Derive the following in free logic.

Note: Consider the fact that one of the main motivations behind free logic is to create a system in which nothing exists necessarily. This is not achieved in classical predicate logic since both  $\exists x \ x=a \ and \ \exists x \ x=x \ are theorems.$  In free logic, these are not theorems, but the formula above is a theorem. Do you think this means that this system of free logic has failed to achieve its aims? Why or why not?

**Answer**: The fact that this is a theorem is consistent with the aims of free logic because the universal quantifier applies only to existing objects. But from this theorem one can not deduce that some object actually exists, since the rule for universal elimination can be executed only on objects that exist.

2. 
$$\forall xFx$$
,  $\exists x \sim Fx \mid P$  (Don't use QE)

Derive the following in free logic.

- 1. ∃xFx A
- 2. | Fa & ∃x x=a H (F∃E)
- 4. ∃x x=a 2, &E
- 5. a=a =l
- 6. ∃x x=x F∃I 4,5
- 7. ∃x x=x 1, 2-6 F∃E

- 1. ∃xFx A
- 2.  $\mid \exists x \ x=a$  H (F $\forall$ I)
- 3. a=a =l
- 4.  $\forall x x=x$  2-3 (F $\forall$ I)

5. 
$$\vdash \forall x(Fx \rightarrow (Fx \lor Gx))$$

- 1. | ∃x x=a H (F∀I)
- 2. | Fa H (→I)
- 3. | Fa v Ga 2, vI
- 4. Fa  $\rightarrow$  (Fa v Ga) 2-3,  $\rightarrow$ I
- 5.  $\forall x(Fx \rightarrow (Fx \lor Gx))$  1,4  $F \forall I$

Derive the following in free logic.

6. 
$$\forall x(Fx \rightarrow Gx), \exists xFx \mid \exists xGx$$

- 1.  $\forall x(Fx \rightarrow Gx)$
- Α

2. ∃xFx

- Α
- 3. | Fa & ∃x x=a
- H (F∃E)

4. ∃x x=a

3, &E

5. Fa

3, &E

6.  $Fa \rightarrow Ga$ 

1,4 F∀E

7. Ga

5,6 →E

8. ∃xGx

4,7 F∃I

9. ∃x**G**x

2, 3-8 F∃E

1. | ∃x◊~x=x

- H (~I)
- H (F∃E)

2, &E

4. ~□a=a

3, Dual, DN

5. □a=a

Ν

6. P & ~P

4,5 Con

7. P & ~P

1, 2-6 F∃E

8. ~∃x◊~x=x

1-7, ~I

Derive the following in free logic.

2. □∃x x=a

3.  $\Box(\forall x Fx \rightarrow (\exists x x=a \rightarrow Fa))$ 

4. 
$$\Box \forall x Fx \rightarrow \Box (\exists x \ x=a \rightarrow Fa)$$

5.  $\Box(\exists x \ x=a \rightarrow Fa)$ 

Α

N (rule of F∀E)

1.3 →E

## 9. $\Box \forall x(Fx \rightarrow Gx), \Diamond \exists xFx \vdash \Diamond \exists xGx \text{ (Hint: Use result from 6 above)}$

1. 
$$\Box \forall x (Fx \rightarrow Gx)$$

3. 
$$\Box((\forall x(Fx \rightarrow Gx) \rightarrow (\exists xFx \rightarrow \exists xGx)))$$
 N (theorem derived in problem 6)

4.  $\Box \forall x(Fx \rightarrow Gx) \rightarrow \Box(\exists xFx \rightarrow \exists xGx)$ 

5.  $\Box(\exists x Fx \rightarrow \exists x Gx)$ 

 $\Box$ (~ $\exists$ xGx  $\rightarrow$  ~ $\exists$ xFx) 7.

 $\neg \neg \exists xGx \rightarrow \neg \neg \exists xFx$ 8.

□~∃xFx 9.

10. **~**◊∃xFx

11. P & ~P

12. ~□~∃xGx

Α

Α

3, K

1,4 →E

H~I

5, Trans

7, K

6,8 →E

9, Dual, DN

Con 2,10

6-11, ~I

12, Dual

This turned out just to be a propositonal modal logic proof. No quantifier rules are involved.

