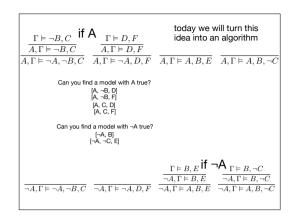


 $\frac{\Gamma \vDash \neg B, C}{A, \Gamma \vDash \neg B, C} \quad \text{if A} \quad \frac{\Gamma \vDash D, F}{A, \Gamma \vDash D, F} \quad \text{today we will turn this idea into an algorithm} \\ \frac{\Gamma \vDash \neg B, C}{A, \Gamma \vDash \neg A, \neg B, C} \quad \frac{\Gamma \vDash D, F}{A, \Gamma \vDash \neg A, D, F} \quad \frac{\Gamma \vDash A, B, E}{A, \Gamma \vDash A, B, E} \quad \frac{\Lambda}{A, \Gamma \vDash A, B, \neg C} \\ \text{Can you find a model with A true?} \\ \text{Can you find a model with ¬A true?} \\ \frac{\Gamma \vDash B, E}{\neg A, \Gamma \vDash A, B, E} \quad \frac{\text{if }}{\neg A, \Gamma \vDash B, \neg C} \quad \frac{\Lambda}{\neg A, \Gamma \vDash B, \neg C} \\ \frac{\Gamma \vDash B, E}{\neg A, \Gamma \vDash A, B, E} \quad \frac{\neg A, \Gamma \vDash B, \neg C}{\neg A, \Gamma \vDash A, B, E} \quad \frac{\neg A, \Gamma \vDash B, \neg C}{\neg A, \Gamma \vDash A, B, C} \\ \frac{\neg A, \Gamma \vDash A, B, C}{\neg A, \Gamma \vDash A, B, E} \quad \frac{\neg A, \Gamma \vDash A, B, C}{\neg A, \Gamma \vDash A, B, E} \quad \frac{\neg A, \Gamma \vDash A, B, C}{\neg A, \Gamma \vDash A, B, C}$

$$\frac{\Gamma \vDash \neg B, C}{A, \Gamma \vDash \neg B, C} \text{ if A } \underset{\Gamma \vDash D, F}{\prod_{A, \Gamma \vDash D, F}} \underset{A, \Gamma \vDash A, D, F}{\text{today we will turn this idea into an algorithm}} \\ \frac{\Gamma \vDash \neg B, C}{A, \Gamma \vDash \neg A, \neg B, C} \frac{A, \Gamma \vDash D, F}{A, \Gamma \vDash A, D, F} \xrightarrow{A, \Gamma \vDash A, B, E} \overline{A, \Gamma \vDash A, B, \neg C} \\ \text{Can you find a model with A true?} \\ \text{Find all valuations } \Gamma \text{ (consistent sets of literals) such that } \\ \Gamma \text{ satisfies the CNF:} \\ \Gamma \vDash \neg A, \neg B, C \quad \Gamma \vDash \neg A, D, F \quad \Gamma \vDash A, B, E \quad \Gamma \vDash A, B, \neg C \\ \text{and no smaller set of literals satisfies this CNF.} \\ \text{Can you find a model with } \neg A \text{ true?} \\ \frac{\Gamma \vDash B, E \text{ if } \neg A}{\neg A, \Gamma \vDash B, E} \xrightarrow{\neg A, \Gamma \vDash B, \neg C} \neg A, \Gamma \vDash B, \neg C} \\ \frac{\neg A, \Gamma \vDash B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg C}{\neg A, \Gamma \vDash A, B, \neg C} \\ \frac{\neg A, \Gamma \vDash A, B, \neg$$



cs	[Or xs Or xs <- cs, not (A 'elem' xs)]	[Or (delete (neg A) xs) I Or xs <- cs, not (A `elem` xs)	
¬A∨C∨D			
¬B∨F∨D			
¬B∨¬F∨¬C			
¬D∨¬B			
B∨¬C∨¬A	How can we simplify if we choose to make A True?		
B∨F∨C			
B∨¬F∨¬D			
A∨E			
A∨F			
¬FvCv¬E			
$A \lor \neg C \lor \neg E$			

cs	[Or xs Or xs <- cs, not (A `elem` xs)]	[Or (delete (neg A) xs) I Or xs <- cs, not (A `elem` xs)
¬A∨C∨D	¬A∨CvD	
¬B∨F∨D	¬B∨F∨D	
$\neg B \lor \neg F \lor \neg C$	¬B∨¬F∨¬C	
¬D∨¬B	∃B∀∃B	
$B \vee \neg C \vee \neg A$	Bv¬C∨¬A	
B∨F∨C	B∨F∨C	
B∨¬F∨¬D	B∨¬F∨¬D	
A∨E		
A∨F		
$\neg F \lor C \lor \neg E$	¬F∨C∨¬E	
A∨¬C∨¬E		

