

Variable-Shape Linear Algebra: Mathematical Foundations and High-Performance Implementation

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Abstract

Variable-Shape Linear Algebra (VSLA) treats *dimension* as intrinsic data rather than a rigid constraint. This paper makes four concrete contributions: (1) formalization of VSLA through equivalence classes of finite-dimensional vectors modulo trailing-zero padding; (2) construction of two semiring instantiations—convolution and Kronecker products—with complete algebraic characterization; (3) asymptotic complexity analysis showing FFT-accelerated convolution achieves $\mathcal{O}(mnd_{\max} \log d_{\max})$ for matrix-vector operations compared to $\mathcal{O}(mnd_{\max}^2)$ for naive approaches; (4) an open-source C99 library with Python bindings (implementation in progress). Unlike existing ragged tensor frameworks (TensorFlow Ragged, PyTorch NestedTensors), VSLA provides mathematically rigorous semiring structures with provable algebraic identities, enabling principled dimension-aware computation for adaptive AI architectures, multi-resolution signal processing, and scientific computing applications.

Keywords: Variable-shape tensors, semiring algebra, automatic differentiation, high-performance computing, adaptive neural networks

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1 Context and Motivation

1.1 The Dimension Problem

Traditional linear algebra fixes dimensions m, n *a priori*. Contemporary challenges—adaptive neural networks, multi-resolution signal analysis, dynamic meshes—demand structures whose shapes evolve in real time.

Running Example: Consider training a convolutional neural network where filter widths adapt dynamically based on input complexity. A standard 3×3 convolution kernel $K_1 = [1, -1, 2]$ might expand to $K_2 = [1, -1, 2, 0, 1]$ for high-resolution features. Traditional frameworks require manual padding: $K'_1 = [1, -1, 2, 0, 0]$ before operations, losing semantic information and incurring unnecessary computation on artificial zeros.

Existing approaches fall short:

- **TensorFlow Ragged Tensors:** Handle variable-length sequences but lack rigorous algebraic structure and semiring properties.
- **PyTorch NestedTensors:** Provide dynamic shapes but without mathematical guarantees or efficient sparse representations.
- **Manual zero-padding:** Obscures mathematical structure, wastes computation, and lacks provable algebraic identities.

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1.2 The VSLA Solution

VSLA incorporates the shape directly into every algebraic object through mathematically rigorous equivalence classes. Operations such as addition or convolution implicitly coerce operands to a common dimension while preserving sparsity and algebraic properties. In our example, $K_1 \oplus K_2 = [2, -2, 4, 0, 1]$ automatically, with provable semiring laws and efficient sparse computation.

1.3 Roadmap

This paper proceeds as follows: §2 establishes mathematical preliminaries; §3–§?? develop two semiring models with complete proofs; §??–§7 bridge theory to implementation; §8–§9 provide empirical validation and context. Appendix contains detailed proofs and API specifications.

2 Mathematical Preliminaries

Key Definitions

Dimension-aware vector: An equivalence class $[(d, v)]$ where $d \in \mathbb{N}$ is the logical dimension and $v \in \mathbb{R}^d$ is the data vector.

Zero-padding equivalence: $(d_1, v) \sim (d_2, w)$ iff their extensions to $\max(d_1, d_2)$ dimensions are equal.

Shape-semiring: A semiring S with degree function $\deg : S \rightarrow \mathbb{N}$ satisfying $\deg(x+y) \leq \max(\deg x, \deg y)$ and $\deg(xy) = \deg x + \deg y$.

Variable-shape operation: An operation that automatically promotes operands to compatible shapes before computation.

Table 1: Notation Table

Symbol	Meaning
D	Set of dimension-aware vectors
$[(d, v)]$	Equivalence class of vector $v \in \mathbb{R}^d$
$\deg x$	Logical dimension/degree of element x
$\iota_{m \rightarrow n}$	Zero-padding map from \mathbb{R}^m to \mathbb{R}^n
\oplus, \otimes_c	Addition and convolution in Model A
\oplus, \otimes_K	Addition and Kronecker product in Model B
d_{\max}	Maximum degree in a matrix or operation
$\mathcal{O}(\cdot)$	Asymptotic complexity bound

3 Mathematical Foundations

3.1 The Dimension-Aware Space

Definition 3.1 (Dimension-Aware Vectors). *Define the graded set*

$$D_e := \bigsqcup_{d \geq 0} \{d\} \times \mathbb{R}^d,$$

where $\mathbb{R}^0 := \{[]\}$ denotes the empty vector.

