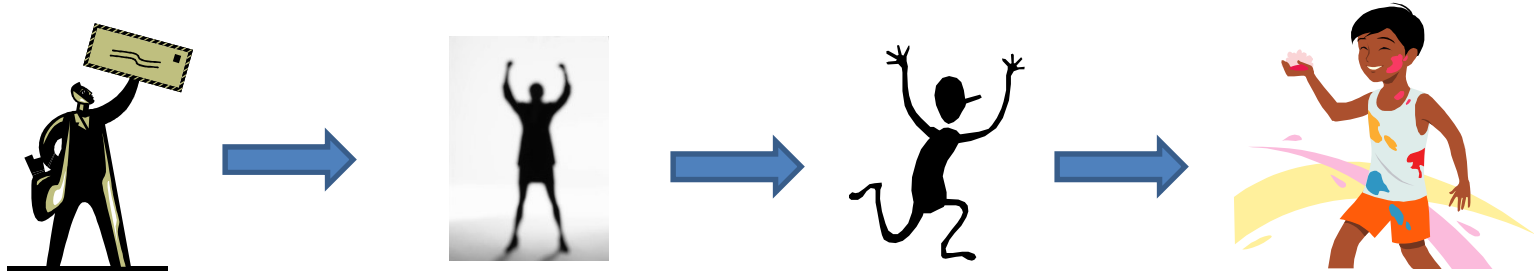




Small World Phenomenon



Six Degrees of Separation



- Milgram asked randomly chosen “starter” individuals to each try **forwarding a letter to a designated “target”** person.
- The target’s name, address, occupation and some personal information are provide.
- Each participant could only advance the letter by forwarding it to a single acquaintance that he or she knew.
- A third of the letters arrived at the target.
- In a median of **six steps**.



- Killworth and Bernard studied the strategies that people employ for choosing how to forward a message toward a target.
 - a mixture of primarily **geographic** and **occupational** features being used, with different features being favored depending on the **characteristics** of the target in relation to the sender.



Existence of short paths



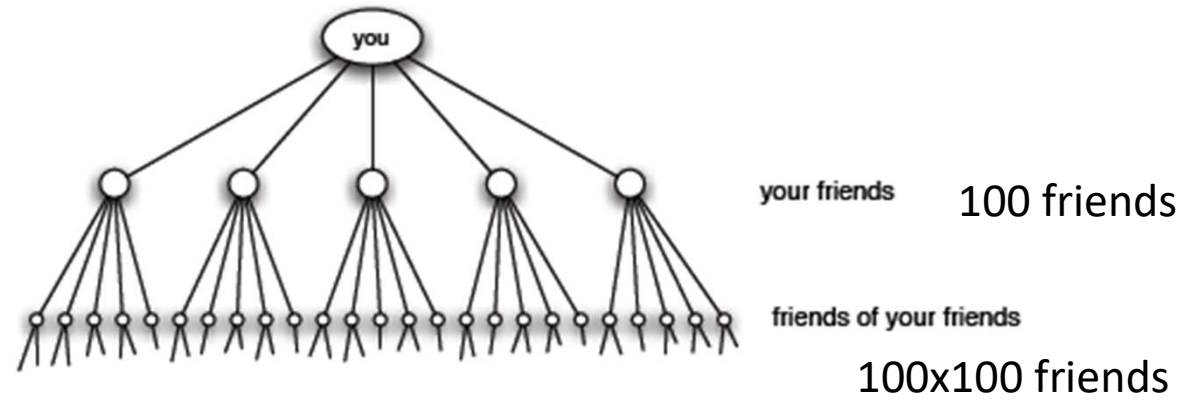


- Two striking facts
 - Short paths are there in abundance.
 - People, acting without any sort of global “map” of the network, are effective at collectively finding these short paths.
- Do you expect the same happen on social networks ?

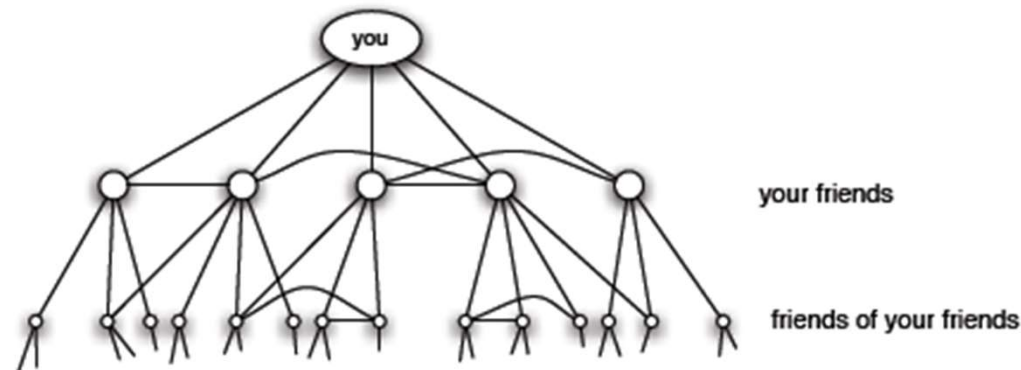


Structure and Randomness

Each friend is
linked to 100
new people



Triadic closure



Highly clustered individuals, so
why do shortest paths exist ?



