

# ER DATA MODEL CONCEPTS

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# Database

## Relational Database

A relational database is a type of database that stores and provides access to data points that are related to one another.



Figure 1: Relational Database



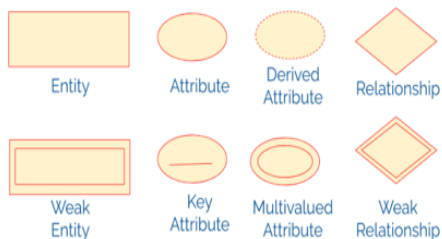




# ERD Notations

There are several Notations for ER Diagrams. Most widely used ones are:

## Chen's Notation:



## Crow's Foot Notation:

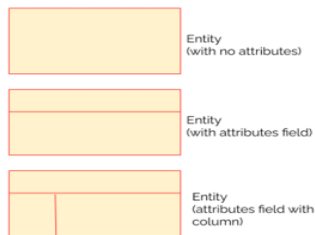


Figure 3: Chen's Notation and Crow's Foot Notation

An entity is an object in the real world with an independent existence that can be differentiated from other objects. An entity might be:

- An object with physical existence (e.g., a lecturer, a student, a car).
- An object with conceptual existence (e.g., a course, a job, a position).



## Student

| Roll_no | Student_name | Age | Mobile_no  |
|---------|--------------|-----|------------|
| 1       | Andrew       | 18  | 7089117222 |
| 2       | Angel        | 19  | 8709054568 |
| 3       | Priya        | 20  | 9864257315 |
| 4       | Analisa      | 21  | 9847852156 |

Figure 4: Entity Table

# Entity Type

The entity type is a collection of the entity having similar attributes.

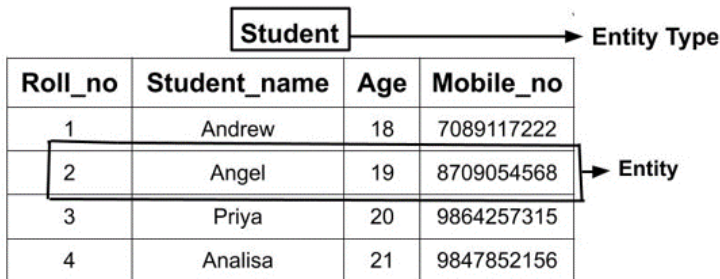


Figure 5: Entity Type

# Types of Entity Type

**Strong Entity Type:** Strong entity are those entity types which has a key attribute. The primary key helps in identifying each entity uniquely.

**Weak Entity Type:** Weak entity type doesn't have a key attribute. Weak entity type can't be identified on its own. It depends upon some other strong entity for its distinct identity.

## Example

There can be a room only if building exists.

# Entity Set

Entity Set is a collection of entities of the same entity type.

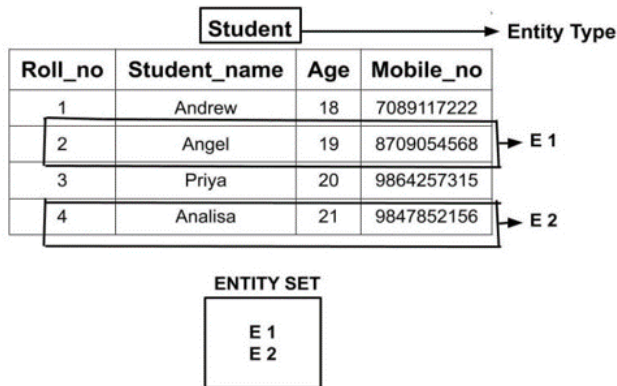


Figure 6: Entity Set

# Attributes

Each entity is described by a set of attributes.

