ch[1];
44 }
45 void Node::set(Node *t, bool _d)
46 {
47     ch[_d] = t; t->p = this;
48 }
49 void Node::Update()
50 {
51 }
52 }
53 void Node::Relax()
54 {
55     if (this == Null) return;
56 }
57 }

```

1.5 Binary Heap

双射堆, $\text{ind}[v]$ 表示标号为 v 的节点在堆中的位置

```

1 const int MAX_V = 100000 + 10;
2 struct Heap
3 {
4     int tot;
5     int a[MAX_V], h[MAX_V], ind[MAX_V];
6     void exchange(int i, int j) {
7         swap(h[i], h[j]);

```

```

8         swap(ind[h[i]], ind[h[j]]);
9     }
10    inline int val(int x) {
11        return a[h[x]];
12    }
13    void fixUp(int x) {
14        if (x / 2 && val(x / 2) < val(x))
15            exchange(x, x / 2), fixUp(x / 2);
16    }
17    void fixDown(int x) {
18        int p = x * 2; if (p > tot) return;
19        if (p < tot && val(p + 1) > val(p)) ++ p;
20        if (val(p) > val(x))
21            exchange(p, x), fixDown(p);
22    }
23    void Update(int i, int x) {
24        a[i] = x;
25        fixUp(ind[i]);
26        fixDown(ind[i]);
27    }
28    int top() {
29        return h[1];
30    }
31    void pop() {
32        exchange(1, tot);
33        -- tot;
34        fixDown(1);
35    }
36    void insert(int i, int x) {
37        ++ tot;
38        h[tot] = i;
39        ind[i] = tot;
40        a[i] = x;
41        fixUp(tot);
42    }
43 } H;

```

1.6 Leftist Tree

没写 delete 操作，注意初始化内存池和 null 节点

```

1 struct Node
2 {
3     int dis, val;
4     Node *ch[2];
5 } *null;
6
7 Node* merge(Node *u, Node *v)
8 {
9     if (u == null) return v;
10    if (v == null) return u;
11    if (u->val < v->val) swap(u, v);
12    u->ch[1] = merge(u->ch[1], v);
13    if (u->ch[1]->dis > u->ch[0]->dis)
14        swap(u->ch[1], u->ch[0]);
15    u->dis = u->ch[1]->dis + 1;
16    return u;
17 }
18
19 Node* newNode(int w)
20 {
21     Node *t = totNode++;
22     t->ch[0] = t->ch[1] = null;

```

```

23     t->val = w; t->dis = 0;
24     return t;
25 }

```

1.7 Treap

包含 build, insert 和 erase , 执行时注意初始化内存池和 null 节点

```

1  struct Node *null;
2
3  struct Node
4  {
5      int key, val, size;
6      Node *ch[2];
7      Node() {
8          key = INT_MAX;
9          val = size = 0;
10     }
11     Node(int __val) {
12         size = 1;
13         val = __val;
14         key = bigRand();
15         ch[0] = ch[1] = null;
16     }
17     int bigRand() {
18         return rand() * RAND_MAX + rand();
19     }
20     void update() {
21         size = ch[0]->size + ch[1]->size + 1;
22     }
23 };
24
25 struct Treap
26 {
27     Node *root;
28     Treap() {
29         root = null;
30     }
31     void rot(Node *&t, int d) {
32         Node *p = t->ch[d]; t->ch[d] = p->ch[! d]; p->ch[! d] = t;
33         t->update(); p->update();
34         t = p;
35     }
36
37     void insert(Node *&t, int x) {
38         if (t == null) {
39             t = new Node(x);
40             return;
41         }
42         int dir = x >= t->val;
43         insert(t->ch[dir], x);
44         if (t->ch[dir]->key < t->key)
45             rot(t, dir);
46         else
47             t->update();
48     }
49
50     void erase(Node *&t, int x) {
51         if (t == null)
52             return;
53         if (t->val == x) {
54             int dir = t->ch[1]->key < t->ch[0]->key;
55             if (t->ch[dir] == null) {

```



```

56         delete t;
57         t = null;
58         return;
59     }
60     rot(t, dir);
61     erase(t->ch[! dir], x);
62     t->update();
63     return;
64 }
65 bool dir = x > t->val;
66 erase(t->ch[dir], x);
67 t->update();
68 }
69
70 void insert(int x) {
71     insert(root, x);
72 }
73
74 void erase(int x) {
75     erase(root, x);
76 }
77 };

```

1.8 Segment Tree

包含建树和区间操作样例，没有写具体操作

