

# BST Questions

lecture 1

Is A Binary Search Tree

Is Balanced Tree

Largest Bst Subtree

Size, Sum, Max, Min, Find In Bst

lecture 2

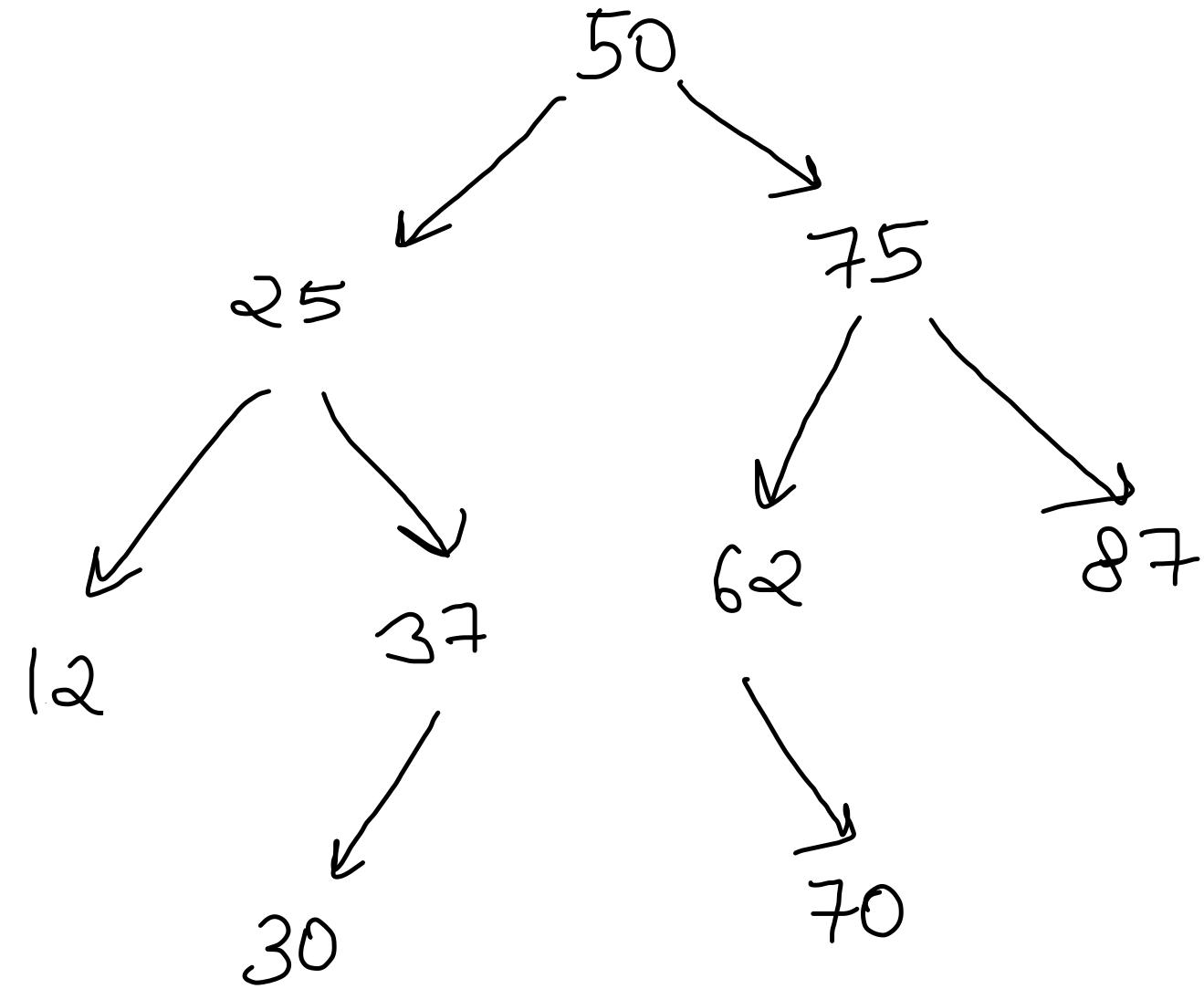
Binary Search Tree - Introduction

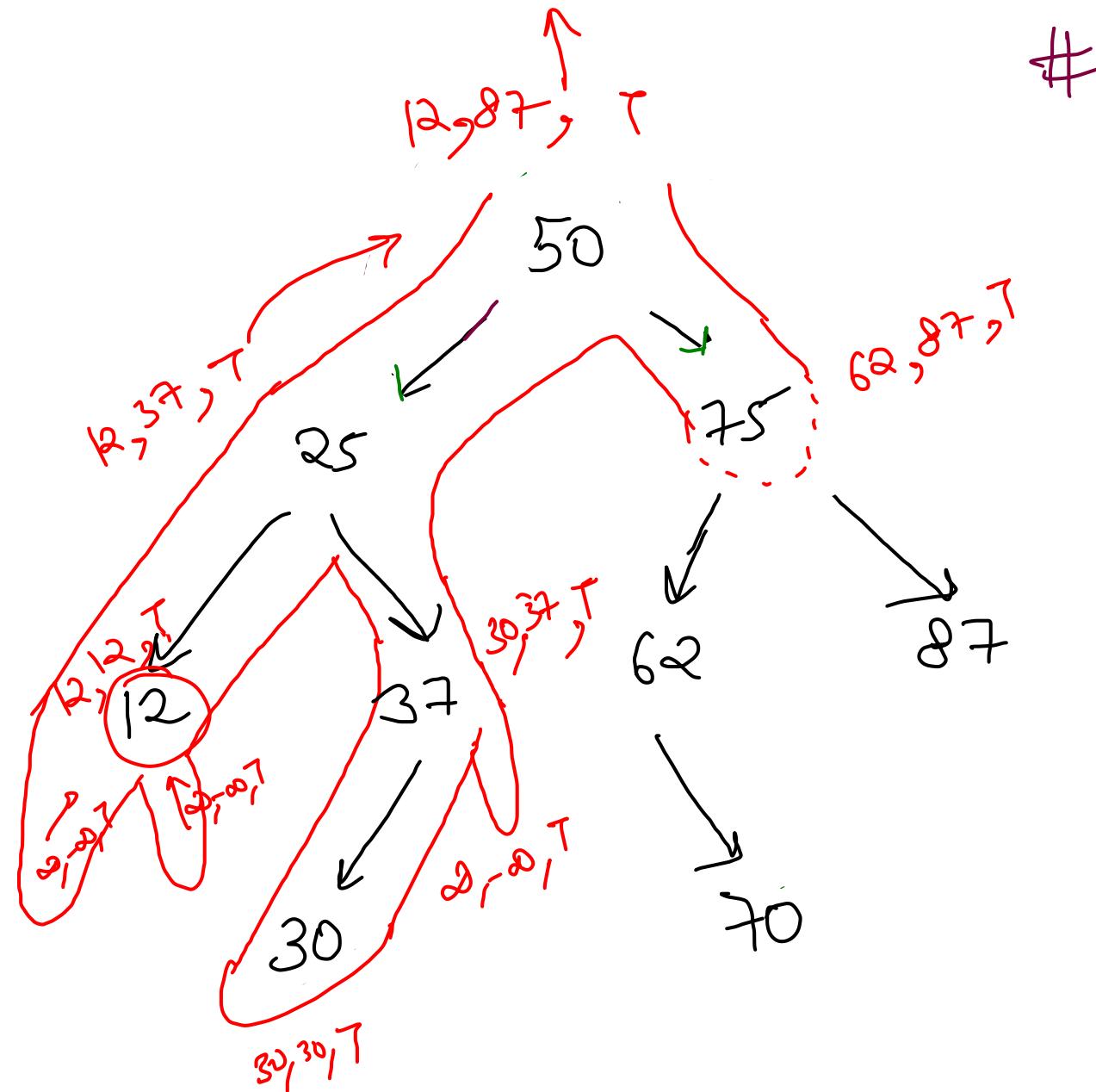
Binary Search Tree - Constructor

Add Node To Bst

lecture 3

- </> Remove Node From Bst
- </> Replace With Sum Of Larger
- </> Lca Of Bst
- </> Print In Range
- </> Target Sum Pair In Bst
- ▣ Target Sum Pair - Bst - Alternate Approaches
- ▣ Iterable And Iterator





# Binary Search Tree # unique nodes

① left nodes < root node < right nodes

isBinaryTree ABST?

1st approach

```

public static isBSTPair isBST(Node root){
    if(root == null) return new isBSTPair();

    isBSTPair left = isBST(root.left);
    isBSTPair right = isBST(root.right);

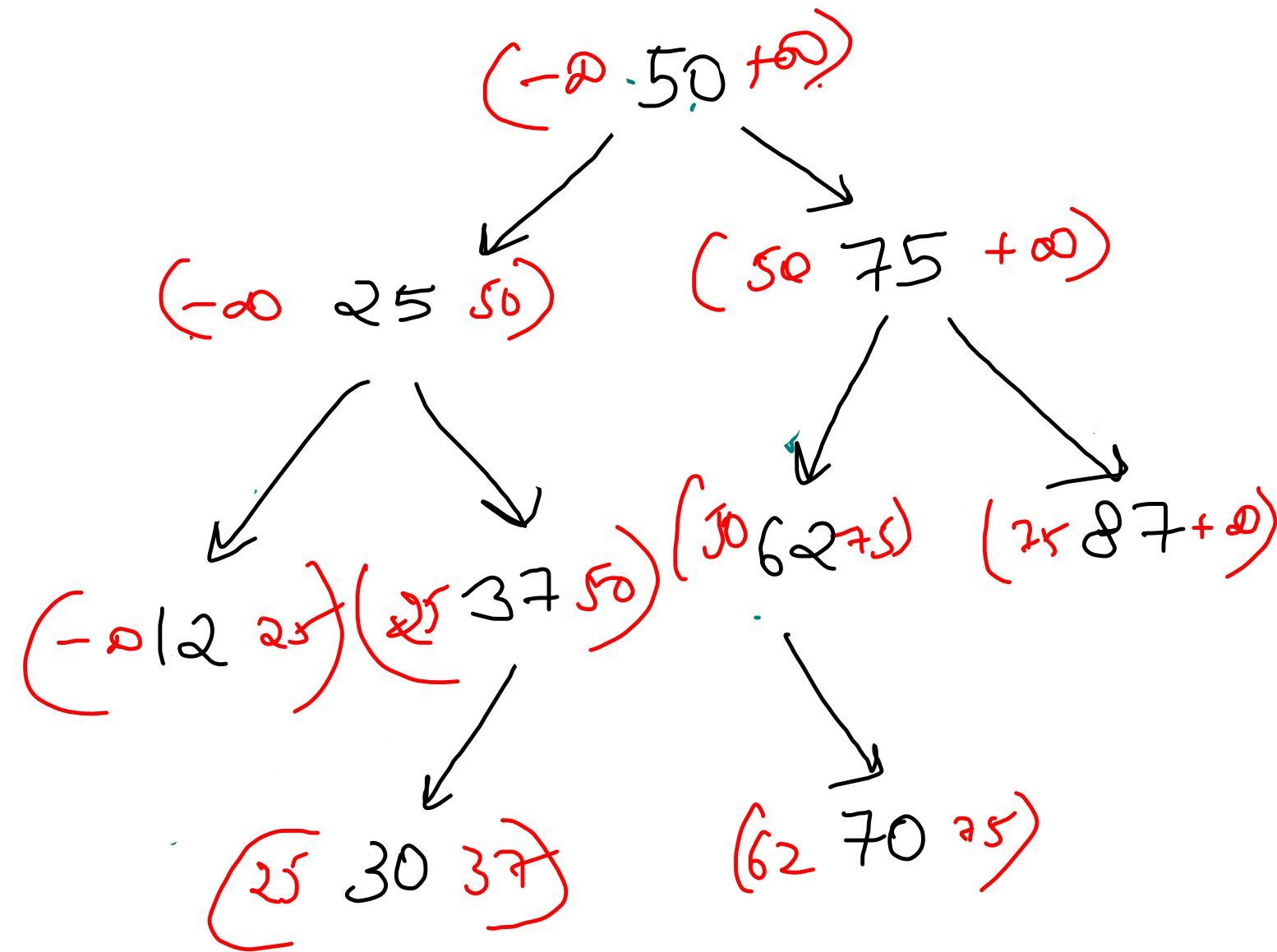
    isBSTPair curr = new isBSTPair();
    if(left.max < root.data && root.data < right.min
        && left.isBST && right.isBST){
        curr.isBST = true;
    } else {
        curr.isBST = false;
    }

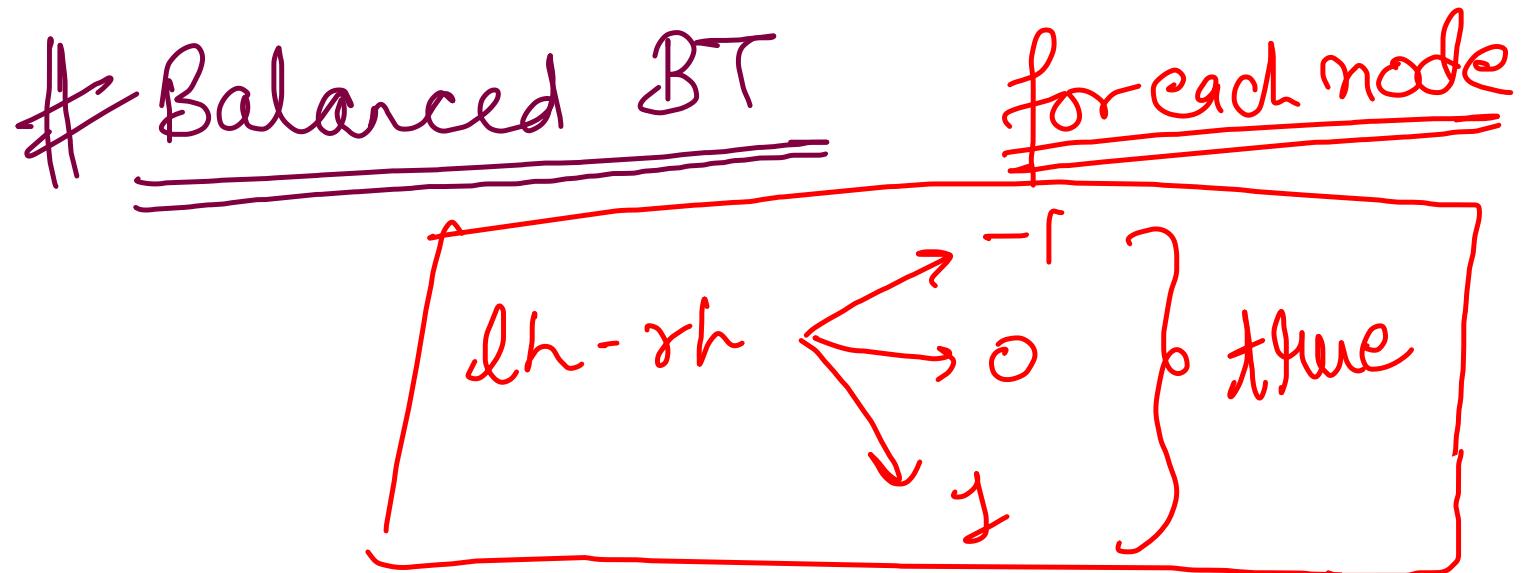
    curr.min = Math.min(root.data, Math.min(left.min, right.min));
    curr.max = Math.max(root.data, Math.max(left.max, right.max));
    return curr;
}

