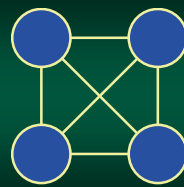




## Graph Theory

### Part 6

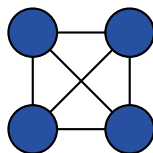


## Graphs

Nope, not the same as charts

## Graphs

- Lists are just a special case of another structure - the *graph*
- Graphs are the basis for all of computer science
- Computer science is not about chips, processors, etc...
- ... this is just implementation technology



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## Motivation

- Several real-life problems can be converted to problems on graphs
- They are one of the pervasive data structures used in computer science
- They are useful tool for modeling real-world problems
- Allows us to abstract details and focus on the problem

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## Where are Graphs Used?

- The easy answer is: *everywhere*
- In computer science
  - state machines
  - mazes and networks
- Other fields
  - chemistry
  - physics
  - government

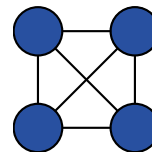
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## Terminology

- The terminology for graphs is a bit different from trees and linked lists
- Rather, it is more generalized
  - "nodes" are called *"vertices"*
  - "branches" or "links" are called *"edges"*



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## Formal Definition

- A graph  $G = (V, E)$  is defined by a pair of two sets
  - a finite set  $V$  of items called vertices
  - a finite set  $E$  of vertex pairs called edges

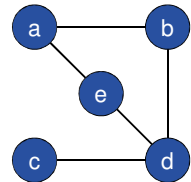
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## Example Graph

- Set of vertices  $V$ 
  - a
  - b
  - c
  - d
  - e



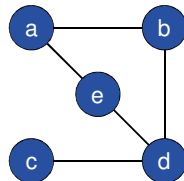
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## Example Graph

- Set of edges  $E$ 
  - (a, b)
  - (a, e)
  - (b, d)
  - (c, d)
  - (d, e)



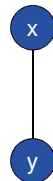
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## Adjacent and Incident

- If two vertices  $x$  and  $y$  share an edge  $(x, y)$ , they are said to be *adjacent*
- The edge  $(x, y)$  is called *incident* on vertices  $x$  and  $y$



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## Directed Graph

- A *directed graph (digraph)* is a graph where each edge has a source and target vertex
- This is the basis most of the data structures used today:
  - trees
  - linked lists

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## Undirected Graphs

- Undirected graphs have edges that link both vertices together
- So, the edge has the same meaning for both directions (set rather than tuple)
- Examples:
  - mathematical equality: if  $a = b$  then  $b = a$
  - marriage: if Jane is married to Joe, Joe is married to Jane

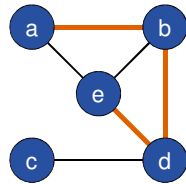
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## Paths

- A **path** is a sequence of vertices  $v_1, v_2, \dots, v_k$  such that consecutive vertices  $v_n$  and  $v_{n+1}$  are adjacent
- This can represent
  - a physical path
  - logical connection
  - etc...



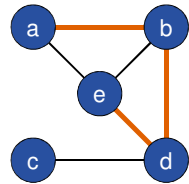
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## Paths

- In a **simple path**, all edges of a path are distinct
- The **length** of a path is measured by either the total number of edges or vertices



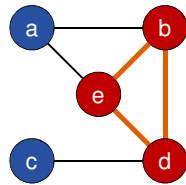
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## Cycles

- Unlike linked lists, graphs can have loops
- A **cycle** is a sequence of vertices  $v_1, v_2, \dots, v_k$  such that consecutive vertices  $v_n$  and  $v_{n+1}$  are adjacent and  $v_1 = v_k$



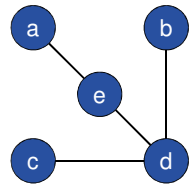
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## Cycles

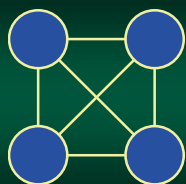
- An **acyclic graph** contains no cycles
- An acyclic directed graph is often called a **DAG** (Directed Acyclic Graph)



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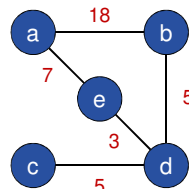


## Types of Graphs

Broad categories for a broad topic

## Weighted Graphs

- Weighted graph** has values on each edge
- These values can be anything – and are defined by the ADT using the graph



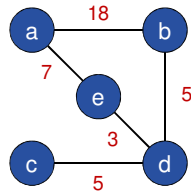
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## Weighted Graph

- Typical uses
  - networks – for finding the fastest speed
  - driving – fastest route
  - etc...
- Example:
  - minimum path from **a** to **b** is:  $a \rightarrow e \rightarrow d \rightarrow b$



