Informatics 1 Functional Programming Lecture 7

Function properties

Philip Wadler
University of Edinburgh

Compare and contrast

Sum, Product, Concat

```
sum :: [Int] -> Int
sum = foldr (+) 0

product :: [Int] -> Int
product = foldr (*) 1

concat :: [[a]] -> [a]
concat = foldr (++) []
```

Part I

Fold, right and left

Fold right

```
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr f u [] = u
foldr f u (x:xs) = x 'f' (foldr f v xs)
 foldr (++) [] ["abc", "def", "ghi"]
  foldr (++) [] ("abc" : ("def" : ("ghi" : [])))
 "abc" ++ foldr (++) [] ("def" : ("ghi" : []))
=
  "abc" ++ ("def" ++ foldr (++) [] ("ghi" : []))
=
 "abc" ++ ("def" ++ ("ghi" ++ foldr (++) [] []))
  "abc" ++ ("def" ++ ("ghi" ++ []))
```

Fold left

```
foldl :: (b -> a -> b) -> b -> [a] -> b
foldl f u [] = u
foldl f u (x:xs) = foldl f (u 'f' x) xs
 foldl (++) [] ["abc", "def", "qhi"]
 foldl (++) [] ("abc" : ("def" : ("ghi" : [])))
=
 foldl (++) ([] ++ "abc") ("def" : ("qhi" : []))
=
 foldl (++) (([] ++ "abc") ++ "def") ("ghi" : [])
=
 foldl (++) ((([] ++ "abc") ++ "def") ++ "qhi") []
=
  ((([] ++ "abc") ++ "def") ++ "qhi")
```

Fold right, non-empty list

```
foldr1 :: (a -> a -> a) -> [a] -> a
foldr1 f [x] = x
foldr1 f (x:xs) = x 'f' (foldr1 f xs)
 foldr1 ('max') [3, 1, 4, 2]
 foldr1 ('max') (3 : (1 : (4 : (2 : []))))
=
 3 'max' foldr1 ('max') (1 : (4 : (2 : [])))
=
 3 'max' (1 'max' foldr1 ('max') (4 : (2 : [])))
=
 3 'max' (1 'max' (4 'max' foldr1 ('max') (2 : [])))
=
 3 'max' (1 'max' (4 'max' 2))
```

Fold left, non-empty list

```
foldl1 :: (a \rightarrow a \rightarrow a) \rightarrow [a] \rightarrow a
foldl1 f (x:xs) = foldl f x xs
  fold11 ('max') [3, 1, 4, 2]
=
  fold11 ('max') (3 : (1 : (4 : (2 : []))))
=
  foldl ('max') 3 (1 : (4 : (2 : [])))
  foldl ('max') (3 'max' 1) (4 : (2 : []))
=
  foldl ('max') ((3 'max' 1) 'max' 4) (2 : [])
=
  foldl ('max') (((3 'max' 1) 'max' 4) 'max' 2)[]
  (((3 'max' 1) 'max' 4) 'max' 2)
```

Part II

Append

Append

```
(++) :: [a] -> [a] -> [a]
[] ++ ys = ys
(x:xs) ++ ys = x : (xs ++ ys)
 "abc" ++ "de"
  ('a' : ('b' : ('c' : []))) ++ ('d' : ('e' : []))
 'a' : (('b' : ('c' : [])) ++ ('d' : ('e' : [])))
=
  'a' : ('b' : (('c' : []) ++ ('d' : ('e' : []))))
=
  'a' : ('b' : ('c' : ([] ++ ('d' : ('e' : [])))))
  'a' : ('b' : ('c' : ('d' : ('e' : []))))
 "abcde"
```

Append

```
(++) :: [a] -> [a] -> [a]
[] ++ ys = ys
(x:xs) ++ ys = x : (xs ++ ys)
 "abc" ++ "de"
 'a' : ("bc" ++ "de")
 'a' : ('b' : ("c" ++ "de"))
=
  'a' : ('b' : ('c' : ("" ++ "de")))
=
 'a' : ('b' : ('c' : "de"))
 "abcde"
```

Properties of operators

- There are a few key properties about operators: *associativity*, *identity*, *commutativity*, *distributivity*, *zero*, *idempotence*. You should know and understand these properties.
- When you meet a new operator, the first question you should ask is "Is it associative?" The second is "Does it have an identity?"
- Associativity is our friend, because it means we don't need to worry about parentheses. The program is easier to read.
- Associativity is our friend, because it is key to writing programs that run twice as fast on dual-core machines, and a thousand times as fast on machines with a thousand cores.

