

Note: Only the first three problems are solved completely below. Solutions to 7-10 are provided on this [lecture video](#) as well as [these corresponding slides](#). Solution outlines to 11-13 are below.

1. $P \rightarrow Q \vdash \sim Q \rightarrow \sim P$

1.	$P \rightarrow Q$	A
2.	$\sim Q$	H \rightarrow I
3.	P	H \sim I
4.	Q	1,3 \rightarrow E
5.	$Q \& \sim Q$	2,4 $\&$ I
6.	$\sim P$	3-5, \sim I
7.	$\sim Q \rightarrow \sim P$	2-6, \rightarrow I

2. $P \& Q \vdash \sim(P \rightarrow \sim Q)$

1.	$P \& Q$	A
2.	$P \rightarrow \sim Q$	H
3.	P	1, $\&$ E
4.	$\sim Q$	2,3 \rightarrow E
5.	Q	1, $\&$ E
6.	$Q \& \sim Q$	4,5 $\&$ I
7.	$\sim(P \rightarrow \sim Q)$	2-6, \sim I

3. $(P \ \& \ Q) \rightarrow (R \vee S), \sim(R \vee S) \vdash \sim(P \ \& \ Q)$

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|----|---------------------------------------|---------------------|
| 1. | $(P \ \& \ Q) \rightarrow (R \vee S)$ | A |
| 2. | $\sim(R \vee S)$ | A |
| 3. | $P \ \& \ Q$ | H |
| 4. | $R \vee S$ | 1,3 \rightarrow E |
| 5. | $(R \vee S) \ \& \ \sim(R \vee S)$ | 2,4 $\&$ I |
| 6. | $\sim(P \ \& \ Q)$ | 3-5, \sim I |

11. $\sim(P \rightarrow Q) \vdash P \ \& \ \sim Q$

Outline of proof.

Since you are trying to prove a conjunction, consider proving each conjunct independently by \sim I.

For the first leg, hypothesize $\sim P$. From here it's not clear what to do, but you know you need a contradiction, and the only formula you can contradict is the assumption, which means that you want to prove $(P \rightarrow Q)$.

Proceed by \rightarrow I, hypothesizing P . This gives you a contradiction to work with, and so you can derive Q by hypothesizing $\sim Q$ for \sim I. This will allow you to show $(P \rightarrow Q)$ which contradicts the assumption, giving you $\sim\sim P$ and thus P .

Proceed similarly to prove $\sim Q$, by hypothesizing Q for \sim I and deriving $(P \rightarrow Q)$ once again.

12. $S \leftrightarrow T, T \vee S \vdash \sim(T \rightarrow \sim S)$

Outline of proof.

Since you have a disjunction, you will want to use \vee E to derive the conclusion twice, each time by \sim I.

So you first show $T \rightarrow \sim(T \rightarrow \sim S)$, by hypothesizing T for \rightarrow I and then $(T \rightarrow \sim S)$ for \sim I.

By involving the biconditional, you will arrive at a contradiction in a few steps.

Repeat the procedure for the second disjunct S , then draw the conclusion by \vee E.

13. $Q \rightarrow R \vdash \sim Q \vee R$

Outline of proof.

Hypothesize $\sim(\sim Q \vee R)$ for \sim I.

The problem now is to generate a contradiction which takes some creativity.

If you look at your hypothesis, you'll see it affords you the opportunity to prove Q by hypothesizing $\sim Q$ and using \vee I to derive a contradiction.

But if you have Q , you have R by \rightarrow E, and then you can use \vee I again to contradict the original hypothesis.