```

1 struct Tree
2 {
3     int l, r;
4     Tree *ch[2];
5     Tree() {}
6     Tree(int _l, int _r, int *sqn) {
7         l = _l; r = _r;
8         if (l + 1 == r)
9             return;
10        int mid = l + r >> 1;
11        ch[0] = new Tree(l, mid, sqn);
12        ch[1] = new Tree(mid, r, sqn);
13    }
14
15    void insert(int p, int x) {
16        if (p < l || p >= r)
17            return;
18        //some operations
19        if (l + 1 == r)
20            return;
21        ch[0]->insert(p, x);
22        ch[1]->insert(p, x);
23    }
24
25    int query(int _l, int _r, int x) {
26        if (_r <= l || _l >= r)
27            return 0;
28        if (_l <= l && _r >= r)
29            // return information in [l, r)
30            //merge ch[0]->query(_l, _r, x), ch[1]->query(_l, _r, x) and return
31        }
32    };

```

1.9 Heavy-Light Decomposition

包含 BFS 剖分过程，线段树部分视题目而定

```
1 struct Tree()
2 {
3
4 };
5
6 int father[MAX_N], size[MAX_N], depth[MAX_N];
7 int bfsOrd[MAX_N], pathId[MAX_N], ordInPath[MAX_N], sqn[MAX_N];
8 Tree *root[MAX_N];
9
10 void doBfs(int s)
11 {
12     int *que = bfsOrd;
13     int qh = 0, qt = 0;
14     father[s] = -1; depth[s] = 0;
15
16     for(que[qt++] = s; qh < qt; ) {
17         int u = que[qh++];
18         foreach(iter, adj[u]) {
19             int v = *iter;
20             if (v == father[u])
21                 continue;
22             father[v] = u;
23             depth[v] = depth[u] + 1;
24             que[qt++] = v;
25         }
26     }
27 }
28
29 void doSplit()
30 {
31     for(int i = N - 1; i >= 0; -- i) {
32         int u = bfsOrd[i];
33         size[u] = 1;
34         foreach(iter, adj[u]) {
35             int v = *iter;
36             if (v == father[u])
37                 continue;
38             size[u] += size[v];
39         }
40     }
41
42     memset(pathId, -1, sizeof pathId);
43     for(int i = 0; i < N; ++ i) {
44         int top = bfsOrd[i];
45         if (pathId[top] != -1)
46             continue;
47
48         int cnt = 0;
49         for(int u = top; u != -1; ) {
50             sqn[cnt] = val[u];
51             ordInPath[u] = cnt;
52             pathId[u] = top;
53             ++ cnt;
54
55             int next = -1;
56             foreach(iter, adj[u]) {
57                 int v = *iter;
58                 if (v == father[u])
59                     continue;
60                 if (next < 0 || size[next] < size[v])
```

```

61         next = v;
62     }
63     u = next;
64 }
65
66     root[top] = new Tree(0, cnt, sqn);
67 }
68 }
69
70 void prepare()
71 {
72     doBfs(0);
73     doSplit();
74 }

```

1.10 KD Tree

读入 N 个点，输出距离每个点的最近点。

```

1  const int MAX_N = 100000 + 10;
2  const int MAX_NODE = 200000 + 10;
3  const LL INF = 2000000000000000000LL;
4
5  int N;
6
7  struct Point
8  {
9      int x, y, id;
10 };
11
12 LL dis(const Point &a, const Point &b)
13 {
14     return 1LL * (a.x - b.x) * (a.x - b.x) + 1LL * (a.y - b.y) * (a.y - b.y);
15 }
16
17 struct Node
18 {
19     Point p;
20     int maxX, minX, maxY, minY;
21     int l, r, d;
22     Node *ch[2];
23 };
24
25 LL ret;
26 LL ans[MAX_N];
27 Node *root;
28 Point p[MAX_N], queryPoint;
29 Node *totNode, nodePool[MAX_NODE];
30
31 int cmpx(const Point &a, const Point &b)
32 {
33     return a.x < b.x;
34 }
35 int cmpy(const Point &a, const Point &b)
36 {
37     return a.y < b.y;
38 }
39
40 Node* newNode(int l, int r, Point p, int deep)
41 {
42     Node *t = totNode++;
43     t->l = l; t->r = r;
44     t->p = p; t->d = deep;

```

