

MPTEE: Bringing Flexible and Efficient Memory Protection to Intel SGX

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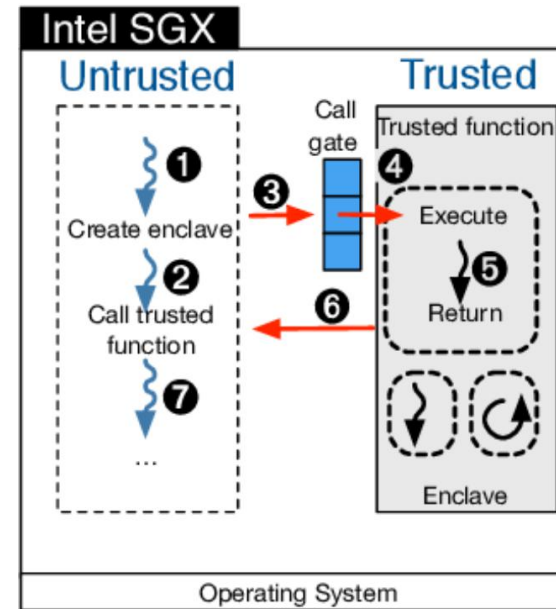
Background

➤ Inter SGX

- Software Guard Extension
- Enclave Cache Page

➤ Intel MPX

- Memory Protection Extensions
- Four bound registers (BND0 ~ BND3)



Motivation

➤ **Need of Flexible memory protection**

- Dynamically update memory-page permissions
- Least-privilege principle & Security

➤ **Existing System**

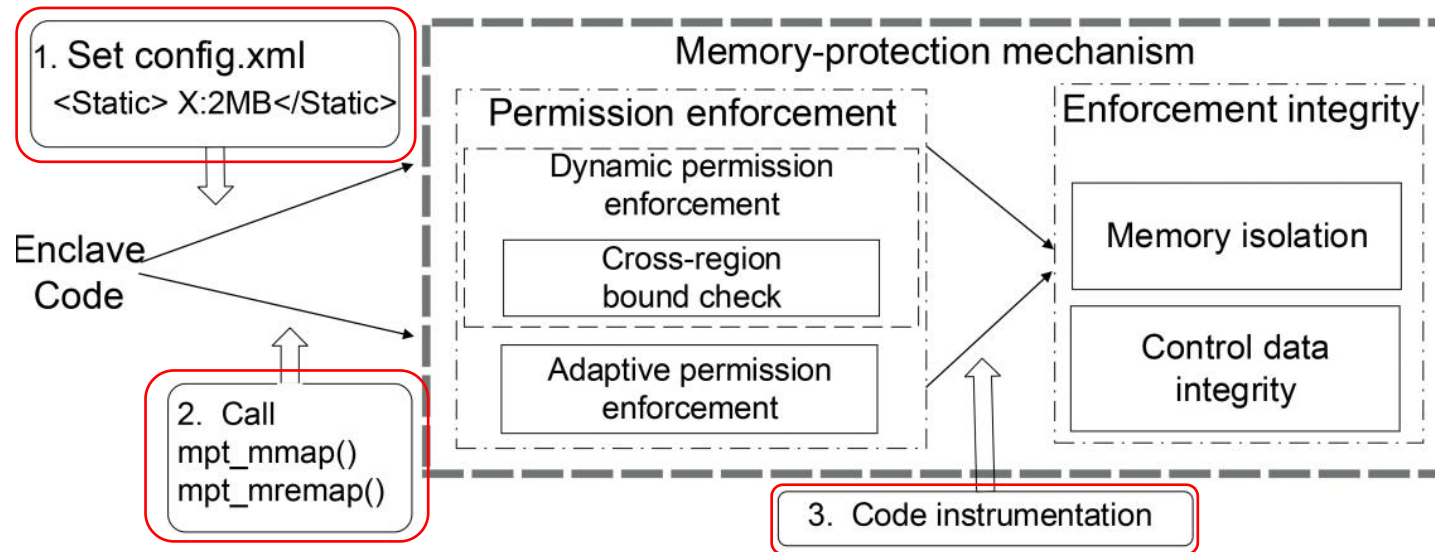
- SGX: SECS & Signature-Verification mechanism
- SGX-Shield: R15 Register & SGXCrypter: OS page table
- MPX: Just four registers