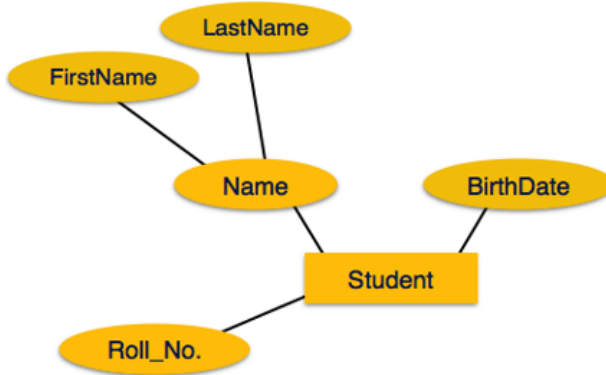
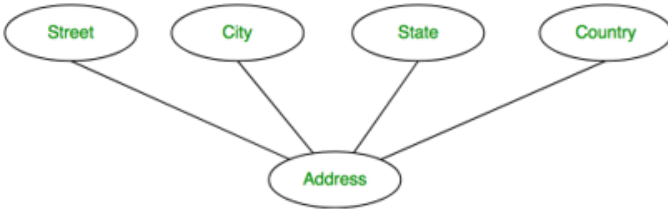


Figure 7: Attributes



# Composite Attributes

A composite attribute, **not to be confused with a composite key**, is an attribute that can be further subdivided to yield additional attributes. For example, the attribute ADDRESS can be subdivided into street, city, state, and zip code.



### Figure 9: Composite Attributes

# Multi-valued Attributes

Multivalued attributes are attributes that can have many values. For instance, a person may have several college degrees, and a household may have several different phones, each with its own number.

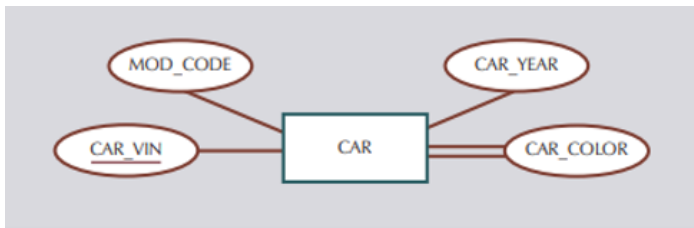


Figure 10: Multi-valued Attributes



# Derived Attributes

A derived attribute is an attribute whose value is calculated (derived) from other attributes.

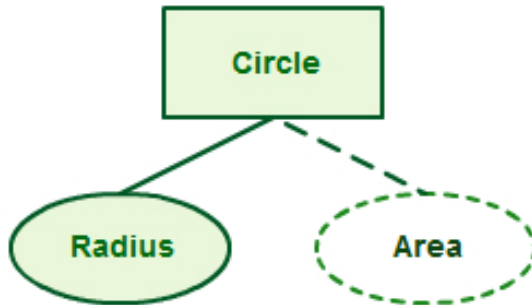


Figure 11: Derived Attributes

# Null

A null is a special symbol, independent of data type, which means either unknown or inapplicable. It does not mean zero or blank.

- Can represent:
  - An unknown attribute value
  - A known, but missing, attribute value
  - A “not applicable” condition
- Can create problems when functions such as COUNT, AVERAGE and SUM are used.
- Can create logical problems when relational tables are linked.

# Null

Student

| <u>Std_ID</u> | First Name | Last Name | Email  |
|---------------|------------|-----------|--|
| 1109          | John       | Price     | <a href="mailto:jp@gmail.com">jp@gmail.com</a> |
| 1401          | Bran       | NULL      | <a href="mailto:bw@gmail.com">bw@gmail.com</a> |
| 1756          |            | Natan     | <a href="mailto:cn@gmail.com">cn@gmail.com</a> |
| 1945          | Tom        | Black     | <a href="mailto:tb@gmail.com">tb@gmail.com</a> |
| 1987          | Nick       | Norris    | <a href="mailto:nn@gmail.com">nn@gmail.com</a> |



This is NOT NULL



This is NULL

Figure 12: Null

## Keys



The key is an attribute or a group of attributes whose values can be used to uniquely identify an individual entity in an entity set.