Properties of append

```
prop_append_assoc :: [Int] -> [Int] -> Bool
prop_append_assoc xs ys zs =
    xs ++ (ys ++ zs) == (xs ++ ys) ++ zs

prop_append_ident :: [Int] -> Bool
prop_append_ident xs =
    xs ++ [] == xs && xs == [] ++ xs

prop_append_cons :: Int -> [Int] -> Bool
prop_append_cons x xs =
    [x] ++ xs == x : xs
```

Infix vs prefix notation

```
(++) :: [a] -> [a] -> [a]
[] ++ ys = ys
(x:xs) ++ ys = x : (xs ++ ys)
prop_append_assoc :: [Int] -> [Int] -> [Int] -> Bool
prop_append_assoc xs ys zs =
  xs ++ (ys ++ zs) == (xs ++ ys) ++ zs
                         VS
append :: [a] -> [a] -> [a]
append [] ys = ys
append (x:xs) ys = x : append xs ys
prop_append_assoc :: [Int] -> [Int] -> Bool
prop_append_assoc xs ys zs =
  append xs (append ys zs) == append (append xs ys) zs
```

Efficiency

```
(++) :: [a] -> [a] -> [a]
[] ++ ys = ys
(x:xs) ++ ys = x : (xs ++ ys)
 "abc" ++ "de"
 'a' : ("bc" ++ "de")
 'a' : ('b' : ("c" ++ "de"))
=
 'a' : ('b' : ('c' : ("" ++ "de")))
 'a' : ('b' : ('c' : "de"))
=
 "abcde"
```

Computing xs ++ ys takes about n steps, where n is the length of xs.

A useful fact

```
-- prop_sum.hs
import Test.QuickCheck

prop_sum :: Int -> Property
prop_sum n = n >= 0 ==> sum [1..n] == n * (n+1) 'div' 2

[melchior]dts: ghci prop_sum.hs
GHCi, version 6.8.3: http://www.haskell.org/ghc/ :? for help
*Main> quickCheck prop_sum
+++ OK, passed 100 tests.
*Main>
```

Associativity and Efficiency: Left vs. Right

Let n_1, n_2, n_3, n_4 be the lengths of xs_1, xs_2, xs_3, xs_4 .

Associated to the left

$$((xs_1 ++ xs_2) ++ xs_3) ++ xs_4$$

computing takes

$$n_1 + (n_1 + n_2) + (n_1 + n_2 + n_3)$$

steps. If we have m lists of length n, it takes about m^2n steps.

Associated to the right

$$xs_1 ++ (xs_2 ++ (xs_3 ++ xs_4))$$

computing takes

$$n_1 + n_2 + n_3$$

steps. If we have m lists of length n, it takes about mn steps. When m = 1000, the first is a thousand times slower than the second!

Associativity and Efficiency: Sequential vs. Parallel

Sequential:

$$(((((((x_1+x_2)+x_3)+x_4)+x_5)+x_6)+x_7)+x_8)$$

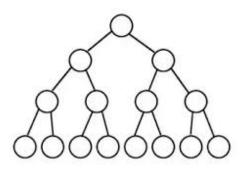
Summing 8 numbers takes 7 steps. If we have m numbers it takes m-1 steps.

Parallel:

$$((x_1+x_2)+(x_3+x_4))+((x_5+x_6)+(x_7+x_8))$$

Summing 8 numbers takes 3 steps.

Full Binary Tree



If we have m numbers it takes $\log_2 m$ steps. When m = 1000, the first is a hundred times slower than the second!

Map-Reduce

