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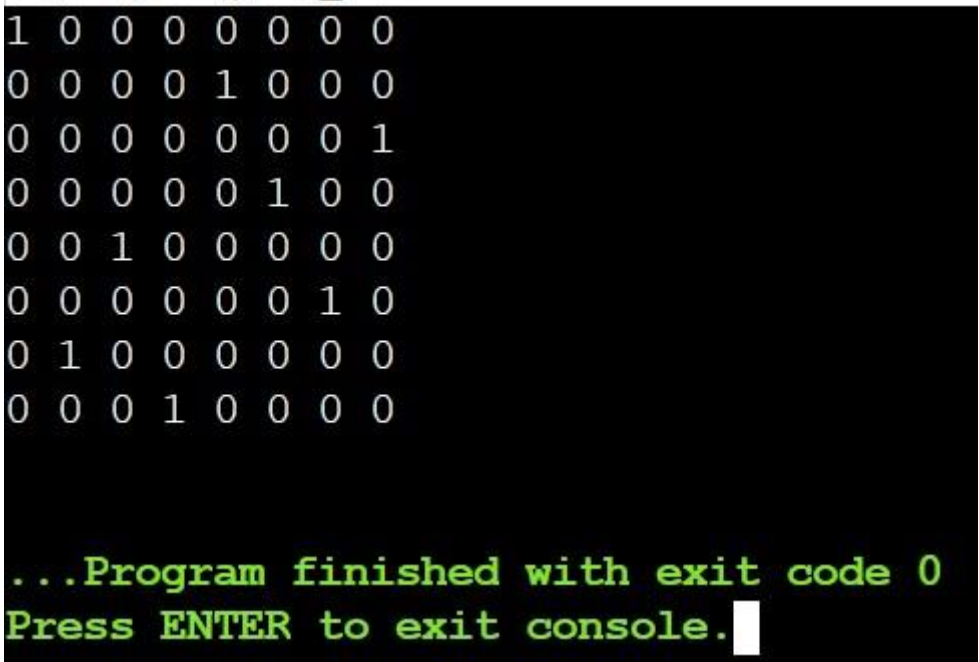
Department of Master of Computer Application

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| Experiment | 8 |
| Aim | To implement Backtracking algorithm (8 queens problem) |
| Objective | 1) Understand the problem of placing 8 queens on an 8x8 chessboard such that no two queens threaten each other (8 queens problem). 2) Understand the Time complexity of N queen problem |
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| Algorithm and Explanation of the technique Used | <pre>void placeQueens(vector<vector<int>>& board, int row, int col) { if (row == N) { printSolution(board); return; } for (int c = 0; c < N; ++c) { if (isSafe(board, row, c)) { board[row][c] = 1; placeQueens(board, row + 1, c); board[row][c] = 0; } } }</pre> <pre>void solveNQueens() { vector<vector<int>> board(N, vector<int>(N, 0)); placeQueens(board, 0, 0); }</pre> <p>The time complexity of the backtracking solution for the 8 queens problem is $O(N!)$.</p> |
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**Program
Code**

```
public class NQueens {  
  
    static final int N = 8;  
    static boolean isSafe(int[][] board, int  
row, int col) {  
        for (int i = 0; i < row; ++i)  
            if (board[i][col] == 1)  
                return false;  
        for (int i = row, j = col; i >= 0 && j  
>= 0; --i, --j)  
            if (board[i][j] == 1)  
                return false;  
        for (int i = row, j = col; i >= 0 && j <  
N; --i, ++j)  
            if (board[i][j] == 1)  
                return false;  
        return true;  
    }  
  
    static boolean solveNQueens(int[][] board,  
int row) {  
        if (row == N)  
            return true;  
        for (int col = 0; col < N; ++col) {  
            if (isSafe(board, row, col)) {  
                board[row][col] = 1;  
                if (solveNQueens(board, row +  
1))  
                    return true;  
                board[row][col] = 0;  
            }  
        }  
        return false;  
    }  
  
    static void printSolution(int[][] board) {  
        for (int i = 0; i < N; ++i) {  
            for (int j = 0; j < N; ++j)  
                System.out.print(board[i][j] + "  
");  
            System.out.println();  
        }  
    }  
  
    public static void main(String[] args) {  
        int[][] board = new int[N][N];  
        if (solveNQueens(board, 0))  
            printSolution(board);  
        else  
            System.out.println("Solution does  
not exist.");  
    }  
}
```

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| | <code>}</code> |
| Output |  <pre>1 0 0 0 0 0 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 1 0 0 0 0 0 1 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 1 0 0 1 0 0 0 0 0 0 0 0 0 1 0 0 0 0 ...Program finished with exit code 0 Press ENTER to exit console.</pre> |
| Justification of the complexity calculated | <p>The time complexity of the backtracking solution for the 8 queens problem is $O(N!)$.</p> <p>In the worst-case scenario, the backtracking algorithm explores all possible permutations of placing 8 queens on an 8x8 chessboard, which is $N!$ ($8!$ in this case). This is because for the first row, you have 8 choices, for the second row you have 7 choices (one queen has been placed in the first row), for the third row you have 6 choices, and so on, resulting in a total of $8! = 40320$ possible permutations to check.</p> <p>Therefore, the time complexity of the backtracking solution for the 8 queens problem is $O(N!)$, where N is the size of the chessboard.</p> |

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| Conclusion | <p>The backtracking algorithm for the 8 queens problem offers significant advantages in solving combinatorial optimization challenges, particularly in scenarios where exhaustive search is feasible and efficient. By systematically exploring the solution space and leveraging constraints to prune branches, the algorithm efficiently identifies valid arrangements of queens on the chessboard, ensuring that no two queens threaten each other. This approach finds application in various fields, such as puzzle-solving, constraint satisfaction, and combinatorial optimization tasks, where it enables the rapid identification of solutions while minimizing computational complexity.</p> |
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