```

45     t->maxX = t->minX = p.x;
46     t->maxY = t->minY = p.y;
47     return t;
48 }
49
50 void updateInfo(Node *t, Node *p)
51 {
52     t->maxX = max(t->maxX, p->maxX);
53     t->maxY = max(t->maxY, p->maxY);
54     t->minX = min(t->minX, p->minX);
55     t->minY = min(t->minY, p->minY);
56 }
57
58 Node* build(int l, int r, int deep)
59 {
60     if (l == r) return NULL;
61     if (deep & 1) sort(p + l, p + r, cmpx);
62     else sort(p + l, p + r, cmpy);
63     int mid = (l + r) >> 1;
64     Node *t = newNode(l, r, p[mid], deep & 1);
65     if (l + 1 == r) return t;
66     t->ch[0] = build(l, mid, deep + 1);
67     t->ch[1] = build(mid + 1, r, deep + 1);
68     if (t->ch[0]) updateInfo(t, t->ch[0]);
69     if (t->ch[1]) updateInfo(t, t->ch[1]);
70     return t;
71 }
72
73 void updateAns(Point p)
74 {
75     ret = min(ret, dis(p, queryPoint));
76 }
77
78 LL calc(Node *t, LL d)
79 {
80     LL tmp;
81     if (d) {
82         if (queryPoint.x >= t->minX && queryPoint.x <= t->maxX) tmp = 0;
83         else tmp = min(abs(queryPoint.x - t->maxX), abs(queryPoint.x - t->minX));
84     } else {
85         if (queryPoint.y >= t->minY && queryPoint.y <= t->maxY) tmp = 0;
86         else tmp = min(abs(queryPoint.y - t->maxY), abs(queryPoint.y - t->minY));
87     }
88     return tmp * tmp;
89 }
90
91 void query(Node *t)
92 {
93     if (t == NULL) return;
94     if (t->p.id != queryPoint.id) updateAns(t->p);
95     if (t->l + 1 == t->r) return;
96     LL dl = t->ch[0] ? calc(t->ch[0], t->d) : INF;
97     LL dr = t->ch[1] ? calc(t->ch[1], t->d) : INF;
98     if (dl < dr) {
99         query(t->ch[0]);
100         if (ret > dr) query(t->ch[1]);
101     } else {
102         query(t->ch[1]);
103         if (ret > dl) query(t->ch[0]);
104     }
105 }
106
107 void solve()

```

```

108 {
109     scanf("%d", &N);
110     for(int i = 0; i < N; ++ i) {
111         scanf("%d%d", &p[i].x, &p[i].y);
112         p[i].id = i;
113     }
114     totNode = nodePool;
115     root = build(0, N, 1);
116
117     for(int i = 0; i < N; ++ i) {
118         queryPoint = p[i];
119         ret = INF;
120         query(root);
121         ans[p[i].id] = ret;
122     }
123     for(int i = 0; i < N; ++ i)
124         printf("%I64d\n", ans[i]);
125 }
126
127 int main()
128 {
129     int T; scanf("%d", &T);
130     for( ; T --; )
131         solve();
132     return 0;
133 }

```

1.11 Manacher

len[i] means the max palindrome length centered i/2

eg: cs = "abbacabbaddabbaae"

len = 1 0 1 4 1 0 1 0 9 0 1 0 1 4 1 0 1 0 1 4 1 0 1 2 1 0 1 0

```

1 void palindrome(char cs[], int len[], int n) {
2     for (int i = 0; i < n * 2; ++i) {
3         len[i] = 0;
4     }
5     for (int i = 0, j = 0, k; i < n * 2; i += k, j = max(j - k, 0)) {
6         while (i - j >= 0 && i + j + 1 < n * 2 && cs[(i - j) / 2] == cs[(i + j + 1) /
7             2])
8             j++;
9         len[i] = j;
10        for (k = 1; i - k >= 0 && j - k >= 0 && len[i - k] != j - k; k++) {
11            len[i + k] = min(len[i - k], j - k);
12        }
13    }
14 }

```

1.12 Z Algorithm

传入字符串 s 和长度 N, next[i]=LCP(s, s[i..N-1])

```

1 void z(char *s, int *next, int N)
2 {
3     int j = 0, k = 1;
4     while (j + 1 < N && s[j] == s[j + 1]) ++ j;
5     next[0] = N - 1; next[1] = j;
6     for(int i = 2; i < N; ++ i) {
7         int far = k + next[k] - 1, L = next[i - k];
8         if (L < far - i + 1) next[i] = L;
9         else {
10            j = max(0, far - i + 1);

```

```

11         while (i + j < N && s[j] == s[i + j]) ++ j;
12         next[i] = j; k = i;
13     }
14 }
15 }

```

1.13 Aho-Corasick Automaton

包含建 trie 和构造自动机的过程