```

## 2<sup>nd</sup> approach

```
public static boolean isBST(Node root, int min, int max){  
    if(root == null) return true;  
  
    if(root.data > min && root.data < max){  
        boolean lres = isBST(root.left, min, root.data);  
        boolean rres = isBST(root.right, root.data, max);  
        return lres && rres;  
    }  
    return false;  
}
```





```

public static balPair isBalanced(Node root){
    if(root == null) return new balPair();

    balPair left = isBalanced(root.left);
    balPair right = isBalanced(root.right);

    balPair curr = new balPair();

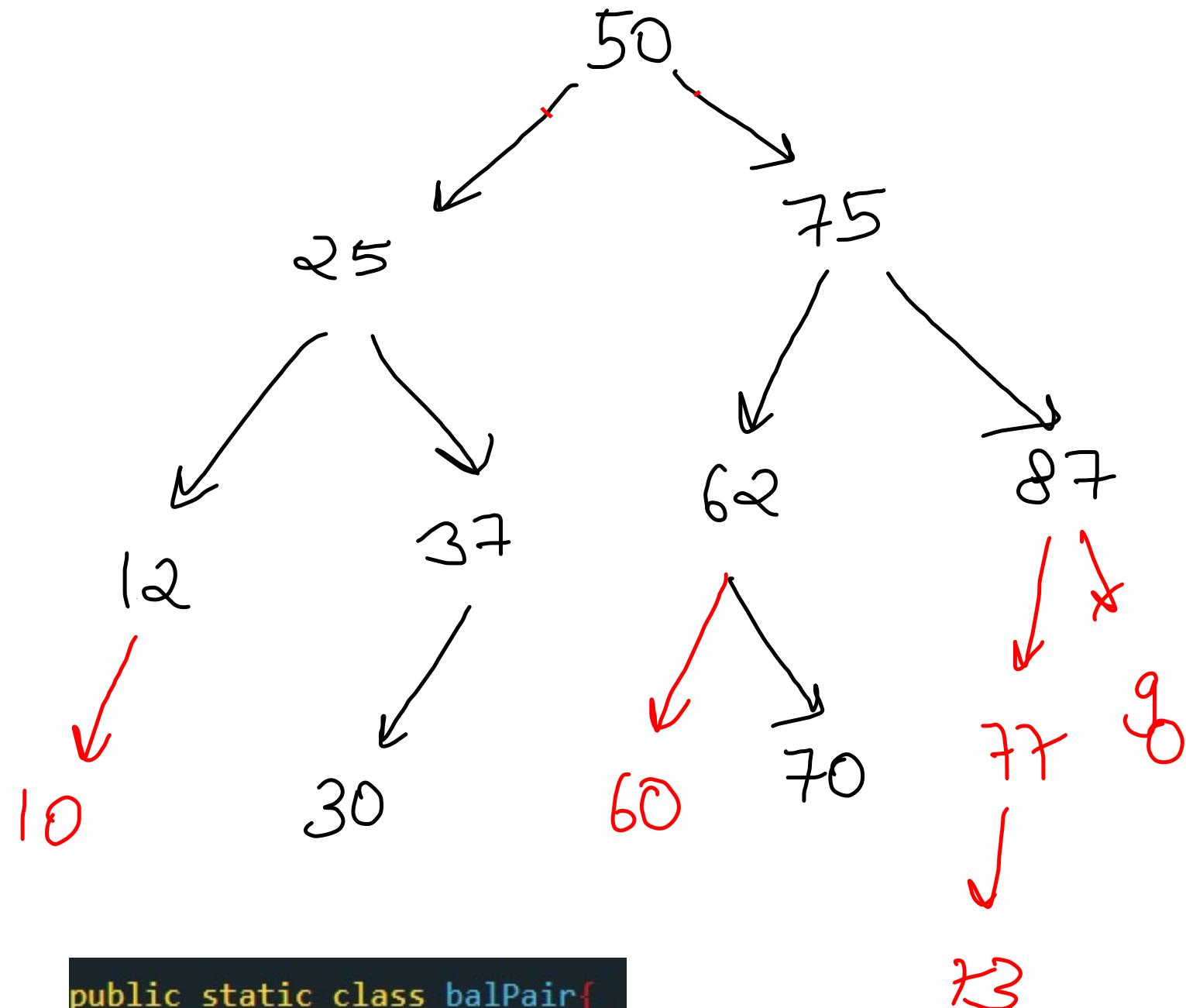
    if(left.isBal && right.isBal && Math.abs(left.height - right.height) <= 1){
        curr.isBal = true;
    } else {
        curr.isBal = false;
    }

    curr.height = Math.max(left.height, right.height) + 1;
    return curr;
}

```

TC  $\Rightarrow O(N)$

SC  $\Rightarrow O(h)$



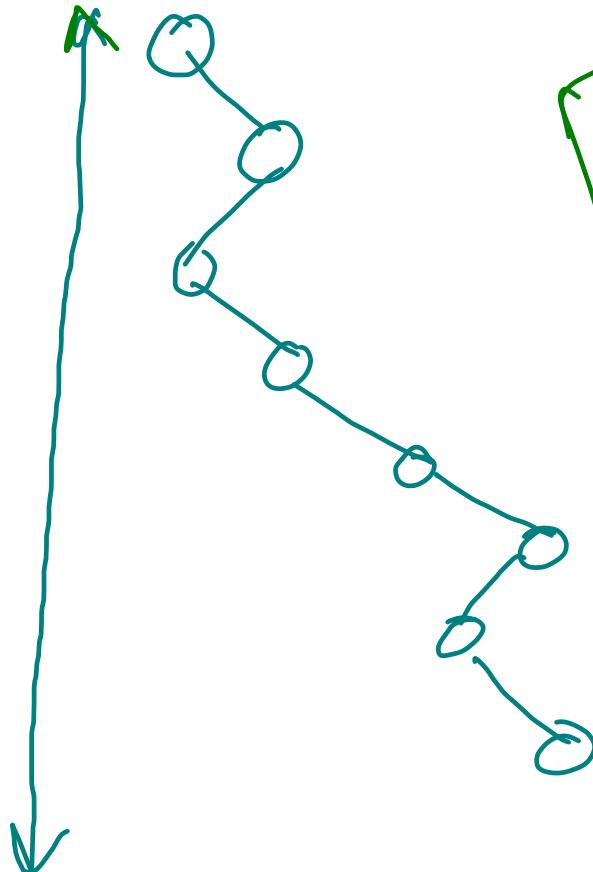
```

public static class balPair{
    int height = 0;
    boolean isBal = true;
}

```

## nodes to height formula

N nodes



Skewed tree  
 $h_{max} = N$

Recursion (Space)