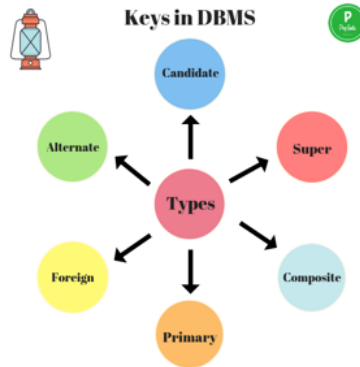


Figure 13: Keys

# Keys

- A **candidate key** is a simple or composite key that is unique and minimal. It is unique because no two rows in a table may have the same value at any time. It is minimal because every column is necessary in order to attain uniqueness.
- A **composite key** is composed of two or more attributes, but it must be minimal.
- The **primary key** is a candidate key that is selected by the database designer to be used as an identifying mechanism for the whole entity set. It must uniquely identify tuples in a table and not be null.

# Keys

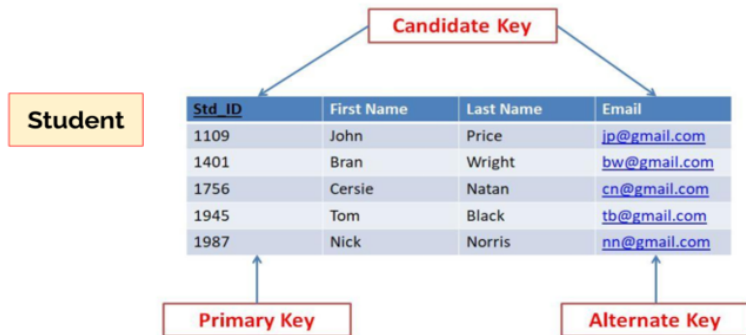


Figure 14: Candidate Key, Primary Key, Alternate Key

# Keys

- A **secondary key** is an attribute used strictly for retrieval purposes (can be composite), for example: Phone and Last Name.
- **Alternate keys** are all candidate keys not chosen as the primary key.
- A **foreign key** (FK) is an attribute in a table that references the primary key in another table OR it can be null.

# Keys

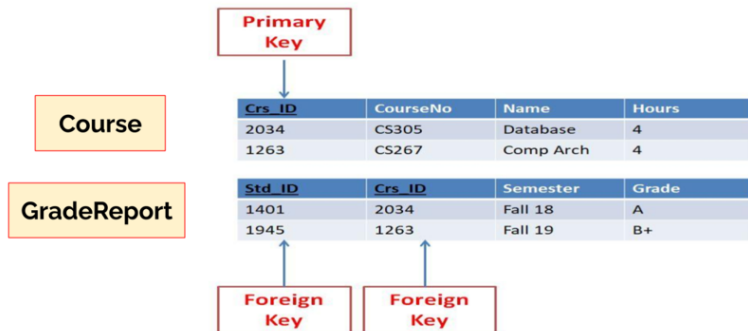


Figure 15: Primary Key, Foreign Key



# Keys

**STUDENT**

| STUD_NO | STUD_NAME | STUD_PHONE | STUD_STATE | STUD_COUNT<br>RY | STUD_AGE |
|---------|-----------|------------|------------|------------------|----------|
| 1       | RAM       | 9716271721 | Haryana    | India            | 20       |
| 2       | RAM       | 9898291281 | Punjab     | India            | 19       |
| 3       | SUJIT     | 7898291981 | Rajsthan   | India            | 18       |
| 4       | SURESH    |            | Punjab     | India            | 21       |

**Table 1**

**STUDENT\_COURSE**

| STUD_NO | COURSE_NO | COURSE_NAME       |
|---------|-----------|-------------------|
| 1       | C1        | DBMS              |
| 2       | C2        | Computer Networks |
| 1       | C2        | Computer Networks |

**Table 2**

**Figure 16: Keys**

# Relationships

Relationship is an association between entities. The entities that participate in a relationship are also known as participants, and each relationship is identified by a name that describes the relationship. The relationship name is an active or passive verb.

