

Binary search

Input: An array A of elements in sorted order, and an element x .

Output: Return the position of x if it exists; otherwise output nil.

Idea: Set q = middle item. Check if $x \leq A[q]$,
and search either to left or right of q as appropriate

$q =$	1	2	3	4	5	6	7	8	9	10	11	12	13	14
$A[q]$	4	7	10	15	19	20	42	54	67	75	81	87	93	96

BSearch(A, p, r, x) :

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if  $p = r$  then
  if  $A[p] = x$  then return  $p$ 
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else
   $q \leftarrow \lfloor (p + r) / 2 \rfloor$ 
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First call: BSearch($A, 1, n, x$)

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First call: **BSearch**($A, 1, n, x$)

BSearch($A, 1, 14, 67$)

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First call: **BSearch**($A, 1, n, x$)

BSearch($A, 1, 14, 67$)

$q \leftarrow \lfloor (p + r) / 2 \rfloor = 7$

$67 > 42 = A[7]$

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First call: **BSearch**($A, 1, n, x$)

BSearch($A, 8, 14, 67$)

$q \leftarrow \lfloor (p + r) / 2 \rfloor = 11$

$67 \leq 81 = A[11]$

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First call: **BSearch**($A, 1, n, x$)

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$q \leftarrow \lfloor (p + r) / 2 \rfloor = 9$

$67 \leq 67 = A[9]$

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First call: **BSearch**($A, 1, n, x$)

BSearch($A, 8, 9, 67$)

$q \leftarrow \lfloor (p + r) / 2 \rfloor = 8$

$67 > 54 = A[8]$

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First call: **BSearch**($A, 1, n, x$)

BSearch($A, 9, 9, 67$)

$p = r$ AND $67 = A[p] = A[9]$

Success

Binary search

Previous example was of a successful search, in which the search key was in the the array.

What happens if the search key is not in the array?

q =	1	2	3	4	5	6	7	8	9	10	11	12	13	14
A[q]	4	7	10	15	19	20	42	54	67	75	81	87	93	96

BSearch(A, p, r, x) :

```
if p = r then
    if A[p] = x then return p
    else return nil
else
    q ← ⌊(p + r)/2⌋
    if x ≤ A[q] then BSearch(A, p, q, x)
    else BSearch(A, q + 1, r, x)
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First call: **BSearch**(A, 1, n, x)

BSearch(A, 1, 14, 66)

Binary search

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$p = r$ AND $66 \neq A[p] = A[9]$

Failure
Return *nil*

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In fact, algorithm knows where 66 *should* be. In case of failure it could be modified to return locations of predecessor and successor of x in the array.

Analysis of Binary Search

Analysis: Let $T(n)$ be the number of comparisons needed for running the algorithm on at most n elements. (Assume n a power of 2).

Recurrence: A single comparison eliminates half of the array. Therefore, we search for the element in the remaining half, which has size $n/2$. Thus, the **recurrence** counting the number of comparisons is:

$$T(1) = 1. \text{ Otherwise, if } n > 1, T(n) = T\left(\frac{n}{2}\right) + 1.$$

Solve the **recurrence** by iterating it: (Assume n a power of 2)

$$\begin{aligned} T(n) &= T(n/2) + 1 \\ &= (T(n/2^2) + 1) + 1 \\ &= T(n/2^2) + 2 \\ &= \dots \\ &= T(n/2^i) + i \\ &= \dots \\ &= T(1) + \log n = \Theta(\log n) \end{aligned}$$

Proof of correctness
more formal analysis and analysis
when n is not a power of 2 are
developed in practice homework

Note: may sometimes terminate faster than $\Theta(\log n)$, but **worst-case running time** is $\Theta(\log n)$.