Definition 3.2 (Zero-Padding Equivalence). *For $m \leq n$ let $\iota_{m \rightarrow n} : \mathbb{R}^m \rightarrow \mathbb{R}^n$ append $n - m$ trailing zeros. Put*

$$(d_1, v) \sim (d_2, w) \iff \iota_{d_1 \rightarrow n}(v) = \iota_{d_2 \rightarrow n}(w), \quad n := \max(d_1, d_2).$$

Proposition 3.3. *The relation \sim is an equivalence relation, yielding the set $D := D_e / \sim$ of dimension-aware vectors.*

Proof. Reflexivity and symmetry are immediate from Definition 3.2. For transitivity pad to $n := \max(d_1, d_2, d_3)$. \square

3.2 Additive Structure

Theorem 3.4. $(D, +, 0)$ is a commutative monoid where

$$[(d_1, v)] + [(d_2, w)] := [(n, \iota_{d_1 \rightarrow n}(v) + \iota_{d_2 \rightarrow n}(w))], \quad n := \max(d_1, d_2), \quad 0 := [(0, [])].$$

Proof. Well-definedness follows from Proposition 3.3. Associativity and commutativity inherit from \mathbb{R}^n . \square

4 Model A: The Convolution Semiring

4.1 Convolution Product

Definition 4.1. For $v \in \mathbb{R}^{d_1}$ and $w \in \mathbb{R}^{d_2}$ define the discrete convolution

$$(v * w)_k := \sum_{i+j=k+1} v_i w_j, \quad k = 0, \dots, d_1 + d_2 - 2.$$

Put

$$[(d_1, v)] \otimes_c [(d_2, w)] := \begin{cases} 0, & d_1 d_2 = 0, \\ [(d_1 + d_2 - 1, v * w)], & \text{otherwise.} \end{cases}$$

Theorem 4.2. $(D, +, \otimes_c, 0, 1)$ is a commutative semiring with $1 := [(1, [1])]$.

Proof. We verify the semiring axioms:

Associativity of \otimes_c : For $a, b, c \in D$ with representatives $[(d_1, u)], [(d_2, v)], [(d_3, w)]$, we need $(a \otimes_c b) \otimes_c c = a \otimes_c (b \otimes_c c)$. By definition, $(u * v) * w$ and $u * (v * w)$ both equal

$$\sum_{i+j+k=n+2} u_i v_j w_k$$

when expanding the convolution index arithmetic. Thus both products have degree $d_1 + d_2 + d_3 - 2$ and identical coefficients.

Commutativity of \otimes_c : The convolution $(u * v)_k = \sum_{i+j=k+1} u_i v_j = \sum_{j+i=k+1} v_j u_i = (v * u)_k$ by symmetry of the index condition.

Distributivity: For $a, b, c \in D$, we have $a \otimes_c (b + c) = a \otimes_c b + a \otimes_c c$ since convolution distributes over pointwise addition: $u * (v + w) = u * v + u * w$ coefficientwise.

Identity elements: The zero element $0 = [(0, [])]$ satisfies $0 \otimes_c a = 0$ by the first case in the definition. The one element $1 = [(1, [1])]$ satisfies $(1 * v)_k = v_k$ for all k , making it the multiplicative identity. \square

Theorem 4.3 (Polynomial Isomorphism). *The map $\Phi([(d, v)]) := \sum_{i=0}^{d-1} v_{i+1} x^i$ is a semiring isomorphism $D \cong \mathbb{R}[x]$.*

Theorem 4.4 (Completion). *Equip D with the norm $\|[(d, v)]\|_1 := \sum_{i=1}^d |v_i|$. The Cauchy completion of D is isomorphic to the power-series ring $\mathbb{R}[[x]]$.*

Proof. The isomorphism Φ from Theorem 4.3 extends to the completion. Every Cauchy sequence (f_n) in D with respect to $\|\cdot\|_1$ corresponds to a sequence of polynomials $(\Phi(f_n))$ that converges coefficientwise. Since $\|f_n - f_m\|_1 = \sum_i |\text{coeff}_i(\Phi(f_n)) - \text{coeff}_i(\Phi(f_m))|$, Cauchy sequences in D map to coefficient-wise Cauchy sequences in $\mathbb{R}[x]$. The completion consists of formal power series $\sum_{i=0}^{\infty} a_i x^i$ where the sequence of partial sums is Cauchy in the ℓ^1 norm. This is precisely the ring $\mathbb{R}[[x]]$ of convergent power series, and Φ extends to a ring isomorphism between the completions. \square

5 Model B: The Kronecker Semiring

5.1 Kronecker Product

Definition 5.1. For $v \in \mathbb{R}^{d_1}$, $w \in \mathbb{R}^{d_2}$, let

$$v \otimes_K w := (v_1 w_1, \dots, v_1 w_{d_2}, v_2 w_1, \dots, v_{d_1} w_{d_2}).$$

Define

$$[(d_1, v)] \otimes_K [(d_2, w)] := [(d_1 d_2, v \otimes_K w)].$$

Theorem 5.2. $(D, +, \otimes_K, 0, 1)$ is a non-commutative semiring.

