



Methodology of CS

Introduction to algorithms
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What is an algorithm?

How to use addition(+), subtraction(-), multiplication(*), and division(/)?

How to draw one line and visit all edges in a graph?

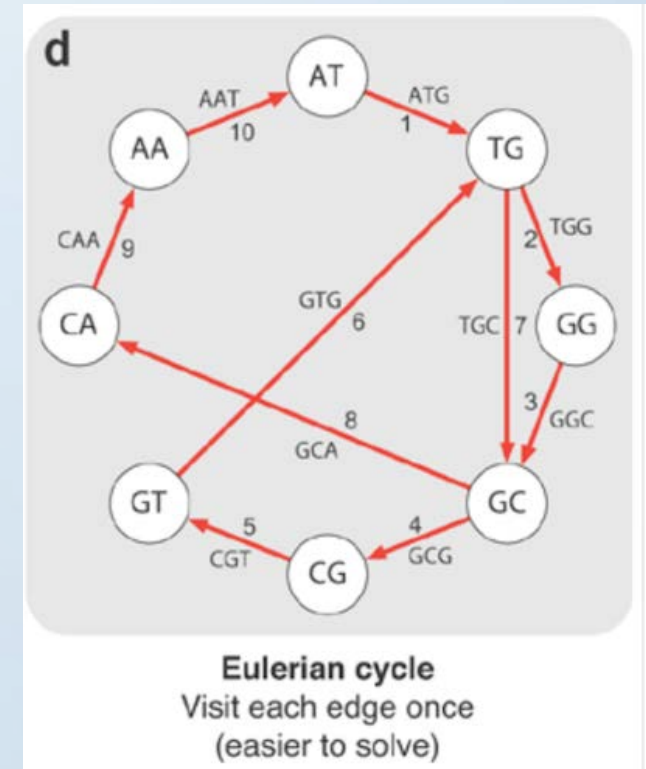
Handwritten polynomial long division showing steps 1 to 4:

① multiply: $2x-1$ multiplied by $7x^3 + x^2$ gives $14x^4 - 7x^3$.

② subtract: $14x^4 - 5x^3 - 11x^2 - 11x + 8$ minus $(14x^4 - 7x^3)$ gives $2x^3 - 11x^2 - 11x + 8$.

③: The next step is to subtract $2x^3 - 11x^2$ from the result.

④ this into this: The result is $-10x^2 - 11x + 8$.



What is an algorithm?

- a set of rules that precisely defines a sequence of operations

addition(+), subtraction(-),
multiplication(*), and division(/)

Euclid's Algorithm to obtain
greatest common divisor (gcd)

Sort Algorithms

Eulerian Path Algorithm

Shortest Path Algorithms
(Dijkstra, SPFA, Floyd)

How to design a delivery route
for a courier in S.F. Express?

How to play chess
with AlphaGo?

How to analyze
big data?

How to evaluate an algorithm?

Computational Complexity

Computational Complexity

Time Complexity

- Time needed for a sequence of operations

Rough
Estimation

Space Complexity

- Space needed for data structures used in a program

-
- Time Efficiency
 - Running time is always measured to evaluate an algorithm.
 - It also depends on CPU and other factors.