$O(h)$

$O(\text{height or depth})$

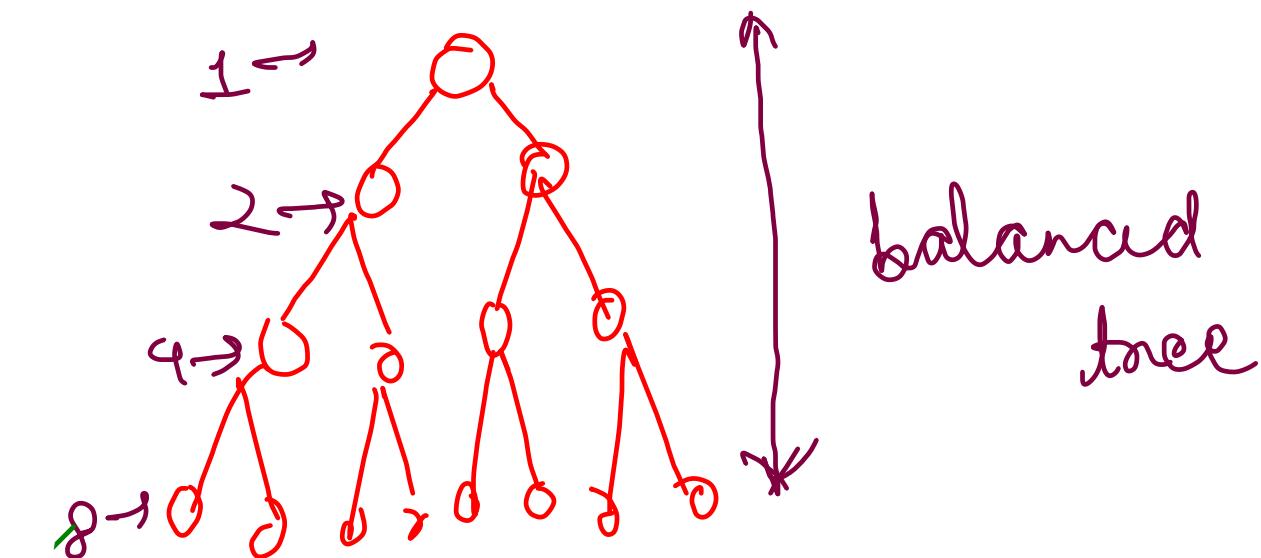
skewed  $\Theta(n)$

balanced  $\Theta(\log n)$

$\Theta(\log_2 N) \rightarrow$  Best case

$\Theta(\log N) \rightarrow$  Avg case

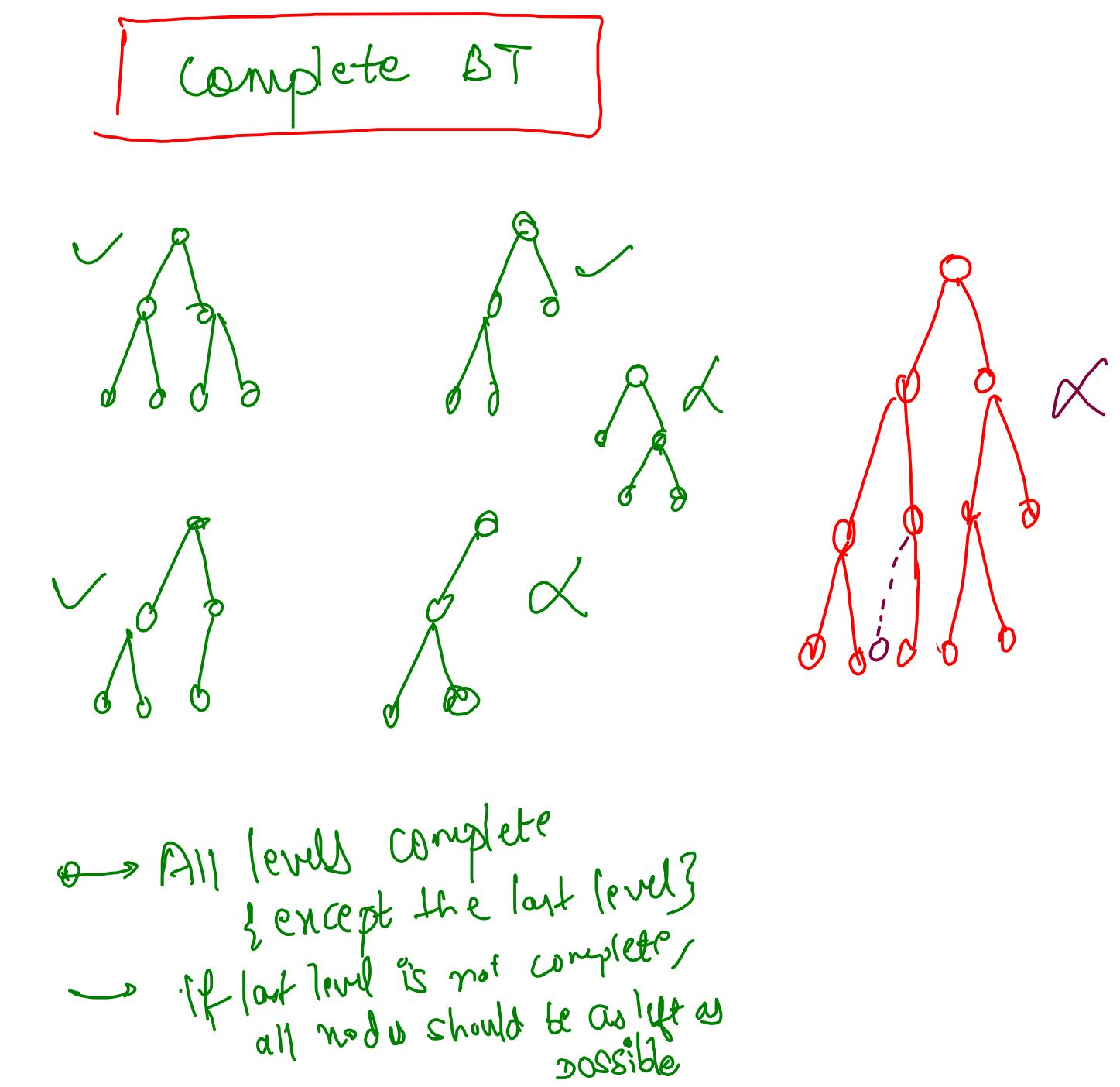
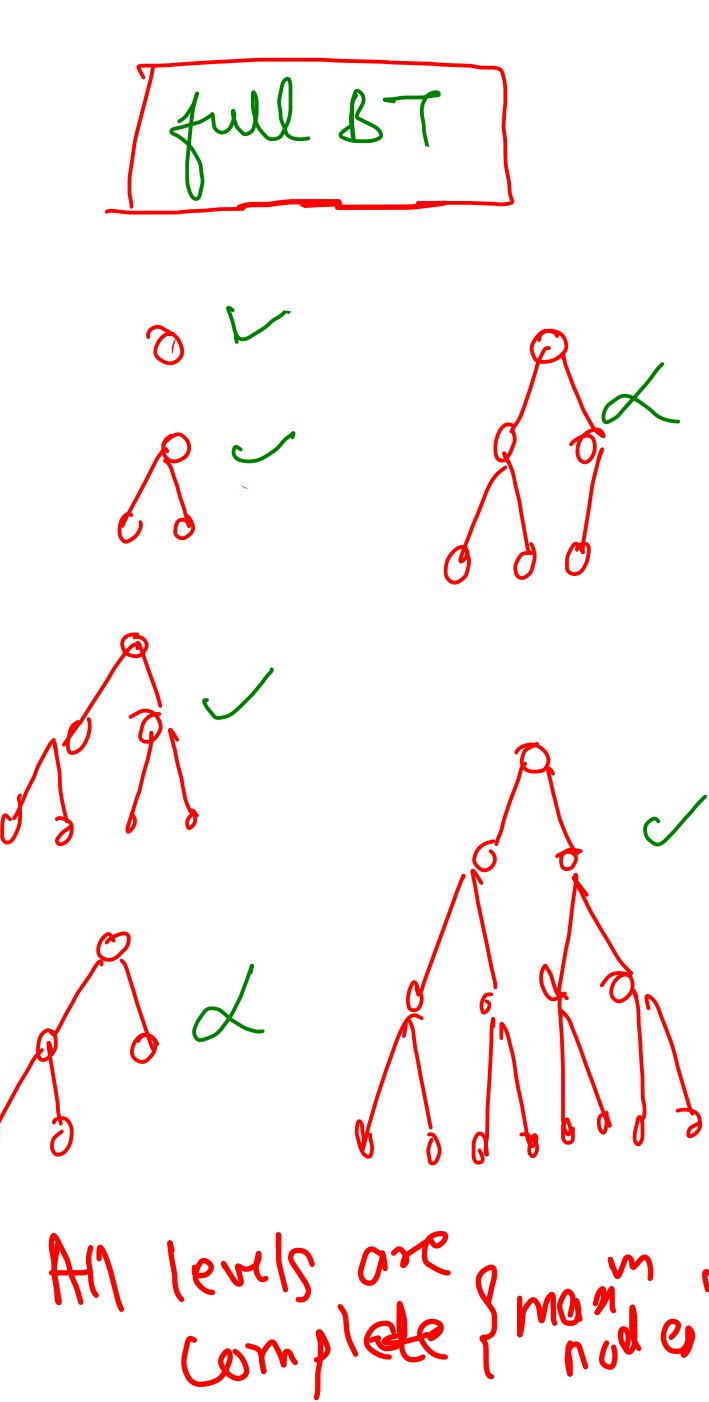
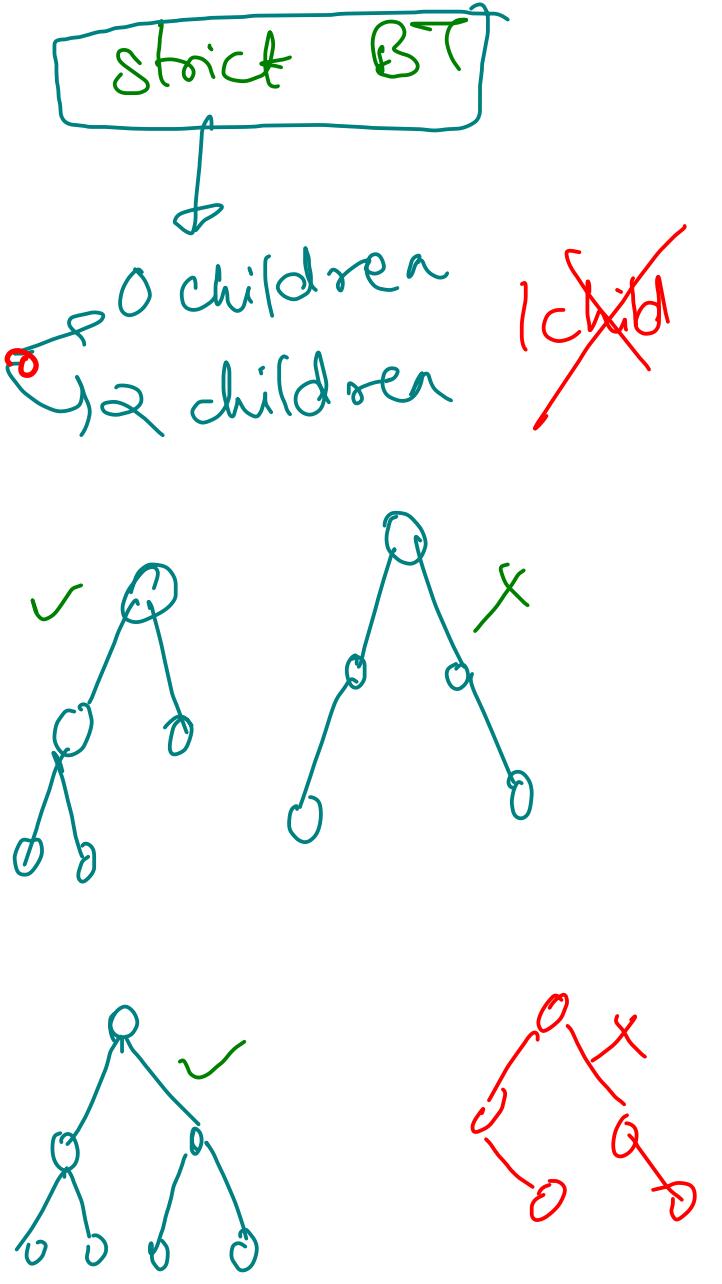
$\Theta(N) \rightarrow$  Worst case