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## Connected and Unconnected

- A **connected** graph
  - has a path from every vertex to all other vertex
  - so, everything is connected somehow
- An **unconnected** graph
  - at least one vertex exists in which no path exists to another vertex
  - so, there are 2+ sub-graphs that are unlinked

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## Connected and Unconnected

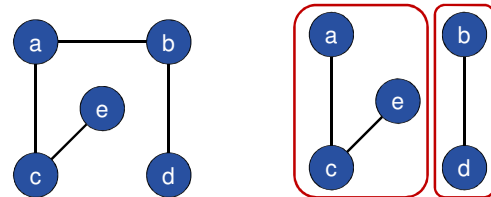
- The **connected component** is the maximum connected subgraph of a given graph
- If the graph is connected, then the whole graph is one single connected component

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## Connected and Unconnected



Connected

Unconnected

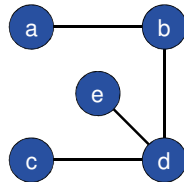
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## Trees

- Tree** is defined as a connected acyclic graph
- Forest** is a collection of trees
- Rooted tree** is a tree graph that selects an arbitrary vertex as the "root"

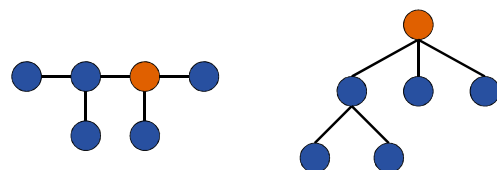


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## Trees



Tree

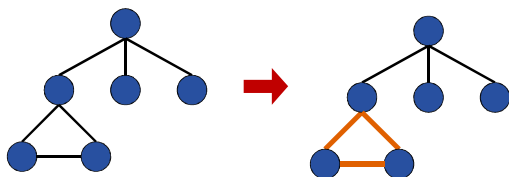
Same Tree

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## Trees



**NOT a Tree**

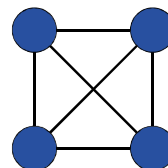
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## Complete Graphs

- A *complete graph* is one in which all pairs of vertices are adjacent
- In other words, every vertex is connected to every other vertex



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## Complete Graphs

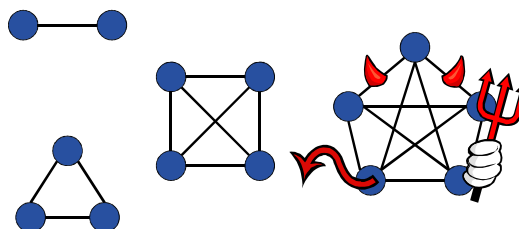
- The number of edges in a complete graph...
  - if  $n$  is the total number of vertices, each vertex is incident to  $n - 1$  edges
  - we can compute  $n \times (n - 1)$  edges, but this would count each edge twice!
  - so, the number of edges =  $n \times (n - 1) / 2$
- So, for a non-complete graph...
  - number of edges  $< n \times (n - 1) / 2$

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## Example Complete Graphs



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## Birth of Graph Theory

... and, later, computer science!

## Birth of Graph Theory

- Graph theory was created due to a perplexing question troubling mathematician Leonhard Euler
- Lived in Königsberg (now "Kaliningrad" in Russia)
  - city occupied 2 islands plus areas on both banks
  - 7 bridges over the Pregel River



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## Birth of Graph Theory

- People wondered if they could leave home...
  - cross every bridge exactly once
  - and return to their starting point
- This is known as the *Konigsberg Bridge Problem*

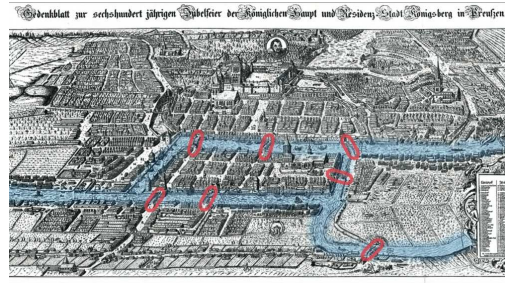


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## Konigsberg Bridge Problem

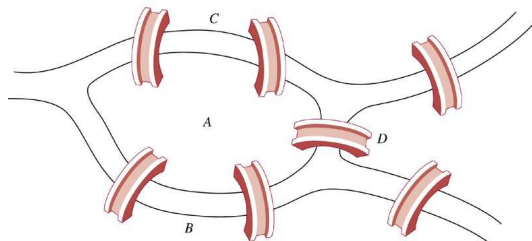


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## Konigsberg Bridge Problem



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## Konigsberg Bridge Problem