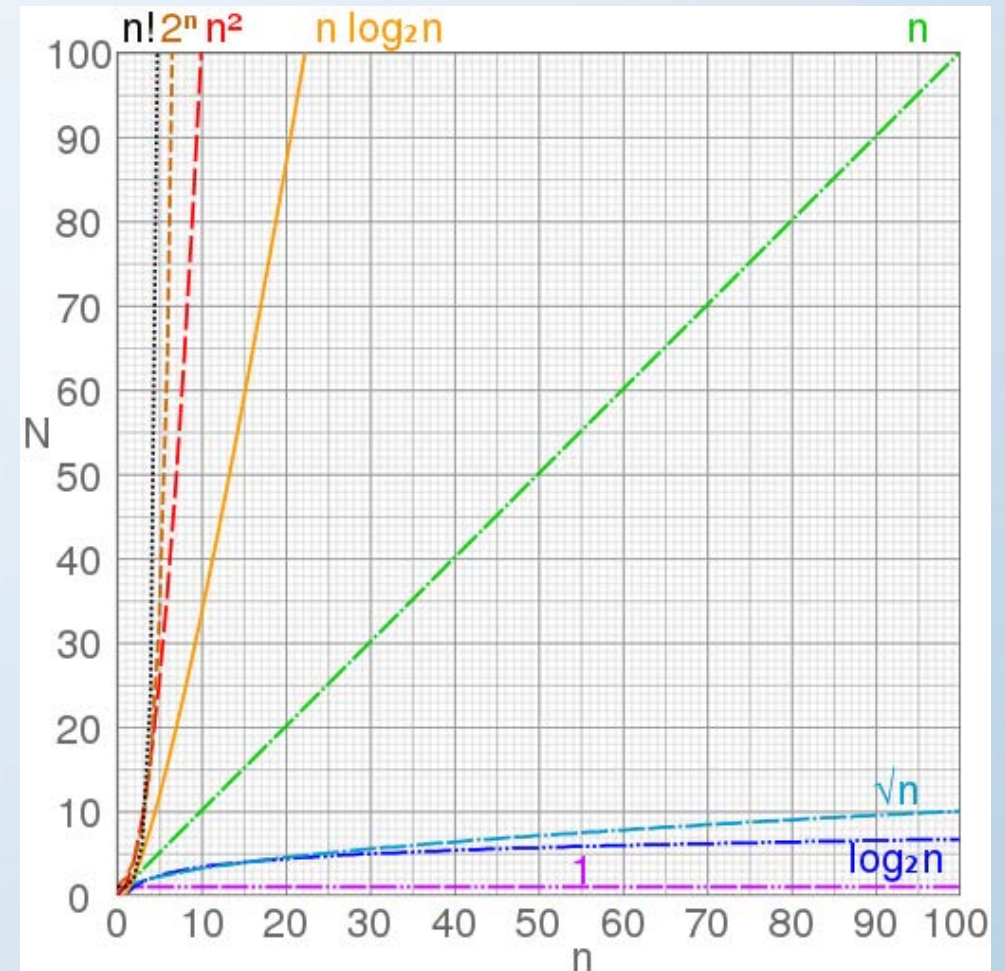
Specific
Evaluation

- Space Efficiency
- Memory is always measured to evaluate an algorithm.
- It also depends on programming language and other factors.

Computational Complexity

Big O Notation

- $O(1)$ Constant Complexity
- $O(\log_2 n)$ Logarithmic Complexity
- $O(n)$ Linear Complexity
- $O(n \log_2 n)$ Linearithmic Complexity
- $O(n^2)$ $O(n^3)$... Polynomial Complexity
- $O(a^n)$ Exponential Complexity
- $O(n!)$ Factorial Complexity



Big O Notation evaluates the magnitude of complexity of an algorithm based on input size n

Core Ideas

Coefficients $O(kn)$ ❌

Constant $O(n + c)$ ❌

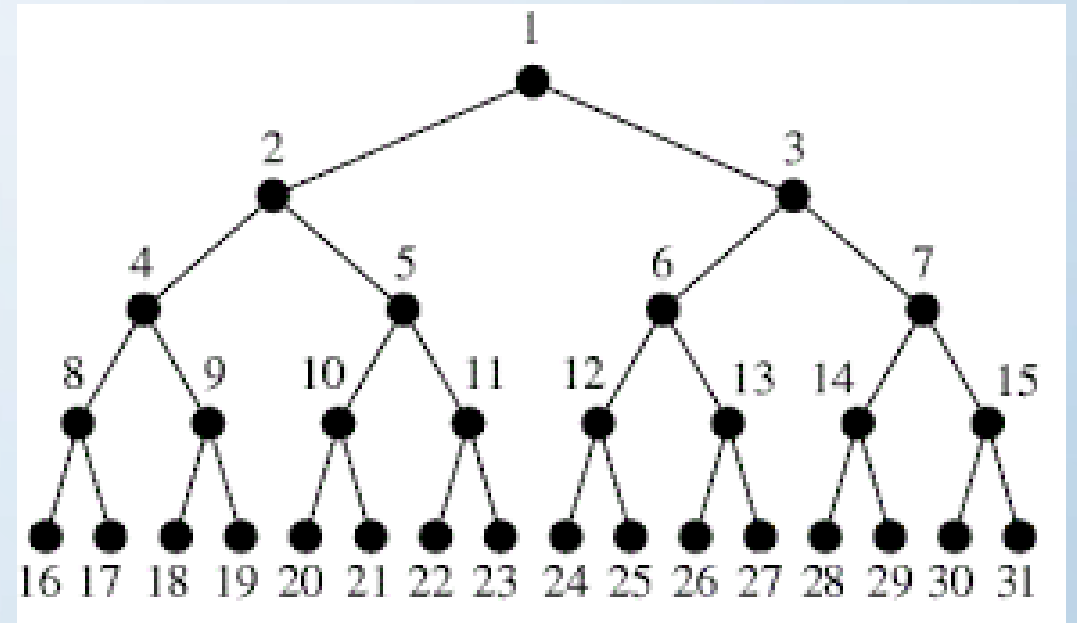


Some Examples

```
count = 0
for i in range(n):
    for j in range(i, n):
        count += 1
```