## Example

- a STUDENT takes a CLASS,
- a PROFESSOR teaches a CLASS,
- a DEPARTMENT employs a PROFESSOR,
- a DIVISION is managed by an EMPLOYEE and
- an AIRCRAFT is flown by a CREW.

# Weak Relationship

A weak, or non-identifying, relationship exists if the primary key of the related entity does not contain a primary key component of the parent entity.

## Course

| <u>Crs_ID</u> | CourseNo | Name      | Hours |
|---------------|----------|-----------|-------|
| 2034          | CS305    | Database  | 4     |
| 1263          | CS267    | Comp Arch | 4     |

## Class

| <u>Class_ID</u> | Crs_ID | Section | Time                      |
|-----------------|--------|---------|---------------------------|
| 10012           | 2034   | 1       | Friday, 8:00-10:00 a.m.   |
| 10017           | 1263   | 3       | Wednesday, 2:30-3:45 p.m. |

Figure 17: Weak Relationship

# Strong Relationship

A strong, or identifying, relationship exists when the primary key of the related entity contains the primary key component of the parent entity.

## Course

| <u>Crs_ID</u> | CourseNo | Name      | Hours |
|---------------|----------|-----------|-------|
| 2034          | CS305    | Database  | 4     |
| 1263          | CS267    | Comp Arch | 4     |

## Class

| <u>Crs_ID</u> | <u>Section</u> | Time                      |
|---------------|----------------|---------------------------|
| 10012         | 1              | Friday, 8:00-10:00 a.m.   |
| 10017         | 3              | Wednesday, 2:30-3:45 p.m. |

Figure 18: Strong Relationship

# One to One (1 : 1) Relationship

A one to one (1 : 1) relationship is the relationship of one entity to only one other entity, and vice versa. It should be rare in any relational database design.

## Example

In nature, each element has one formula and each formula represents one element.



Figure 19: One to One (1 : 1) Relationship

# One to Many ( $1 : M$ ) Relationship

A one to many ( $1 : M$ ) relationship should be the norm in any relational database design and is found in all relational database environments.

## Example

In business, each staff works in one department and in each department, many staff can be working.



**Figure 20:** One to Many ( $1 : M$ ) Relationship

# Many to Many ( $M : N$ ) Relationship

One entity from A can be associated with more than one entity from B and vice versa.

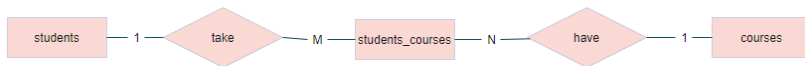
## Example

In university, each student can attend many courses and in each course, there can be many students.



**Figure 21:** Many to Many ( $M : N$ ) Relationship

# Breaking Up $M : N$ Relationship



**Figure 22:** Breaking Up  $M : N$  Relationship



# Unary Relationship (Recursive)

A unary relationship, also called recursive, is one in which a relationship exists between occurrences of the same entity set. In this relationship, the primary and foreign keys are the same, but they represent two entities with different roles.

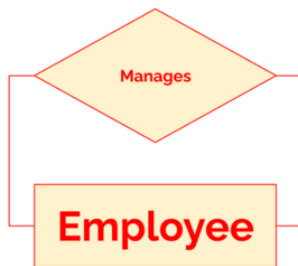


Figure 23: Unary Relationship (Recursive)

# Unary Relationship (Recursive)

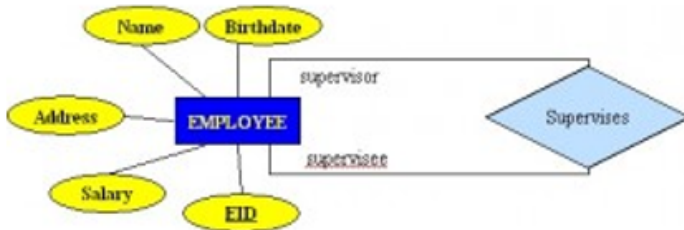


Figure 24: Unary Relationship (Recursive)

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# Integrity Rules and Constraints