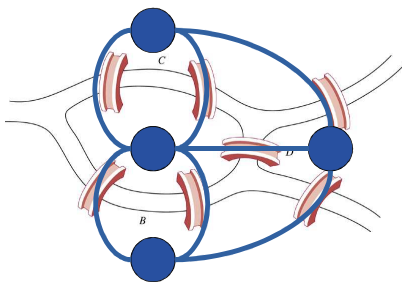
- The problem reduces down to a number of points and links to each point
- From this, Euler created the first graph and began the study of their properties

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## Konigsberg Bridge Problem

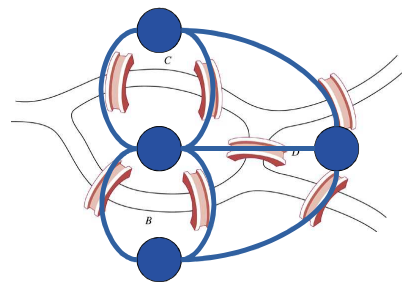


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## Konigsberg Bridge Problem



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## The Solution to Konigsberg

- In 1736, Euler proved that no such traversal exists
- From his work on Konigsberg, a path is named after him
- An *Eulerian circuit*, in a graph...
  - is cycle containing all the edges in the graph
  - and only traversing each edge once

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## The Solution to Konigsberg

- Euler proved:
  - a graph may have an Eulerian circuit **if and only if** there are no vertices with an odd number of edges touching them
- Konigsberg Bridge Problem
  - 4 vertices, all with an odd number of edges
  - Sorry people of Konigsberg, there is no solution!

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## Alan Turing

- Mathematician, logician & cryptographer
- *Father of Computer Science*
  - Highest award in Computer Science is the **Turing Award**
  - Developed Turing Machines



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## Major Work: Turing Machines

- Invented in **1937**
- Logical model – not an actual computer or machine
- Based on 2 graphs (and sets on each of the edge)
- One graph is simple array, but the other could be anything
- From this, he proved programming



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## Major Work: Turing Test

- Used in artificial intelligence
- Consists of a human operator *texting* a human **or** computer
- If the operator can't ascertain if it is a computer or human, the computer is "intelligent"
- No computer has passed it



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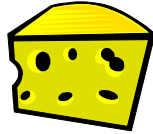


Real World  
Examples

The origin and the usage

## Real World Examples

- How many layers does a computer chip need so that wires in the same layer don't cross?
- How can the season of a sports league be scheduled into the minimum number of weeks?



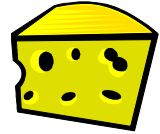
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## Real World Examples

- In what order should a traveling salesman visit cities to minimize travel time?
- Can we color the regions of every map using four colors so that neighboring regions receive different colors?



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## Real World Examples

- How can we lay cable at minimum cost to make every network reachable from every other?
- What is the fastest route from the national capital to each state capital?
- How can  $n$  jobs be filled by  $n$  people with maximum total utility?



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## The London Underground Subway



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## Maze Traversal

- One example of where a graph is useful is a maze traversal
- Basically, any maze can be represented with a graph
- ... and this is not so much different to how networks actually work
- ... a source must find a destination through various vertices

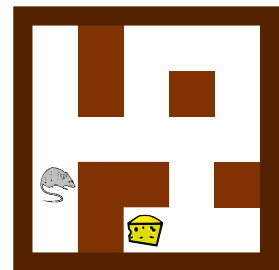
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## Maze Traversal

- This is a simple maze – though not to the mouse!
- We can help him find the cheese if we convert this to a graph



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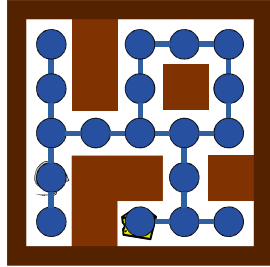
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## Maze Traversal

- The empty spaces are vertices
- The bordering ones are connected with an edge
- Is this a directed graph?



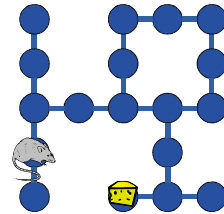
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## Maze Traversal

- So, to help the mouse, we can get to depth-first search on the maze
- If we find it, we can print off the vertices post-order



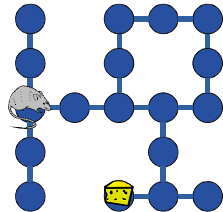
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## Maze Traversal

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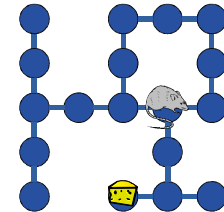
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## Maze Traversal

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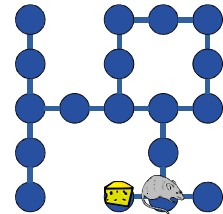
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## Maze Traversal

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