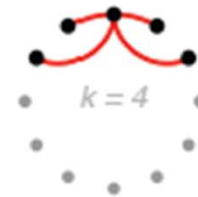
Watts-Strogatz model (Ring)

- Two types of links
 - Homophily
 - We connect to others who are like ourselves.
 - Many triangles
 - Weak ties
 - The links to acquaintances that connect us to parts of the network that would otherwise be far apart.

We start with a ring of n vertices



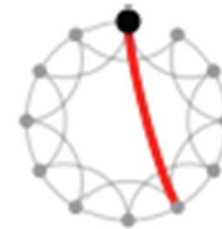
where each vertex is connected to its k nearest neighbors



like so.



With probability p , we reconnect this edge to a vertex chosen uniformly at random over the entire ring, with duplicate edges forbidden. Otherwise, we leave the edge in place.





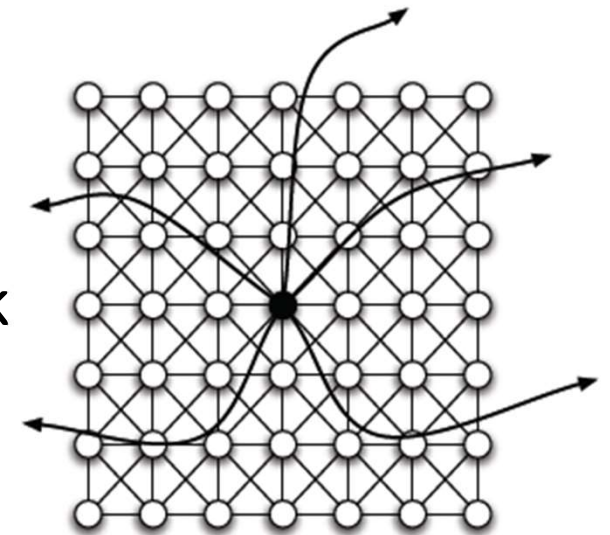
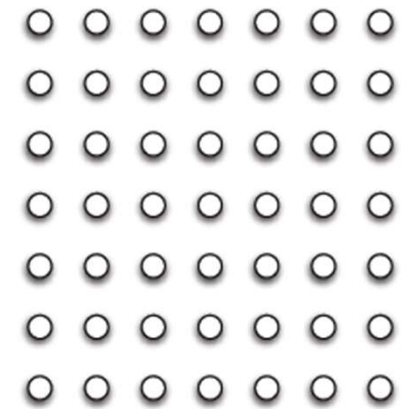
Watts-Strogatz model (Grid)

Everyone lives on a two-dimensional grid

- Two nodes are one grid-step apart if they are directly adjacent to each other.

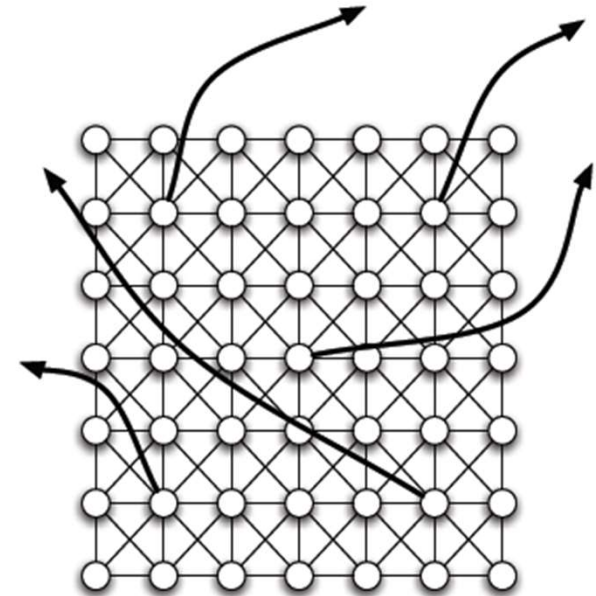
Two types of links

- **Homophily** : each node form a link to all other nodes that lie within a radius of up to **r** grid steps away.
- **Weak ties** : each node also forms a link to **k** other nodes selected uniformly at random from the grid (very far apart).





