Project parser example

Symbol table: empty

String buffer: empty

Instruction buffer:

- Our goal in the parser is to read a program such as shown on the left, and turn it into code that can be executed on the VM.
- The parser has several tasks
 - 1. Decide what statement is being parsed (there is one per line)
 - 2. Add declarations of variables and labels to a *symbol table* which is used in processing variables and jumps
 - 1. For jumps, determine where in the program the jump goes to
 - 2. For variables, determine where in storage the variable is

Symbol table: empty

String buffer: empty

Instruction buffer:

- The presence of functions (subroutines) in our code complicates life
- A function or subroutine is started by a gosublabel label statement and ended by a return statement
- A function will only contain a single return
- Variables declared in a function are
 - Only visible in the function (there scope is the function)
 - At runtime they will be put on the data memory, but we'll talk about that later
- Labels declared in a function are only visible on the function

Symbol table: empty

String buffer: empty

Instruction buffer:

- Handling functions/subroutines will be the most complicated part of the program
- Strings are tracked using a *string buffer*
- Generated code is placed into instruction buffer to be written to the output file when the entire program has been processed
 - Statements cannot be written immediately because, for example, a jump may be to a label that hasn't been declared yet (see, e.g., L2 to the left)
 - Thus code for the jump cannot be finalized until the target of the jump is known.
- Variables are always declared before being used.

start

exit 0

end

deciscal A pushi 4 gosub L1 jump L2 gosublabel L1 declscal A popscal A pushscal A return label L2 popscal A prints exit pgm

Symbol table: empty

String buffer: empty

The *instruction buffer* contains pointers to *Stmt* objects. A program always starts with a start statement.

Stmt is an abstract class (discussed later this week) that all statement types inherit from.

Instruction buffer:

OP_START_PROGRAM,? A program always starts with a *start*

statement

- All variables reside on the data memory, and in the code written to the output file for the start statement must contain the number of variables the *data memory* will hold for the program outside of those needed for functions
 - This count is not known until all variables are declared
 - Therefore, we leave that field in the OP_START_PROGRAM statement undefined

The *start* statement causes a Start statement object to be created and pointed to by the instruction buffer. The operand will be the size of data memory needed to hold outer scope variables, a value that is not yet known.

Symbol table: A, <0,1>

String buffer: empty

Instruction buffer:

 When a variable is declared, it is put into the symbol table, along with its position in the data memory and its length

OP_START_PROGRAM,?

- all variables are either scalar ints with a length of 1 or an array with a length that is the number of elements in the array
- The position for some variable is one plus the location at the end of the previously declared variable

The *declscal A* statement causes and entry for A to be created in the current symbol table. Since this is the first declared variable, its location will be 0.

Symbol table: A, <0,1>

String buffer: empty

statement buffer:

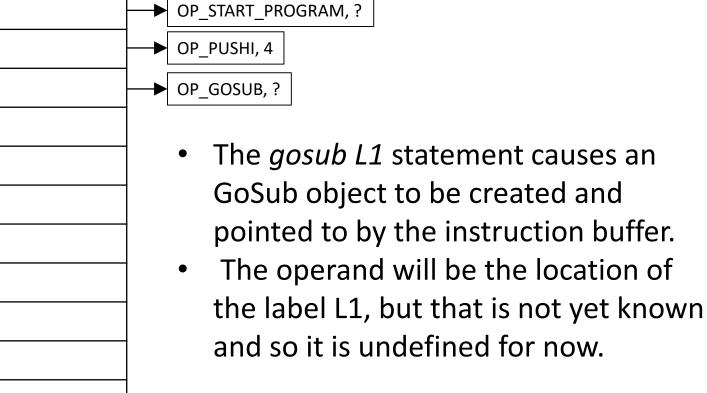
 The pushi 4 statement causes a Push statement object to be created and pointed to by the buffer. The operand is 4, the value to be pushed.

OP_START_PROGRAM,?

OP_PUSHI, 4

Symbol table: A, <0,1>

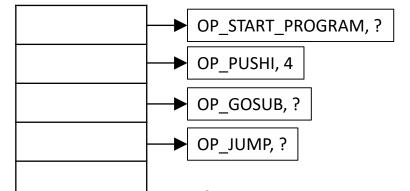
String buffer: empty



Symbol table: A, <0,1>

String buffer: empty

statement buffer:



The *jump L2* statement causes an Jump object to be created and pointed to by the instruction buffer.