➤ **Challenge**

- Limited hardware support
- Strong adversary

Architecture

➤ MPTEE



➤ Permission Enforcement

- Dynamic
- Adaptive

➤ Enforcement Integrity (EI)

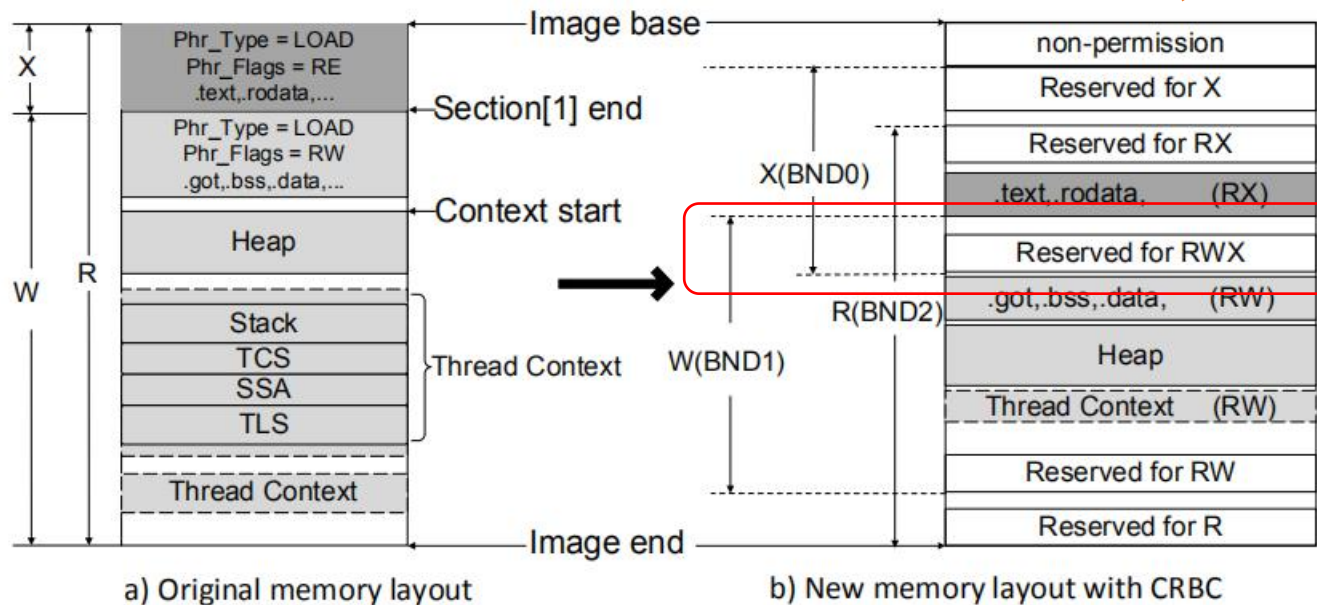
- Memory isolation
- Control-data integrity (CDI)

A memory-protection mechanism

Flexible and **Efficient** enforcement of memory-page permissions in SGX.

Permission Enforcement

➤ Cross-Region Bound Check*



Permission	Bound range
non-perm.	(ImageBase, bnd0.lb)
X	(bnd0.lb, bnd2.lb)
RX	(bnd2.lb, bnd1.lb)
RWX	(bnd1.lb, bnd0.ub)
RW	(bnd0.ub, bnd1.ub)
R	(bnd1.ub, bnd2.ub)

Overlapping

➤ Unique memory layout of enclaves

- Memory sections that have same permissions are adjacent

➤ Cross-region bounds

- ❶ BND0 ~ BND2 (Three Registers)
- ❷ X/R/RX/RW/RWX/No (Six permissions)
- ❸ RWX - RW: bnd0.ub decrease

