

Department of Data Science and Knowledge Engineering

Logic 2020/2021 Resit Exam Questions

— Do not turn this page before the official start of the exam! —

First Name, Surname:
Student ID:
Program: Bachelor Data Science and Knowledge Engineering
Course code: KEN1530
Examiners: dr. Otti D'Huys, dr. ir. ing. Nico Roos
Date/time: July 1 st , 2021, 9:30-11:30h
Format: Closed book exam.
Allowed aides: Pens.

Instructions to students:

General instructions:

- The exam consists of 10 questions on 12 pages (excluding the 1 cover page(s)).
- Answer every question on a separate piece of paper. Do not mix the answers on different exam
 questions.
- Number each page of answers you submit in the top left corner.
- Ensure that you properly motivate your answers.
- Do not use red pens, and write in a readable way. Answers that cannot be read easily cannot be graded and may therefore lower your grade.
- You are not allowed to have a communication device within your reach, nor to wear or use a watch.
- If you think a question is ambiguous, or even erroneous, explain this in detail in your answer.
- If you have not registered for the exam, your answers will not be graded, and thus handled as invalid.

For on-site exams:

- Fill in your name and student ID number on every page of answers you submit.
- Please make sure that communication devises and watches are not within reach.
- The use of pencils is not allowed.

For online proctored exams:

- The scan / photographs of your submitted answers must be readable.
- Reserve space for your ID in the top right corner of every page of answers you submit.

Success!

The following table will be filled by the examiner:

Question:	1	2	3	4	5	6	7	8	9	10	Total
Points:	10	10	15	10	10	10	10	5	10	10	100
Score:											

Question 1 (10 points)

- Alice likes running or Bob likes cycling.
- If Chris likes dancing, then Bob does not like cycling.
- It is not the case that (Alice likes running and Chris likes dancing).
- If Bob likes cycling and Alice does not like running, then Chris likes dancing.

Determine the answers of the following questions using a truth table. (Draw the full table!) Answer with: Yes, No or Unknown!

- Does Alice like running?
- Does Bob like cycling?
- Does Chris like dancing?

Use the atomic propositions a, b and c, use the truth-value 0 representing false and the truth-value 1 representing true, and enumerate the valuations of a, b and c as binary numbers from 0 to 7.

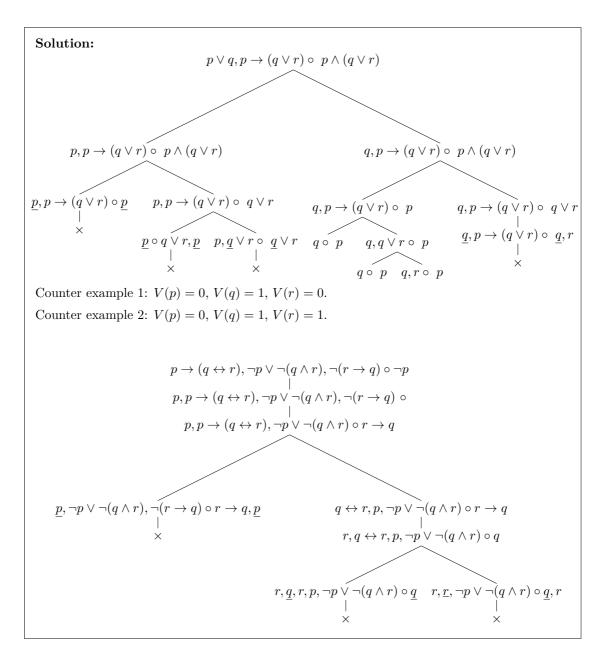
Solı	ıtio	n:														
	a	b	c	$a \lor b$	c	\rightarrow	_	b	_	$(a \wedge c)$	(b	\wedge	\neg	a)	\rightarrow	c
	0	0	0	0		1	1		1	0		0	1		1	
	0	0	1	0		1	1		1	0		0	1		1	
	0	1	0	1 1		1	0		1	0		1	1		0	
	0	1	1	1 1		0	0		1	0		1	1		1	
\Rightarrow	1	0	0	1		1	1		1	0		0	0		1	
	1	0	1	1		1	1		0	1		0	0		1	
\Rightarrow	1	1	0	1 1		1	0		1	0		0	0		1	
	1	1	1	1 1		0	0		0	1		0	0		1	

- Does Alice like running? yes
- Does Bob like cycling? unknown
- Does Chris like dancing? no

Question 2 (10 points)

Investigate the validity of the following inferences with the aid of a semantic tableau. If an inference is invalid, give **all** the counterexamples.