```

1
2 struct acNode
3 {
4     int id;
5     acNode *ch[26], *fail;
6 } *totNode, *root, nodePool[MAX_V];
7
8 acNode* newNode()
9 {
10     acNode *now = totNode++;
11     now->id = 0; now->fail = 0;
12     memset(now->ch, 0, sizeof now->ch);
13     return now;
14 }
15
16 void acInsert(char *c, int id)
17 {
18     acNode *cur = root;
19     while (*c) {
20         int p = *c - 'A'; //change the index
21         if (! cur->ch[p]) cur->ch[p] = newNode();
22         cur = cur->ch[p];
23         ++ c;
24     }
25     cur->id = id;
26 }
27
28 void getFail()
29 {
30     acNode *cur;
31     queue<acNode*> Q;
32     for(int i = 0; i < 26; ++ i)
33         if (root->ch[i]) {
34             root->ch[i]->fail = root;
35             Q.push(root->ch[i]);
36         } else root->ch[i] = root;
37     while (! Q.empty()) {
38         cur = Q.front(); Q.pop();
39         for(int i = 0; i < 26; ++ i)
40             if (cur->ch[i]) {
41                 cur->ch[i]->fail = cur->fail->ch[i];
42                 Q.push(cur->ch[i]);
43             } else cur->ch[i] = cur->fail->ch[i];
44     }
45 }

```

1.14 Suffix Array

对于串 a 求 SA , 长度为 N , M 为元素值范围 , $height[i] = LCP(suf[rank[i]], suf[rank[i]-1])$

```

1 const int MAX_N = 1000000 + 10;
2

```

```

3  int rank[MAX_N], height[MAX_N];
4
5  int cmp(int *x, int a, int b, int d)
6  {
7      return x[a] == x[b] && x[a + d] == x[b + d];
8  }
9
10 void doubling(int *a, int N, int M)
11 {
12     static int sRank[MAX_N], tmpA[MAX_N], tmpB[MAX_N];
13     int *x = tmpA, *y = tmpB;
14     for(int i = 0; i < M; ++ i) sRank[i] = 0;
15     for(int i = 0; i < N; ++ i) ++ sRank[x[i] = a[i]];
16     for(int i = 1; i < M; ++ i) sRank[i] += sRank[i - 1];
17     for(int i = N - 1; i >= 0; -- i) sa[-- sRank[x[i]]] = i;
18
19     for(int d = 1, p = 0; p < N; M = p, d <= 1) {
20         p = 0; for(int i = N - d; i < N; ++ i) y[p++] = i;
21         for(int i = 0; i < N; ++ i)
22             if (sa[i] >= d) y[p++] = sa[i] - d;
23         for(int i = 0; i < M; ++ i) sRank[i] = 0;
24         for(int i = 0; i < N; ++ i) ++ sRank[x[i]];
25         for(int i = 1; i < M; ++ i) sRank[i] += sRank[i - 1];
26         for(int i = N - 1; i >= 0; -- i) sa[-- sRank[x[y[i]]]] = y[i];
27         swap(x, y); x[sa[0]] = 0; p = 1;
28         for(int i = 1; i < N; ++ i)
29             x[sa[i]] = cmp(y, sa[i], sa[i - 1], d) ? p - 1 : p++;
30     }
31 }
32
33 void calcHeight()
34 {
35     for(int i = 0; i < N; ++ i) rank[sa[i]] = i;
36     int cur = 0;
37     for(int i = 0; i < N; ++ i)
38         if (rank[i]) {
39             if (cur) cur--;
40             for( ; a[i + cur] == a[sa[rank[i] - 1] + cur]; ++ cur);
41             height[rank[i]] = cur;
42         }
43 }

```

1.15 Suffix Automaton

.....保重

```

1  struct State
2  {
3      int val;
4      State *suf, *go[26];
5  } *root, *last;
6
7  State statePool[MAX_N], *curState;
8
9  void extend(int w)
10 {
11     State *p = last, *np = curState++;
12     np->val = p->val + 1;
13     for( ; p && ! p->go[w]; p = p->suf)
14         p->go[w] = np;
15     if (! p)
16         np->suf = root;
17     else {

```

```

18     State *q = p->go[w];
19     if (q->val == p->val + 1)
20         np->suf = q;
21     else {
22         State *nq = curState++;
23         nq->val = p->val + 1;
24         memcpy(nq->go, q->go, sizeof q->go);
25         nq->suf = q->suf;
26         q->suf = np->suf = nq;
27         for( ; p && p->go[w] == q; p = p->suf)
28             p->go[w] = nq;
29     }
30 }
31 last = np;
32 }

```

1.16 *Dancing Links

Comming Soon

Chapter 2

Graph Theory and Network Algorithms

2.1 Dijkstra

求 s 到其他点的最短路