Proposition 5.3. $x \otimes_K y = y \otimes_K x$ iff $\deg x = 1$ or $\deg y = 1$ (i.e. one operand is scalar).

Lemma 5.4 (Scalar-Commutation). If $x = \alpha 1$ with $\alpha \in \mathbb{R}$ then $x \otimes_K y = y \otimes_K x$ for all $y \in D$.

Proof. Both products equal αy by definition. \square

6 Variable-Shape Linear Algebra

6.1 Shape-Semirings and Shape-Matrices

Definition 6.1. A shape-semiring is a semiring S equipped with $\deg: S \rightarrow \mathbb{N}$ such that $\deg(x + y) \leq \max\{\deg x, \deg y\}$ and $\deg(xy) = \deg x \deg y$.

The convolution and Kronecker models are shape-semirings.

Theorem 6.2 (Matrix Product). For an $m \times n$ shape-matrix $A = (a_{ij})$ and an $n \times p$ shape-matrix $B = (b_{jk})$ over a shape-semiring,

$$(AB)_{ik} = \sum_{j=1}^n a_{ij} \otimes b_{jk}$$

exists and yields an $m \times p$ shape-matrix.

6.2 Rank, Spectrum and Complexity

Theorem 6.3 (Complexity). Let $d_{\max} = \max_{i,j} \deg a_{ij}$. Then

- Model A: matrix-vector multiply costs $\mathcal{O}(mn d_{\max} \log d_{\max})$ via FFT.
- Model B: the same task costs $\mathcal{O}(mn d_{\max}^2)$.

7 Implementation Design

7.1 API Mapping

C Library API Mapping

Tensor Creation: // C API

```
vsla_tensor_t* vsla_new(uint8_t rank, const uint64_t shape[],  
                        vsla_model_t model, vsla_dtype_t dtype);
```

// Python wrapper

```
def new(shape: List[int], model: Model, dtype: DType) -> Tensor
```

Variable-Shape Operations: // C API

```
vsla_error_t vsla_add(vsla_tensor_t* out, const vsla_tensor_t* a,  
                     const vsla_tensor_t* b);
```

// Python wrapper

```
def add(x: Tensor, y: Tensor) -> Tensor # automatic promotion
```

Semiring Products: // Model A (convolution)

```
vsla_error_t vsla_conv(vsla_tensor_t* out, const vsla_tensor_t* a,  
                      const vsla_tensor_t* b);
```

// Model B (Kronecker)

```
vsla_error_t vsla_kron(vsla_tensor_t* out, const vsla_tensor_t* a,  
                      const vsla_tensor_t* b);
```

7.2 Memory Model

Memory Layout and Optimization

Equivalence Class Storage: VSLA tensors store only the minimal representative of each equivalence class. A tensor with logical shape (d_1, d_2, \dots, d_k) containing trailing zeros is stored with reduced dimensions, avoiding explicit zero storage.

Capacity Management: Physical memory allocation uses power-of-2 growth policy:

```
capacity[i] = next_pow2(shape[i]) // for each dimension i  
total_size = product(capacity[i]) * sizeof(element_type)
```

Memory Alignment: All tensor data is 64-byte aligned for optimal SIMD and cache performance:

```
void* data = aligned_alloc(64, total_size);
```

Zero-Padding Avoidance: Operations automatically promote shapes without materializing padding zeros. A 3×5 tensor added to a 7×2 tensor conceptually becomes 7×5 , but only non-zero regions are computed.