- $\{p \lor q, p \to (q \lor r), \neg(r \land p) \to \neg q\} \models p \land (q \land r)$
- $\bullet \ \{p \to (q \leftrightarrow r), \neg p \lor \neg (q \land r), \neg (r \to q)\} \models \neg p$



Question 3 (15 points)

Prove by means of natural deduction:

- From the premises $p \to \neg(r \to q)$ derive $p \to (\neg q \land r)$
- From the premise $\neg(\neg p \lor \neg q)$ derive $p \land q$

```
Solution:
 3
 4
 5
 6
 7
 8
 9
                                      \rightarrowI, 5, 9
\perpI, 3, 10
 10
 11
                                      ¬E, 4–11
 12
                              \rightarrowI, 2, 12
 13
 1
                             \vee I, 2
 3
                              \perp I, 1, 3
 4
                              \neg E, 2-4
 5
 6
 7
                             \perp I, 1, 7
 8
                              \neg E, 6-8
 9
 10
                              \wedge I, 5, 9
        p \wedge q
```

Question 4 (10 points)

Consider the following two premises of a syllogism:

No cow is pink.

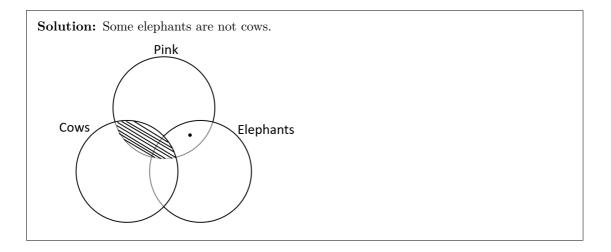
Some elephants are pink.

??

Student name:

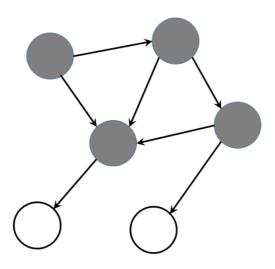
Complete the syllogism with a conclusion that makes it valid.

Demonstrate the validity of your syllogism using the method with Venn diagrams.



Question 5 (10 points)

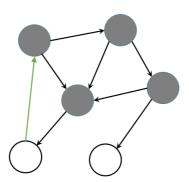
Consider the predicate logic model shown below. The model has a unary predicate P and a binary predicate S. Shaded objects have property P, a \rightarrow from a to b meant that Sab is true in this model.



- (a) State two closed formulas that are true in this model. Your formulas should include both predicates P and S.
- (b) Add one pair to I(S) (one arrow to the model representation), such that the statement $\forall x(Px \to \exists ySyx)$ becomes true.

Solution:

- (a) Examples include $\forall x(\exists ySxy \to Px), \neg \exists x(Px \land Rxx), \forall x \forall y((Syx \land \neg Px) \to Py), \exists y \exists x(Sxy \land Py \land Px)$
- (b) Any arrow arriving at the upper left object makes the formula true. An example is below



Question 6 (10 points)

Are the following formulas a tautology? Show by means of a semantic tableau.

- $\forall x \exists y (Ryx \to Rxx)$
- $\forall x ((\exists y Ryx) \to Rxx)$

Solution:

This statement is a tautology

This tableau has an open branch: we find a countermodel with two objects a and b, with b related to a, but a not related to itself.

Question 7 (10 points)

Prove by means of natural deduction:

- From the premises $\{ \forall x (Px \to \exists y Rxy), \forall x (\exists y Ryx \to \neg Px), \exists x Px \}$ derive $\exists x \neg Px$
- From the premises $\{ \forall x (Ax \to \neg Bx), \exists x (Bx \land Cx) \}$ derive $\exists x (Cx \land \neg Ax)$

```
\forall x (Px \to \exists y Rxy)
                     1
                     2
                              \forall x (\exists y Ryx \to \neg Px)
                              \exists x P x
                     3
                               a \mid Pa
                     4
                                     Pa \rightarrow \exists y Ray
                                                                             \forall E, 1
                     5
                     6
                                     \exists y Ray
                                                                             \rightarrowE, 5, 4
Solution: 7
                                     b \mid Rab
                                           \exists yRyb
                     8
                                                                             ∃I, 7
                     9
                                           \exists yRyb \rightarrow \neg Pb
                                                                             \forall E, 2
                     10
                                           \neg Pb
                                                                             \rightarrowE, 9, 8
                     11
                                           \exists x \neg Px
                                                                             ∃I, 10
                                     \exists x \neg Px
                                                                             \exists E, 6, 7-11
                     12
                               \exists x \neg Px
                                                                             \exists E,\ 3,\ 4\text{--}12
                     13
```

```
Solution:
          \forall x (Ax \to \neg Bx)
 1
 2
          \exists x (Bx \wedge Cx)
 3
          a \mid Ba \wedge Ca
               Ca
                                              \wedge E, 3
 4
                                              \wedge E, 3
 5
               Ba
                     Aa
 6
 7
                     Aa \to \neg Ba
                                             \forall E, 1
                                             \rightarrowE, 6, 7
 8
                     \neg Ba
                     \perp
                                             \perpI, 5, 8
 9
                                              \neg I, 6-9
 10
                \neg Aa
               \neg Aa \wedge Ca
                                             \wedge I, 4, 10
 11
               \exists x (\neg Ax \land Cx)
 12
                                              ∃I, 11
 13
          \exists x (\neg Ax \land Cx)
                                              \exists E,\ 2,\ 3\text{--}12
```