$$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{h-1} = N$$

$$2^h = N$$

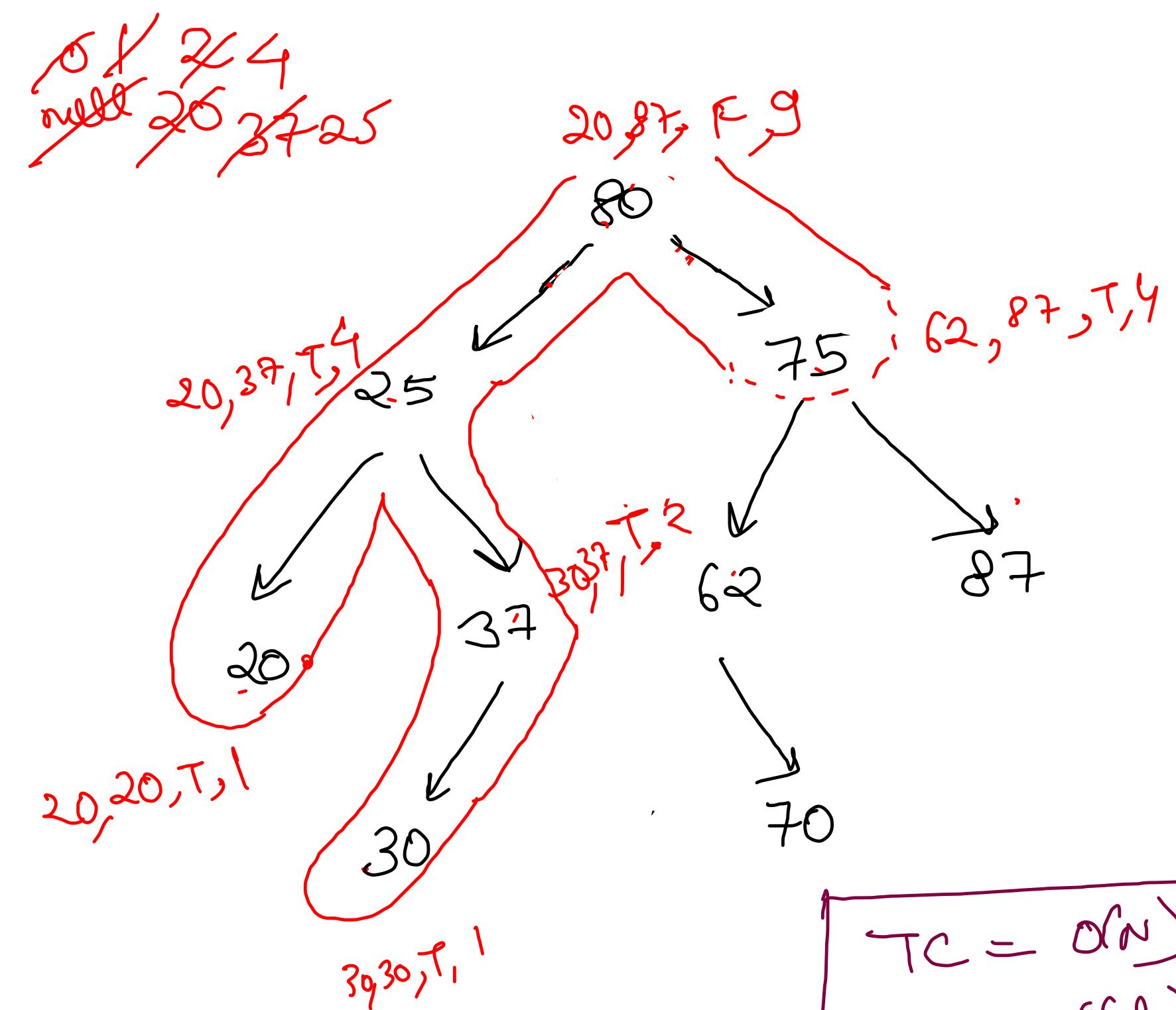
$h = \log_2 N$



All full BT are complete BT  
but complete BT may or may not be full.

# Largest BST subtree

```
public static class isBSTPair{  
    int max = Integer.MIN_VALUE;  
    int min = Integer.MAX_VALUE;  
    boolean isBST = true;  
    int count = 0;  
}  
  
static int maxCount = 0;  
static Node largestBST = null;  
  
public static isBSTPair largestBST(Node root){  
    if(root == null) return new isBSTPair();  
  
    isBSTPair left = largestBST(root.left);  
    isBSTPair right = largestBST(root.right);  
  
    isBSTPair curr = new isBSTPair();  
    curr.count = left.count + right.count + 1;  
  
    if(left.max < root.data && root.data < right.min  
        && left.isBST && right.isBST){  
        curr.isBST = true;  
  
        if(curr.count > maxCount){  
            maxCount = curr.count;  
            largestBST = root;  
        }  
    }  
    else {  
        curr.isBST = false;  
    }  
  
    curr.min = Math.min(root.data, Math.min(left.min, right.min));  
    curr.max = Math.max(root.data, Math.max(left.max, right.max));  
  
    return curr;  
}
```

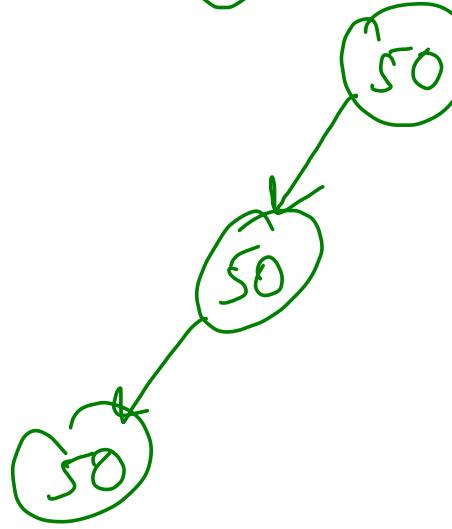


$$TC = O(n)$$

$$SC = O(h)$$

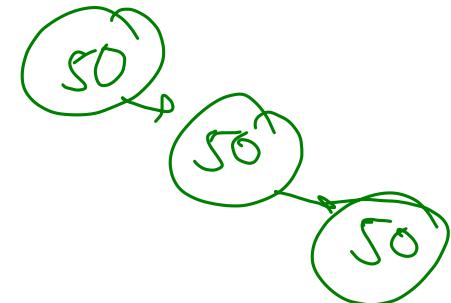
# ~~BST nodes duplicate~~

①

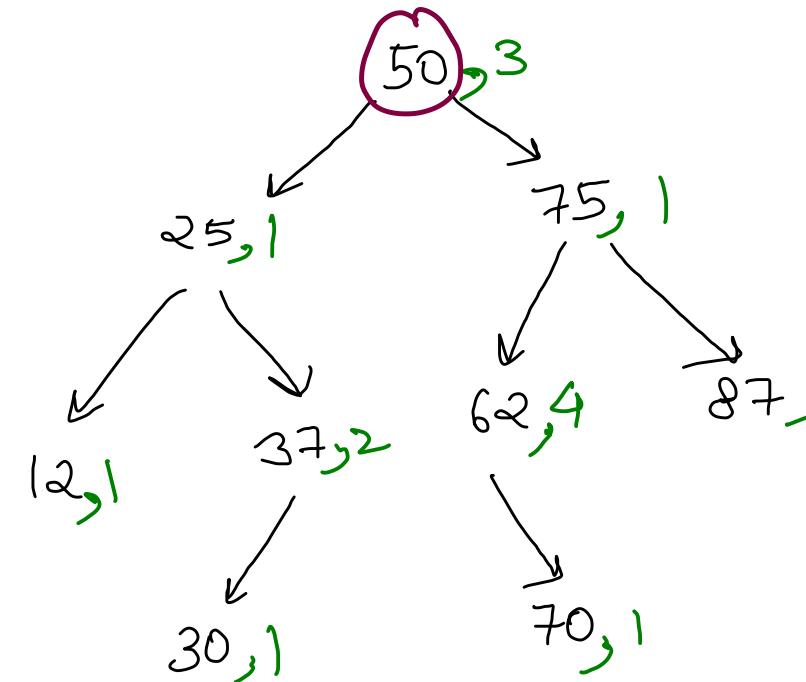


②

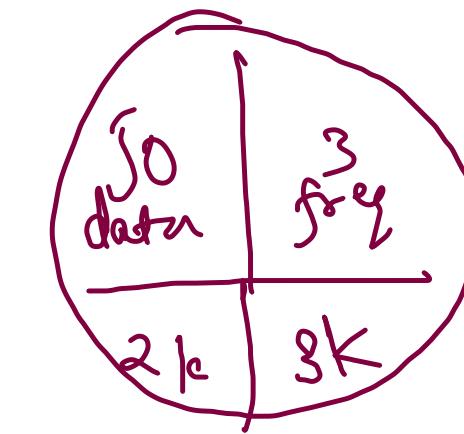
$<$  left,  $\geq$  right



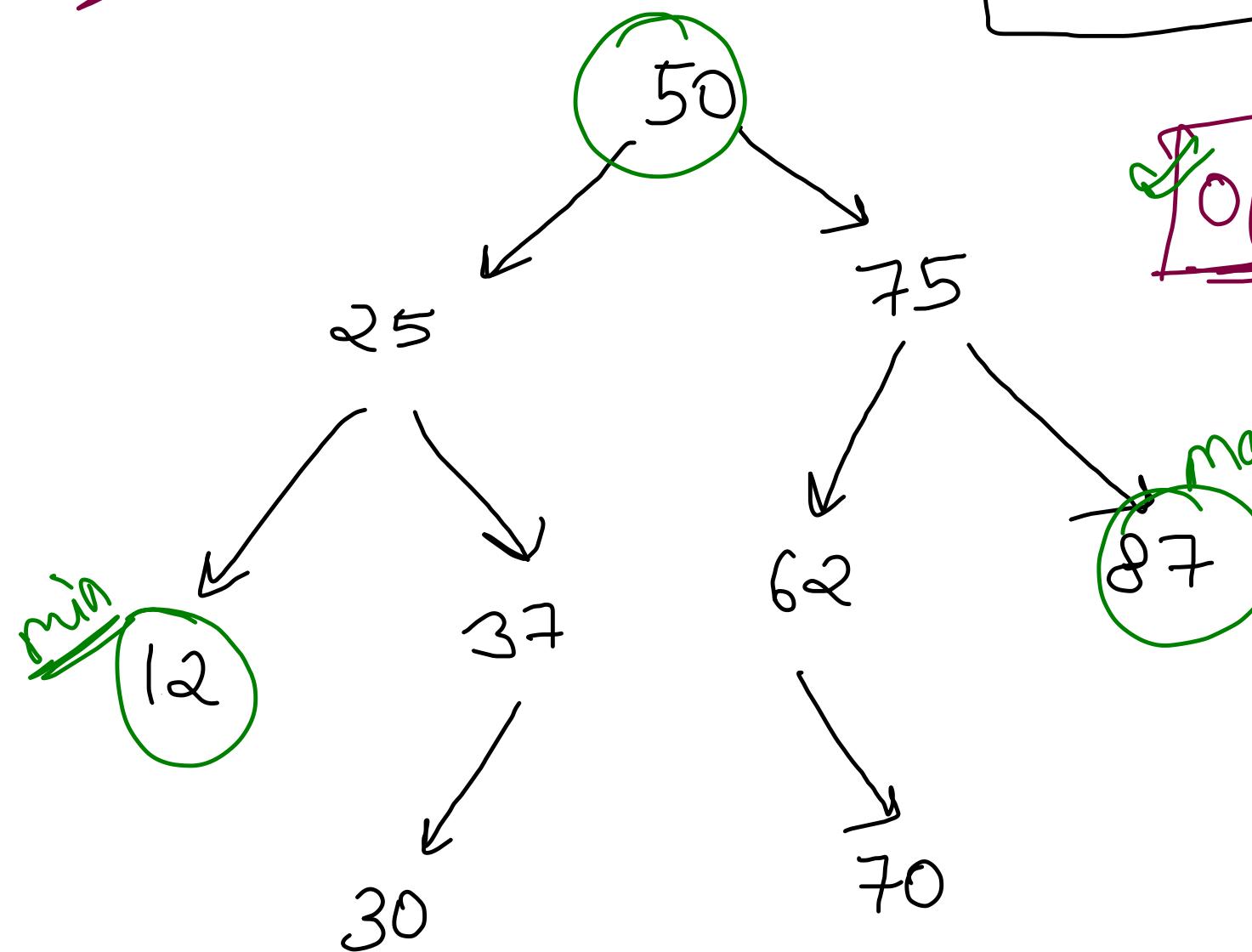
③ frequency



Node  
 $\xrightarrow{\hspace{1cm}}$   
 $\xrightarrow{\hspace{1cm}}$  int freq



Max & Min



Time  $O(N)$   $\Leftrightarrow \{ \text{size}, \text{sum} \rightarrow \text{same as BT} \}$