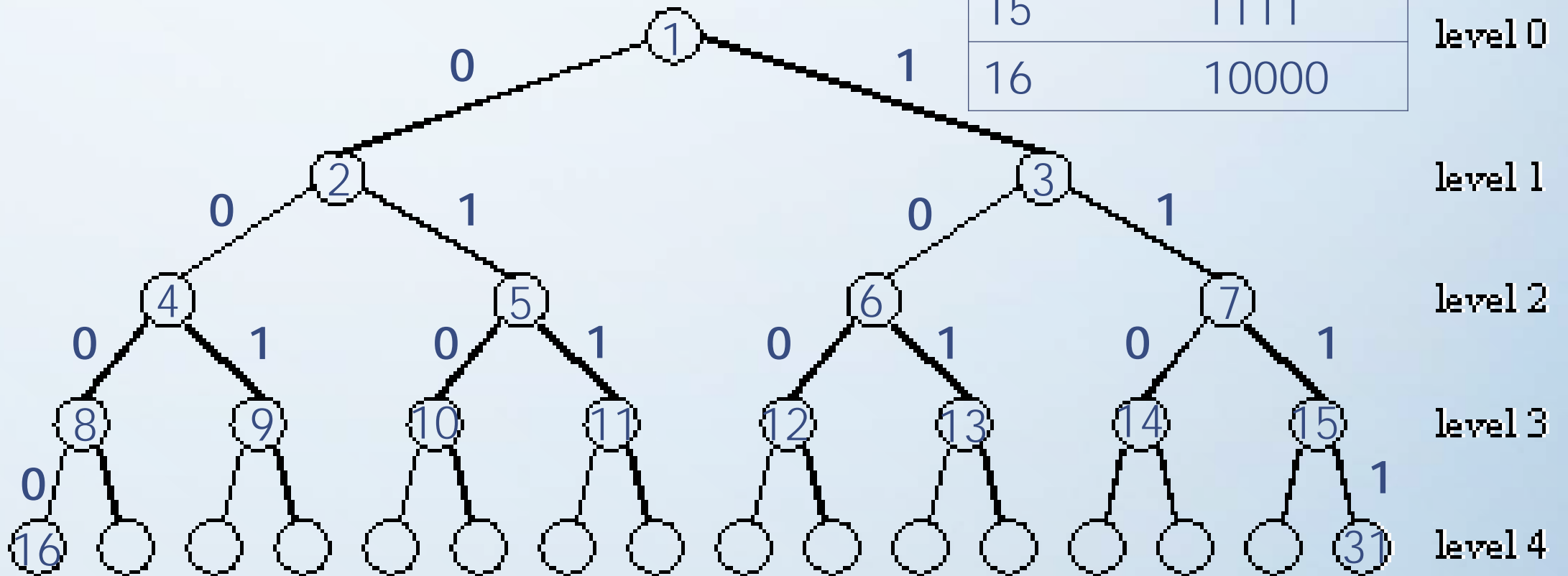
Time: $O(n^2)$
Space: $O(1)$

Look for a number in this perfect binary tree with n nodes from the root.



A Perfect Binary Tree

Decimal	Binary
5	101
14	1110
15	1111
16	10000



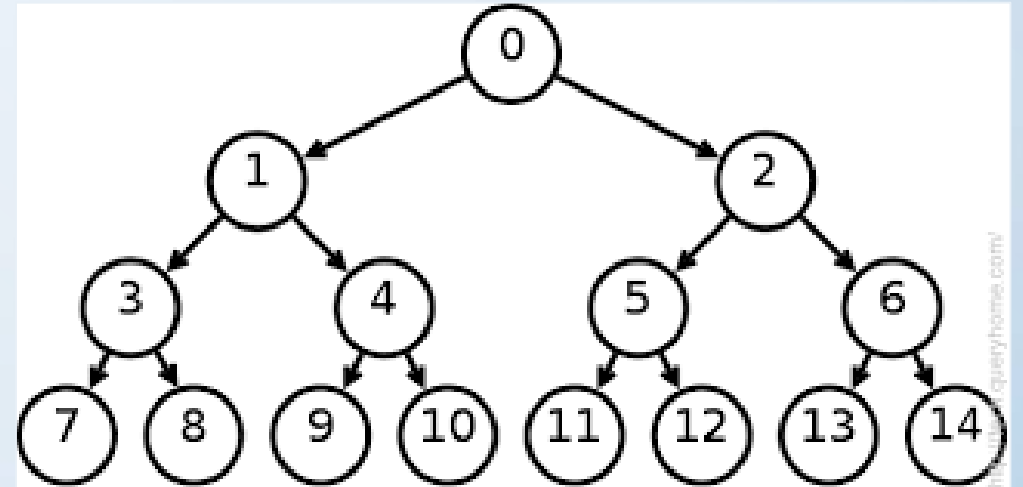
The depth of a perfect binary tree with n nodes is $\text{int}(\log_2 n)$