## Domain Integrity

Domain restricts the values of attributes in the relation and is a constraint of the relational model. We need more specific ways to state what data values are or are not allowed and which format is suitable for an attribute. For example, the Employee ID (EID) must be unique or the employee Birthdate is in the range [Jan 1, 1950, Jan 1, 2000].

| ID   | NAME     | SEMENSTER       | AGE |
|------|----------|-----------------|-----|
| 1000 | Tom      | 1 <sup>st</sup> | 17  |
| 1001 | Johnson  | 2 <sup>nd</sup> | 24  |
| 1002 | Leonardo | 5 <sup>th</sup> | 21  |
| 1003 | Kate     | 3 <sup>rd</sup> | 19  |
| 1004 | Morgan   | 8 <sup>th</sup> | A   |

Not allowed. Because AGE is an integer attribute

Figure 26: Domain Integrity

# Entity Integrity

To ensure entity integrity, it is required that every table have a primary key. Neither the PK nor any part of it can contain null values. This is because null values for the primary key mean we cannot identify some rows.

## EMPLOYEE

| EMP_ID | EMP_NAME | SALARY |
|--------|----------|--------|
| 123    | Jack     | 30000  |
| 142    | Harry    | 60000  |
| 164    | John     | 20000  |
|        | Jackson  | 27000  |

Not allowed as primary key can't contain a NULL value

Figure 27: Entity Integrity

# Referential Integrity

Referential integrity requires that a foreign key must have a matching primary key or it must be null. This constraint is specified between two tables (parent and child); it maintains the correspondence between rows in these tables.

| Crs_ID | CourseNo | Name      | Hours |
|--------|----------|-----------|-------|
| 2034   | CS305    | Database  | 4     |
| 1263   | CS267    | Comp Arch | 4     |

| Std_ID | Crs_ID | Semester | Grade |
|--------|--------|----------|-------|
| 1401   | 2034   | Fall 18  | A     |
| 1945   | 1263   | Fall 19  | B+    |
| 2654   | 2015   | Fall 19  | B     |

This value is not allowed because this value is not defined as a primary key in the course table.

Figure 28: Referential Integrity



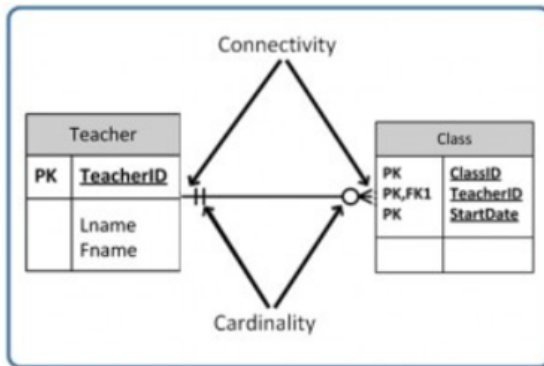








# Cardinality and Connectivity

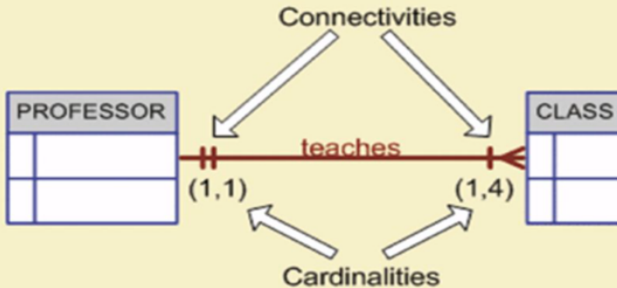


**Figure 29:** Cardinality and Connectivity

# Cardinality and Connectivity

**FIGURE 4.7**

**Connectivity and cardinality in an ERD**





# Types of Relationship

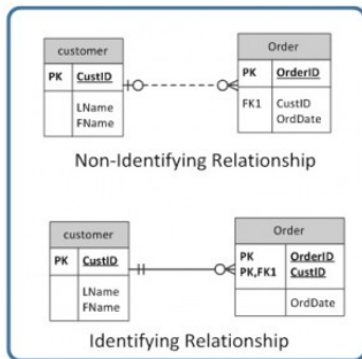
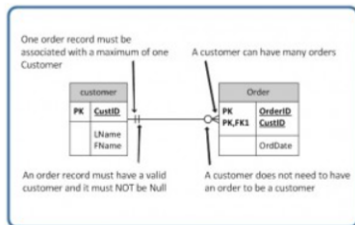