cs	[Or xs I Or xs <- cs, not (A 'elem' xs)]	[Or (delete (neg A) xs) I Or xs <- cs, not (A `elem` xs)]	
¬A∨C∨D	¬A∨C∨D	C∨D	
¬B∨F∨D	¬B∨F∨D	¬B∨F∨D	
$\neg B \lor \neg F \lor \neg C$	¬B∨¬F∨¬C	¬Bv¬Fv¬C	
¬D∨¬B	¬D∨¬B	¬D∨¬B	
B∨¬C∨¬A	B∨¬C∨¬A	B∨¬C	
B∨F∨C	B∨F∨C	B√F√C	
B∨¬F∨¬D	B∨¬F∨¬D	Bv¬Fv¬D	
A∨E			
A∨F			
$\neg F \lor C \lor \neg E$	¬F∨C∨¬E	¬FvCv¬E	
$A \lor \neg C \lor \neg E$			

```
CS
                                                                                                                  cs << A
                     import Data.List(elem, delete)
data Literal a = N a | P a deriving (Eq, Show)
 ¬A∨C∨D
                                                                                                                        C∨D
-BV-FVD neg :: Literal a = N a | P a derivir

-BV-FVD neg :: Literal a -> Literal a

neg (P a) = N a

-BV-FV-C neg (N a) = P a

-DV-B data Clause a = Or [Literal a]

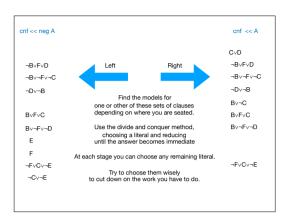
data Form a = And [Clause a]
                                                                                                                    \neg B \lor F \lor D
                                                                                                                 ¬B∨¬F∨¬C
                                                                                                                      ¬D∨¬B
B∨¬C∨¬A
                     (<<) :: Eq a => [Clause a] -> Literal a -> [Clause a] Bv \neg C -- reduces the clauses with x set true Bv \vdash V C
B√F√C
                    cs << x = [ Or (delete (neg x) xs)
| Or xs <- cs, not (x `elem` xs) ]
B∨¬F∨¬D
                                                                                                                  B∨¬F∨¬D
A∨E
A∨F
 ¬F∨C∨¬E
                                                                                                                  ¬F∨C∨¬E
A∨¬C∨¬E
```

```
cs << neg A
                                                       cs << A
 \neg A \lor C \lor D
                                                        C∨D
 ¬B∨F∨D
                                                         ¬B∨F∨D
 \neg B \lor \neg F \lor \neg C
                                                         ¬B∨¬F∨¬C
  \neg D \lor \neg B
                                                         \neg D \lor \neg B
 B∨¬C∨¬A
 B√F√C
                                                        B√F√C
 B∨¬F∨¬D
                                                         B∨¬F∨¬D
 A∨E
 A∨F
 \neg F \lor C \lor \neg E
                                                         \neg F \lor C \lor \neg E
 A \lor \neg C \lor \neg E
cs << x = [ Or (delete (neg x) xs)
                 | Or xs <- cs, not (x `elem` xs) ]
```