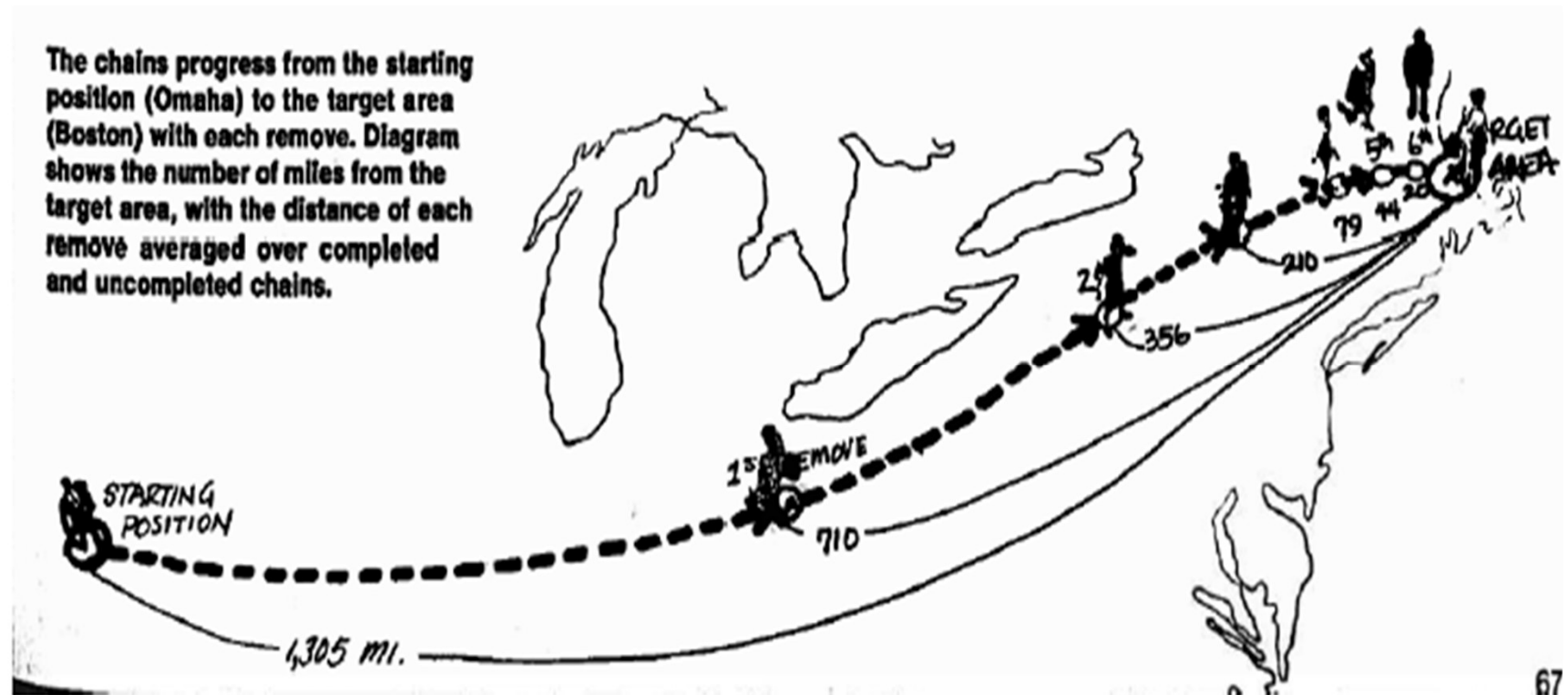
- The combinations
 - weak ties
 - reach many people in a small no. steps
 - mainly homophilous links and few weak ties
 - many triangles
 - If only one out of every k nodes is allowed to have a single random friend, this is enough to make the world small.