```
1 int used[MAX_N], dis[MAX_N];
2 void dijkstra(int s) {
3     fill(dis, dis + N, INF); dis[s] = 0;
4     priority_queue<pair<int, int>> que;
5     que.push(make_pair(-dis[s], s));
6     while (!que.empty()) {
7         int u = que.top().second; que.pop();
8         if (used[u]) continue;
9         used[u] = true;
10        foreach(e, E[u])
11            if (dis[u] + e->w < dis[e->t]) {
12                dis[e->t] = dis[u] + e->w;
13                que.push(make_pair(-dis[e->t], e->t));
14            }
15    }
16 }
```

2.2 *Minimum Directed Spanning Tree

Comming Soon

2.3 KuhnMunkres

求完备匹配的最大权匹配，建好的完全图用 $w[i][j]$ 存储，点数为 N

```
1 const int MAX_N = 200 + 10;
2 const int INF = 1000000000;
3
4 int N, flag;
5 int w[MAX_N][MAX_N];
6 int fx[MAX_N], fy[MAX_N], lx[MAX_N], ly[MAX_N], slk[MAX_N], mat[MAX_N];
7
8 int DFS(int x) {
9     fx[x] = flag;
10    for(int i = 1; i <= N; ++i)
11        if (lx[x] + ly[i] != w[x][i])
12            slk[i] = min(slk[i], lx[x] + ly[i] - w[x][i]);
13    else if (fy[i] != flag) {
14        fy[i] = flag;
15        if (mat[i] < 0 || DFS(mat[i])) {
16            mat[i] = x;
17            return true;
18        }
19    }
```

```

19     }
20     return false;
21 }
22
23 int KM() {
24     for(int i = 1; i <= N; ++ i) {
25         mat[i] = -1;
26         fx[i] = 0; fy[i] = 0;
27         ly[i] = 0; lx[i] = -INF;
28         for(int j = 1; j <= N; ++ j)
29             lx[i] = max(lx[i], w[i][j]);
30     }
31     for(int now = 1; now <= N; ++ now) {
32         ++ flag; for(int i = 1; i <= N; ++ i) slk[i] = INF;
33         while (! DFS(now)) {
34             int p(INF); for(int i = 1; i <= N; ++ i)
35                 if (fy[i] != flag) p = min(p, slk[i]);
36             for(int i = 1; i <= N; ++ i) {
37                 if (fx[i] == flag) lx[i] -= p;
38                 if (fy[i] == flag) ly[i] += p;
39                 slk[i] = INF;
40             }
41             ++ flag;
42         }
43     }
44     long long ans = 0;
45     for(int i = 1; i <= N; ++ i) ans += lx[i], ans += ly[i];
46     return ans;
47 }

```

2.4 Maximum Flow

iSAP 算法求 S 到 T 的最大流，点数为 cntN，边表存储在 *E[] 中

```

1 struct Edge
2 {
3     int t, c;
4     Edge *n, *r;
5 } *E[MAX_V], edges[MAX_M], *totEdge;
6
7 Edge* makeEdge(int s, int t, int c)
8 {
9     Edge *e = totEdge++;
10    e->t = t; e->c = c; e->n = E[s];
11    return E[s] = e;
12 }
13
14 void addEdge(int s, int t, int c)
15 {
16     Edge *p = makeEdge(s, t, c), *q = makeEdge(t, s, 0);
17     p->r = q; q->r = p;
18 }
19
20 int maxflow()
21 {
22     static int cnt [MAX_V];
23     static int h [MAX_V];
24     static int que [MAX_V];
25     static int aug [MAX_V];
26     static Edge *cur [MAX_V];
27     static Edge *prev [MAX_V];
28     fill(h, h + cntN, cntN);
29     memset(cnt, 0, sizeof cnt);

```

```

30     int qt = 0, qh = 0; h[T] = 0;
31     for(que[qt++] = T; qh < qt; ) {
32         int u = que[qh++];
33         ++ cnt[h[u]];
34         for(Edge *e = E[u]; e; e = e->n)
35             if (e->r->c && h[e->t] == cntN) {
36                 h[e->t] = h[u] + 1;
37                 que[qt++] = e->t;
38             }
39     }
40     memcpy(cur, E, sizeof E);
41     aug[S] = INF; Edge *e;
42     int flow = 0, u = S;
43     while (h[S] < cntN) {
44         for(e = cur[u]; e; e = e->n)
45             if (e->c && h[e->t] + 1 == h[u])
46                 break;
47         if (e) {
48             int v = e->t;
49             cur[u] = prev[v] = e;
50             aug[v] = min(aug[u], e->c);
51             if ((u = v) == T) {
52                 int by = aug[T];
53                 while (u != S) {
54                     Edge *p = prev[u];
55                     p->c -= by;
56                     p->r->c += by;
57                     u = p->r->t;
58                 }
59                 flow += by;
60             }
61         } else {
62             if (!-- cnt[h[u]]) return flow;
63             h[u] = cntN;
64             for(e = E[u]; e; e = e->n)
65                 if (e->c && h[u] > h[e->t] + 1)
66                     h[u] = h[e->t] + 1, cur[u] = e;
67             ++ cnt[h[u]];
68             if (u != S) u = prev[u]->r->t;
69         }
70     }
71     return flow;
72 }

```

2.5 Minimum Cost Maximum Flow

注意图的初始化，费用和流的类型依题目而定

