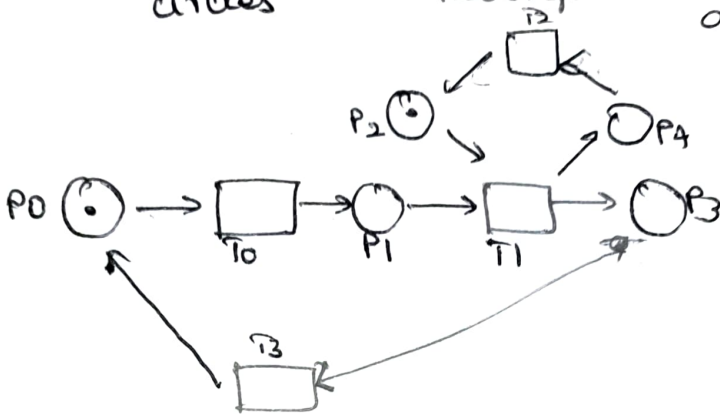


Petri-nets

• It consists of four elements:

- (i) Places \downarrow circles
- (ii) Transitions \downarrow Rectangles
- (iii) Edges \downarrow Directed arrows
- (iv) Tokens \downarrow Solid circles

Eg:



4 places: $P0 \rightarrow P4$ 3 transitions: $T0 \rightarrow T3$. Edges connect places to transitions and vice versa.
 $P0$ and $P2$ have a single token represented by dot.

\Rightarrow Petri net may contain cycles. Here it has 2 cycles.

$\{P0, T0, P1, T1, P3, T3\}$ and $\{P2, T1, P4, T2\}$

* Tokens represent the state of the petri nets.

There are rules by which petri nets change its state.

"A petri net" changes from one state to the next state when a transition fires. We define input, output for edges in the petri-net. ($In(T0) = P0$, $Out(T0) = P1$, $Out(T1) = P4, P3$)

Firing rules for a transition:

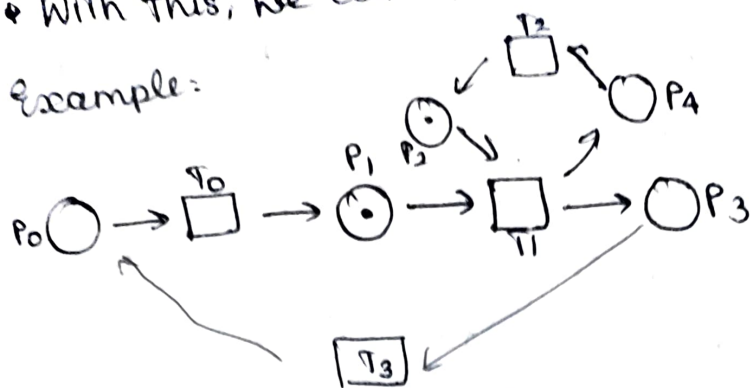
* A transition is able to fire when there is atleast one token on each of the input of transition.

* When a transition happens, it removes one token from input and would place it on output

• Every transition will get active at a point. If it is not eligible we will disable it.

With this, we can model non-deterministic behaviour

Example:



(Next state of previous example)

(T_1 could not do as P_1 did not have token)

From this state, we can fire T_1 (as both P_1, P_2 has tokens)

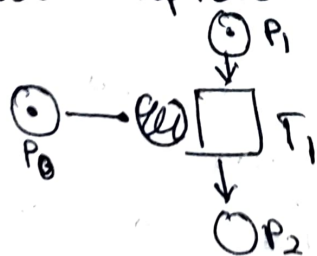
Once you have two transition states ready to fire, you can go in any order. After this state, we will reach the first state or second state. (T_2 after T_3) (T_3 after T_2), (T_3 after T_0 after T_2)

Extended edge types:

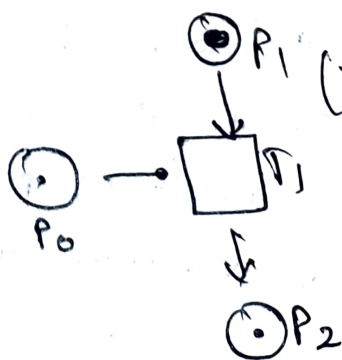
We will see the extensions to the petri net structure, two of it are read edges, inhibitor edge

* Read edges (arrow) can be drawn only from place to transitions unlike the normal one.