Decentralized Search

- Milgram small-world experiment

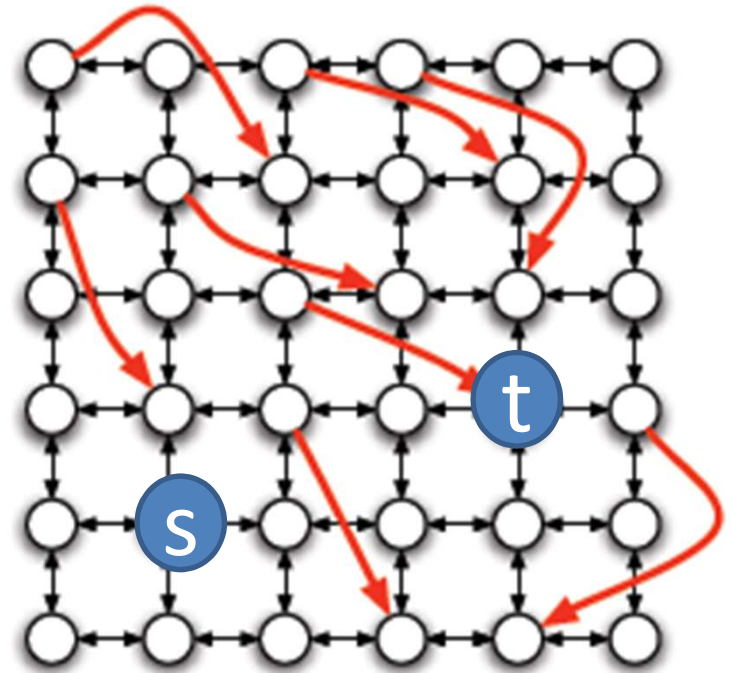






A model for decentralized search

- Starting with the **grid-based model** of Watts and Strogatz
- Node **s** has to forward a message to node **t** passing it along edges of the network.
- **s** knows the **location of t** but **s** does not know the **random edges** out of any nodes other than itself.
- Each intermediate nodes has this partial information as well.





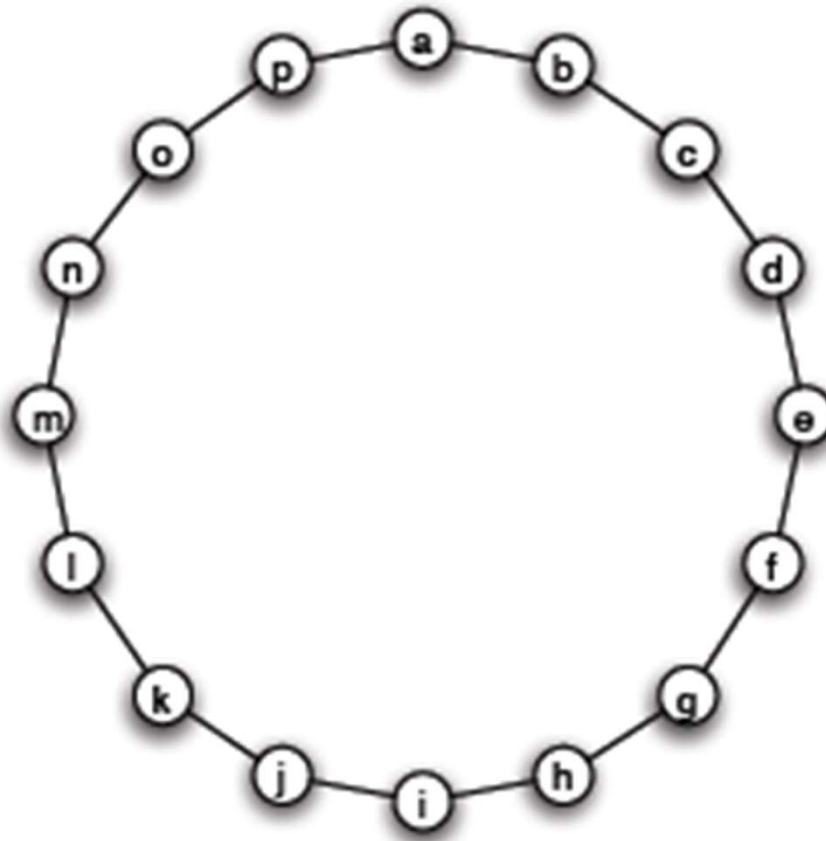
A model for decentralized search

- We will evaluate different search procedures according to their delivery time.
- Require a large number of steps to reach a target — much larger than the true length of the shortest path.