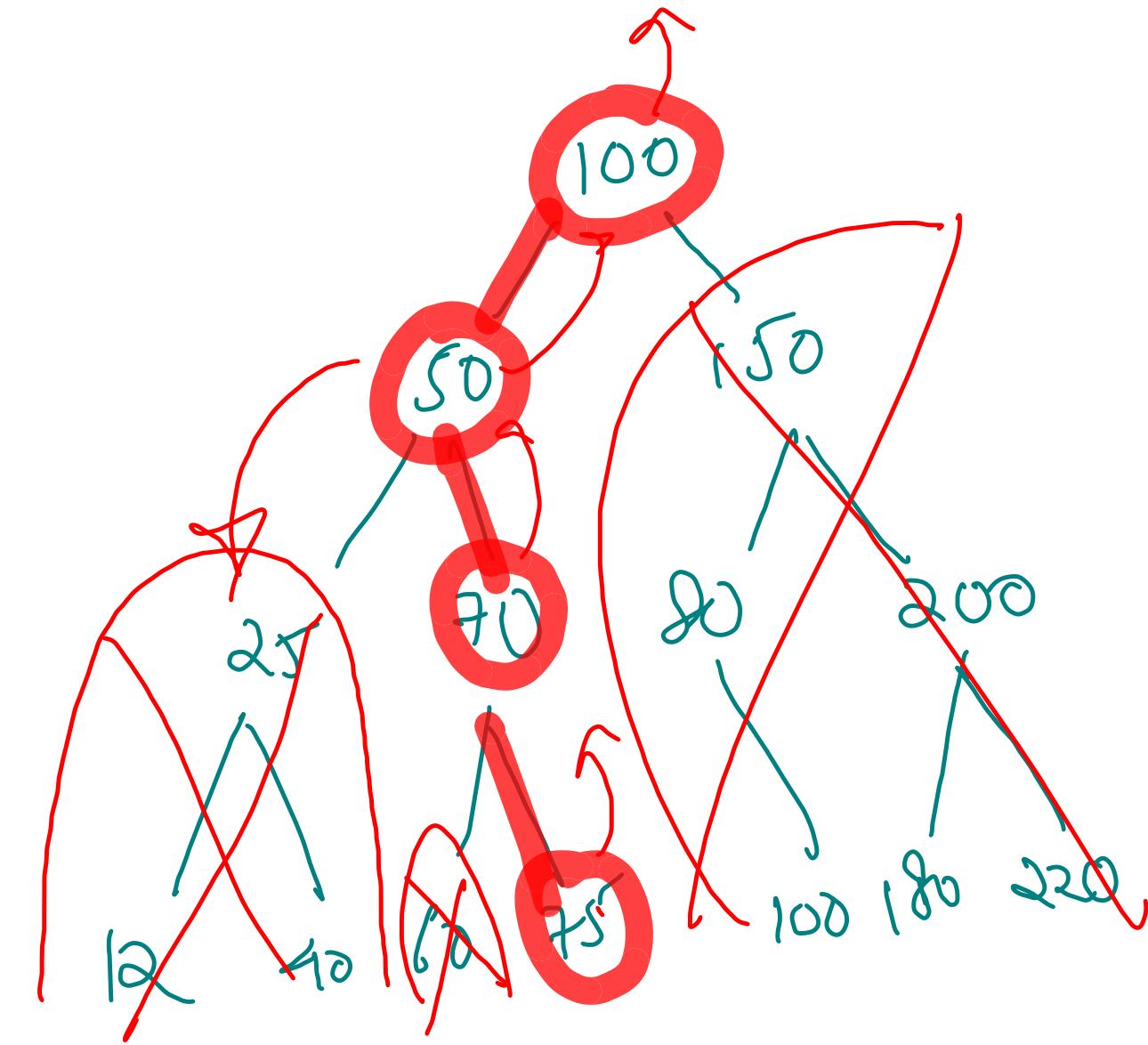
Move as left as possible

move as right as possible

```
public static int max(Node node) {  
    if(node.right == null)  
        return node.data;  
    return max(node.right);  
}  
  
public static int min(Node node) {  
    if(node.left == null)  
        return node.data;  
    return min(node.left);  
}
```

# Find/Search in BST

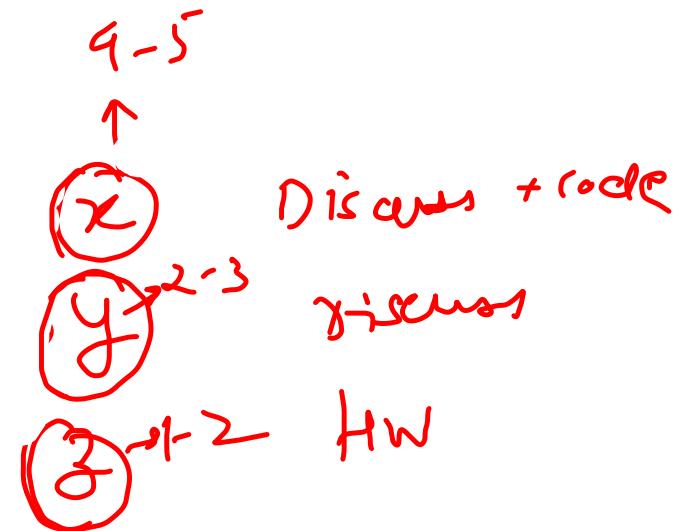
```
public static boolean find(Node node, int data){  
    if(node == null) return false;  
  
    if(node.data == data) return true;  
  
    if(data < node.data)  
        return find(node.left, data);  
  
    return find(node.right, data);  
}
```



BST search  
 $\rightarrow O(h) = \begin{cases} \text{Balanced BST} \\ O(\log_2 N) \end{cases}$

## # Tentative Level 1 Schedule

- ① 23rd October :- BST Lecture 2 (Saturday)
- ② 24th October :- BST Lecture 3 (Sunday Morning)  
HashMap Lecture 1 (Sunday Evening)
- ③ 28th October :- HashMap Lecture 2 (Thursday)
- ④ 30th October :- Heap Lecture 1 & 2 (Saturday Morning & Evening)
- ⑤ 31st October :-  
    Morning class → Remaining Questions.  
    ( Binary Search, + Other questions)  
    Evening class off
- ⑥ 4th November :- Diwali off (Thursday)
- ⑦ 6th November :- Bhaiya Dhuji off (Morning and/or Evening class)
- ⑧ 7th Nov - 14th Nov : Graph Lectures Resumes { 6 lectures }
- ⑨ 18th Nov (Thursday) & 20th Nov (Saturday AM & PM) :- Level 1 revision off
- ⑩ 21st Nov (Sunday) :- Level 2 start



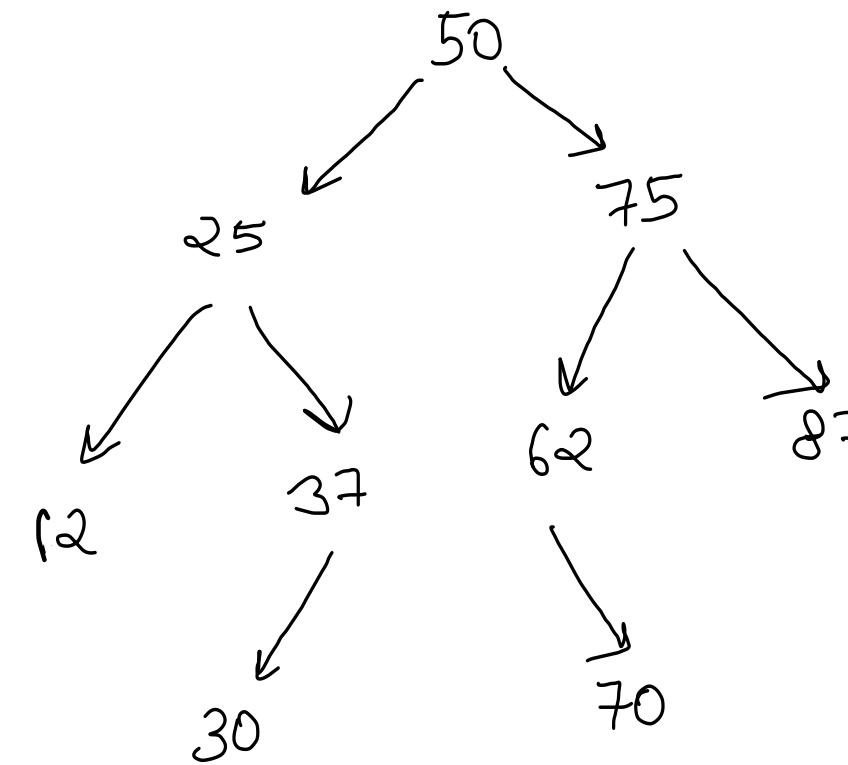
② Property :→ Inorder traversal of BST will always give sorted order.

### Inorder Traversal

```
public static void inorder(Node root){  
    if(root == null) return;  
  
    ① → inorder(root.left);  
    System.out.println(root.data);  
    ③ → inorder(root.right);  
}
```



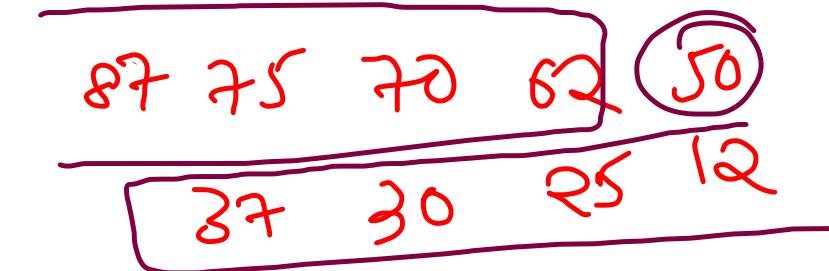
Increasing order



DFS | Euler traversal

### Reverse-Inorder traversal

```
public static void revInorder(Node root){  
    if(root == null) return;  
  
    revInorder(root.right);  
    System.out.println(root.data);  
    revInorder(root.left);  
}
```



Decreasing order

## ~~Construction~~

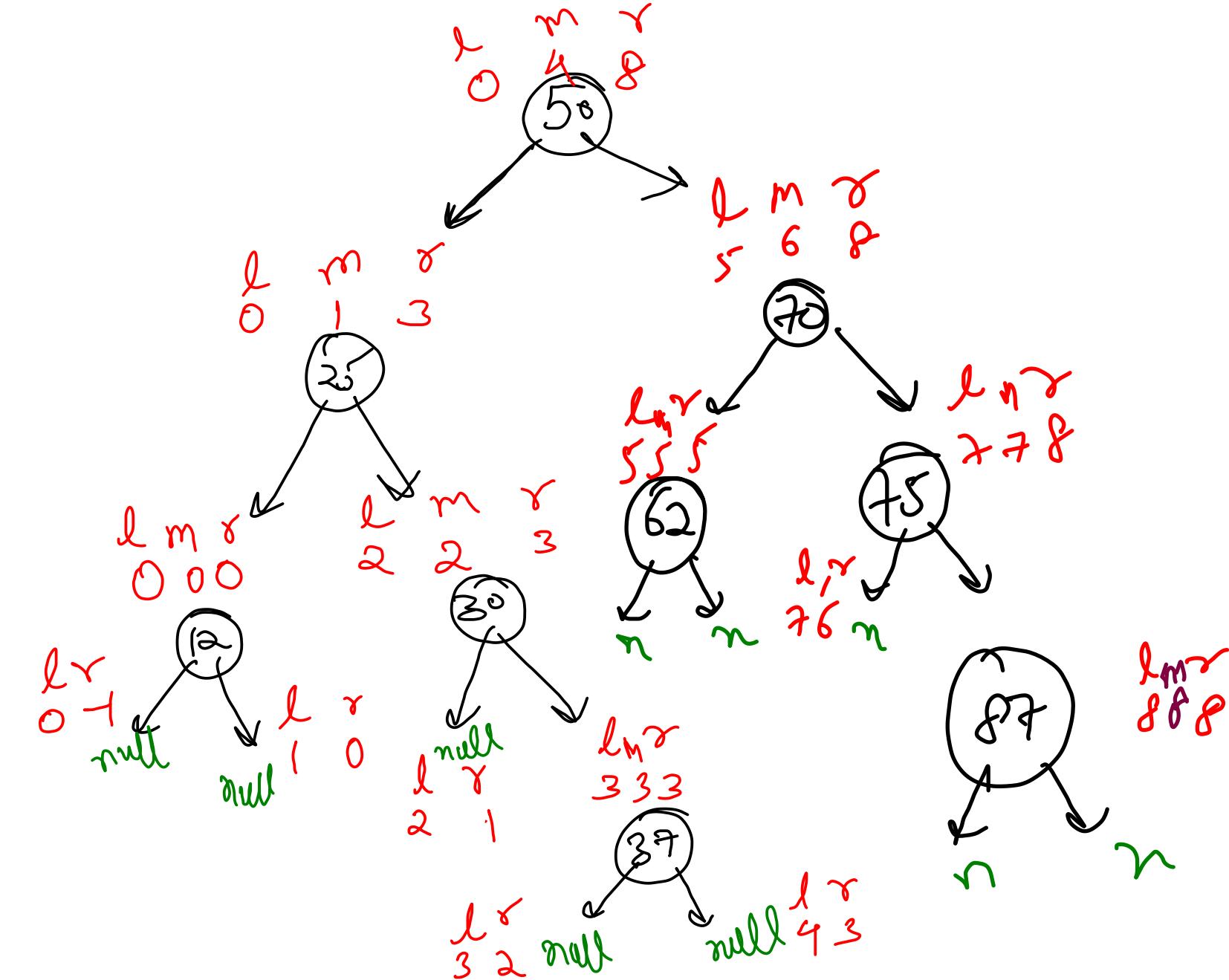
$\{12, 25, 30, 37, 50, 62, 70, 75, 87\}$