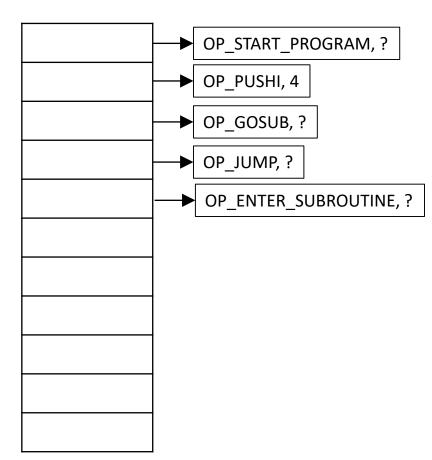
The operand will be the location of the label L2, but that is not yet known and so it is undefined for now.

exit 0

end

Symbol table: **A, <0,1> L1 <4, 0>**

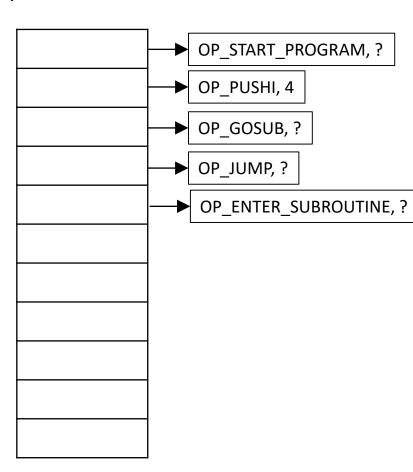
String buffer: empty



- The gosublabel L1
 statement causes an
 GoSubLabel object to be
 created and pointed to by
 the instruction buffer.
- The label L1 is added to the outer scope symbol table, along with its location (4) and its length (zero) since it doesn't occupy any space in the data memory.
- The operand will be the size of the data memory needed for the subroutine, which is not yet known, and so is undefined.

Symbol table: A, <0,1> L1 <4, 0> A <1, 0>

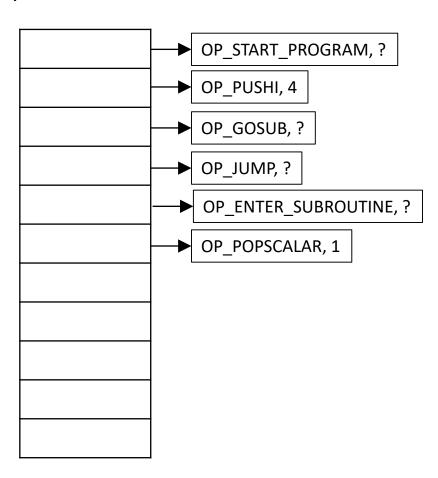
String buffer: empty



- The *declscal A* statement causes an entry for A in the symbol table.
- Since we are in a new scope (the subroutine) it is ok for there to be a name that is the same as a name in the outer scope.
- Since it is a scalar the length is 0.
- It is the first variable declared in the subroutine, so its location will be 1 plus the location of the last storage declared in the outer scope.
- We will get rid of the subroutine declarations at the end of the subroutine, so we need to mark where the subroutine table begins.

Symbol table: A, <0,1> L1 <4, 0> A <1, 0>

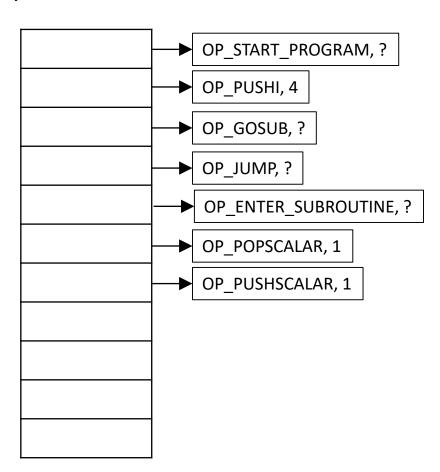
String buffer: empty



- The popscal A
 statement causes a
 PopScalar object to be
 created and pointed to
 by the instruction
 buffer.
- The operand will be the location of A in the data memory.
- Since variables are always declared before being used, A can be looked up in the symbol table and its location (1) determined immediately.