Some Examples

```
count = 0
for i in range(n):
    for j in range(i, n):
        count += 1
```

Time: $O(n^2)$
Space: $O(1)$

Look for a number in this perfect binary tree with n nodes from the root.



Time: $O(\log n)$
Space: $O(n \log n)$

A Task For You

Fibonacci Numbers



Algorithm A

```
f = [1, 1]
for i in range(2, n):
    f.append(f[i-1] + f[i-2])
print(f[n-1])
```

Time: $O(n)$
Space: $O(n)$

```
f = [1, 1]
for i in range(2, n):
    f[i%2] += f[(i+1)%2]
print(f[(n-1)%2])
```

Time: $O(n)$
Space: $O(1)$

Algorithm B

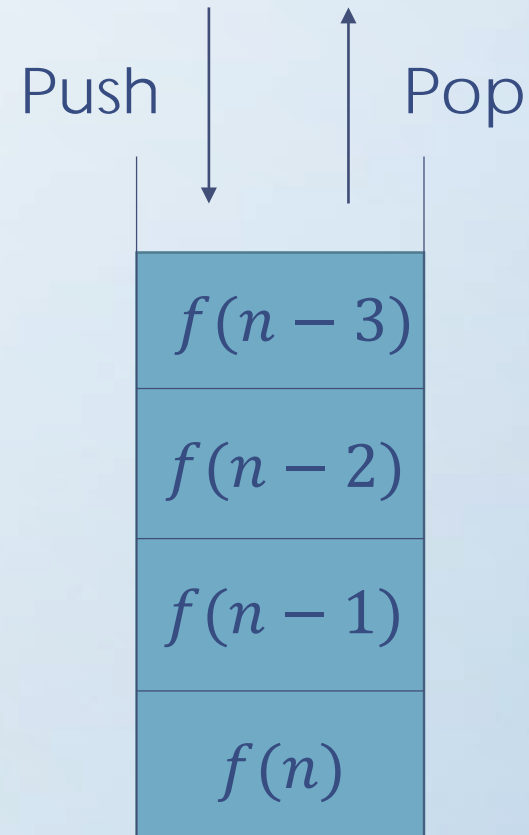
```
def f(n):  
    if n <= 2:  
        return 1  
    return f(n-1) + f(n-2)  
print(f(n))
```

Stack Overflow!

An exponential time algorithm

Time: $O(f(n))$

Space: $O(f(n))$



Algorithm C

```
sqrt5 = 5**0.5  
print(1/sqrt5*((1+sqrt5)/2)**n-  
      ((1-sqrt5)/2)**n))
```

Time: $O(1)$
Space: $O(1)$

$$f_n = \frac{1}{\sqrt{5}} \cdot \left[\left(\frac{1 + \sqrt{5}}{2} \right)^n - \left(\frac{1 - \sqrt{5}}{2} \right)^n \right]$$


```
In [3]: n = 100
```

```
In [4]: f = [1,1]
        for i in range(2,n):
            f.append(f[i-1]+f[i-2])
        print(f[n-1])
```

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```
In [5]: f = [1,1]
        for i in range(2,n):
            f[i%2] += f[(i+1)%2]
        print(f[(n-1)%2])
```

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```
In [6]: sqrt5 = 5**0.5
        print(int(1/sqrt5*((1+sqrt5)/2)**n-((1-sqrt5)/2)**n))
```

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The Best Algorithm?

```
f = [1,1]
for i in range(2,n):
    f[i%2] += f[(i+1)%2]
print(f[(n-1)%2])
```

Time: $O(n)$

Space: $O(1)$

A Better Algorithm

Fast Matrix Exponentiation

Matrix Multiplication &
Divide and Conquer

Time: $O(\log n)$
Space: $O(1)$

Thank you!