7.3 Algorithm Complexity

Complexity Analysis:

- **Model A:** $\mathcal{O}(mnd_1 d_{\max} \log d_{\max})$ via FFT convolution
- **Model B:** $\mathcal{O}(mnd_1 d_{\max}^2)$ for naive Kronecker products
- **Memory:** $\mathcal{O}(mnd_{\max})$ with sparse storage avoiding materialized zeros

Algorithm 1 FFT-Accelerated Convolution (Model A)

Require: Tensors $A \in \mathbb{R}^{m \times d_1}$, $B \in \mathbb{R}^{d_1 \times n}$ with $\deg(A_{ij}), \deg(B_{jk}) \leq d_{\max}$

Ensure: Result tensor $C \in \mathbb{R}^{m \times n}$ with $C_{ik} = \sum_j A_{ij} \otimes_c B_{jk}$

```
1: for  $i = 1$  to  $m$  do
2:   for  $k = 1$  to  $n$  do
3:     conv_sum  $\leftarrow 0$ 
4:     for  $j = 1$  to  $d_1$  do
5:       Pad  $A_{ij}$  and  $B_{jk}$  to length  $2d_{\max}$ 
6:        $\hat{A} \leftarrow \text{FFT}(A_{ij})$ ,  $\hat{B} \leftarrow \text{FFT}(B_{jk})$ 
7:        $\hat{C} \leftarrow \hat{A} \odot \hat{B}$  {pointwise multiply}
8:       conv_sum  $\leftarrow$  conv_sum +  $\text{IFFT}(\hat{C})$ 
9:     end for
10:     $C_{ik} \leftarrow$  conv_sum
11:  end for
12: end for
```

8 Theoretical Analysis

Table 2: Theoretical Performance Analysis: VSLA vs. Existing Approaches

Operation	Zero-Padding	Ragged Tensors	VSLA
Vector Add	$\mathcal{O}(n_{\max})$	$\mathcal{O}(\sum n_i)$	$\mathcal{O}(\sum n_i)$
Matrix-Vector	$\mathcal{O}(mn_{\max})$	$\mathcal{O}(m \sum n_i)$	$\mathcal{O}(m \sum n_i)$
Convolution (FFT)	$\mathcal{O}(n_{\max} \log n_{\max})$	N/A	$\mathcal{O}(n_{\max} \log n_{\max})$
Memory Usage	$\mathcal{O}(mn_{\max})$	$\mathcal{O}(m \sum n_i)$	$\mathcal{O}(m \sum n_i)$

Analysis: VSLA provides theoretical advantages over zero-padding when $\sum n_i \ll mn_{\max}$ (sparse case). Empirical validation pending full implementation.

9 Related Work

Ragged Tensor Frameworks:

- **TensorFlow RaggedTensors** [9]: Handle variable-length sequences but lack algebraic structure and formal semiring properties.
- **PyTorch NestedTensors** [10]: Support dynamic shapes with view-based optimizations but without mathematical foundations.
- **JAX vmap** [11]: Vectorization over varying shapes, but requires manual padding management.

Semiring Algebra in Computing:

- **GraphBLAS** [12]: Semiring-based sparse linear algebra for graphs, but fixed-dimension matrices.
- **Differentiable Programming** [13]: Automatic differentiation over semirings, complementary to VSLA’s shape flexibility.

Novelty: VSLA’s equivalence-class formulation provides rigorous algebraic foundations missing in existing variable-shape frameworks, enabling provable optimization and correctness guarantees.

10 Gradient Support and Integration

Automatic Differentiation: VSLA operations are differentiable with custom vector-Jacobian products (VJPs):

PyTorch Integration Example

```
class VSLAAdd(torch.autograd.Function):
    @staticmethod
    def forward(ctx, x, y):
        ctx.save_for_backward(x, y)
        return vsla_add(x, y) # automatic shape promotion

    @staticmethod
    def backward(ctx, grad_output):
        x, y = ctx.saved_tensors
        # Gradients respect original shapes
        grad_x = grad_output[:x.shape[0], :x.shape[1]]
        grad_y = grad_output[:y.shape[0], :y.shape[1]]
        return grad_x, grad_y

# Usage in neural networks
x = VSLATensor([1, 2, 3]) # shape (3,)
y = VSLATensor([4, 5, 6, 7]) # shape (4,)
z = VSLAAdd.apply(x, y) # shape (4,), z = [5,7,9,7]
loss = z.sum()
loss.backward() # gradients flow correctly
```

JAX Custom Call Integration: Similar integration possible via `jax.custom_call` with XLA primitives for GPU acceleration.

11 Applications

- **Adaptive AI Architectures:** mixture-of-experts with dynamic specialist widths.
- **Multi-Resolution Signal Processing:** wavelets, adaptive filters, compression.
- **Scientific Computing:** adaptive mesh refinement, multigrid, domain decomposition.

12 Future Research Directions

- Categorical formulation of VSLA as a semiring-enriched category.
- Sub-quadratic tensor algorithms and parallel implementations.
- Integration with automatic differentiation and quantum computing.

13 Conclusion

Dimension-aware computation replaces brittle padding with algebraic rigor. VSLA unifies flexible data shapes with efficient algorithms, promising advances across adaptive AI, signal processing and beyond.

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