

For each of these statements, say if it is true or false. Give a proof or provide a counterexample for your answer.

(In the actual exam, you will be provided with a short list of undecidable/NP-hard languages.)

- (1) There is a 2-state NFA for the language  $(01)^*$ .
- (2) There is a 2-state DFA for the language  $(01)^*$ .
- (3) If  $L$  is regular over  $\Sigma = \{0, 1\}$ , then  $L'$  is also regular, where

$$L' = \{x \mid x \in L \text{ and } x \text{ starts and ends with the same symbol}\}.$$

- (4) The language  $L = \{wxw^R x^R : x, w \in \Sigma^*\}$  is context-free over alphabet  $\Sigma = \{a, b\}$ .
- (5) If  $L_1$  and  $L_2$  are regular languages, then the following language is context-free:

$$L = \{xy \mid x \in L_1, y \in L_2, \text{ and } |x| = |y|\}.$$

- (6) The grammar  $S \rightarrow Sa \mid a$  is  $LR(0)$ .
- (7) The following language is decidable:

$$L = \{\langle R \rangle \mid \text{Regular expression } R \text{ generates at least one string of even length.}\}$$

- (8) The following language is decidable:

$$L = \{\langle G_1, G_2 \rangle \mid \text{CFG } G_2 \text{ generates some string that CFG } G_1 \text{ does not generate.}\}$$

- (9) The following language is decidable:

$$L = \{\langle M, k \rangle \mid \text{TM } M \text{ accepts at most } k \text{ inputs.}\}$$

- (10) The following language is NP-complete:

$$L = \{\langle \phi, \psi \rangle \mid \text{Boolean formulas } \phi \text{ and } \psi \text{ share a common satisfying assignment.}\}$$