Analysis of Decentralized Search

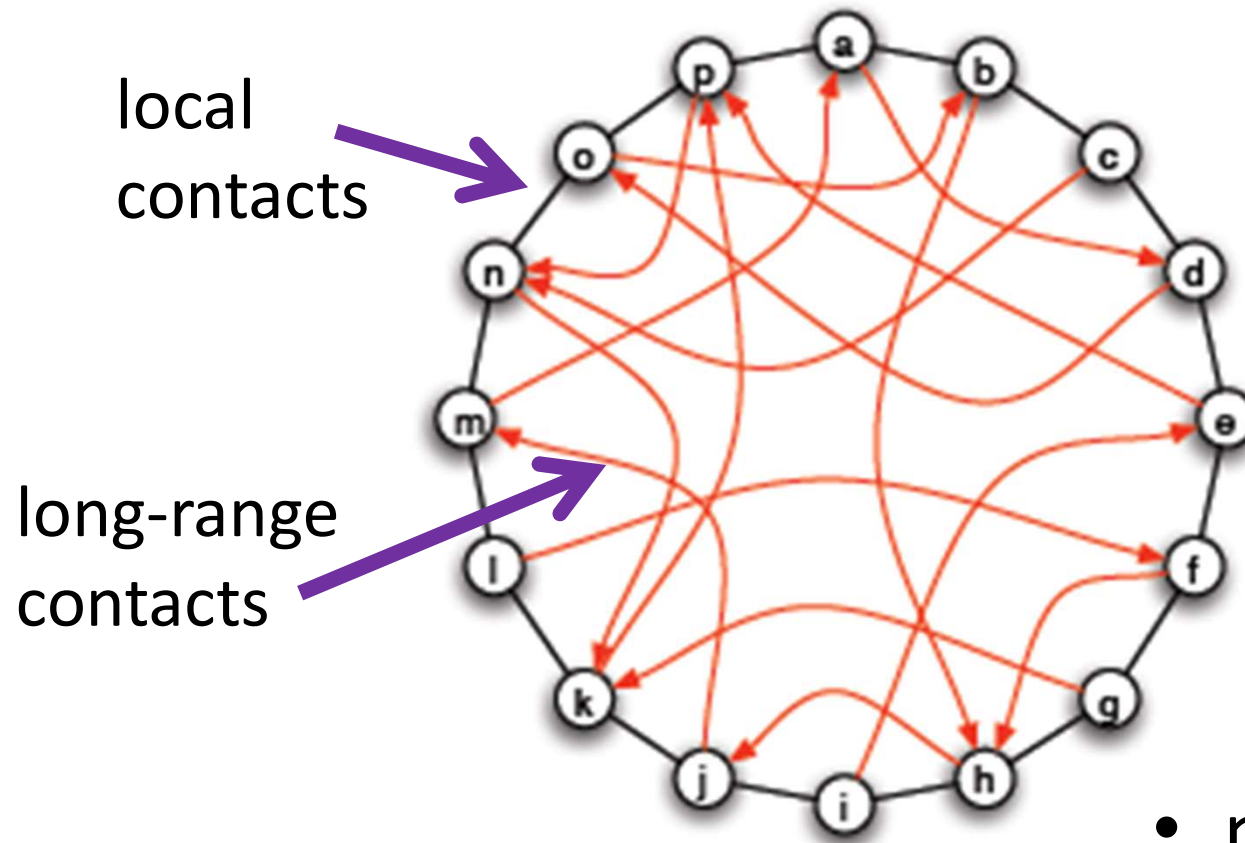
- A set of n nodes arranging on a 1D ring