```

1  int flow, cost;
2
3  struct Edge
4  {
5      int t, c, w;
6      Edge *n, *r;
7  } *totEdge, edges[MAX_M], *E[MAX_V];
8
9  Edge* makeEdge(int s, int t, int c, int w)
10 {
11     Edge *e = totEdge++;
12     e->t = t; e->c = c; e->w = w; e->n = E[s];
13     return E[s] = e;
14 }
15

```

```

16 void addEdge(int s, int t, int c, int w)
17 {
18     Edge *st = makeEdge(s, t, c, w), *ts = makeEdge(t, s, 0, -w);
19     st->r = ts; ts->r = st;
20 }
21
22 int SPFA()
23 {
24     static int que[MAX_V];
25     static int aug[MAX_V];
26     static int in[MAX_V];
27     static int dist[MAX_V];
28     static Edge *prev[MAX_V];
29     int qh = 0, qt = 0;
30
31     int u, v;
32     fill(dist, dist + cntN, INF); dist[S] = 0;
33     fill(in, in + cntN, 0); in[S] = true;
34     que[qt++] = S; aug[S] = INF;
35     for( ; qh != qt; ) {
36         u = que[qh]; qh = (qh + 1) % MAX_N;
37         for(Edge *e = E[u]; e; e = e->n) {
38             if (! e->c) continue;
39             v = e->t;
40             if (dist[v] > dist[u] + e->w) {
41                 dist[v] = dist[u] + e->w;
42                 aug[v] = min(aug[u], e->c);
43                 prev[v] = e;
44                 if (! in[v]) {
45                     in[v] = true;
46                     if (qh != qt && dist[v] <= dist[que[qh]]) {
47                         qh = (qh - 1 + MAX_N) % MAX_N;
48                         que[qh] = v;
49                     } else {
50                         que[qt] = v;
51                         qt = (qt + 1) % MAX_N;
52                     }
53                 }
54             }
55         }
56         in[u] = false;
57     }
58
59     if (dist[T] == INF) return false;
60     cost += dist[T] * aug[T];
61     flow += aug[T];
62     for(u = T; u != S; ) {
63         prev[u]->c -= aug[T];
64         prev[u]->r->c += aug[T];
65         u = prev[u]->r->t;
66     }
67     return true;
68 }
69
70 int minCostFlow()
71 {
72     flow = cost = 0;
73     while(SPFA());
74     return cost;
75 }

```

2.6 Strongly Connected Component

N 个点的图求 SCC, totID 为时间标记, top 为栈顶, totCol 为强联通分量个数, 注意初始化

```
1  int totID, totCol;
2  int col[MAX_N], low[MAX_N], dfn[MAX_N];
3  int top, stack[MAX_N], instack[MAX_N];
4
5  int tarjan(int u)
6  {
7      low[u] = dfn[u] = ++ totID;
8      instack[u] = true; stack[++ top] = u;
9
10     int v;
11     foreach(it, adj[u]) {
12         v = it->first;
13         if (dfn[v] == -1)
14             low[u] = min(low[u], tarjan(v));
15         else if (instack[v])
16             low[u] = min(low[u], dfn[v]);
17     }
18
19     if (low[u] == dfn[u]) {
20         do {
21             v = stack[top--];
22             instack[v] = false;
23             col[v] = totCol;
24         } while(v != u);
25         ++ totCol;
26     }
27     return low[u];
28 }
29
30 void solve()
31 {
32     totID = totCol = top = 0;
33     fill(dfn, dfn + N, 0);
34     for(int i = 0; i < N; ++ i)
35         if (! dfn[i])
36             tarjan(i);
37 }
```

2.7 *2-SAT

Comming Soon

Chapter 3

Number Theory

3.1 Chinese Remainder Theorem

包括扩展欧几里得，求逆元，和保证除数互质条件下的 CRT