Note, however, that line 3 was only possible because the theorem was derived in free logic.

In free logic, every theorem introduced involving quantifiers must be a theorem in free logic. Hence, e.g., you can not introduce

 $\Box$  ( $\forall x Fx \rightarrow Fa$ )

since this is not a theorem in free logic.

Derive the following in free logic.

- 1. ∃y□(Fy & ∃x y=x)
- 2. □ (Fa & ∃x a=x) & ∃y y=a
- 3.  $\Box$ ((Fa &  $\exists$ x a=x)  $\rightarrow$   $\exists$ yFy)
- 4. □(Fa & ∃x a=x) → □∃yFy
- 5. □(Fa & ∃y a=y)
- 6. □∃yFy
- **7**. □∃yFy

- Α
- H F∃E
- N (rule of F∃I in free logic)
  - 3, K
  - 2, &E
  - 4,5 →E
  - 1, 2-6 F∃E

11.  $\exists x \Box Fx, \forall x \Box (Fx \rightarrow Gx) \models \exists x \Box Gx$ 

- 1. ∃x□Fx
- 2.  $\forall x \square (Fx \rightarrow Gx)$
- 3. | □Fa & ∃x x=a
- 4. ∃x x=a
- 5. □(Fa → Ga)
- 6. □Fa
- 7. □Fa → □Ga
- 8. □Ga
- 9. ∃x□Gx
- 10.∃x□Gx

- Α
- Α
  - H (F∃E)
  - 3, &E
  - 2,4 F∀E
  - 3, &E
  - 5, K
  - $6,7 \rightarrow E$
  - 4,8 F∃I
  - 1, 3-9 F∃E

Derive the following in free logic.

# 12. $\exists x(Fx \lor Gx) \mid \exists xFx \lor \exists xGx$

1. ∃x(Fx v Gx) A		Α
2.	(Fa v Ga) & ∃x x=a	H (F∃E)
3.	Fa v Ga	2, &E
4.	∃x <b>x=</b> a	2, &E
5.	<sub> </sub> Fa	$H \rightarrow I$
6.	∃xFx	4,6 F∃I
7.	∃xFx v ∃xGx	6, vI
8.	Fa → (∃xFx v ∃xGx)	5-7 →I
9.	<sub> </sub> Ga	$H \rightarrow I$
10.	∃xGx	4,9 F∃I
11.	∃xFx v ∃xGx	10, vI
12.	Ga → (∃xFx v ∃xGx)	5-7 →I
13. ∃xFx v ∃xGx		3,8,12 vE
14. ∃xFx v ∃xGx		1, 2-13 (F∃E)

Derive the following in free logic.

## 13. $\exists x \exists y Gxy \mid \forall x \forall y \ x=y \rightarrow \forall z Gzz$

```
1. ∃х∃уСху
                                                Α
                                                H \rightarrow I
2. |\forall x \forall y x=y|
3.
                                                \mathsf{H}\,\mathsf{F}\forall\mathsf{I}
     ∃x x=a
4.
       ∃yGby & ∃x x=b
                                                H F∃E
       ∃yGby
5.
                                                4, &E
          Gbc & ∃y y=c
                                                H F∃E
6.
7.
          ∃x x=b
                                                4, &E
          ∀y b=y
                                                2,7 F∀E
8.
                                                                          We could have
9.
                                                6, &E
                                                                         gone to Gbb here,
          ∃у у=с
                                                                          but for the ∀FE
10.
          b=c
                                                8,9 F∀E
                                                                         step on line 14 we
11.
          Gbc
                                                6, &E
                                                                        would then require
12.
                                                10, 11 =E<sup>1</sup>
          Gcc
                                                                          ∃y y=b, so you
13.
                                                2,3 F∀E
          ∀у а=у
                                                                        would need to take
14.
                                                9,13 F∀E
          a=c
                                                                          a step to switch
                                                                        the variable from x
15.
                                                12,14 = E
          Gaa
                                                                           to y on line 7.
16.
                                                5, 6-15 F∃E
        Gaa
17.
      Gaa
                                                1, 4-16 F∃E
18.  ∀zGzz
                                                3-17 ∀I
19. \forall x \forall y \ x=y \rightarrow \forall zGzz
                                                2-18 →I
```

Derive the following in free logic.