- directed edges to two adjacent nodes



- A set of n nodes arranging on a 1D ring

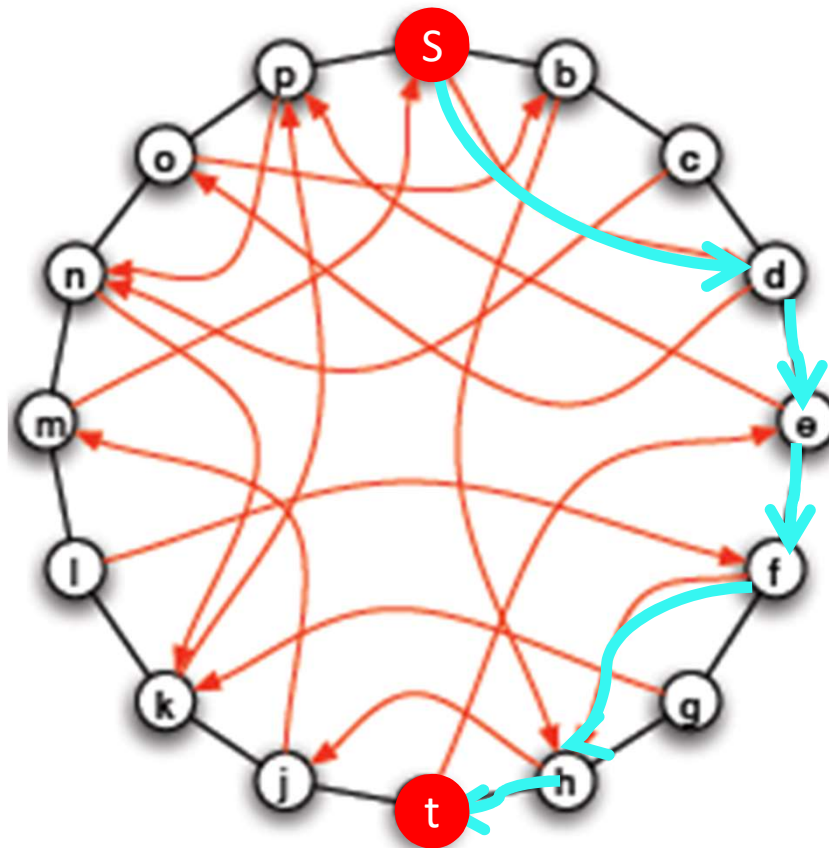


- directed edges to two adjacent nodes
- a single directed edge to some other node
- $\text{prob} \propto d(v, w)^{-1}$