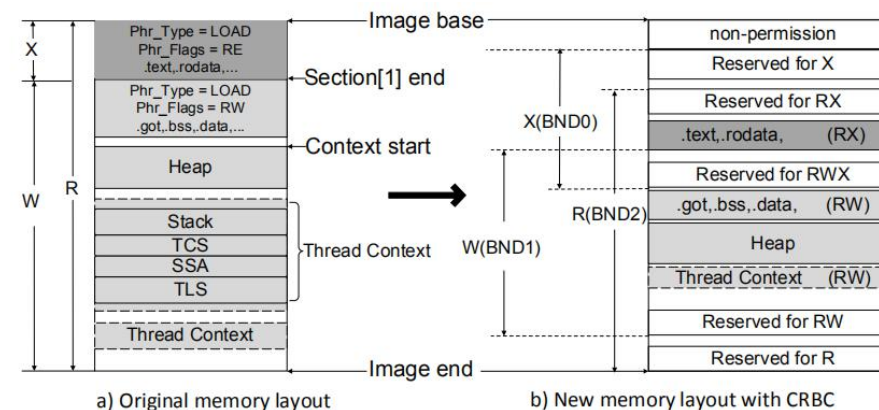
Permission Enforcement

➤ Cross-Region Bound Check*

API	Description
<code>void *mpt_mmap(size, flags)</code>	Acquires a memory buffer, which is at least <i>size</i> bytes. It is restricted as specified in <i>flags</i> .
<code>void *mpt_mremap(pages, size, flags)</code>	Changes <i>old_pages</i> to <i>flags</i> . If the <i>pages</i> are located at a region boundary, it changes the permissions; otherwise, it allocates new memory buffer and copy content over.
<code>void *mpt_munmap(pages)</code>	Frees an acquired memory buffer, the freed memory will be sealed and swapped out.
<code>void *mpt_write(pages, size, content)</code>	Writes <i>content</i> to the mapped region. The function will not be bound-checked.

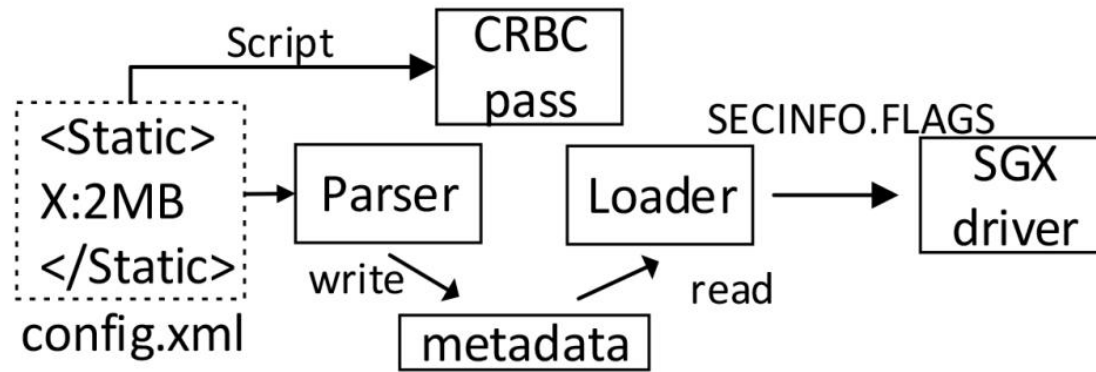
➤ Four APIs

- `mpt_mmap()` for acquire memory
- `mpt_mremap()` for change permission



Permission Enforcement

➤ Adaptive Permission Enforcement



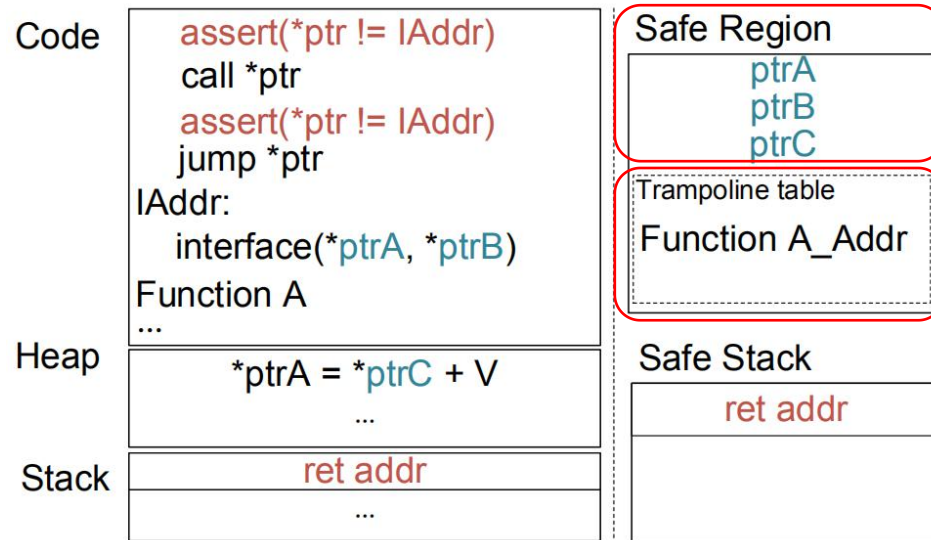
➤ Setting the static region in **config.xml**