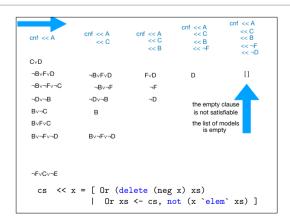
```
cs [ Or xs | xs <- cs,
not (neg A `elem` xs) ]
                         cs << neg A
                                       cs << A
¬A∨C∨D
                                      C√D
¬B∨F∨D
          ¬B∨F∨D
                                       ¬B∨F∨D
 ¬Bv¬Fv¬C ¬Bv¬Fv¬C
                                       ¬B∨¬F∨¬C
                                       ¬D∨¬B
¬D∨¬B
          ¬D∨¬B
B∨¬C∨¬A
                                       B∨¬C
B√F√C
          B∨F∨C
                                       B√F√C
B∨¬F∨¬D
         B∨¬F∨¬D
                                       B∨¬F∨¬D
A∨E
          A∨E
A∨F
          A∨F
¬FvCv¬E ¬FvCv¬E
                                       ¬F∨C∨¬E
Av¬Cv¬E Av¬Cv¬E
cs << x = [ Or (delete (neg x) xs)
           | Or xs <- cs, not (x `elem` xs) ]
```

```
[ Or xs | xs <- cs,
not (neg A 'elem' xs) ]
                                                    cs << A
                                  cs << neg A
¬A∨C∨D
                                                   C∨D
                                                    ¬B∨F∨D
 \neg B \lor F \lor D
             ¬B∨F∨D
                                  \neg B \lor F \lor D
 ¬Bv¬Fv¬C ¬Bv¬Fv¬C
                                  \neg B \lor \neg F \lor \neg C
                                                    ¬B∨¬F∨¬C
 \neg D \lor \neg B
              ¬D∨¬B
                                  \neg D \lor \neg B
                                                    \neg D \lor \neg B
B∨¬C∨¬A
                                                    B∨¬C
B<sub>V</sub>F<sub>V</sub>C
             B∨F∨C
                                 B√F√C
                                                   B∨F∨C
B∨¬F∨¬D
            B∨¬F∨¬D
                                 B∨¬F∨¬D
                                                   B∨¬F∨¬D
A∨E
              A∨E
                                   Ε
A∨F
              A∨F
¬FvCv¬E ¬FvCv¬E
                                 ¬F∨C∨¬E
                                                    ¬F∨C∨¬E
Av¬Cv¬E Av¬Cv¬E
                                  ¬C∨¬E
cs << x = [ Or (delete (neg x) xs) ]
                | Or xs <- cs, not (x `elem` xs) ]
```

```
cs << neg A cs << A
    CS
                                                                          C∨D
              if we choose a literal x then we can consider
              the models in which x is true and
                                                             \neg B \lor F \lor D \neg B \lor F \lor D
 \neg B \lor F \lor D
 \neg \mathsf{B} \lor \neg \mathsf{F} \lor \neg \mathsf{C} \quad \text{ the models in which neg x is true}
                                                              ¬Bv¬Fv¬C ¬Bv¬Fv¬C
                                                              ¬Dv¬B ¬Dv¬B
 \neg \mathsf{D} \vee \neg \mathsf{B}
                [ A:m|m<-models (cnf << A)]
B∨¬C∨¬A
                                                                           B∨¬C
B√F√C
                [ neg A : m | m <- models (cnf << neg A) ]
                                                            B∨F∨C
                                                                          B<sub>V</sub>F<sub>V</sub>C
                                                              Bv¬Fv¬D Bv¬Fv¬D
B∨¬F∨¬D
                                                              Е
A∨E
A∨F
                                                              ¬FvCv¬E ¬FvCv¬E
¬F∨C∨¬E
                                                              \neg C \lor \neg E
A \lor \neg C \lor \neg E
cs << x = [ Or (delete (neg x) xs)
                   | Or xs <- cs, not (x `elem` xs) ]
```



cnf << neg A	cnf << neg A << E	cnf << neg A << E << F	cnf << neg A << E << F << C	
¬B∨F∨D	¬B∨F∨D			
¬B∨¬F∨¬C	¬Bv¬Fv¬C	¬B∨¬C	¬В	
¬D∨¬B	¬D∨¬B	¬D∨¬B	¬D∨¬B	
B√F√C	B∨F∨C			
B∨¬F∨¬D	B∨¬F∨¬D	B∨¬D	B∨¬D	
E				
F	F			
¬F∨C∨¬E	¬F∨C	С	empty clause is not satisfiable	
¬C∨¬E	¬C	¬C	the list of models is empty	



```
*Main> :set -W

*Main> :load DPLL.hs
[1 of 1] Compiling Main ( DPLL.hs, interpreted )

Tutorial6.hs:26:1: warning: [-Wincomplete-patterns]
   Pattern match(es) are non-exhaustive
   In an equation for 'models':
        Patterns not matched:
        []
        (Or []:_)
|
26 | models (Or (x : xs): cs) =
```

```
models :: Eq a => Form a -> [[Literal a]]
-- returns the list of satisfying valuations
models [] = undefined -- null form
models (Or [] : _ )) = undefined -- null clause
models (Or (x : xs): cs )) =
    [ x : m | m <- models (cs << x) ]
++
    [ neg x : m | m <- models (Or xs: (cs << neg x)) ]
```

```
models :: Eq a => Form a -> [[Literal a]]
-- returns the list of satisfying valuations
models [] = undefined -- null form
models (Or [] : _ )) = undefined -- null clause
models (Or [x] : cs ) = undefined -- unit clause
models (Or (x : xs): cs )) =
    [ x : m | m <- models (cs << x) ]
++
    [ neg x : m | m <- models (Or xs: (cs << neg x)) ]</pre>
```