$l=0$        $r=n-1$

```
public TreeNode construct(int[] nums, int l, int r){  
    if(l > r) return null;  
  
    int mid = (l + r) / 2;  
  
    TreeNode root = new TreeNode(nums[mid]);  
  
    root.left = construct(nums, l, mid - 1);  
    root.right = construct(nums, mid + 1, r);  
    return root;  
}  
  
public TreeNode sortedArrayToBST(int[] nums) {  
    TreeNode root = construct(nums, 0, nums.length - 1);  
    return root;  
}
```

$O(N)$  Time Comp

$O(1) = O(\log_2 N)$  Recursion

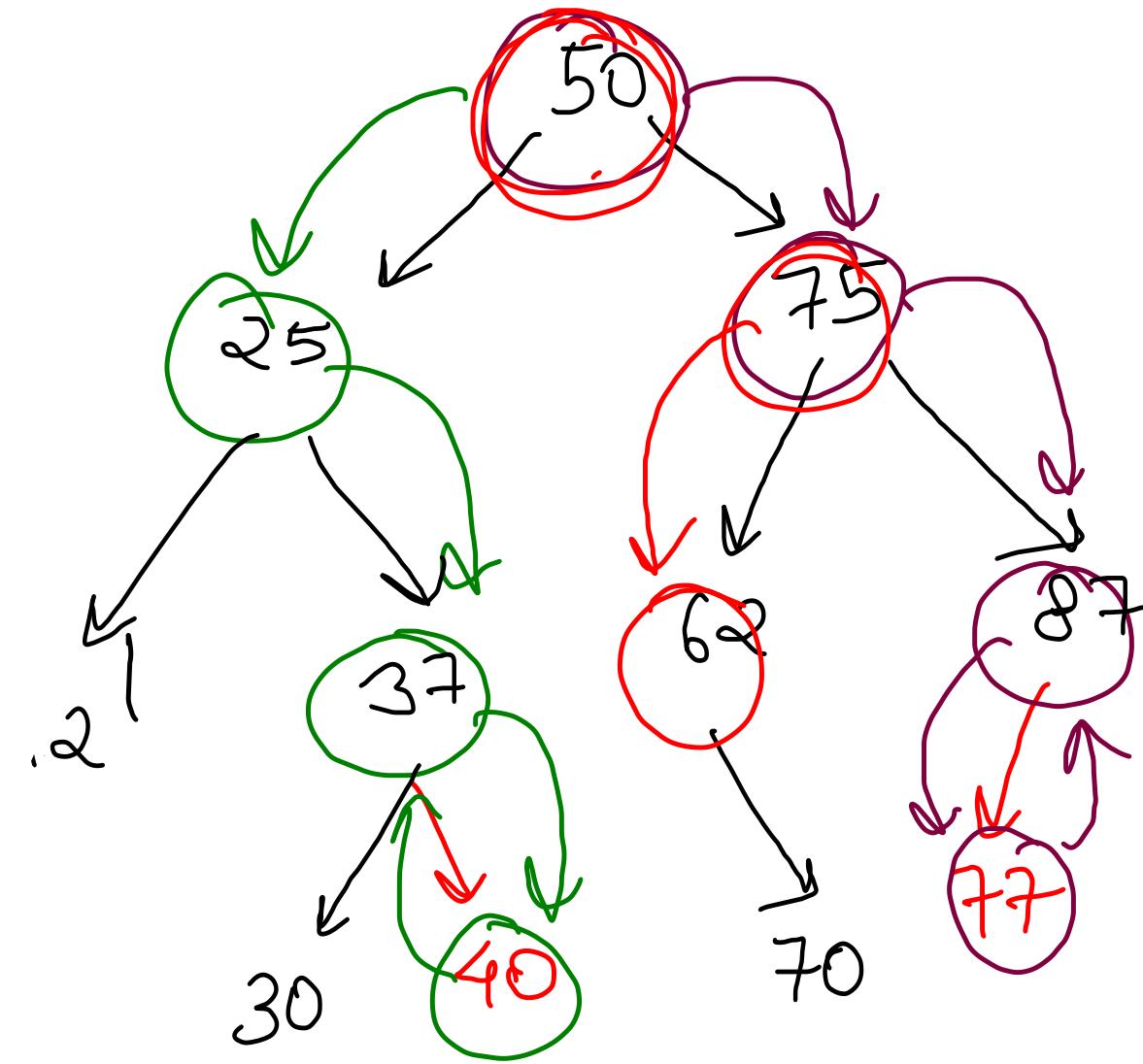


## # Insert Nodes / Add Nodes

~~O(1)~~  
Time

$\Rightarrow \begin{cases} O(\log_2 n) & \text{balanced} \\ O(n) & \text{skewed} \end{cases}$

```
public static Node add(Node node, int data) {  
    if(node == null){  
        Node newNode = new Node(data, null, null);  
        return newNode;  
    }  
  
    if(node.data == data) return node;  
  
    else if(node.data < data)  
        node.right = add(node.right, data);  
    else  
        node.left = add(node.left, data);  
  
    return node;  
}
```



## # Remove node in BST

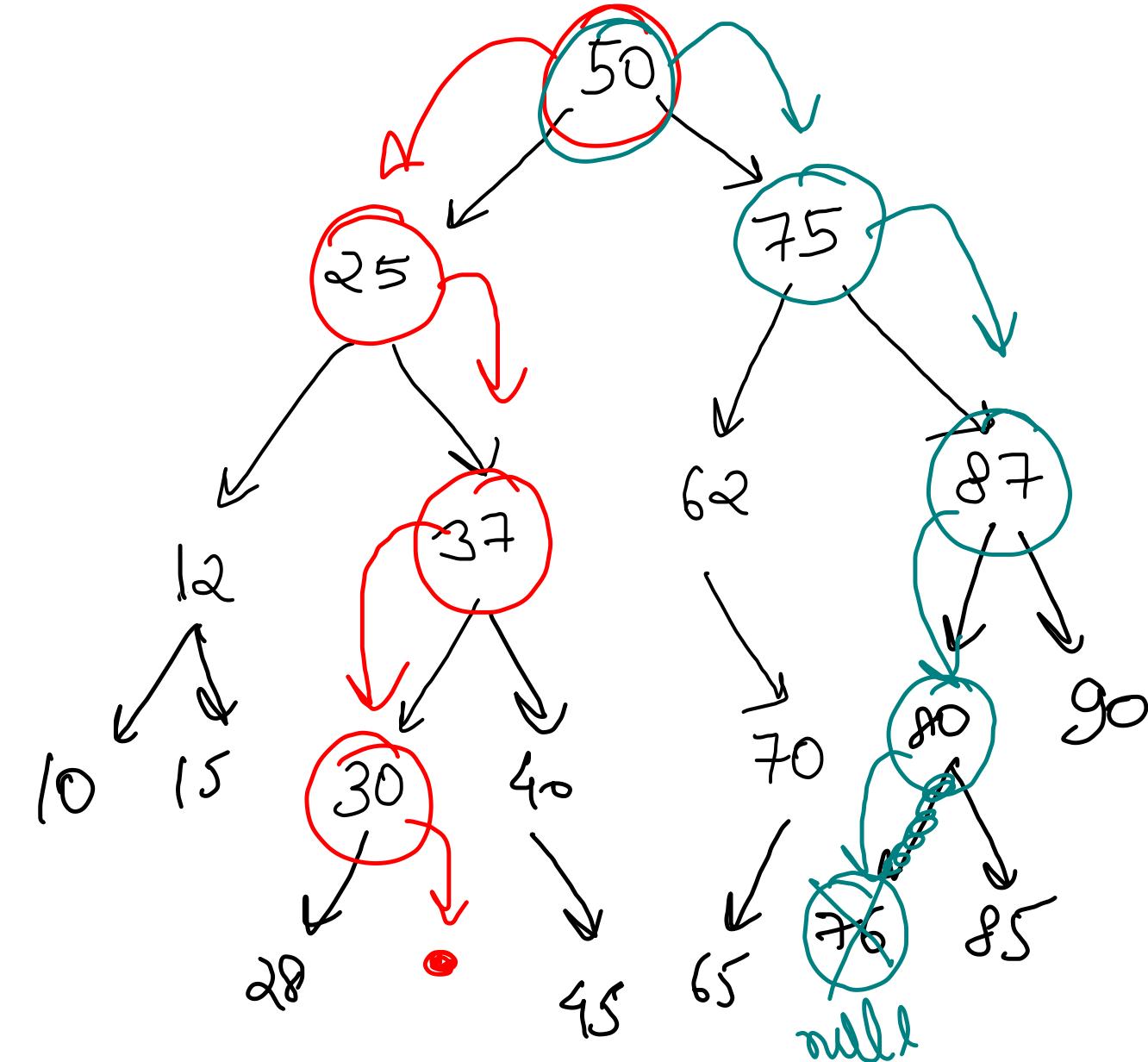
→ Node does not exist

→ Leaf node (0 child)

→ Single child (1 child)

→ 2 children node

```
if (node == null) return null;  
if (node.data == data) {  
    leaf node: return null;  
    } else if (node.data > data)  
        node.left = remove(node.left, data)  
    else {  
        node.right = remove(node.right, data)  
    }
```



## # Remove node in BST

```

public static Node remove(Node node, int data) {
    if(node == null){
        // if node does not exist
        return null;
    }

    if(node.data == data){

        if(node.left == null && node.right == null){
            // leaf node
            return null;
        } else if(node.left == null){
            // only right child
            return node.right;
        } else if(node.right == null){
            // only left child
            return node.left;
        } else {
            // 2 child
            int inorderPredecessor = max(node.left);
            node.data = inorderPredecessor;
            node.left = remove(node.left, inorderPredecessor);
        }
    } else if(node.data < data)
        node.right = remove(node.right, data);
    else
        node.left = remove(node.left, data);

    return node;
}

