# Algorithms for Big Data

Fall Semester 2019

Exercise Set 14

Here is a formalization of MPC model (one of many possible, equivalent):

- Input size N, distributed among machines.
- Machine memory is  $S = N^{\alpha}$  for some  $0 < \alpha < 1$ .
- Machines are numbered with unique ID's, 1 ..  $\frac{N}{S}$ .
- After each round machines send messages addressed to other machines. Each machine can send  $\mathcal{O}(S)$  atomic messages in total and receive  $\mathcal{O}(S)$  atomic messages in total.

## Exercise 1:

Maximum computation: input array x[1 .. N]. Output:  $\max\{x[i]\}$  (on a single machine), in time  $\mathcal{O}(\frac{1}{\alpha})$ .

# Exercise 2:

Broadcasting: as an input one machine has a message m of size  $\mathcal{O}(S)$ . Output: all machines have m. Show  $\mathcal{O}(\frac{1}{\alpha})$  algorithm.

#### Exercise 3:

Reason that broadcasting cannot be done faster, that there is no  $o(\frac{1}{\alpha})$  algorithm.

#### Exercise 4:

Prefix sums: input array x[1 ... N]. Output: array y[1 ... N] where y[i] = x[1] + ... + x[i]. Time:  $\mathcal{O}(\frac{1}{\alpha})$ .

#### Exercise 5:

Offsets: input array x[1 ... N] and S values  $a_1, ..., a_S$ . Output: values  $j_1, ..., j_S$  where  $j_k$  is the position of  $a_k$  in sorted x[1 ... N]. Time:  $\mathcal{O}(\frac{1}{\alpha})$ .

# Exercise 6:

Pivot: input array x[1 ... N] and S values  $a_1, ..., a_{S-1}$ . Output: reshuffle x so that some prefix of machines holds all the values from x smaller than  $a_1$ , then next batch of machines holds all values from x between  $a_1$  and  $a_2$ , etc. Time:  $\mathcal{O}(\frac{1}{\alpha})$ .

### Exercise 7:

Sorting: input array x[1 ... N]. Output: x sorted. Time:  $\mathcal{O}(\frac{1}{\alpha^2})$ . Idea:

- Pick sample of size S.
- Use it as a pivot.
- Show that whp subproblems are of size  $\widetilde{\mathcal{O}}(\frac{N}{\sqrt{S}})$ .
- Recurse on subproblems.