```
1 LL x, y;
2 void exGcd(LL a, LL b)
3 {
4     if (b == 0) {
5         x = 1;
6         y = 0;
7         return;
8     }
9     exGcd(b, a % b);
10    LL k = y;
11    y = x - a / b * y;
12    x = k;
13 }
14
15 LL inversion(LL a, LL b)
16 {
17     exGcd(a, b);
18     return (x % b + b) % b;
19 }
20
21 LL CRT(vector<LL> m, vector<LL> a)
22 {
23     int N = m.size();
24     LL M = 1, ret = 0;
25     for(int i = 0; i < N; ++i)
26         M *= m[i];
27
28     for(int i = 0; i < N; ++i) {
29         ret = (ret + (M / m[i]) * a[i] % M * inversion(M / m[i], m[i])) % M;
30     }
31     return ret;
32 }
```

3.2 Pollard's Rho and Miller-Rabbin

大数分解和素性判断

```
1 typedef long long LL;
2
3 LL modMul(LL a, LL b, LL P)
4 {
5     LL ret = 0;
6     for( ; a; a >>= 1) {
7         if (a & 1) {
```

```

8         ret += b;
9         if (ret >= P) ret -= P;
10    }
11    b <<= 1;
12    if (b >= P) b -= P;
13 }
14 return ret;
15 }
16
17 LL modPow(LL a, LL u, LL P)
18 {
19     LL ret = 1;
20     for( ; u >>= 1, a = modMul(a, a, P))
21         if (u & 1) ret = modMul(ret, a, P);
22     return ret;
23 }
24
25 int millerRabin(LL N)
26 {
27     if (N == 2) return true;
28     LL t = 0, u = N - 1, x, y, a;
29     for( ; ! (u & 1); ++ t, u >>= 1) ;
30     for(int k = 0; k < 10; ++ k) {
31         a = rand() % (N - 2) + 2;
32         x = modPow(a, u, N);
33         for(int i = 0; i < t; ++ i, x = y) {
34             y = modMul(x, x, N);
35             if (y == 1 && x > 1 && x < N - 1) return false;
36         }
37         if (x != 1) return false;
38     }
39     return true;
40 }
41
42 LL gcd(LL a, LL b)
43 {
44     return ! b ? a : gcd(b, a % b);
45 }
46
47 LL pollardRho(LL N)
48 {
49     LL i = 1, x = rand() % N;
50     LL y = x, k = 2, d = 1;
51     do {
52         d = gcd(x - y + N, N);
53         if (d != 1 && d != N) return d;
54         if (++ i == k) y = x, k <<= 1;
55         x = (modMul(x, x, N) - 1 + N) % N;
56     } while (y != x);
57     return N;
58 }
59
60 void getFactor(LL N)
61 {
62     if (N < 2) return;
63     if (millerRabin(N)) {
64         //do some operations
65         return;
66     }
67     LL x = pollardRho(N);
68     getFactor(x);
69     getFactor(N / x);
70 }

```

3.3 *Baby-step Giant-step

Comming soon

Chapter 4

Algebraic Algorithms

4.1 *Linear Equations in Z_m

Comming Soon

4.2 *Linear Equations in R

Comming Soon

4.3 *Fast Fourier Transform

Comming Soon

Chapter 5

Computational Geometry

5.1 Basic Operations

平面几何基本操作，之后的几个都需要先敲它

```
1 #include <cstdio>
2 #include <cstring>
3 #include <algorithm>
4 #include <iostream>
5 #include <climits>
6 #include <numeric>
7 #define foreach(e,x) for(__typeof(x.begin()) e=x.begin();e!=x.end();++e)
8 #define REP(i,n) for(int i=0;i<n;++i)
9 using namespace std;
10
11 const double EPS = 1e-8;
12 inline int sign(double a) {
13     return a < -EPS ? -1 : a > EPS;
14 }
15
16 struct Point {
17     double x, y;
18     Point() {}
19 }
20 Point(double _x, double _y) :
21     x(_x), y(_y) {}
22
23 Point operator+(const Point&p) const {
24     return Point(x + p.x, y + p.y);
25 }
26 Point operator-(const Point&p) const {
27     return Point(x - p.x, y - p.y);
28 }
29 Point operator*(double d) const {
30     return Point(x * d, y * d);
31 }
32 Point operator/(double d) const {
33     return Point(x / d, y / d);
34 }
35 bool operator<(const Point&p) const {
36     int c = sign(x - p.x);
37     if (c)
38         return c == -1;
39     return sign(y - p.y) == -1;
40 }
41 double dot(const Point&p) const {
42     return x * p.x + y * p.y;
43 }
44 double det(const Point&p) const {
45     return x * p.y - y * p.x;
```