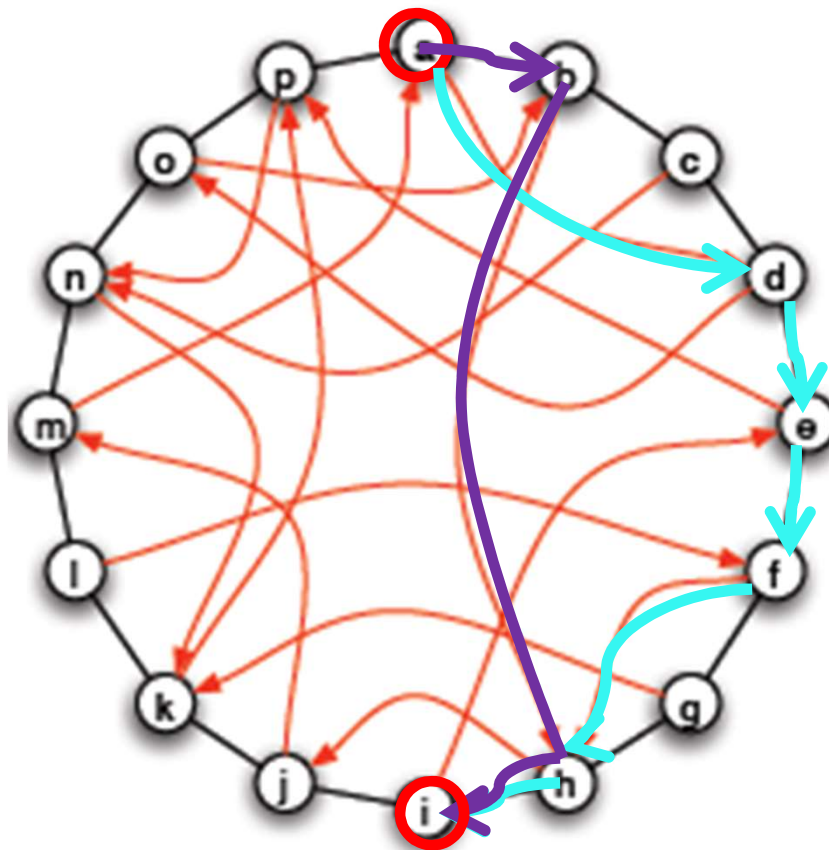
Myopic Search

- start node $s \rightarrow$ target node t





the shortest path from a to i



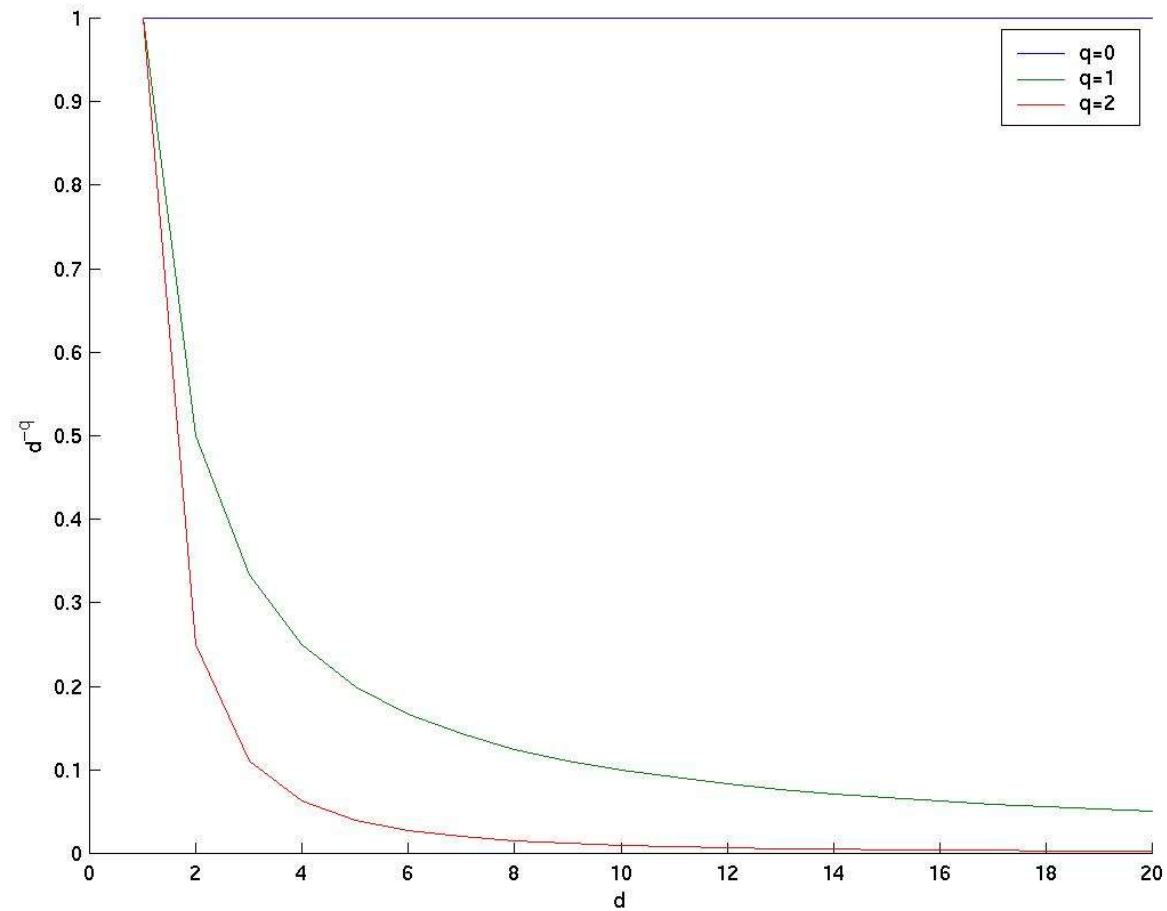


Generalizing the network model

- Nodes on a grid and each node has edges to each other node within **r** grid steps.
- Each of its **k** random edges is generated in a way that decays with distance
 - Let $d(v, w)$ denote the number of grid steps between nodes v and w .
 - Probability of an edge is proportional to $d(v, w)^{-q}$



Graph of d^{-q}

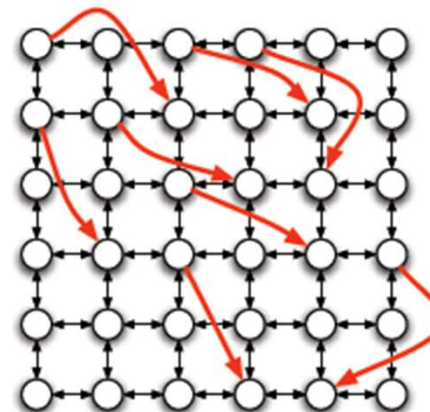
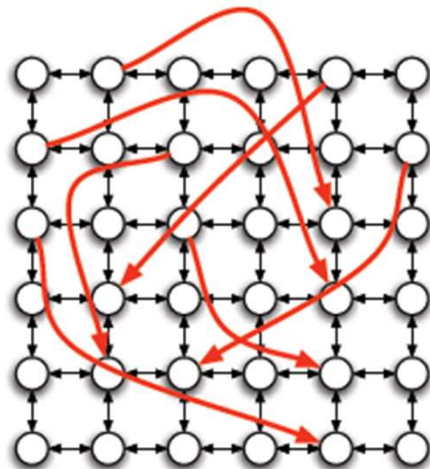




Generalizing the network model

- $q=0$, the links are chosen uniformly at random
- Small q , the long-range links are “too random”
- Large q , the long-range links are “not random enough,”