➤ `<Permission>:<Region size>`

➤ CRBC pass: Reduce the performance overhead

Enforcement Integrity

➤ Memory Isolation



➤ Safe Region

- Protect code and data for permission enforcement
- No-permission region

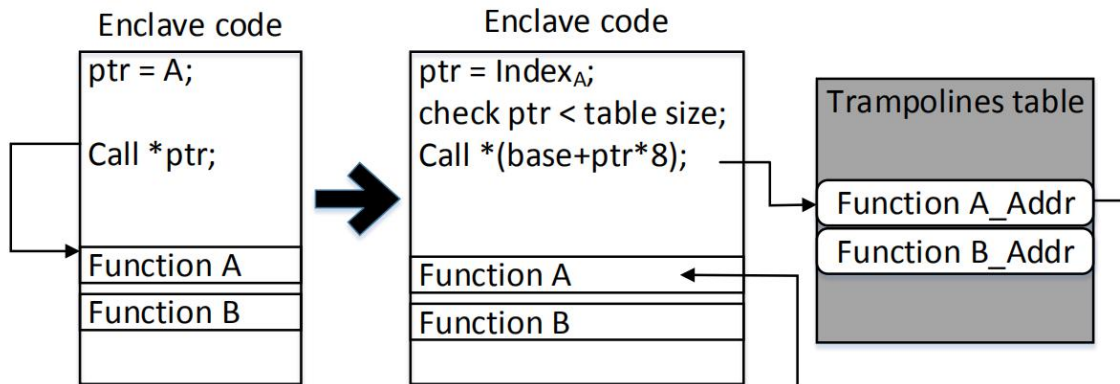
➤ Safe Stack & Trampoline table

- Ensure the integrity of return addresses and function pointers

➤ Safe Region & Trampoline table

Enforcement Integrity

➤ Control-Data Integrity



➤ Trampoline table

- protect pointers for indirect calls and jumps

➤ Static analysis: All functions whose addresses are ever taken

Evaluation

➤ Experiments

- Dell workstation
- Intel Xeon E3-1225v5 CPU & 8G RAM
- Ubuntu 16.04 Server

