Derive the following using any of the first eight non-derived rules as well as the rule of \rightarrow I.

1.
$$P \& Q, Q \rightarrow (S \rightarrow T), P \rightarrow (T \rightarrow R) \mid S \rightarrow R$$

2.
$$Q \rightarrow (R \& S) \vdash P \rightarrow (Q \rightarrow R)$$

4.
$$P \lor T$$
, $P \leftrightarrow T \vdash P \& T$

5.
$$P \rightarrow Q$$
, $R \rightarrow S + (P \lor R) \rightarrow (Q \lor S)$

6.
$$(P \lor Q) \rightarrow (S \& M), (Q \lor P) \vdash (M \& S)$$

1. P & Q, Q
$$\rightarrow$$
 (S \rightarrow T), P \rightarrow (T \rightarrow R) \mid S \rightarrow R

2.
$$Q \rightarrow (S \rightarrow T)$$

3.
$$P \rightarrow (T \rightarrow R)$$

6.
$$T \rightarrow R$$
 3,5 $\rightarrow E$

8.
$$S \rightarrow T$$
 2,7 $\rightarrow E$

11 S
$$\rightarrow$$
 R 4-10, \rightarrow I

2.
$$Q \rightarrow (R \& S) \vdash P \rightarrow (Q \rightarrow R)$$

1.
$$Q \rightarrow (R \& S)$$
 A

S.
$$Q \rightarrow R$$
 3-5 \rightarrow

6.
$$Q \rightarrow R$$
 3-5 $\rightarrow I$
7. $P \rightarrow (Q \rightarrow R)$ 2-6, $\rightarrow I$

3. R v R | R

- 1. R v R A
- 2. $R H (Get R \rightarrow R)$
- 3. $R \rightarrow R$ 2-2, $\rightarrow I$
- 4. R 1, 3, 3 vE

4. $P \vee T$, $P \leftrightarrow T \vdash P \& T$

1.
$$P \vee T$$
 A
2. $P \leftrightarrow T$ A
3. P H
4. $P \to T$ 2, $\leftrightarrow E$
5. T 3,4 $\to E$
6. $P \& T$ 3,5 &I
7. $P \to (P \& T)$ 3-6, $\to I$
8. T H
9. $T \to P$ 2, $\leftrightarrow E$
10. P 8,9 $\to E$
11. $P \& T$ 8,10 &I
12. $T \to (P \& T)$ 8-11 $\to I$

13. P & T 1, 7, 12 vE

5.
$$P \rightarrow Q$$
, $R \rightarrow S \vdash (P \lor R) \rightarrow (Q \lor S)$

1.
$$P \rightarrow Q$$
 A
2. $R \rightarrow S$ A
3. $P \lor R$ H
4. $P \lor P$ H
5. $Q \lor S$ 5, $V \lor P \rightarrow (Q \lor S)$ 4-6 $\rightarrow I$
8. $R \lor P \rightarrow (Q \lor S)$ 4-6 $\rightarrow I$
9. $Q \lor S$ 9 $\lor I$
11. $Q \lor S$ 9 $\lor I$
12. $Q \lor S$ 3,7,11 $\lor E$
13. $Q \lor R$ 3-12, $\rightarrow I$

Strategy:

Hypothesize P v R for →I

then derive

 $P \rightarrow (Q \vee S)$ and

 $R \rightarrow (Q \vee S)$

By two more applications of \rightarrow I.

6. $(P \lor Q) \rightarrow (S \& M), (Q \lor P) \vdash (M \& S)$

1. $(P \lor Q) \rightarrow (S \& M)$

l) A

2. (Q v P)

Α

3. Q

Н

4. P v Q

3, vI

5. $Q \rightarrow (P \lor Q)$

3-4, →I

6. P

Н

7. | P v Q

6, vI

8. $P \rightarrow (P \vee Q)$

6-7, →I

9. P v Q

2, 5, 8 vE

10. S & M

1,9 →E

11. S

10, &E

12. M

10, &E

13. M&S

11,12 &I

Strategy:

Derive

(P v Q)

from

(Q v P)

using vE.

7. R v ~R | (R v R) v (~R v R)

2. R

5.
$$R \rightarrow ((R \lor R) \lor (\sim R \lor R))$$
 2-4, $\rightarrow I$

9.
$$\sim R \rightarrow ((R \lor R) \lor (\sim R \lor R))$$
 6-8, $\rightarrow I$

Strategy:

Α

Н

2, vI

3, vI

Η

7, vI

1,5,9 vE

6, vI

Demonstrate the entire disjunction twice, once from R, and

again from ~R for vE.

Philosophy 160 HW 4 Derivations using →I Solutions