Symbol table: A, <0,1> L1 <4, 0> A <1, 0>

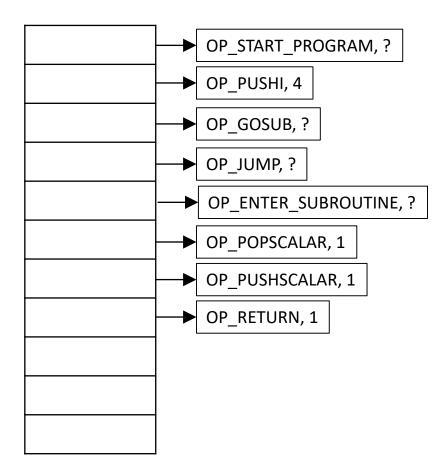
String buffer: empty



- The pushscal A
 statement causes a
 PushScalar statement
 object to be created
 and pointed to by the
 instruction buffer.
- The operand will be the location of A.
 Since variables are always declared before being used, A can be looked up in the symbol table and its location (1) determined.

Symbol table: A, <0,1> L1 <4, 0>

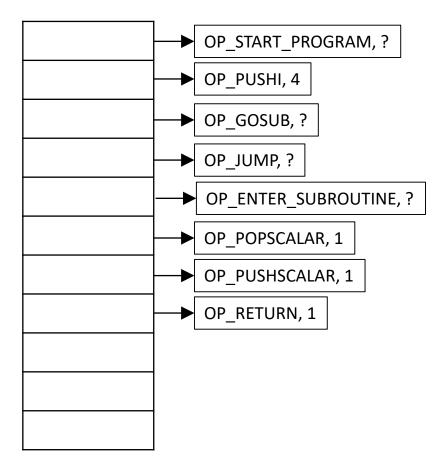
String buffer: empty



- The Return statement causes a Return statement object to be created and pointed to by the instruction buffer.
- The operand of the OP_ENTER_SUBROUTINE statement is the combined length all variables declared in the subroutine and gotten from the symbol table. This allows storage space in the data memory to be allocated for the variables.
- The symbol table for the subroutine is popped or deleted. The operand is not used and is undefined.

Symbol table: A, <0,1> L1 <4, 0> L2 <8, 0>

String buffer: empty

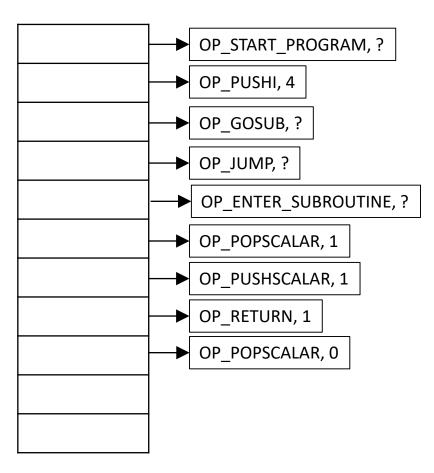


- The label L2 statement causes a symbol table entry for the label to be created.
- Its location is the location of the next instruction in the instruction buffer following the label, 8 in this case.
- Its length is 0.

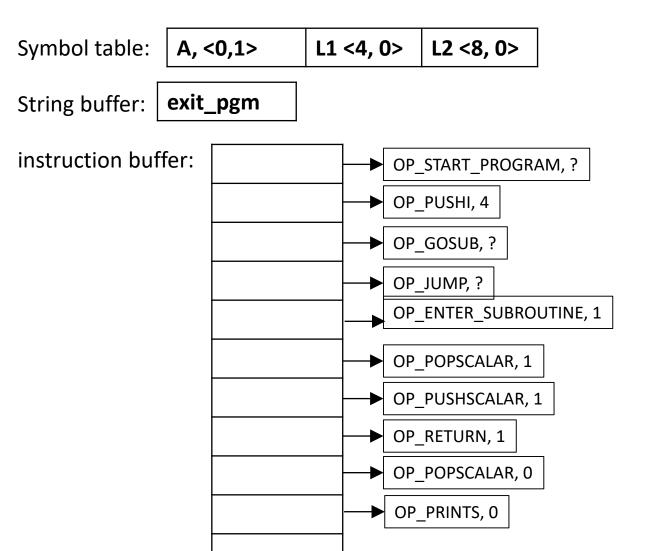
Symbol table: A, <0,1> L1 <4, 0> L2 <8, 0>