```

46     }
47     double alpha() const {
48         return atan2(y, x);
49     }
50     double distTo(const Point&p) const {
51         double dx = x - p.x, dy = y - p.y;
52         return hypot(dx, dy);
53     }
54     double alphaTo(const Point&p) const {
55         double dx = x - p.x, dy = y - p.y;
56         return atan2(dy, dx);
57     }
58     void read() {
59         scanf("%lf%lf", &x, &y);
60     }
61     double abs() {
62         return hypot(x, y);
63     }
64     double abs2() {
65         return x * x + y * y;
66     }
67     void write() {
68         cout << "(" << x << ", " << y << ")" << endl;
69     }
70 };
71
72 #define cross(p1,p2,p3) ((p2.x-p1.x)*(p3.y-p1.y)-(p3.x-p1.x)*(p2.y-p1.y))
73
74 #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
75
76 Point isSS(Point p1, Point p2, Point q1, Point q2) {
77     double a1 = cross(q1,q2,p1), a2 = -cross(q1,q2,p2);
78     return (p1 * a2 + p2 * a1) / (a1 + a2);
79 }
80
81 double calcArea(const vector<Point>&ps) {
82     int n = ps.size();
83     double ret = 0;
84     for (int i = 0; i < n; ++i) {
85         ret += ps[i].det(ps[(i + 1) % n]);
86     }
87     return ret / 2; //maybe need abs(ret)
88 }

```

5.2 Convex Hull

Montone Chain

```

1  vector<Point> convexHull(vector<Point> ps) {
2      int n = ps.size();
3      if (n <= 1)
4          return ps;
5      sort(ps.begin(), ps.end());
6      vector<Point> qs;
7      for (int i = 0; i < n; qs.push_back(ps[i++])) {
8          while (qs.size() > 1 && crossOp(qs[qs.size()-2], qs.back(), ps[i]) <= 0)
9              qs.pop_back();
10     }
11     for (int i = n - 2, t = qs.size(); i >= 0; qs.push_back(ps[i--])) {
12         while (qs.size() > t && crossOp(qs[qs.size()-2], qs.back(), ps[i]) <= 0)
13             qs.pop_back();
14     }
15     qs.pop_back();

```

```

16     return qs;
17 }

```

5.3 Convex Diameter

```

1  double convexDiameter(const vector<Point>&ps) {
2      int n = ps.size();
3      int is = 0, js = 0;
4      for (int i = 1; i < n; ++i) {
5          if (ps[i].x > ps[is].x)
6              is = i;
7          if (ps[i].x < ps[js].x)
8              js = i;
9      }
10     double maxd = ps[is].distTo(ps[js]);
11     int i = is, j = js;
12     do {
13         if ((ps[(i + 1) % n] - ps[i]).det(ps[(j + 1) % n] - ps[j]) >= 0)
14             (++j) %= n;
15         else
16             (++i) %= n;
17         maxd = max(maxd, ps[i].distTo(ps[j]));
18     } while (i != is || j != js);
19     return maxd;
20 }

```

5.4 Convex Cut

```

1  vector<Point> convexCut(const vector<Point>&ps, Point q1, Point q2) {
2      vector<Point> qs;
3      int n = ps.size();
4      for (int i = 0; i < n; ++i) {
5          Point p1 = ps[i], p2 = ps[(i + 1) % n];
6          int d1 = crossOp(q1, q2, p1), d2 = crossOp(q1, q2, p2);
7          if (d1 >= 0)
8              qs.push_back(p1);
9          if (d1 * d2 < 0)
10             qs.push_back(isSS(p1, p2, q1, q2));
11     }
12     return qs;
13 }

```

Chapter 6

Classic Problems

6.1 Nim Game

对于 N 堆石子，每人轮流取

- 1) 每堆石子个数为 1，根据奇偶性直接判胜负
- 2) 有一堆个数大于 1，先手必胜（直接根据奇偶性调整留一个还是取光）
- 3) Nim 游戏：每堆个数任意，xor 和为 0 则为必败态，否则必胜，证明略
- 4) Moore's Nim K 游戏：从最少 1 堆最多 K 堆中取任意数量的石子，结论：把所有堆的石子个数按二进制表示，如果任意一位一的个数总和都为 $K+1$ 的倍数则为必败态，否则为必胜态，显然 Nim 是 $K=1$ 的特殊情况。
- 5) anti-nim 游戏：取到最后一个石子为败，结论：必胜态当且仅当：1) 所有堆石子数都为 1 且游戏的 SG 值为 0, 2) 存在某堆石子数大于 1 且游戏的 sg 值不为 0