Note: the topic of this question (epistemic logic) has been replaced in 2021/2022 by Hoare logic and tableau proofs for PDL. Please ignore this question.

Question 8 (5 points)

Translate the following sentence into a formula of epistemic logic, using appropriate translation keys.

• Bob does not know whether Chris knows whether Alice likes sailing.

Use the modal operators \square_A , \square_B and \square_C , and the atomic proposition a.

Solution:

• $\neg(\Box_B(\Box_C a \lor \Box_C \neg a) \lor \Box_B \neg(\Box_C a \lor \Box_C \neg a))$

or

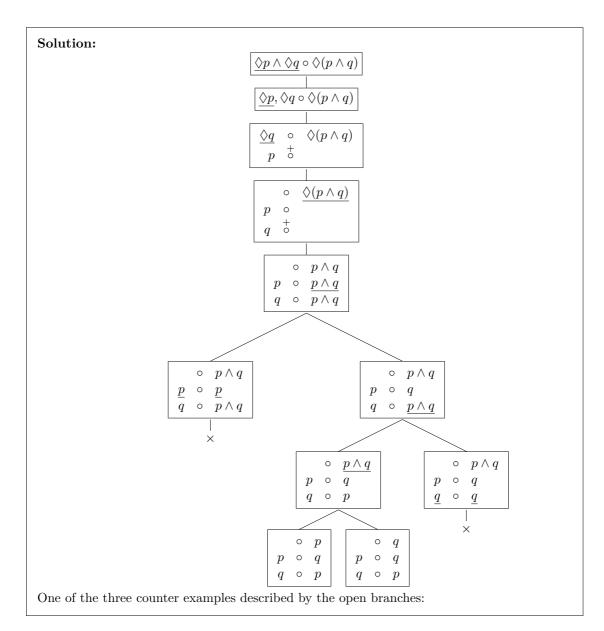
 $\neg \Box_B(\Box_C a \vee \Box_C \neg a) \wedge \neg \Box_B \neg (\Box_C a \vee \Box_C \neg a)$

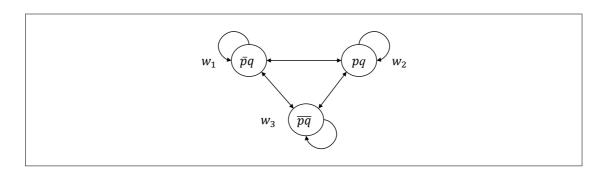
Note: the topic of this question (epistemic logic) has been replaced in 2021/2022 by Hoare logic and tableau proofs for PDL. Please ignore this question.

Question 9 (10 points)

Investigate the validity of the following inferences with the aid of a semantic tableau. If an inference is invalid, give a counterexample.

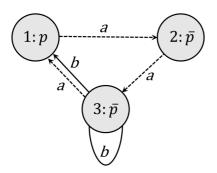
- $\Diamond p \wedge \Diamond q \models \Diamond (p \wedge q)$
- $\neg \Diamond (p \lor q) \models \Box \neg p \land \Box \neg q$





Question 10 (10 points)

Consider the labelled transition system below, with 3 states and actions a and b.



- (a) In which states are the following formula true?
 - $\langle a^* \rangle p$.
 - $[b] \neg p$
- (b) Give all the elements of the relation defined by the action $(?\neg p; a \cup b)^*$.

Solution:

- (a) In which states are the following formula true?
 - $\langle a^* \rangle p$: true in all states.
 - $[b]\neg p$: true in states 1 and 2.
- (b) The elements of the relation defined by the action $(?\neg p; a\cup b)^*$ are (1,1),(2,2),(3,3),(2,3),(3,1),(2,1)