

$$\begin{pmatrix} 5 & 6 & 3 \\ 0 & 0.8 & -0.6 \\ 0 & 0 & 2.5 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 8 \\ 1.4 \\ 1.5 \end{pmatrix},$$

thereby obtaining the desired answer

$$x = \begin{pmatrix} -1.4 \\ 2.2 \\ 0.6 \end{pmatrix}$$

by computing first x_3 , then x_2 , and finally x_1 .

Computing an LU decomposition

We have now shown that if we can create an LUP decomposition for a nonsingular matrix A , then forward and back substitution can solve the system $Ax = b$ of linear equations. Now we show how to efficiently compute an LUP decomposition for A . We start with the case in which A is an $n \times n$ nonsingular matrix and P is absent (or, equivalently, $P = I_n$). In this case, we factor $A = LU$. We call the two matrices L and U an **LU decomposition** of A .

We use a process known as **Gaussian elimination** to create an LU decomposition. We start by subtracting multiples of the first equation from the other equations in order to remove the first variable from those equations. Then, we subtract multiples of the second equation from the third and subsequent equations so that now the first and second variables are removed from them. We continue this process until the system that remains has an upper-triangular form—in fact, it is the matrix U . The matrix L is made up of the row multipliers that cause variables to be eliminated.

Our algorithm to implement this strategy is recursive. We wish to construct an LU decomposition for an $n \times n$ nonsingular matrix A . If $n = 1$, then we are done, since we can choose $L = I_1$ and $U = A$. For $n > 1$, we break A into four parts:

$$\begin{aligned} A &= \left(\begin{array}{c|ccc} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{array} \right) \\ &= \begin{pmatrix} a_{11} & w^T \\ v & A' \end{pmatrix}, \end{aligned}$$

where v is a column $(n - 1)$ -vector, w^T is a row $(n - 1)$ -vector, and A' is an $(n - 1) \times (n - 1)$ matrix. Then, using matrix algebra (verify the equations by

simply multiplying through), we can factor A as

$$\begin{aligned} A &= \begin{pmatrix} a_{11} & w^T \\ v & A' \end{pmatrix} \\ &= \begin{pmatrix} 1 & 0 \\ v/a_{11} & I_{n-1} \end{pmatrix} \begin{pmatrix} a_{11} & w^T \\ 0 & A' - vw^T/a_{11} \end{pmatrix}. \end{aligned} \quad (28.8)$$

The 0s in the first and second matrices of equation (28.8) are row and column $(n-1)$ -vectors, respectively. The term vw^T/a_{11} , formed by taking the outer product of v and w and dividing each element of the result by a_{11} , is an $(n-1) \times (n-1)$ matrix, which conforms in size to the matrix A' from which it is subtracted. The resulting $(n-1) \times (n-1)$ matrix

$$A' - vw^T/a_{11} \quad (28.9)$$

is called the **Schur complement** of A with respect to a_{11} .

We claim that if A is nonsingular, then the Schur complement is nonsingular, too. Why? Suppose that the Schur complement, which is $(n-1) \times (n-1)$, is singular. Then by Theorem D.1, it has row rank strictly less than $n-1$. Because the bottom $n-1$ entries in the first column of the matrix

$$\begin{pmatrix} a_{11} & w^T \\ 0 & A' - vw^T/a_{11} \end{pmatrix}$$

are all 0, the bottom $n-1$ rows of this matrix must have row rank strictly less than $n-1$. The row rank of the entire matrix, therefore, is strictly less than n . Applying Exercise D.2-8 to equation (28.8), A has rank strictly less than n , and from Theorem D.1 we derive the contradiction that A is singular.

Because the Schur complement is nonsingular, we can now recursively find an LU decomposition for it. Let us say that

$$A' - vw^T/a_{11} = L'U',$$

where L' is unit lower-triangular and U' is upper-triangular. Then, using matrix algebra, we have

$$\begin{aligned} A &= \begin{pmatrix} 1 & 0 \\ v/a_{11} & I_{n-1} \end{pmatrix} \begin{pmatrix} a_{11} & w^T \\ 0 & A' - vw^T/a_{11} \end{pmatrix} \\ &= \begin{pmatrix} 1 & 0 \\ v/a_{11} & I_{n-1} \end{pmatrix} \begin{pmatrix} a_{11} & w^T \\ 0 & L'U' \end{pmatrix} \\ &= \begin{pmatrix} 1 & 0 \\ v/a_{11} & L' \end{pmatrix} \begin{pmatrix} a_{11} & w^T \\ 0 & U' \end{pmatrix} \\ &= LU, \end{aligned}$$

thereby providing our LU decomposition. (Note that because L' is unit lower-triangular, so is L , and because U' is upper-triangular, so is U .)

Of course, if $a_{11} = 0$, this method doesn't work, because it divides by 0. It also doesn't work if the upper leftmost entry of the Schur complement $A' - vw^T/a_{11}$ is 0, since we divide by it in the next step of the recursion. The elements by which we divide during LU decomposition are called **pivots**, and they occupy the diagonal elements of the matrix U . The reason we include a permutation matrix P during LUP decomposition is that it allows us to avoid dividing by 0. When we use permutations to avoid division by 0 (or by small numbers, which would contribute to numerical instability), we are **pivoting**.