Read edges won't remove the token from where it reads represented by \rightarrow



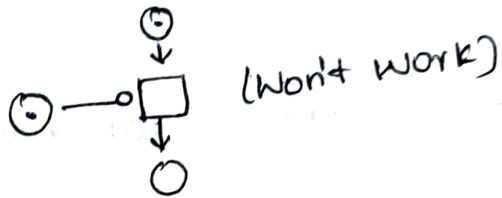
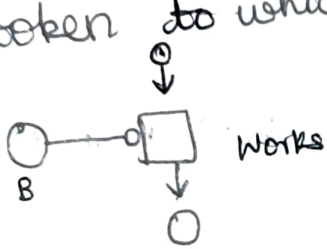
Firing T_1



(Token is removed from place but not from place)

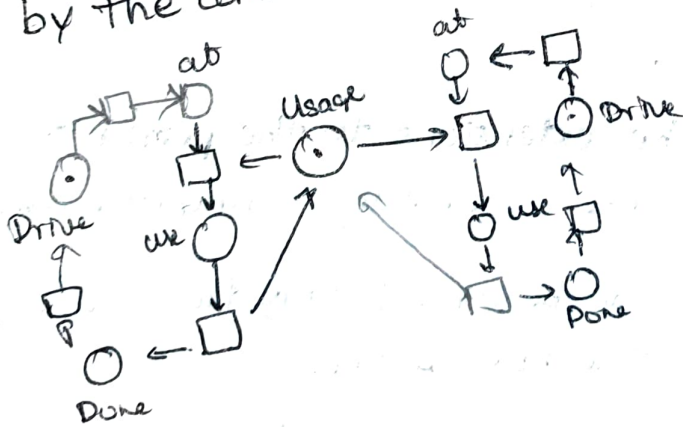
Inhibitor edge: (Opposite of normal head edges)

- * The logical operation is and (presence of token) in all the input places of a transition.
- + In inhibitor edges, the not logic is used i.e. presence of a token inhibits the firing of the token to which it is connected. ($\circ \rightarrow$)



eg: Mutual exclusion:

- * Two concurrent process cannot take place simultaneously. So we declare two petri nets connected by the conditional statement.

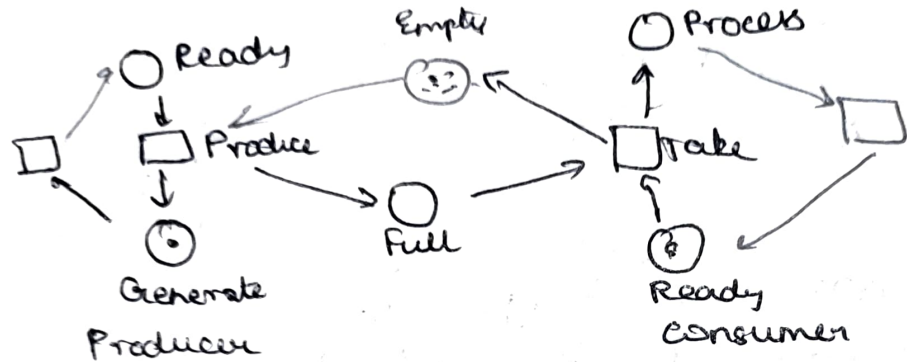


Resource Allocation:

- * It is possible that a single place may contain multiple tokens at one time.
- In these kinds of problems, place represents the number of available units of the resource.
- * If there is no token then you need to wait till you get to have a unit in the place.

* The number of units in a system can be variable

eg: Classical consumer-producer problem



When it is empty,
producer will work.
When it is full, consumer
will start to work.