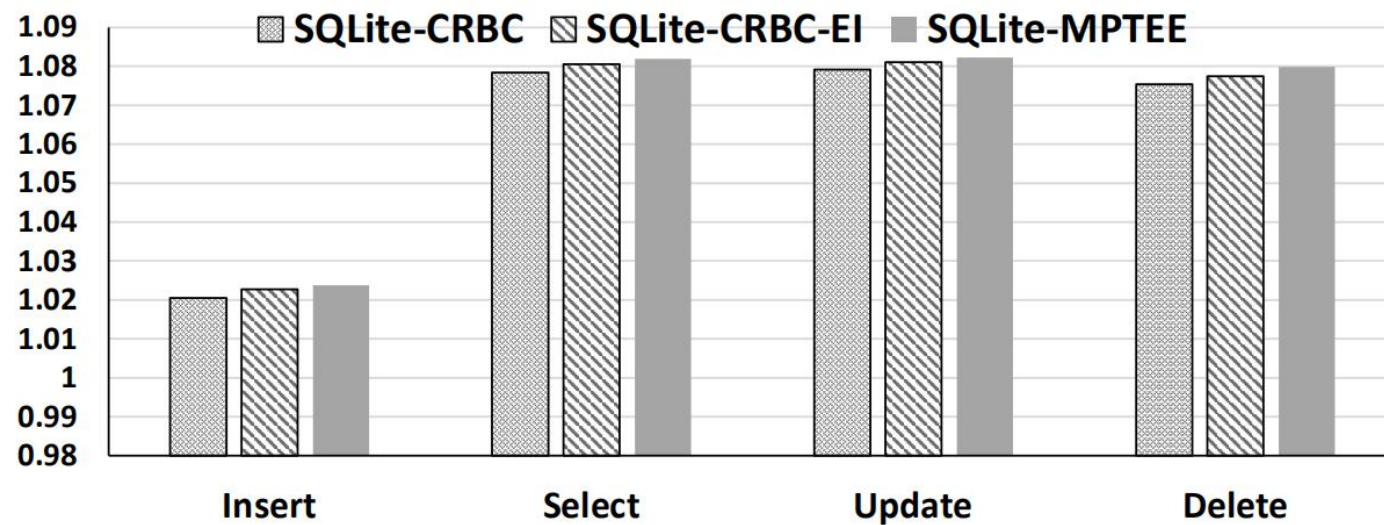
➤ Evaluation

- Performance
- Utility Analysis
- Security Analysis

Evaluation

SQLite is a C-language library
Implements small and fast SQL database engine.

➤ Performance

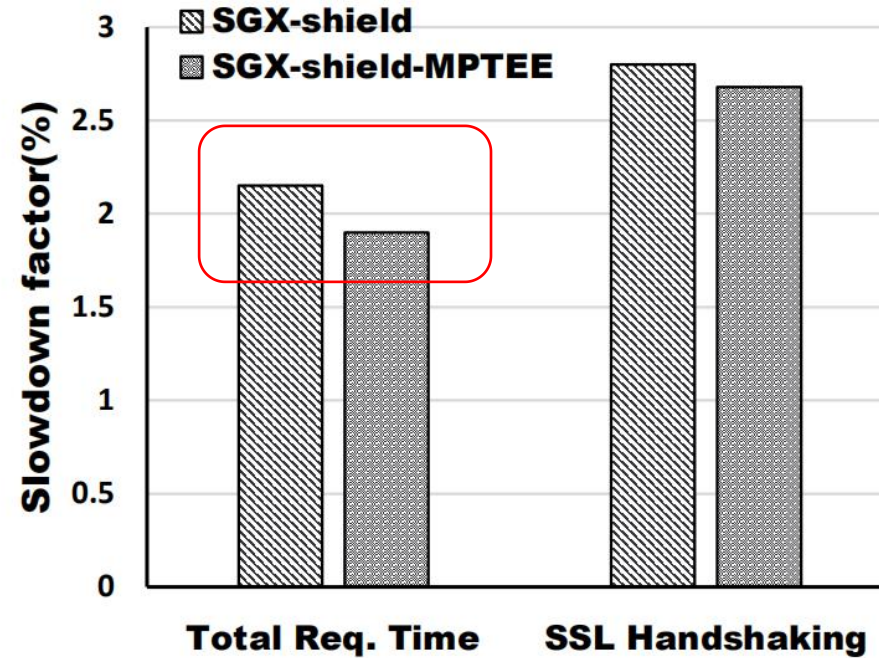


Low overhead imposed by MPTEE

- CRBC imposes 2%–7% (avg 6.6%) overhead
- EI imposes smaller than 1%
- More frequent memory accesses leading to more bound checks.

Evaluation

➤ Performance

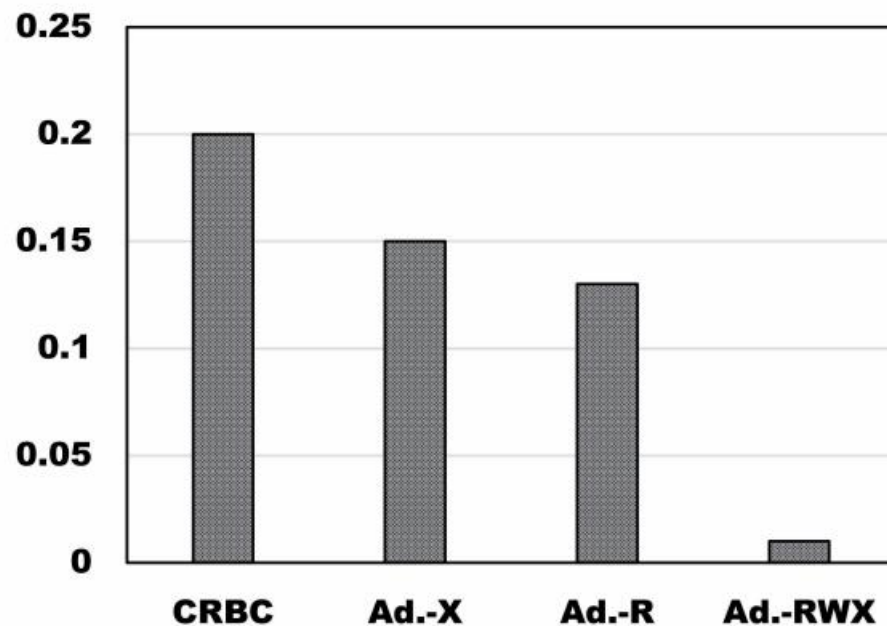


Requesting HTML files from the HTTPS server

- SGX-Shield-MPTEE has lower overhead 1.9% (with no adaptive permission)
- SGX-Shield is 2.2% & Only provide one permission