An important class of matrices for which LU decomposition always works correctly is the class of symmetric positive-definite matrices. Such matrices require no pivoting, and thus we can employ the recursive strategy outlined above without fear of dividing by 0. We shall prove this result, as well as several others, in Section 28.3.

Our code for LU decomposition of a matrix A follows the recursive strategy, except that an iteration loop replaces the recursion. (This transformation is a standard optimization for a “tail-recursive” procedure—one whose last operation is a recursive call to itself. See Problem 7-4.) It assumes that the attribute $A.rows$ gives the dimension of A . We initialize the matrix U with 0s below the diagonal and matrix L with 1s on its diagonal and 0s above the diagonal.

LU-DECOMPOSITION(A)

```

1   $n = A.rows$ 
2  let  $L$  and  $U$  be new  $n \times n$  matrices
3  initialize  $U$  with 0s below the diagonal
4  initialize  $L$  with 1s on the diagonal and 0s above the diagonal
5  for  $k = 1$  to  $n$ 
6       $u_{kk} = a_{kk}$ 
7      for  $i = k + 1$  to  $n$ 
8           $l_{ik} = a_{ik}/u_{kk}$            //  $l_{ik}$  holds  $v_i$ 
9           $u_{ki} = a_{ki}$                //  $u_{ki}$  holds  $w_i^T$ 
10     for  $i = k + 1$  to  $n$ 
11         for  $j = k + 1$  to  $n$ 
12              $a_{ij} = a_{ij} - l_{ik}u_{kj}$ 
13 return  $L$  and  $U$ 
```

The outer **for** loop beginning in line 5 iterates once for each recursive step. Within this loop, line 6 determines the pivot to be $u_{kk} = a_{kk}$. The **for** loop in lines 7–9 (which does not execute when $k = n$), uses the v and w^T vectors to update L and U . Line 8 determines the elements of the v vector, storing v_i in l_{ik} , and line 9 computes the elements of the w^T vector, storing w_i^T in u_{ki} . Finally, lines 10–12 compute the elements of the Schur complement and store them back into the ma-

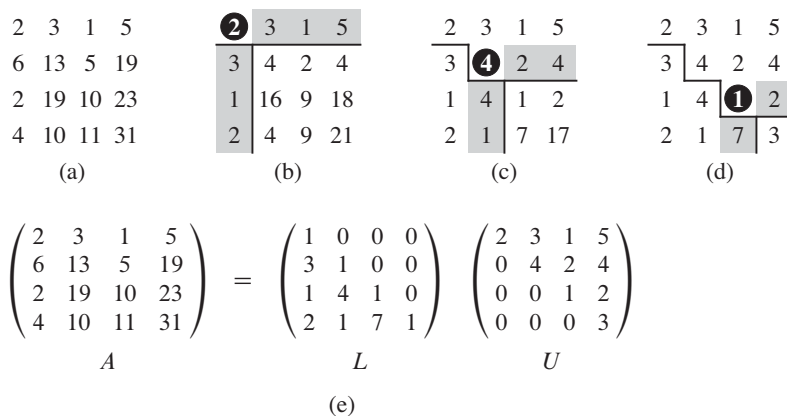


Figure 28.1 The operation of LU-DECOMPOSITION. (a) The matrix A . (b) The element $a_{11} = 2$ in the black circle is the pivot, the shaded column is v/a_{11} , and the shaded row is w^T . The elements of U computed thus far are above the horizontal line, and the elements of L are to the left of the vertical line. The Schur complement matrix $A' - vw^T/a_{11}$ occupies the lower right. (c) We now operate on the Schur complement matrix produced from part (b). The element $a_{22} = 4$ in the black circle is the pivot, and the shaded column and row are v/a_{22} and w^T (in the partitioning of the Schur complement), respectively. Lines divide the matrix into the elements of U computed so far (above), the elements of L computed so far (left), and the new Schur complement (lower right). (d) After the next step, the matrix A is factored. (The element 3 in the new Schur complement becomes part of U when the recursion terminates.) (e) The factorization $A = LU$.

trix A . (We don't need to divide by a_{kk} in line 12 because we already did so when we computed l_{ik} in line 8.) Because line 12 is triply nested, LU-DECOMPOSITION runs in time $\Theta(n^3)$.

Figure 28.1 illustrates the operation of LU-DECOMPOSITION. It shows a standard optimization of the procedure in which we store the significant elements of L and U in place in the matrix A . That is, we can set up a correspondence between each element a_{ij} and either l_{ij} (if $i > j$) or u_{ij} (if $i \leq j$) and update the matrix A so that it holds both L and U when the procedure terminates. To obtain the pseudocode for this optimization from the above pseudocode, just replace each reference to l or u by a ; you can easily verify that this transformation preserves correctness.

Computing an LUP decomposition

Generally, in solving a system of linear equations $Ax = b$, we must pivot on off-diagonal elements of A to avoid dividing by 0. Dividing by 0 would, of course, be disastrous. But we also want to avoid dividing by a small value—even if A is