Figure 30: Identifying and Non-Identifying Relationship



# Types of Relationship

## Optional Relationship:



## Mandatory Relationship:

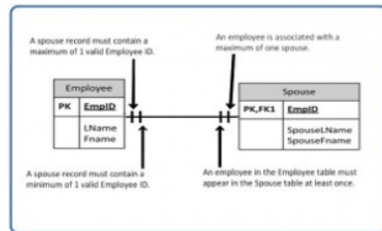
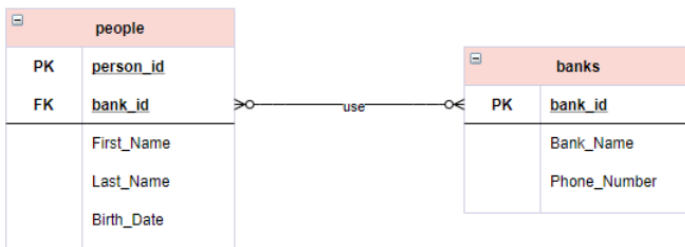


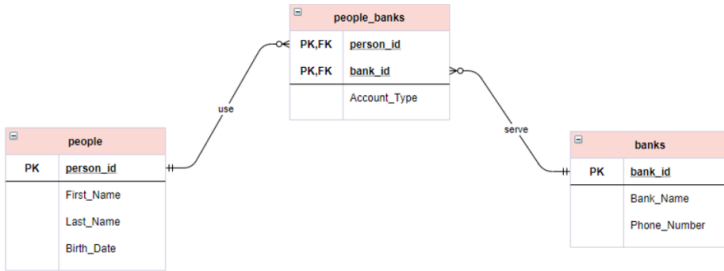
Figure 31: Optional and Mandatory Relationship



# Many to Many Examples



## Many to Many Examples



# Many to Many Examples

*people*

| <u>person_id</u> | <u>bank_id</u> | First_Name | Last_Name | Birth_Date |
|------------------|----------------|------------|-----------|------------|
| 1                | 1              | Jack       | London    | 20.05.1989 |
| 1                | 2              | Jack       | London    | 20.05.1989 |
| 2                | 3              | Harry      | Larry     | 12.07.1974 |
| 3                | 1              | Henry      | Jackson   | 13.02.1990 |
| 2                | 2              | Harry      | Larry     | 12.07.1974 |

*banks*

| <u>bank_id</u> | Bank_Name | Phone_Number |
|----------------|-----------|--------------|
| 1              | X bank    | 3330777      |
| 2              | Y bank    | 5551444      |
| 3              | Z bank    | 6662888      |

# Many to Many Examples

*people*

| person_id | First_Name | Last_Name | Birth_Date |
|-----------|------------|-----------|------------|
| 1         | Jack       | London    | 20.05.1989 |
| 2         | Paul       | Kane      | 30.04.1965 |
| 3         | Harry      | Larry     | 12.07.1974 |
| 4         | Henry      | Jackson   | 13.02.1990 |
| 5         | Mary       | Dune      | 16.01.1980 |

*people\_banks*

| person_id | bank_id | Account_Type                  |
|-----------|---------|-------------------------------|
| 1         | 1       | Checking Account              |
| 1         | 2       | Savings Account               |
| 2         | 1       | Individual Retirement Account |
| 2         | 3       | Checking Account              |
| 3         | 3       | Checking Account              |

*banks*

| bank_id | Bank_Name | Phone_Number |
|---------|-----------|--------------|
| 1       | X bank    | 3330777      |
| 2       | Y bank    | 5551444      |
| 3       | Z bank    | 6662888      |