

\mathcal{V} -Profunctors.

We will see that Feasibility relation is just a special case of \mathcal{V} -profunctors where $\mathcal{V} = \text{Bool}$.

In order to understand the interpretation of co-design problem [formalizing and operating] we need understand pro-functors.

(Intuitively)

Bool-Category		Preorder
Bool-functor		Monotone map
Bool-profunctor		feasibility relation

Profunctors:

\mathcal{V} -Profunctors:

Def: let $\mathcal{V} = (V, \leq, I, \otimes)$ be a quantale (unital commutative) and let \mathcal{X} and \mathcal{Y} be \mathcal{V} -categories. A \mathcal{V} -profunctor from \mathcal{X} to \mathcal{Y} denoted $\Phi: \mathcal{X} \multimap \mathcal{Y}$, is a \mathcal{V} -functor $\Phi: \mathcal{X}^{\text{op}} \times \mathcal{Y} \rightarrow \mathcal{V}$

Because this is the codomain \mathcal{V} is enriched in itself.

- By def of \mathcal{V} -product and \mathcal{X}^{op} definition we can say that,

\mathcal{V} -profunctor is same as a function $\Phi: \text{Ob}(\mathcal{X}) \times \text{Ob}(\mathcal{Y}) \rightarrow \mathcal{V}$ such that for any $x, x' \in \mathcal{X}$ and $y, y' \in \mathcal{Y}$ the following inequality holds in \mathcal{V} :

$$x(x', x) \otimes \Phi(x, y) \otimes \mathcal{Y}(y, y') \leq \Phi(x', y')$$

Proof: \mathcal{V} -profunctor is $\Phi: \mathcal{X}^{\text{op}} \times \mathcal{Y} \rightarrow \mathcal{V}$

so, by definition of \mathcal{V} -product

$$\text{ob}(\mathcal{X}^{\text{op}} \times \mathcal{Y}) = \mathcal{X}^{\text{op}} \times \mathcal{Y} \quad (x, y), (x', y') \text{ are objects in } \mathcal{X}^{\text{op}} \times \mathcal{Y}$$

$$\text{so, } \mathcal{X}^{\text{op}} \otimes \mathcal{Y}((x, y), (x', y')) = \mathcal{X}^{\text{op}}(x, x') \otimes \mathcal{Y}(y, y') \quad \text{--- (i)}$$

$$\mathcal{X}^{\text{op}}(x, x') = x(x', x) \quad [\text{By def of } \mathcal{X}^{\text{op}}] \quad \text{--- (ii)}$$

$$\mathcal{V}(\Phi(x, y), \Phi(x', y')) := \Phi(x, y) \multimap \Phi(x', y') \quad \text{--- (iii)}$$

(Remark 2.89)

Now according to ϕ

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$$d_{\mathcal{X} \times \mathcal{Y}}((x, y), (x', y')) \leq d(\phi(x, y), \phi(x', y'))$$

$$x(x', x) \otimes (y(y', y')) \leq \phi(x, y) \longrightarrow \phi(x', y')$$

Now according to def 2.79

$$\underbrace{\gamma = \text{Bad profunctor}}_{\times} \quad \underbrace{\chi(x', x)}_{\times} \quad \underbrace{\phi(x, y)}_{\times} \quad \underbrace{\gamma(y, y')}_{\times} \leq \underbrace{\phi(x', y')}_{\times}$$

- Feasibility Matrix: ϕ can be expressed as a matrix. The (m,n) th entry is the value of $\phi(m,n) \in B$.

~~Fill out the Bool main.~~

Ex: 4.11 & Ex 4.12 for better understanding ∇ -profunctors

Interpreting the co-design diagram:

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Box : feasibility relation
L.H.S: produces (set of parts , each indicating a resource)
R.H.S: requirements (set " ")

Now, each of these resources (produce or requirements) are taken as pre-orders.

Then we take product pre-order of preorders on the left and similarly for those on the right. The box is then the feasibility relation

Box: () \longrightarrow ()
product of produces product of requirements

Now, feasibility relation says is it possible to get to the source from the target (can we get produce from the requirements)

(Understand movie example)

The collage above statement is valid for the whole