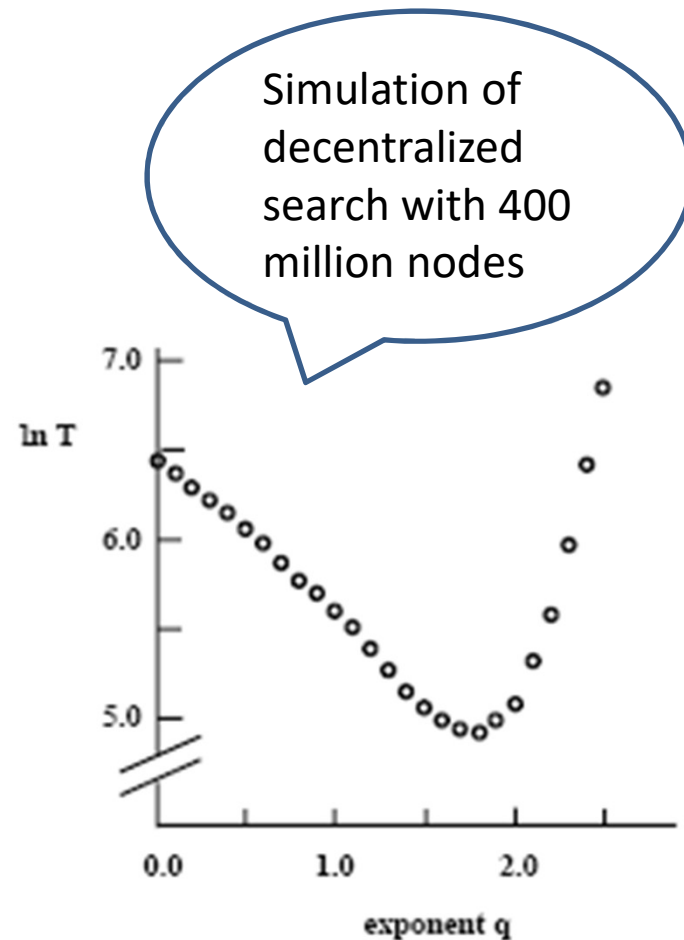
Small q ,
more long
distances
edges



Large q ,
more short
distances
edges



- Decentralized search is most efficient when $q = 2$ (so that random links follow an inverse-square distribution).
- Decentralized search has about the same efficiency on networks of hundreds million nodes across all exponents q between 1.5 and 2.0.





Efficiency of Decentralized Search

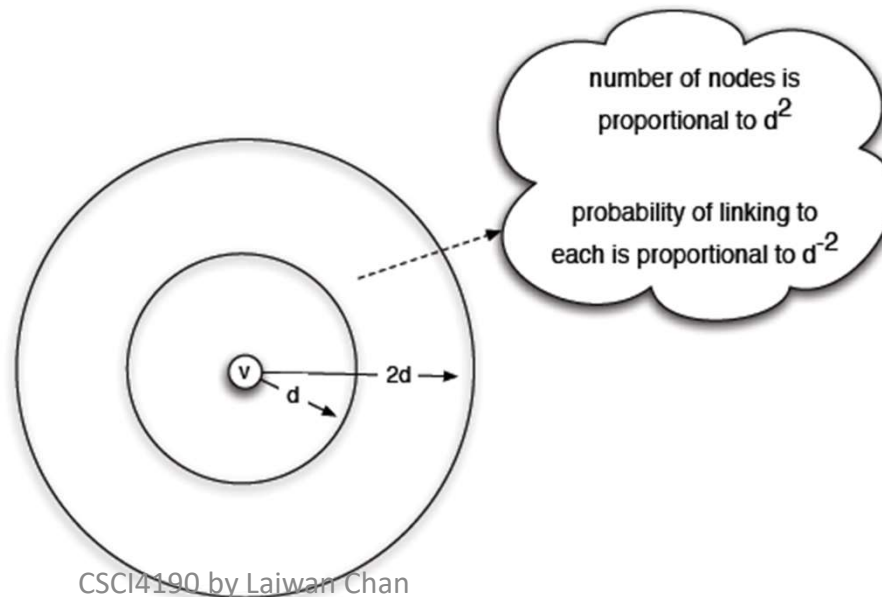
- Kleinberg 2000
 - $n \times n$ grid, $\alpha_0, \alpha_2, \alpha_r$ constants
 - Direct link to nodes within r grid steps away.
 - Prob. of a long range edge : $d(v, w)^{-q}$
- When $q=0$, the expected delivery time is at least $\alpha_0 n^{2/3}$
- When $0 \leq q < 2$, the expected delivery time is at least $\alpha_r n^{(2-q)/3}$
- When $q > 2$, the expected delivery time is at least $\alpha_r n^{(q-2)/(q-1)}$
- When $q = 2$ and $r = 1$, the expected delivery time is at most $\alpha_2 (\log n)^2$



A Rough Calculation Motivating the Inverse-Square Network

- A node v in the network
- Considering the group of nodes lying at distances between d and $2d$ from v
- The probability that a random edge links into any node in this group is approximately independent of the value of d .

When $q=2$,
same probability,
regardless of the distance





Geographic Data on Friendship (Liben-Nowell et al)

- The population density of the LiveJournal network, a blogging site with 500,000 users
- Users provided links to their friends.
- The population density of the users is extremely non-uniform.
- Distance based model
→ Rank-based model