```

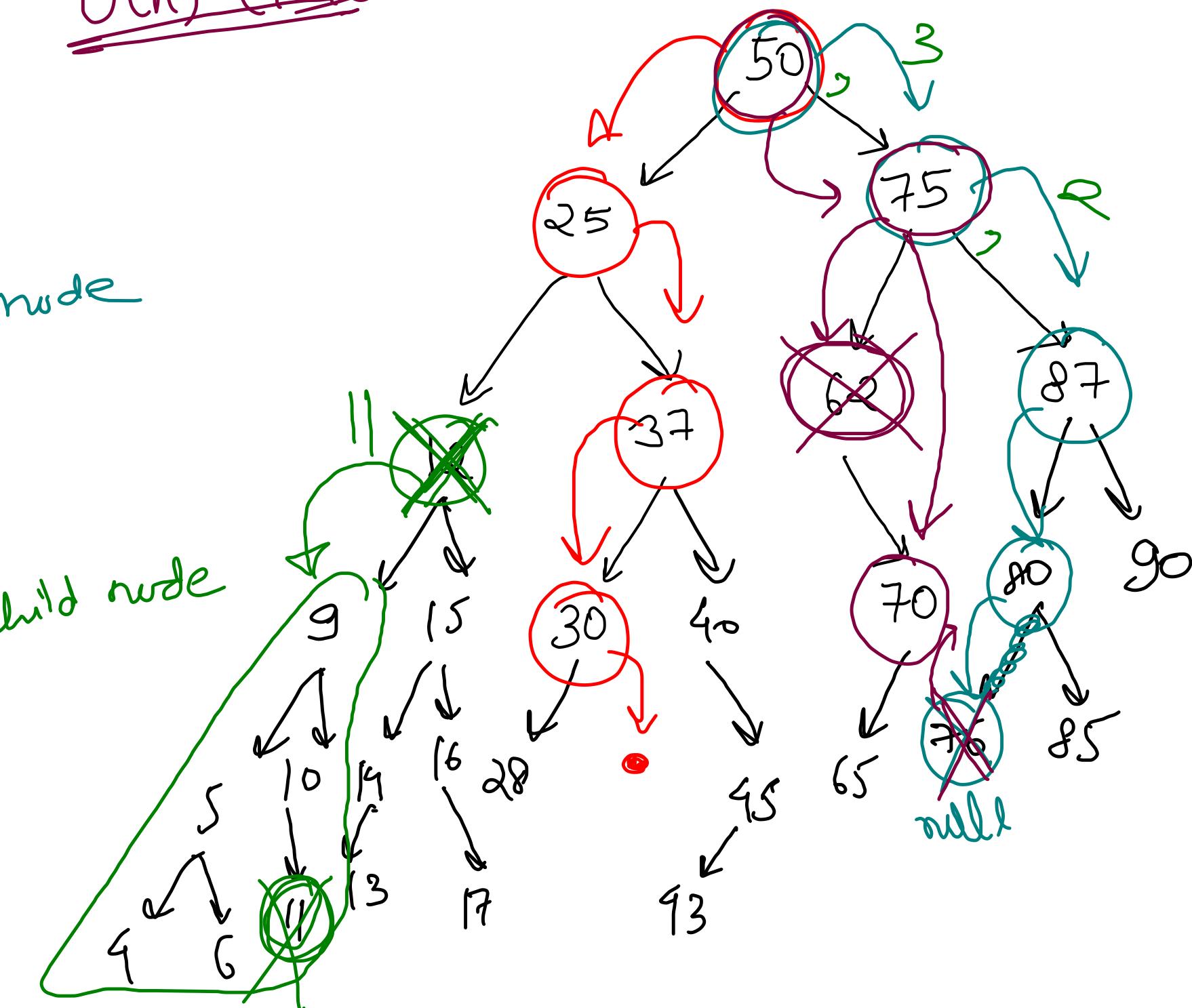
does not exist

leaf node

single child

2 child node

O(lg n) time



Replace with sum of larger

Reverse  
Inorder

87 75 70 62 50 - - -  
87 + 75 + 70 + 62  
+ + + +  
75 70 62

leetcode variation

```
class Solution {
    static int sum = 0;
    public TreeNode bstToGst(TreeNode root) {
        sum = 0;
        return bstToGstHelper(root);
    }

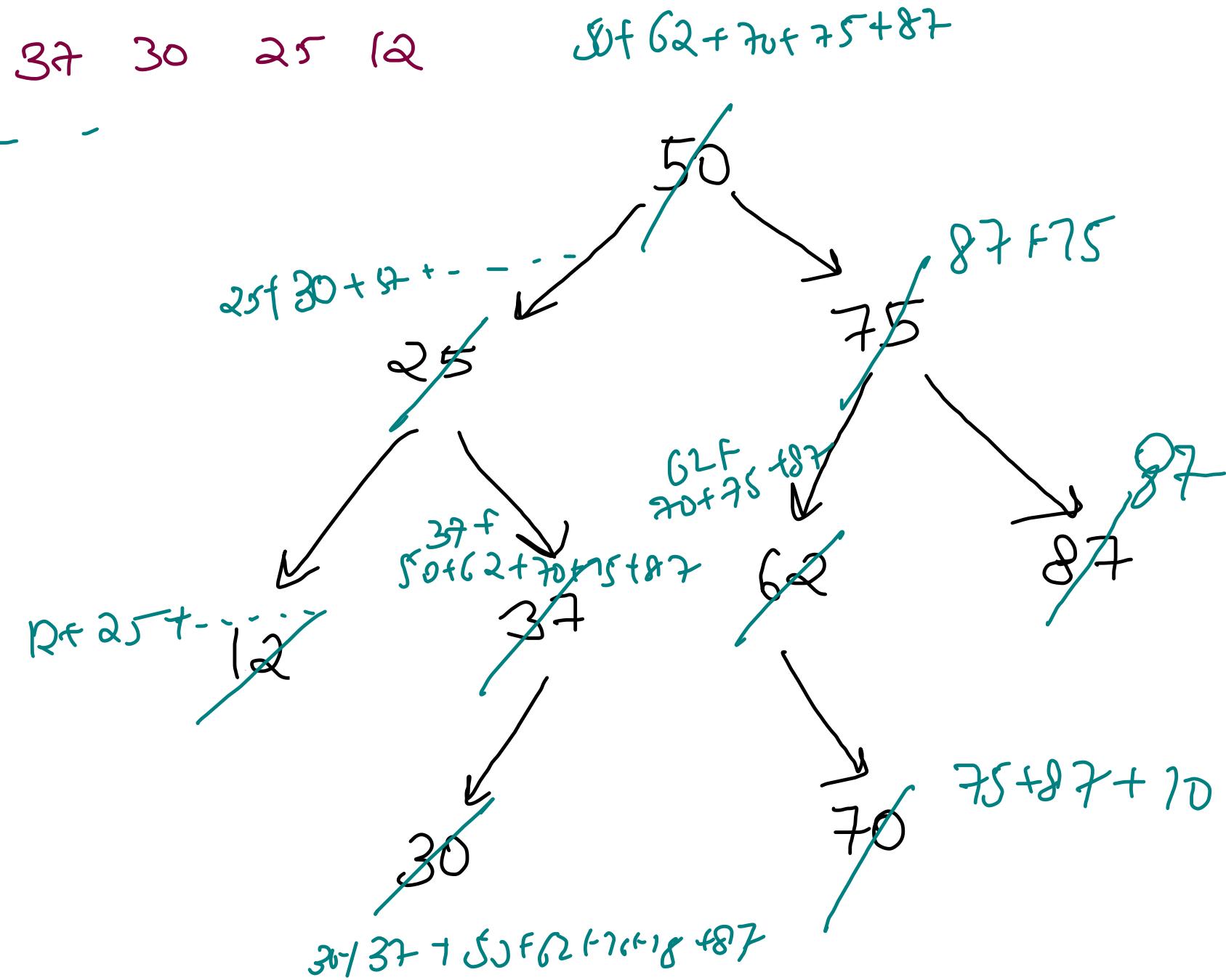
    public TreeNode bstToGstHelper(TreeNode root){
        if(root == null) return null;

        root.right = bstToGstHelper(root.right);

        sum += root.val;
        root.val = sum;

        root.left = bstToGstHelper(root.left);

        return root;
    }
}
```



LCA of BST

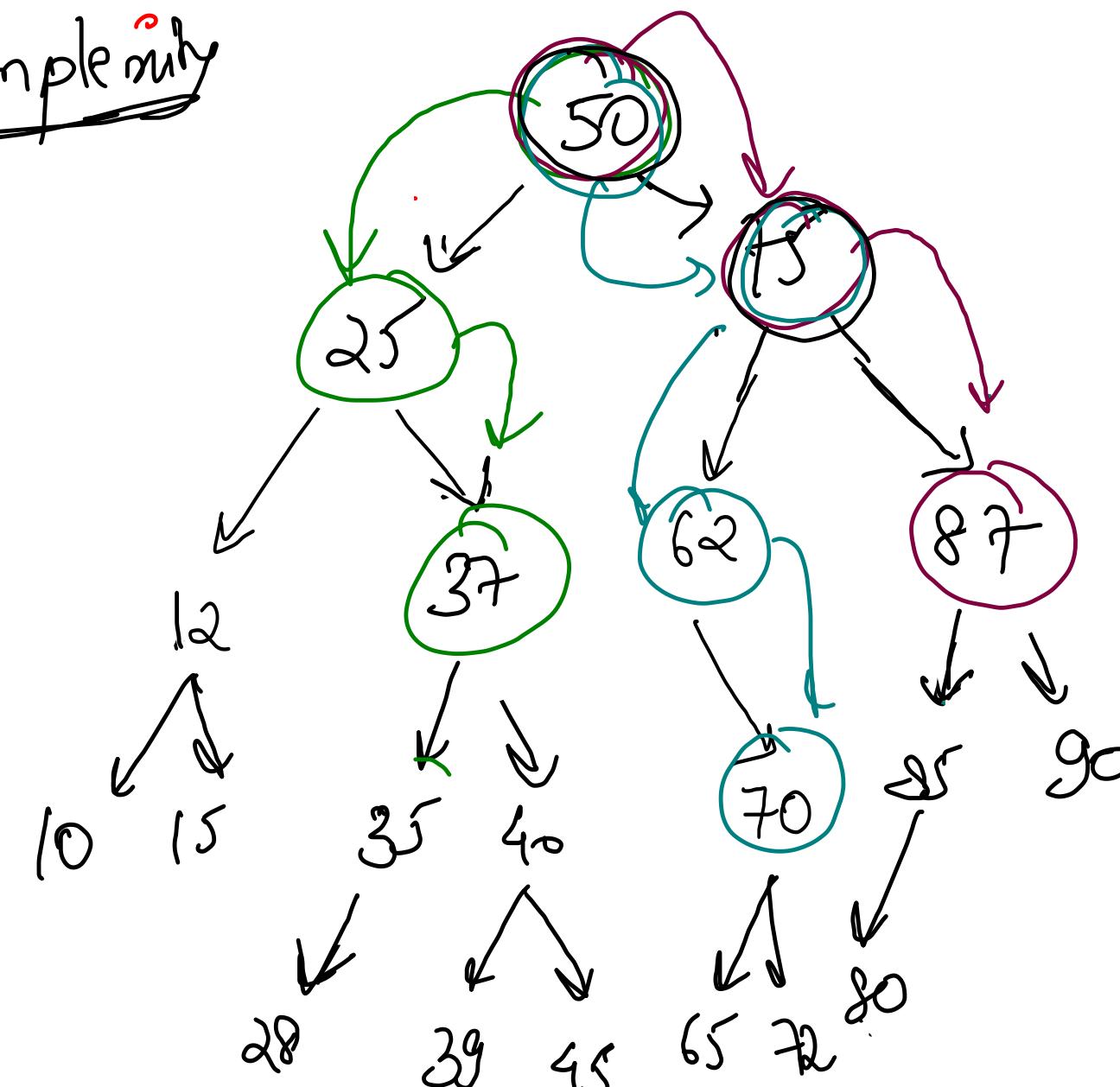
$d_1 < d_2$

\* It will fail if any one node does not exist

```
public static int lca(Node node, int d1, int d2) {
    if(node.data == d1 || node.data == d2){
        return node.data;
    } else if(d1 < node.data && d2 < node.data){
        return lca(node.left, d1, d2);
    } else if(d1 > node.data && d2 > node.data){
        return lca(node.right, d1, d2);
    } else {
        //  $d_1 < node.data < d_2$  or  $d_2 < node.data < d_1$ 
        return node.data;
    }
}
```

$\text{LCA}(a, b) = \text{LCA}(b, a)$  : if  $d_2 > d_1$  swap( $d_1, d_2$ )