Evaluation

➤ Performance



➤ CRBC incur up 20% overhead

➤ Adaptive-RWX incurs almost 0% (hardware)

➤ Bound-checking of R is more frequent than X

Benefits from adaptive permission enforcement

Evaluation

➤ Utility Analysis

- Confirm MPTEE is easy-to-use and practical
- SGX-Shield NRW boundary
 - `mpt_mmap(sgx_end-sgx_start, X)`
- Protecting SGX-Shield code

Evaluation

➤ Threat Model

- Trusted: only Intel SGX itself
- Untrusted: All other software and hardware components
- Vulnerabilities of code running inside SGX: buffer overflows

➤ Security Analysis

- Check-skipping attacks: CDI
- Bound-manipulating attacks: Memory Isolation

Conclusion

➤ **MPTEE**

- Cross-Region Bound Check
- Adaptive Permission Enforcement
- Memory Isolation & Control-data Integrity

➤ **Evaluation**

- Performance
- Utility & Effectiveness.

MPTEE can achieve a flexible and secure memory protection while just imposing a small overhead

Comments

➤ Advantages

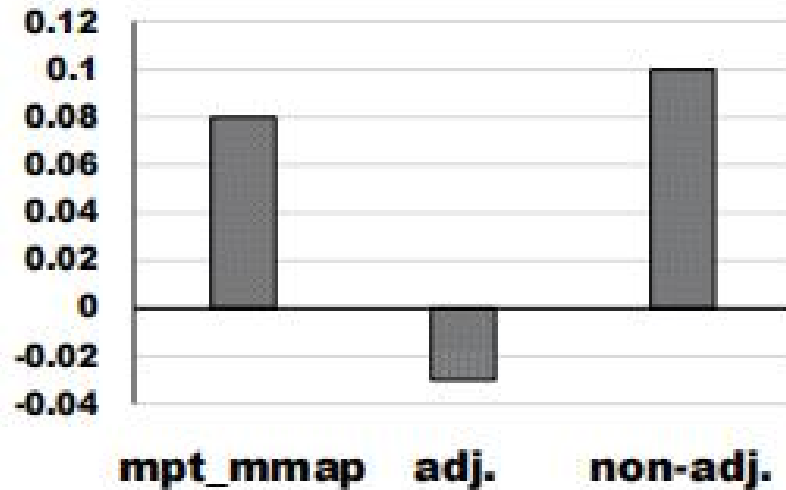
- The MPTEE actually develop a flexible and secure memory-protection mechanism with less hardware support and just impose a small overhead.

➤ Disadvantages

- Intel MPX has been eliminated by Microsoft.
- Argues:
 - As the hardware component, SGX cannot be easily upgraded. Old version SGX will still be the dominating version for a while.
 - For enforcement integrity, MPTEE's design is compatible with new SGX.
- Maintenance burdens

Append

➤ mpt_mremap()



Input : addr, size, flags

Output : new_addr

Data: Metadata of 5 Reserved region, 3 bnd register

```
1 Get cur_region, flags_region, neighbor_region from addr and registers
2 if cur_region == flags_region then
3     new_addr = realloc(addr, size)
4     if realloc failed then
5         if neighbor_region is not reserved || neighbor_region has
           allocated then
6             new_addr = NULL
7         else
8             Get bndX according to flags_region
9             bndX.lb = bndX.lb + size
10            update metadata of cur_region
11            new_addr = realloc(addr, size)
12 else if cur_region is neighbor with flags_region then
13     Get bndX according to flags_region
14     bndX.lb = bndX.lb - size
15     new_addr = addr
16 else
17     Get the dest_heap according to flags
18     new_addr = malloc in dest_heap
19     mpt_write(new_addr, size, addr)
20     mpt_munmap(addr)
21 return new_addr
```