String buffer: empty

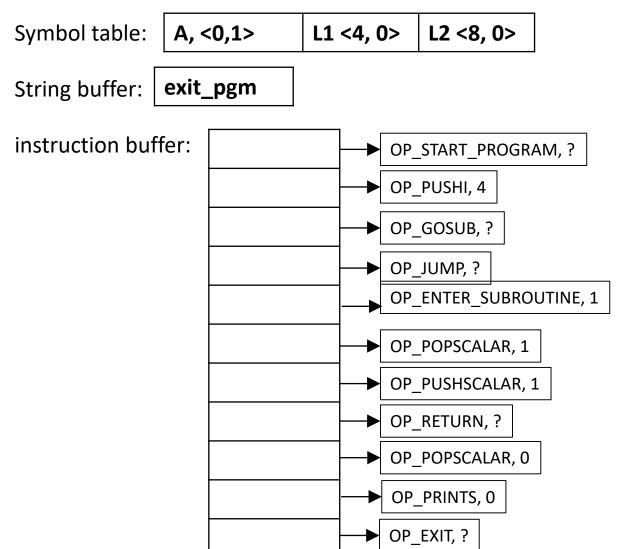
instruction buffer:



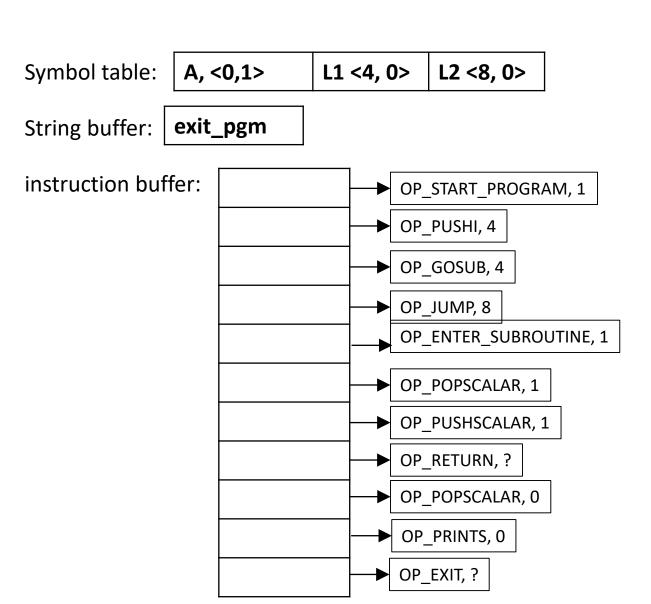
- The popscal A
 statement causes a
 PopScalar statement
 object to be created
 and pointed to by the
 instruction buffer.
- A is looked up in the symbol table, and is now the A declared in the outer scope.
- It's location is the operand of the popscal instruction



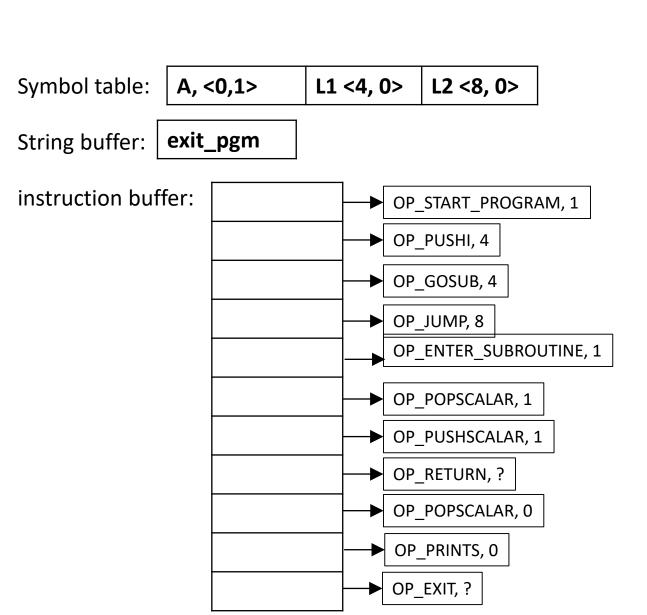
- The prints exit_pgm
 statement causes a
 Prints statement
 object to be created
 and pointed to by the
 instruction buffer.
- The string "exit_pgm"
 is added to the string
 buffer, and its location
 (0) is the operand of
 the instruction.



- The exit statement causes an Exit statement object to be created and pointed to by the instruction BROUTINE, 1 buffer.
- Its operand is unused and left undefined.



- The end statement signals the end of the program text
- All symbols have been defined.
- The parser now goes through all statements with undefined operands that point to labels, and patches them up.
 - Thus, the
 OP_JUMP and
 OP_GOSUB
 statements are
 patched.



- The
 OP_START_PROGRAM
 operand is patched
 with the size of all
 variables declared in
 the outer scope to
 allow the initial stack
 frame to be set up at
 runtime.
- The string table and the instructions are dumped to the .out file to be passed to the VM for execution.
- Now, the parser has completed its work.