$O(\log n)$  time complexity



$$\textcircled{1} \ LCA(35 \ \& \ 40) = 37$$

$$\textcircled{2} \ LCA(87 \ \& \ 80) = 87$$

$$\textcircled{3} \ LCA(87 \ \& \ 70) = 70$$

## Point im Range



37, 50, 62, 70, 75

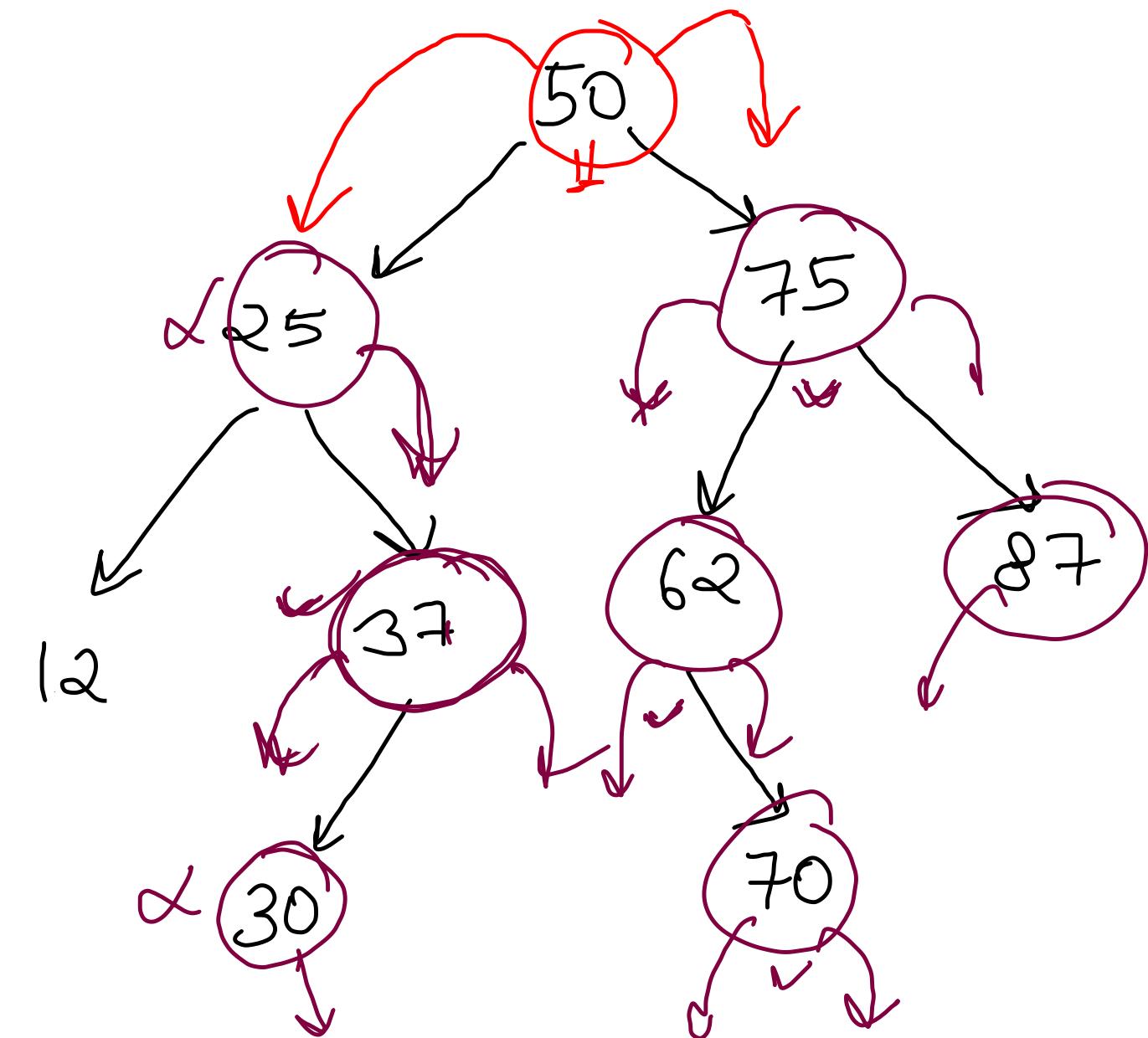
```
public static void pir(Node node, int d1, int d2) {
    if(node == null) return;

    if(!(node.data < d1))
        pir(node.left, d1, d2);

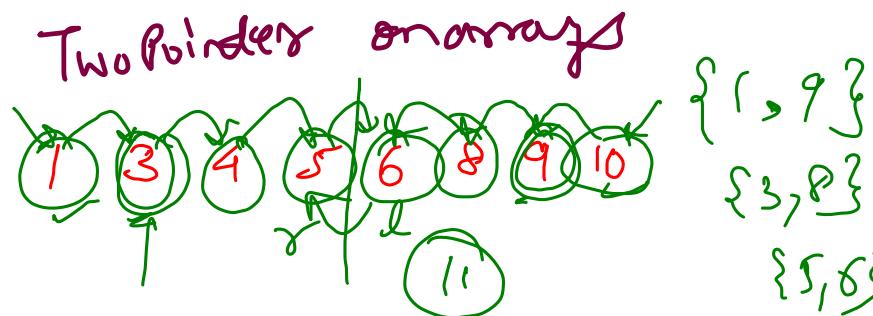
    if(node.data >= d1 && node.data <= d2)
        System.out.println(node.data);

    if(!(node.data > d2))
        pir(node.right, d1, d2);
}
```

Best case:  $O(\log n)$   
Worst case:  $O(n)$



Target sum pair



$\{2, 3, 4, 5, 7, 8, 9, 10, 12, 13, 15, 16, 18, 20\}$

① Approach 1:

Time :  $O(n \times h)$

Space :  $O(h)$  recursion

② Approach 2:

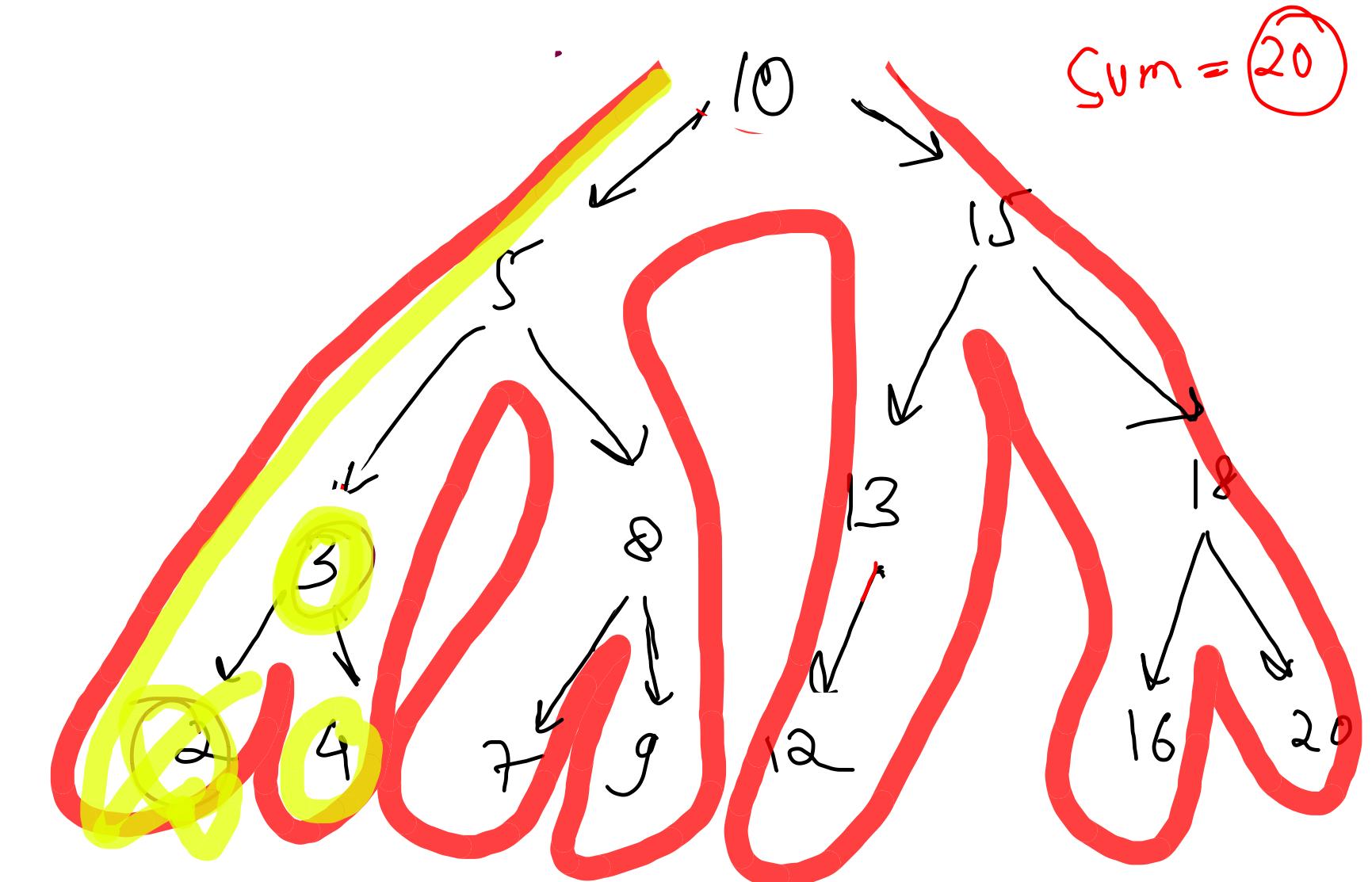
Time :  $O(n)$

Space :  $O(n)$  extra space

Approach 3:

Time :  $\rightarrow O(n)$

Space :  $\rightarrow O(h)$  stack



```

public static boolean find(Node curr, int target){
    if(curr == null) return false;
    if(curr.data == target) return true;
    if(curr.data < target) return find(curr.right, target);
    return find(curr.left, target);
}

public static void approach1Helper(Node curr, int target, Node root){
    if(curr == null) return;
    approach1Helper(curr.left, target, root);
    if(curr.data >= target/2) return;
    int complement = target - curr.data;
    if(find(root, complement) == true){
        System.out.println(curr.data + " " + complement);
    }
    approach1Helper(curr.right, target, root);
}
public static void approach1(Node root, int target){
    approach1Helper(root, target, root);
}

```

## Approach 1

For each node,  
~~find~~ complement.  
Also node < sum/2

Time :  $\rightarrow O(n \times h)$

Space :  $O(h)$

Recursion

```

public static void inorder(Node root, ArrayList<Integer> arr){
    if(root == null) return;

    inorder(root.left, arr);
    arr.add(root.data);
    inorder(root.right, arr);
}

public static void approach2(Node root, int target){
    ArrayList<Integer> arr = new ArrayList<>(); // O(n) space
    inorder(root, arr); // insert elements in sorted order in arraylist - O(n)

    // Two Pointer on arraylist - O(n)
    int left = 0, right = arr.size() - 1;

    while(left < right){
        int sum = arr.get(left) + arr.get(right);
        if(sum == target){
            System.out.println(arr.get(left) + " " + arr.get(right));
            left++; right--;
        } else if(sum < target){
            left++;
        } else {
            right--;
        }
    }
}

```

## Approach 2

- Store inorder in array
- Apply two pointer on array

Time  $\Rightarrow O(n)$  {inorder + 2 pointers}

Space  $\Rightarrow O(n)$  Extra Space  
(array)

Time:  $\rightarrow O(n)$   
Space

```
public static void approach3(Node root, int target){
    Stack<Pair> inorder = new Stack<>();
    inorder.push(new Pair(root, -1));

    Stack<Pair> reverseInorder = new Stack<>();
    reverseInorder.push(new Pair(root, -1));

    int left = iterativeInorder(inorder);
    int right = iterativeReverseInorder(reverseInorder);

    while(left < right){

        if(left + right == target){
            System.out.println(left + " " + right);
            left = iterativeInorder(inorder);
            right = iterativeReverseInorder(reverseInorder);
        } else if(left + right < target){
            left = iterativeInorder(inorder);
        } else {
            right = iterativeReverseInorder(reverseInorder);
        }
    }
}
```

